

Ozyegin University
CS 321 Programming Languages
Sample Problems on Lambda Calculus

1. Apply **one** step of beta reduction on the following lambda terms, if possible. If there is no beta reduction possibility, write NORMAL FORM.

- $(\lambda x. x y z) a b c \Rightarrow \underline{\hspace{2cm} (a y z) b c \hspace{2cm}}$
- $(\lambda w. \lambda p. p w) \Rightarrow \underline{\hspace{2cm} \text{NORMAL FORM} \hspace{2cm}}$
- $(\lambda x. \lambda y. x (y x)) (\lambda w. \lambda v. w) (\lambda z. z) \Rightarrow \underline{\hspace{2cm} (\lambda y. (\lambda w. \lambda v. w) (y (\lambda w. \lambda v. w))) (\lambda z. z) \hspace{2cm}}$
- $(\lambda z. z (\lambda k. k) m) \Rightarrow \underline{\hspace{2cm} \text{NORMAL FORM} \hspace{2cm}}$
- $(\lambda u. (\lambda x. x u) (\lambda y. u)) \Rightarrow \underline{\hspace{2cm} (\lambda u. (\lambda y. u) u) \hspace{2cm}}$
- $(\lambda n. n n) (\lambda n. n n) \Rightarrow \underline{\hspace{2cm} (\lambda n. n n) (\lambda n. n n) \hspace{2cm}}$
- $(\lambda g. (\lambda h. h) g) \Rightarrow \underline{\hspace{2cm} (\lambda g. g) \hspace{2cm}}$

Solution: See <https://vimeo.com/250076161>

2. Reduce the following lambda term to its normal form. If a normal form does not exist, show that the term does not converge.

$$(\lambda x. \lambda y. x (y x)) (\lambda w. \lambda v. w) (\lambda z. z)$$

Solution: See <https://vimeo.com/250078355>

$$\begin{aligned} & (\lambda x. \lambda y. x (y x)) (\lambda w. \lambda v. w) (\lambda z. z) \\ \Rightarrow & (\lambda y. (\lambda w. \lambda v. w) (y (\lambda w. \lambda v. w))) (\lambda z. z) \\ \Rightarrow & (\lambda w. \lambda v. w) ((\lambda z. z) (\lambda w. \lambda v. w)) \\ \Rightarrow & (\lambda w. \lambda v. w) (\lambda w. \lambda v. w) \\ \Rightarrow & (\lambda v. (\lambda w. \lambda v. w)) \end{aligned}$$

3. Reduce the following lambda expression to its normal form.

$$(\lambda f. \lambda m. \lambda p. m (f m p)) (\lambda y. \lambda z. y z) (\lambda w. w)$$

Solution: See <https://vimeo.com/250079704>

$$\begin{aligned} & (\lambda f. \lambda m. \lambda p. m (f m p)) (\lambda y. \lambda z. y z) (\lambda w. w) \\ \Rightarrow & (\lambda m. \lambda p. m ((\lambda y. \lambda z. y z) m p)) (\lambda w. w) \\ \Rightarrow & \lambda p. (\lambda w. w) ((\lambda y. \lambda z. y z) (\lambda w. w) p) \\ \Rightarrow & \lambda p. ((\lambda y. \lambda z. y z) (\lambda w. w) p) \\ \Rightarrow & \lambda p. ((\lambda z. (\lambda w. w) z) p) \\ \Rightarrow & \lambda p. ((\lambda w. w) p) \\ \Rightarrow & \lambda p. p \end{aligned}$$

4. The following are the Church encodings for boolean values and conditional expression. Write down an encoding for logical *and*.

$$\mathbf{true} = \lambda a. \lambda b. a$$

$$\mathbf{false} = \lambda a. \lambda b. b$$

$$\mathbf{if} = \lambda c. \lambda t. \lambda e. c t e$$

$$\mathbf{and} = \underline{\lambda p. \lambda q. p q \mathbf{false}}$$

Solution: See <https://vimeo.com/250079905>

5. Using the encodings below, show that **mult 3 2** is **6**.

$$\begin{aligned}
 \mathbf{0} &= (\lambda f. \lambda x. x) \\
 \mathbf{1} &= (\lambda f. \lambda x. f x) \\
 \mathbf{2} &= (\lambda f. \lambda x. f(fx)) \\
 \mathbf{3} &= (\lambda f. \lambda x. f(f(fx))) \\
 \mathbf{mult} &= \lambda m. \lambda n. \lambda f. \lambda x. m(nf)x
 \end{aligned}$$

Solution:

$$\begin{aligned}
 \mathbf{mult\ 3\ 2} &= (\lambda m. \lambda n. \lambda f. \lambda x. m\ (nf)\ x)\ \mathbf{3\ 2} \\
 &\Rightarrow (\lambda n. \lambda f. \lambda x. \mathbf{3}\ (nf)\ x)\ \mathbf{2} \\
 &\Rightarrow \lambda f. \lambda x. \mathbf{3}\ (\mathbf{2}\ f)\ x \\
 &= \lambda f. \lambda x. \mathbf{3}\ ((\lambda f. \lambda x. f(fx))\ f)\ x && \text{expanded } \mathbf{2} \\
 &= \lambda f. \lambda x. \mathbf{3}\ (\underline{(\lambda f. \lambda x. f(fx))\ f})\ x && \beta\text{-reducing underlined term} \\
 &\Rightarrow \lambda f. \lambda x. \mathbf{3}\ (\lambda x. f(fx))\ x \\
 &= \lambda f. \lambda x. (\lambda f. \lambda x. f(f(fx)))\ (\lambda x. f(fx))\ x && \text{expanded } \mathbf{3} \\
 &= \lambda f. \lambda x. \underline{(\lambda f. \lambda x. f(f(fx)))\ (\lambda x. f(fx))}\ x && \beta\text{-reducing underlined term} \\
 &\Rightarrow \lambda f. \lambda x. (\lambda x. (\lambda x. f(fx))\ ((\lambda x. f(fx))\ ((\lambda x. f(fx))\ x)))\ x \\
 &= \lambda f. \lambda x. (\lambda x. (\lambda x. f(fx))\ ((\lambda x. f(fx))\ (\underline{(\lambda x. f(fx))\ x})))\ x && \beta\text{-reducing underlined term} \\
 &\Rightarrow \lambda f. \lambda x. (\lambda x. (\lambda x. f(fx))\ ((\lambda x. f(fx))\ (f(fx))))\ x \\
 &= \lambda f. \lambda x. (\lambda x. (\lambda x. f(fx))\ (\underline{(\lambda x. f(fx))\ (f(fx))}))\ x && \beta\text{-reducing underlined term} \\
 &\Rightarrow \lambda f. \lambda x. (\lambda x. (\lambda x. f(fx))\ (f(f(f(fx)))))\ x \\
 &= \lambda f. \lambda x. (\lambda x. (\lambda x. f(fx))\ (\underline{f(f(f(fx)))}))\ x && \beta\text{-reducing underlined term} \\
 &\Rightarrow \lambda f. \lambda x. (\lambda x. (f(f(f(f(fx)))))\ x \\
 &= \lambda f. \lambda x. \underline{(\lambda x. (f(f(f(f(fx)))))}\ x && \beta\text{-reducing underlined term} \\
 &\Rightarrow \lambda f. \lambda x. (f(f(f(f(f(fx))))) \\
 &= \mathbf{6}
 \end{aligned}$$