



BITS Pilani
Pilani Campus

CS F214 - Logic in CS

Prolog – Lecture 1

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Source



Learn Prolog Now

<http://www.learnprolognow.org/>

Prolog Interpreters

- Many interpreters available in market

- JIProlog
- SWI-Prolog
- Open Prolog
- Strawberry Prolog
- GNU Prolog
- etc.

We will be following

Today's Lecture



- Introduction to Prolog
 - Facts, Rules and Queries
 - Prolog Syntax
- Unification
- Proof search

Prolog



- *“Programming with Logic”*
- Very different from other programming languages
 - Declarative (not procedural)
 - Recursion (no “for” or “while” loops)
 - Relations (no functions)
 - Unification

Prolog and Web Applications

innovate

achieve

lead

Prolog programs are often smaller

- smallness encourages well written code
- hence, easier to maintain



Source:

<http://www.pathwayslms.com/swipltuts/>

Basic idea of Prolog

- Describe the situation of interest
- Ask a question
- Prolog:
 - logically deduces new facts about the situation we described
 - gives us its deductions back as answers

Consequences



- Think declaratively, not procedurally
 - Challenging
 - Requires a different mindset
- High-level language
 - Not as efficient as, say, C
 - Good for rapid prototyping
 - Useful in many AI applications (knowledge representation, inference)

Knowledge Base 1



woman(mia).
woman(jody).
woman(yolanda).
playsAirGuitar(jody).
party.

?- woman(mia).
yes
?- playsAirGuitar(jody).
yes
?- playsAirGuitar(mia).
no

Knowledge Base 1



woman(mia).
woman(jody).
woman(yolanda).
playsAirGuitar(jody).
party.

?- tattooed(jody).
ERROR: predicate tattooed/1 not defined.
?- party.
yes
?- rockConcert.
no

Knowledge Base 2



happy(yolanda). ← fact

listens2music(mia). ← fact

listens2music(yolanda):- happy(yolanda). ← rule

playsAirGuitar(mia):- listens2music(mia). ← rule

playsAirGuitar(yolanda):- listens2music(yolanda). ← rule

Knowledge Base 2



```
happy(yolanda).  
listens2music(mia).  
listens2music(yolanda):- happy(yolanda).  
playsAirGuitar(mia):- listens2music(mia).  
playsAirGuitar(yolanda):- listens2music(yolanda).
```

head

body

Knowledge Base 2



```
happy(yolanda).  
listens2music(mia).  
listens2music(yolanda):- happy(yolanda).  
playsAirGuitar(mia):- listens2music(mia).  
playsAirGuitar(yolanda):- listens2music(yolanda).
```

```
?- playsAirGuitar(mia).  
yes  
?- playsAirGuitar(yolanda).  
yes  
?-
```

Clauses



```
happy(yolanda).  
listens2music(mia).  
listens2music(yolanda):- happy(yolanda).  
playsAirGuitar(mia):- listens2music(mia).  
playsAirGuitar(yolanda):- listens2music(yolanda).
```

*There are five clauses in this knowledge base:
two facts, and three rules.*

The end of a clause is marked with a full stop.

Predicates



```
happy(yolanda).  
listens2music(mia).  
listens2music(yolanda):- happy(yolanda).  
playsAirGuitar(mia):- listens2music(mia).  
playsAirGuitar(yolanda):- listens2music(yolanda).
```

*There are three predicates in
this knowledge base:*

happy, listens2music, and playsAirGuitar

Knowledge Base 3



happy(vincent).

listens2music(butch).

playsAirGuitar(vincent):- listens2music(vincent), happy(vincent).

playsAirGuitar(butch):- happy(butch).

playsAirGuitar(butch):- listens2music(butch).

Expressing Conjunction



```
happy(vincent).  
listens2music(butch).  
playsAirGuitar(vincent):- listens2music(vincent), happy(vincent).  
playsAirGuitar(butch):- happy(butch).  
playsAirGuitar(butch):- listens2music(butch).
```

The comma “,” expresses conjunction in Prolog

Knowledge Base 3



happy(vincent).

listens2music(butch).

playsAirGuitar(vincent):- listens2music(vincent), happy(vincent).

playsAirGuitar(butch):- happy(butch).

playsAirGuitar(butch):- listens2music(butch).

