

Selection Algorithms

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1 Selection Algorithms

1.0.1 QuickSelect

- Wikipedia Says:

1. In computer science, quickselect is a selection algorithm to find the k th smallest element in an unordered list. It is related to the quicksort sorting algorithm.

1.1 Let's understand QuickSort

1.1.1 Divide and Conquer Idea:

Divide:

The idea is to find a pivot. The pivot is an index q in the array $A[p:r]$, such that all elements of $A[p:q-1]$ are less than or equal to $A[q]$, which is less than or equal to elements of $A[q+1:r]$.

$A[p:q-1]$ and/or $A[q+1:r]$ are/is possibly empty.

Conquer:

Sort the array $A[p:r]$ by recursive calls to Divide

Combine

Merge the sorted arrays $A[p:q-1]$ and $A[q+1:r]$.

Because the sub arrays are sorted, the whole array is already sorted.

Combine is basically not required.

1.1.2 Let's write the Partition Function

Algorithm for partition:

1. input to the algorithm: Array, p = start index for the sub array, r = end index for the sub array
 1. Initialize $i = p - 1$: i is the index that tells that elements from p to i (inclusive) are smaller than $Array[r]$

2. Case b in the diagram beloww: Traverse the array: $j=0$ to $j < r$ if $\text{Array}[j]$ is smaller than or equal to $\text{Array}[r]$: then, move $\text{Array}[j]$ to $\text{Array}[i+1]$ and update i to $i+1$: (Verify in the image below that the loop invariant is satisfied) also increment j to make it ready for the next iteration iterate with new j
3. Loop has ended: j must be equal to r now: because the above loop's termination condition is that. Now $i+1$ is the first element that is greater than or equal to $\text{Array}[r]$ swap $(i+1)$ th position and r th position: At this stage, till i everything is smaller than $\text{Array}[r]$ return $i+1$: $i+1$ is the pivot.

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In [1]: def partition(array, p,r):
    print "*****Working within p:" + str(p) + " and r:" + str(r)
    # initialize data for entering the loop
    i = p - 1
    x = 0
    if r < len(array):
        x = array[r]
    j = p
    while j < r:
        if array[j] <= x:
            i = i + 1
            print "Swap intermediate i:" + str(i) + " and j:" + str(j)
            array[i], array[j] = array[j], array[i]
            j = j + 1
        else:
            j = j + 1
    print "Swap (i + 1):" + str(i+1) + " and j:" + str(j)
    if i < j:
        array[i+1], array[j] = array[j], array[i+1]
    return i+1

def quicksort(A, p , r):
    if p < r:
        q = partition(A, p, r)
        print A[p:q], A[q], A[q+1:r]
        quicksort(A, 0, q-1)
        quicksort(A, q+1, r)
    else:
        return
```

```
In [2]: A = [6, 0, 5, 13, 13, 2, 1, 2]
        quicksort(A, 0, len(A) - 1)
```

```
*****Working within p:0 and r:7
Swap intermediate i:0 and j:1
Swap intermediate i:1 and j:5
Swap intermediate i:2 and j:6
Swap (i + 1):3 and j:7
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[0, 2, 1] 2 [13, 6, 5]
*****Working within p:0 and r:2
Swap intermediate i:0 and j:0
Swap (i + 1):1 and j:2
[0] 1 []
*****Working within p:4 and r:7
Swap intermediate i:4 and j:4
Swap intermediate i:5 and j:5
Swap intermediate i:6 and j:6
Swap (i + 1):7 and j:7
[13, 6, 5] 13 []
*****Working within p:0 and r:6
Swap intermediate i:0 and j:0
Swap intermediate i:1 and j:1
Swap intermediate i:2 and j:2
Swap intermediate i:3 and j:3
Swap (i + 1):4 and j:6
[0, 1, 2, 2] 5 [6]
*****Working within p:0 and r:3
Swap intermediate i:0 and j:0
Swap intermediate i:1 and j:1
Swap intermediate i:2 and j:2
Swap (i + 1):3 and j:3
[0, 1, 2] 2 []
*****Working within p:0 and r:2
Swap intermediate i:0 and j:0
Swap intermediate i:1 and j:1
Swap (i + 1):2 and j:2
[0, 1] 2 []
*****Working within p:0 and r:1
Swap intermediate i:0 and j:0
Swap (i + 1):1 and j:1
[0] 1 []
*****Working within p:5 and r:6
Swap intermediate i:5 and j:5
Swap (i + 1):6 and j:6
[6] 13 []
*****Working within p:0 and r:5
Swap intermediate i:0 and j:0
Swap intermediate i:1 and j:1
Swap intermediate i:2 and j:2
Swap intermediate i:3 and j:3
Swap intermediate i:4 and j:4
Swap (i + 1):5 and j:5
[0, 1, 2, 2, 5] 6 []
*****Working within p:0 and r:4
Swap intermediate i:0 and j:0
Swap intermediate i:1 and j:1

