

DESIGN & ANALYSIS OF ALGORITHM

(CS 617-01) (SP18)

Homework 1

Submitted by

Ashok Kumar Shrestha

Due Date: 30 Jan 2018

Problem 1: page 41, 2-4 (d)

2-4 Inversions

Let $A[1..n]$ be an array of n distinct numbers. If $i < j$ and $A[i] > A[j]$, then the pair (i, j) is called an ***inversion*** of A .

d. Give an algorithm that determines the number of inversions in any permutation on n elements in $\Theta(n \lg n)$ worst-case time. (Hint: Modify merge sort.)

Merge(A,B,C)

$n=0, i=0, j=0, k=0$

While $i < B.length$ && $j < C.length$

if($B[i] < C[j]$)

$A[k] = B[i]$

$i = i+1$

else

$A[k] = C[j]$

$j = j+1$

$n = n + B.length - i$

$k = k+1$

while $i < B.length$:

$A[k] = B[i]$

$i = i+1$

while $j < C.length$:

$A[k] = C[j]$

$j = j+1$

```
return n
```

```
Inversion(A)
```

```
    if (A.length == 1)
        return 0;
    else // A.length > 1
        B = A[1...n/2]
        C = A[n/2+1...n]
        num1 = Inversion(B)
        num2 = Inversion(C)
        num3 = merge(A,B,C)
        return num1+num2+num3
```

Problem 2: Rewrite Insertion Sort in recursive pseudo-code.

Call: Insertion_Sort(A,2,n)

//Pseudocode: Insertion Sort

```
Insertion_Sort(A, n)
```

```
    if (n == 1)
        return;
    else // i <= n
        Insertion_Sort(A, n-1)
        Key = A[ n ]
        j = n - 1
        while ( j > 0 & A[ j ] > key)
            A[ j+1 ] = A[ j ]
            j = j - 1
        A[ j+1 ] = key
        return;
```