?- playsAirGuitar(vincent).

no

?- playsAirGuitar(butch).

yes

?-

Expressing Disjunction

```
happy(vincent).  
listens2music(butch).  
playsAirGuitar(vincent):- listens2music(vincent), happy(vincent).  
playsAirGuitar(butch):- happy(butch).  
playsAirGuitar(butch):- listens2music(butch).
```

```
happy(vincent).  
listens2music(butch).  
playsAirGuitar(vincent):- listens2music(vincent), happy(vincent).  
playsAirGuitar(butch):- happy(butch); listens2music(butch).
```

Prolog & Logic



- Clearly, Prolog has something to do with logic...

	Prolog	Logic
Implication	$A:-B$	$B \rightarrow A$
Conjunction	A,B	$A \wedge B$
Disjunction	$A;B$	$A \vee B$

- Use of inference (modus ponens)
- Negation (?)

Knowledge Base 4



```
woman(mia).  
woman(jody).  
woman(yolanda).  
loves(vincent, mia).  
loves(marsellus, mia).  
loves(pumpkin, honey_bunny).  
loves(honey_bunny, pumpkin).
```

Prolog Variables and Asking Alternatives



```
woman(mia).  
woman(jody).  
woman(yolanda).  
loves(vincent, mia).  
loves(marsellus, mia).  
loves(pumpkin, honey_bunny).  
loves(honey_bunny, pumpkin).
```

```
?- woman(X).  
X=mia;  
X=jody;  
X=yolanda;  
no
```

Knowledge Base 4



```
woman(mia).  
woman(jody).  
woman(yolanda).  
loves(vincent, mia).  
loves(marsellus, mia).  
loves(pumpkin, honey_bunny).  
loves(honey_bunny, pumpkin).
```

```
?- loves(marsellus,X), woman(X).
```

```
X=mia
```

```
yes
```

```
?-
```

Knowledge Base 4



```
woman(mia).  
woman(jody).  
woman(yolanda).  
loves(vincent, mia).  
loves(marsellus, mia).  
loves(pumpkin, honey_bunny).  
loves(honey_bunny, pumpkin).
```

```
?- loves(pumpkin,X), woman(X).
```

```
no
```

```
?-
```


Knowledge Base 5



```
loves(vincent,mia).  
loves(marsellus,mia).  
loves(pumpkin, honey_bunny).  
loves(honey_bunny, pumpkin).  
jealous(X,Y):- loves(X,Z), loves(Y,Z).
```

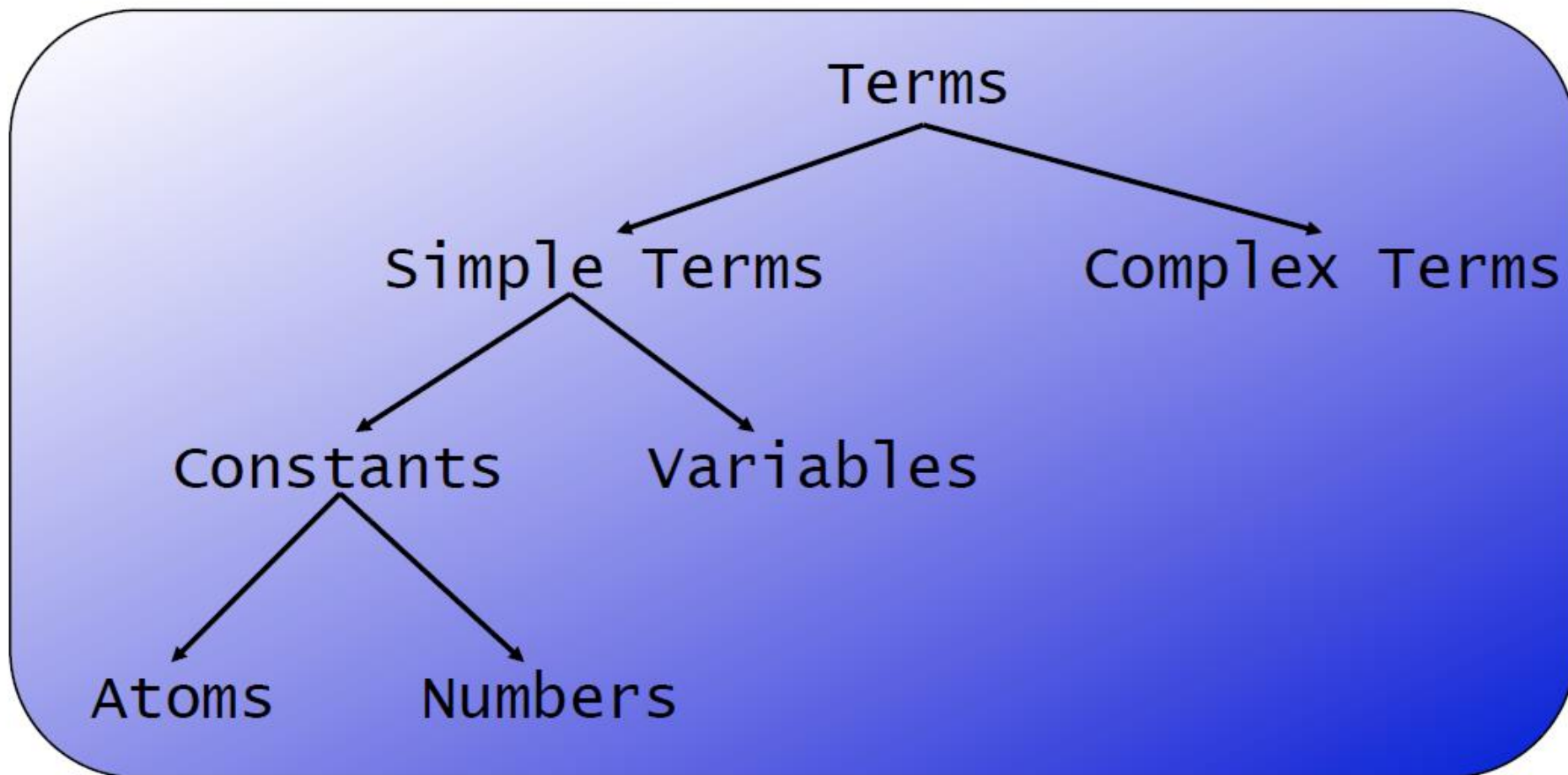
```
?- jealous(marsellus,W).  
W = Vincent  
?-
```