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Swap intermediate i:2 and j:2
Swap intermediate i:3 and j:3
Swap (i + 1):4 and j:4
[0, 1, 2, 2] 5 []
*****Working within p:0 and r:3
Swap intermediate i:0 and j:0
Swap intermediate i:1 and j:1
Swap intermediate i:2 and j:2
Swap (i + 1):3 and j:3
[0, 1, 2] 2 []
*****Working within p:0 and r:2
Swap intermediate i:0 and j:0
Swap intermediate i:1 and j:1
Swap (i + 1):2 and j:2
[0, 1] 2 []
*****Working within p:0 and r:1
Swap intermediate i:0 and j:0
Swap (i + 1):1 and j:1
[0] 1 []

```

In [3]: A

Out[3]: [0, 1, 2, 2, 5, 6, 13, 13]

2 Questions from CLRS:

1. 7.1-1 Using Figure 7.1 as a model, illustrate the operation of PARTITION on the array A = [13, 19, 9, 5, 12, 8, 7, 4, 21, 2, 6, 11]

Answer:

When the function finishes the array looks like the following: [9, 5, 8, 7, 4, 2, 6, 11, 21, 13, 19, 12] where index of 11 is the pivot.

2. 7.1-2 What value of q does Partition function return when all elements in the array A[p..r] have the same value? Modify Partition function so that

$$q = \left\lfloor \frac{p+r}{2} \right\rfloor$$

when all elements in the array A[p..r] have the same value.

Answer: If all the elements in the array A[p..r] are same then Partition function tries to swap each element with itself for the whole array. Partition function would return the last index of the array.

Also look at [StackOverflow link](#)

```

In [1]: def partition2(array, p, r):
        i = p - 1

```

```

q = p
x = array[r]
all_equal = False
if array[q] == x:
    all_equal = True
while q < r:
    if array[q] == x:
        i = i + 1
        q = q + 1
    elif array[q] < x:
        all_equal = False
        i = i + 1
        array[q], array[i] = array[i], array[q]
        q = q + 1
    else:
        all_equal = False
        q = q + 1
# assert j == r
if all_equal == True:
    import math
    return math.floor((p+r)/2)
else:
    array[i+1], array[q] = array[q], array[i+1]
    return i+1

```

In [2]: partition2([2, 2, 2, 2, 2], 0, 4)

Out[2]: 2.0

3. 7.1-3 Give a brief argument that the running time of Partition on a subarray of size n is $\Theta(n)$.

Answer: To show that the running time of Partition on a subarray of size n is $\Theta(n)$, all we need to do is to show that its $\mathcal{O}(n)$ and $\Omega(n)$

So, in the worst case scenario, the Partition algorithm has to traverse the whole array once. This proves that lower bound of Partition algorithm is $\Omega(n)$ In the best case scenario as well Partition function at least needs to increment q for each element in the array. This proves that the algorithm has an upper bound of $\mathcal{O}(n)$

4. 7.1-4 How would you modify QUICKSORT to sort into nonincreasing order?

In [9]: *#Answer to 7.1-4:*

```

#In the partition function just change "<=" to ">="
def partition_increasing_order(array, p, r):
    i = p - 1
    x = array[r]
    q = p
    while q < r:
        if array[q] >= x:
            i = i + 1

```

```

        array[q], array[i] = array[i], array[q]
        q = q + 1
    else:
        q = q + 1
    #assert q == r
    array[i+1], array[q] = array[q], array[i+1]
    return i+1

def quicksort_increasing_order(A, p, r):
    if p < r:
        q = partition_increasing_order(A, p, r)
        quicksort_increasing_order(A, p, q-1)
        quicksort_increasing_order(A, q+1, r)

In [10]: A = [2, 3, 1, 6, 5, 9]
         quicksort_increasing_order(A, 0, 5)

In [11]: A

Out[11]: [9, 6, 5, 3, 2, 1]

```