Syntax of Prolog



- Q: What exactly are facts, rules and queries built out of?
- A: Prolog terms

Prolog terms



- A sequence of characters of upper-case letters, lower-case letters, digits, or underscore, starting with a lowercase letter
 - ✓ *Examples:* **butch**, **big_kahuna_burger**, **playGuitar**
- An arbitrary sequence of characters enclosed in single quotes
 - ✓ *Examples:* **'Vincent'**, **'Five dollar shake'**, **'@\$%'**
- A sequence of special characters
 - ✓ *Examples:* **:**, **;**, **:-**

Numbers



- Integers:
12, -34, 22342
- Floats:
34573.3234, 0.3435

Variables



- A sequence of characters of uppercase letters, lower-case letters, digits, or underscore, starting with either an uppercase letter or an underscore
- Examples:
X, Y, Variable, Vincent, _tag

Complex Terms



- Atoms, numbers and variables are building blocks for **complex terms**
- Complex terms are built out of a **functor** directly followed by a sequence of **arguments**
 - ✓ Arguments are put in round brackets, separated by commas
 - ✓ The functor must be an atom

Examples of complex terms



- Examples we have seen before:
 - ✓ `playsAirGuitar(jody)`
 - ✓ `loves(vincent, mia)`
 - ✓ `jealous(marsellus, W)`
- Complex terms inside complex terms:
 - ✓ `hide(X, father(father(father(butch))))`

Arity



- The number of arguments a complex term has is called its arity
- Examples:
 - woman(mia)** is a term with arity 1
 - loves(vincent,mia)** has arity 2
 - father(father(butch))** arity 1

Arity is important



- You can define two predicates with the same functor but with different arity
- Prolog would treat this as two different predicates!
- In Prolog documentation, arity of a predicate is usually indicated with the suffix "/" followed by a number to indicate the arity



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Unification in Prolog

Unification



- Recall the previous example, where we said that Prolog unifies

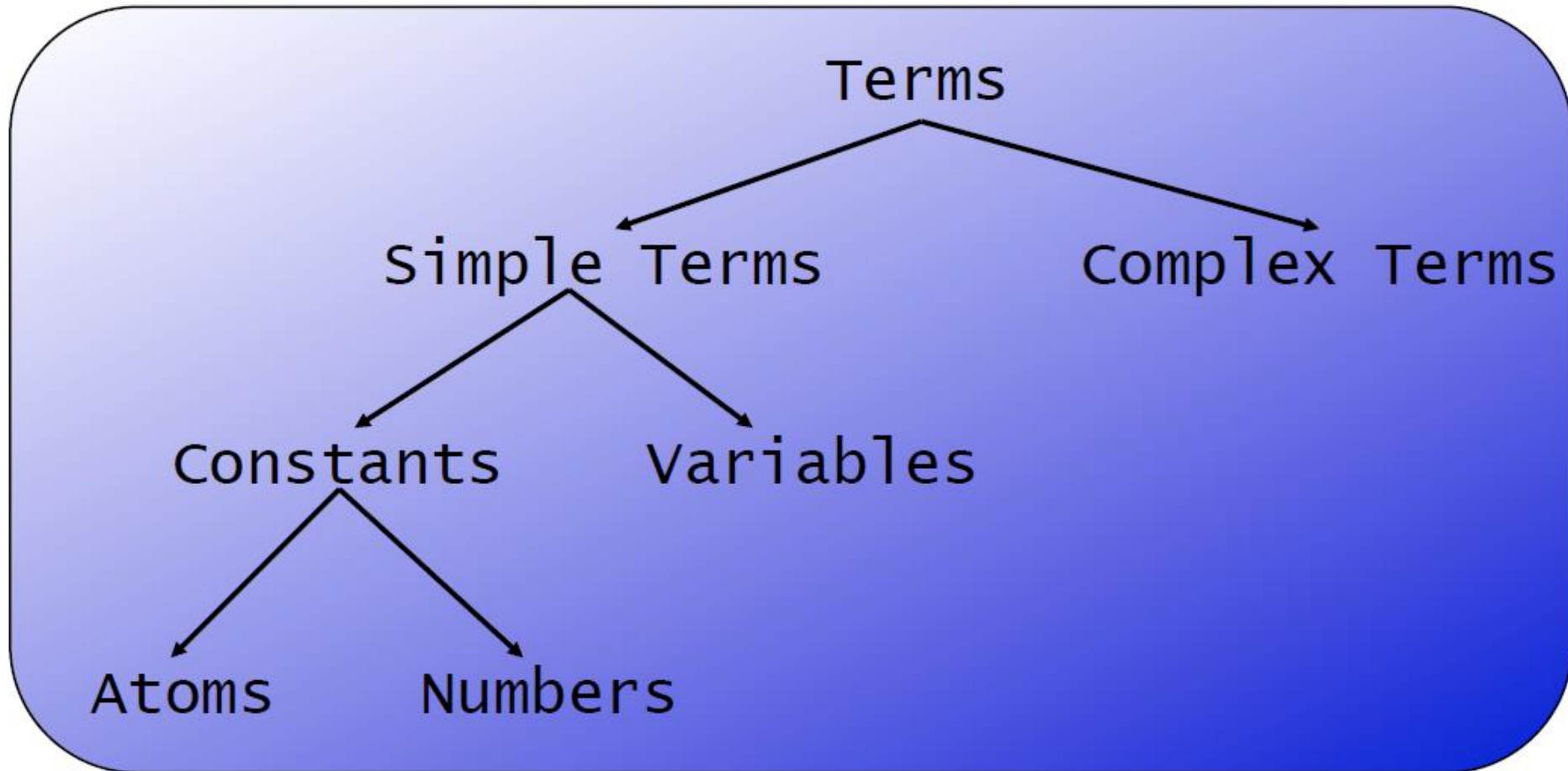
woman(X)

with

woman(mia)

thereby instantiating the variable **X** with the atom **mia**.

Recall Prolog Terms



Unification



- Working definition – two terms unify
 - if they are the same term, or
 - if they contain variables that can be uniformly instantiated with terms in such a way that the resulting terms are equal

Unification



- This means that:
 - **mia** and **mia** unify
 - **42** and **42** unify
 - **woman(mia)** and **woman(mia)** unify
- This also means that:
 - **vincent** and **mia** do not unify
 - **woman(mia)** and **woman(jody)** do not unify

Unification



- What about the terms:
 - **mia** and **X**
 - **woman(Z)** and **woman(mia)**
 - **loves(mia,X)** and **loves(X,vincent)**

Instantiations



- When Prolog unifies two terms, it performs all the necessary instantiations, so that the terms are equal afterwards
- This makes unification a very powerful programming mechanism

Revised Definition

1. If T_1 and T_2 are constants, then T_1 and T_2 unify if they are the same atom, or the same number
2. If T_1 is a variable and T_2 is any type of term, then T_1 and T_2 unify, and T_1 is instantiated to T_2 (and vice versa)
3. If T_1 and T_2 are complex terms then they unify if:
 1. They have the same functor and arity, and
 2. all their corresponding arguments unify, and
 3. the variable instantiations are compatible.

Prolog Unification: =/2



?- mia = mia.

yes

?- mia = vincent.

no

?- mia = X.

X=mia

yes

?-

How will Prolog respond?



?- X=mia, X=vincent.

no

?-

Why? After working through the first goal, Prolog has instantiated X with **mia**, so that it cannot unify it with **Vincent** anymore. Hence the second goal fails.

Example with Complex Terms



?- $k(s(g), Y) = k(X, t(f))$.

$X = s(g)$

$Y = t(f)$

yes

?-

Example with Complex Terms



?- $k(s(g), t(f)) = k(X, t(Y))$.

$X = s(g)$

$Y = f$

yes

?-

One last example



?- loves(X,X) = loves(marsellus,mia).

no

?-

Programming with Unification



vertical (line(point(X,Y), point(X,Z))).
horizontal (line(point(X,Y), point(Z,Y))).

?- vertical(line(point(1,1),point(1,3))).

yes

?- vertical(line(point(1,1),point(3,2))).

no

Programming with Unification



```
vertical ( line(point(X,Y), point(X,Z))).  
horizontal ( line(point(X,Y), point(Z,Y))).
```

```
?- horizontal(line(point(1,1),point(1,Y))).  
Y = 1;  
no  
?- horizontal(line(point(2,3),Point)).  
Point = point(_554,3);  
no
```

Exercise



Which of the following pairs of terms unify? Where relevant, give the variable instantiations that lead to successful unification.

1. $\text{bread} = \text{bread}$
2. $\text{'Bread'} = \text{bread}$
3. $\text{'bread'} = \text{bread}$
4. $\text{Bread} = \text{bread}$
5. $\text{bread} = \text{sausage}$
6. $\text{food}(\text{bread}) = \text{bread}$
7. $\text{food}(\text{bread}) = X$
8. $\text{food}(X) = \text{food}(\text{bread})$
9. $\text{food}(\text{bread}, X) = \text{food}(Y, \text{sausage})$
10. $\text{food}(\text{bread}, X, \text{beer}) = \text{food}(Y, \text{sausage}, X)$
11. $\text{food}(\text{bread}, X, \text{beer}) = \text{food}(Y, \text{kahuna_burger})$
12. $\text{food}(X) = X$
13. $\text{meal}(\text{food}(\text{bread}), \text{drink}(\text{beer})) = \text{meal}(X, Y)$
14. $\text{meal}(\text{food}(\text{bread}), X) = \text{meal}(X, \text{drink}(\text{beer}))$

Exercise



We are working with the following knowledge base:

```
house_elf(dobby).  
witch(hermione).  
witch('McGonagall').  
witch(rita_skeeter).  
magic(X):- house_elf(X).  
magic(X):- wizard(X).  
magic(X):- witch(X).
```

Which of the following queries are satisfied? Where relevant, give all the variable instantiations that lead to success.

1. ?- magic(house_elf).
2. ?- wizard(harry).
3. ?- magic(wizard).
4. ?- magic('McGonagall').
5. ?- magic(Hermione).



Proof Search

Proof Search



- Now that we know about unification, we are in a position to learn how Prolog searches a knowledge base to see if a query is satisfied.
- In other words: we are ready to learn about proof search and search trees

Example



```
f(a).  
f(b).  
g(a).  
g(b).  
h(b).  
k(X):- f(X), g(X), h(X).
```

```
?- k(Y).
```

Example: Search Tree

```
f(a).  
f(b).  
g(a).  
g(b).  
h(b).  
k(X):- f(X), g(X), h(X).
```

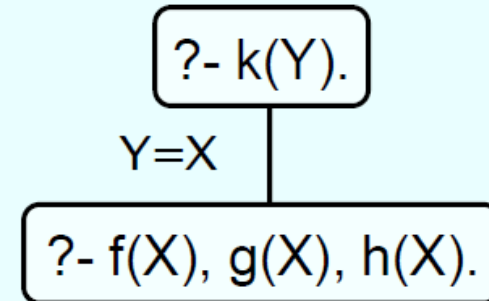
```
?- k(Y).
```

```
?- k(Y).
```

Example: Search Tree

f(a).
f(b).
g(a).
g(b).
h(b).
k(X):- f(X), g(X), h(X).

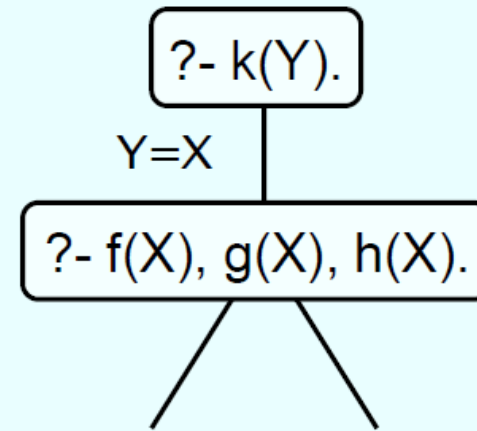
?- k(Y).



Example: Search Tree

f(a).
f(b).
g(a).
g(b).
h(b).
k(X):- f(X), g(X), h(X).

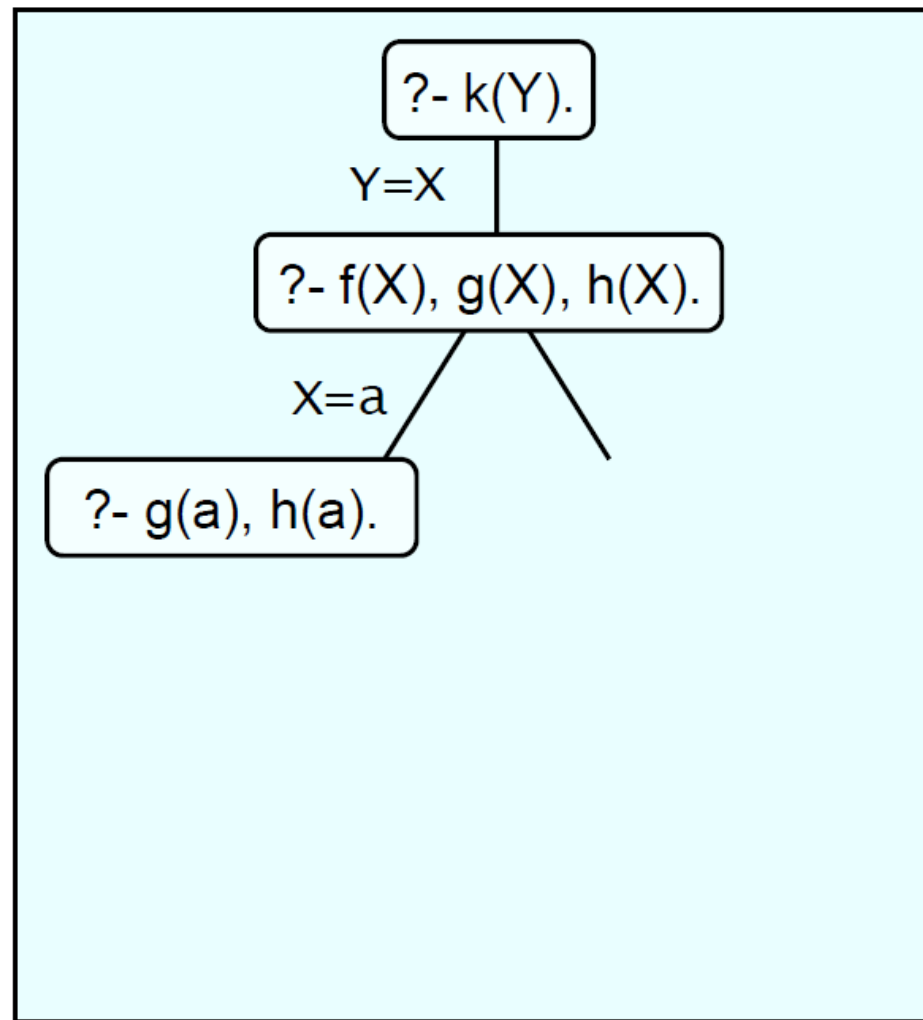
?- k(Y).



Example: Search Tree

$f(a).$
 $f(b).$
 $g(a).$
 $g(b).$
 $h(b).$
 $k(X):- f(X), g(X), h(X).$

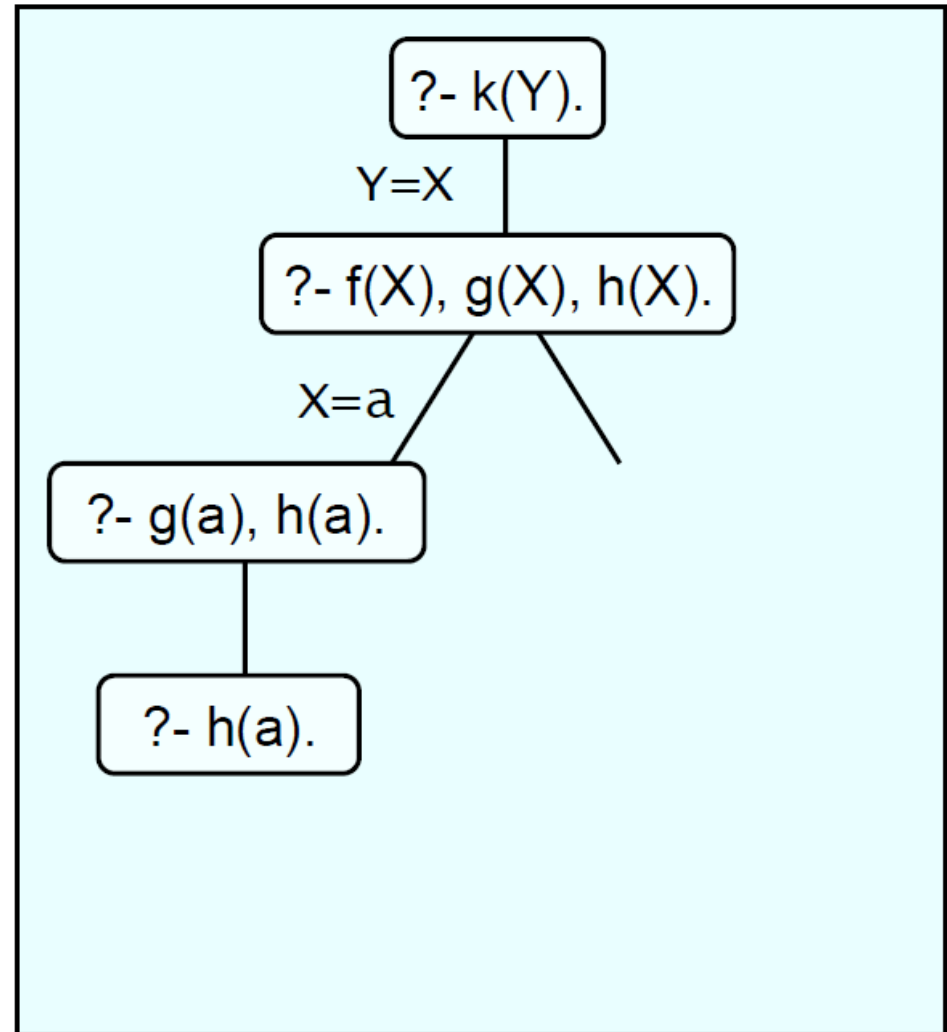
$?- k(Y).$



Example: Search Tree

f(a).
f(b).
g(a).
g(b).
h(b).
k(X):- f(X), g(X), h(X).

?- k(Y).

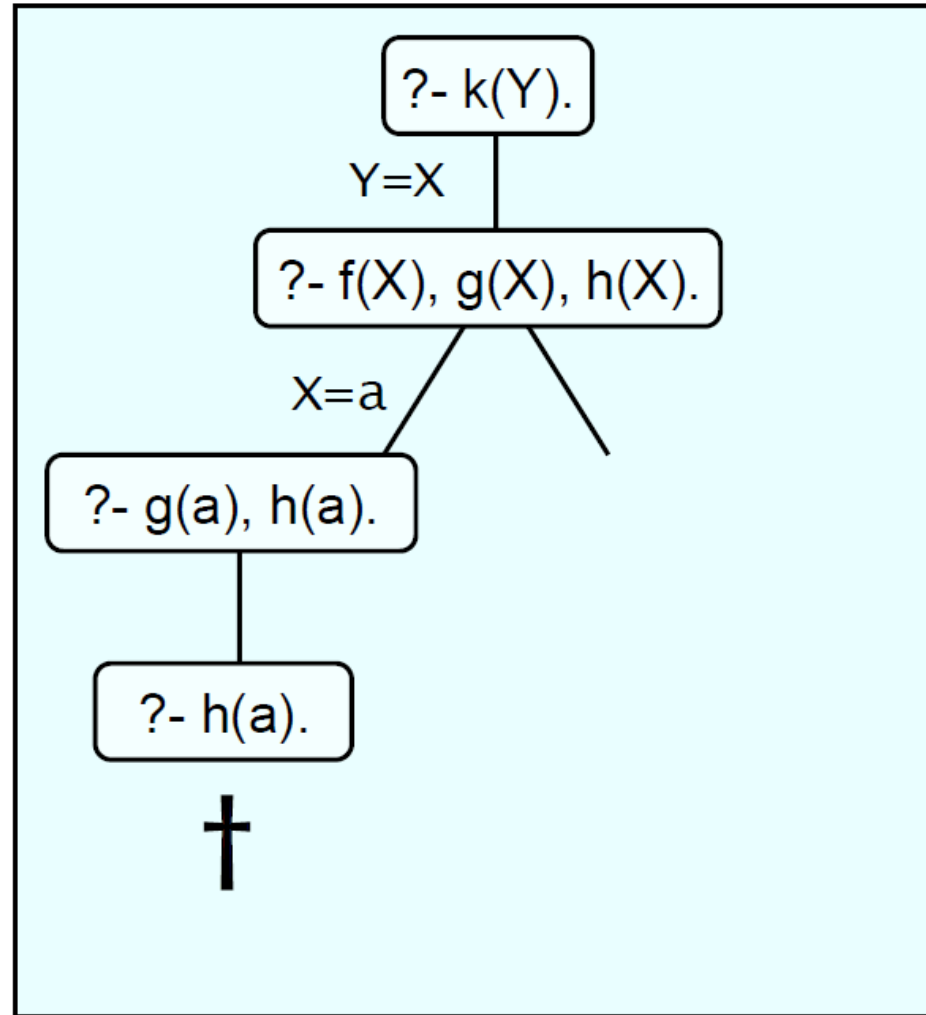


Example: Search Tree



f(a).
f(b).
g(a).
g(b).
h(b).
k(X):- f(X), g(X), h(X).

?- k(Y).

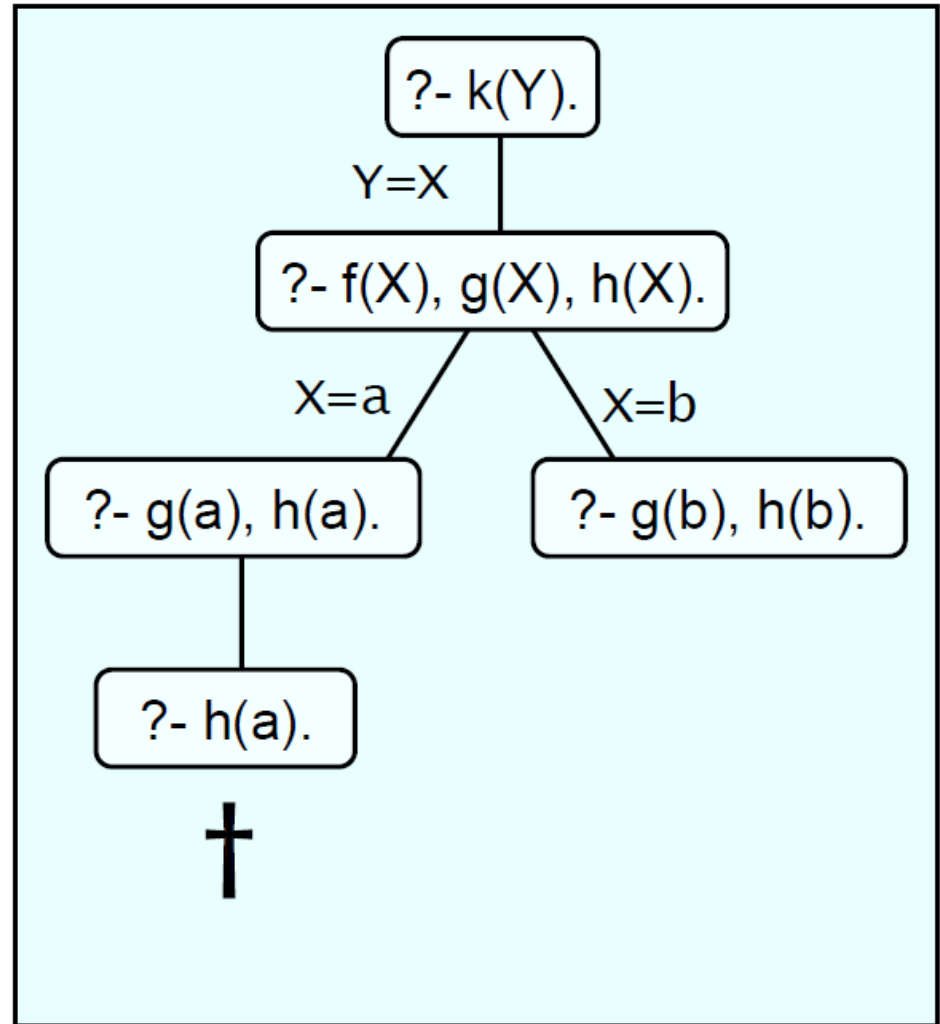


Example: Search Tree



f(a).
f(b).
g(a).
g(b).
h(b).
k(X):- f(X), g(X), h(X).

?- k(Y).

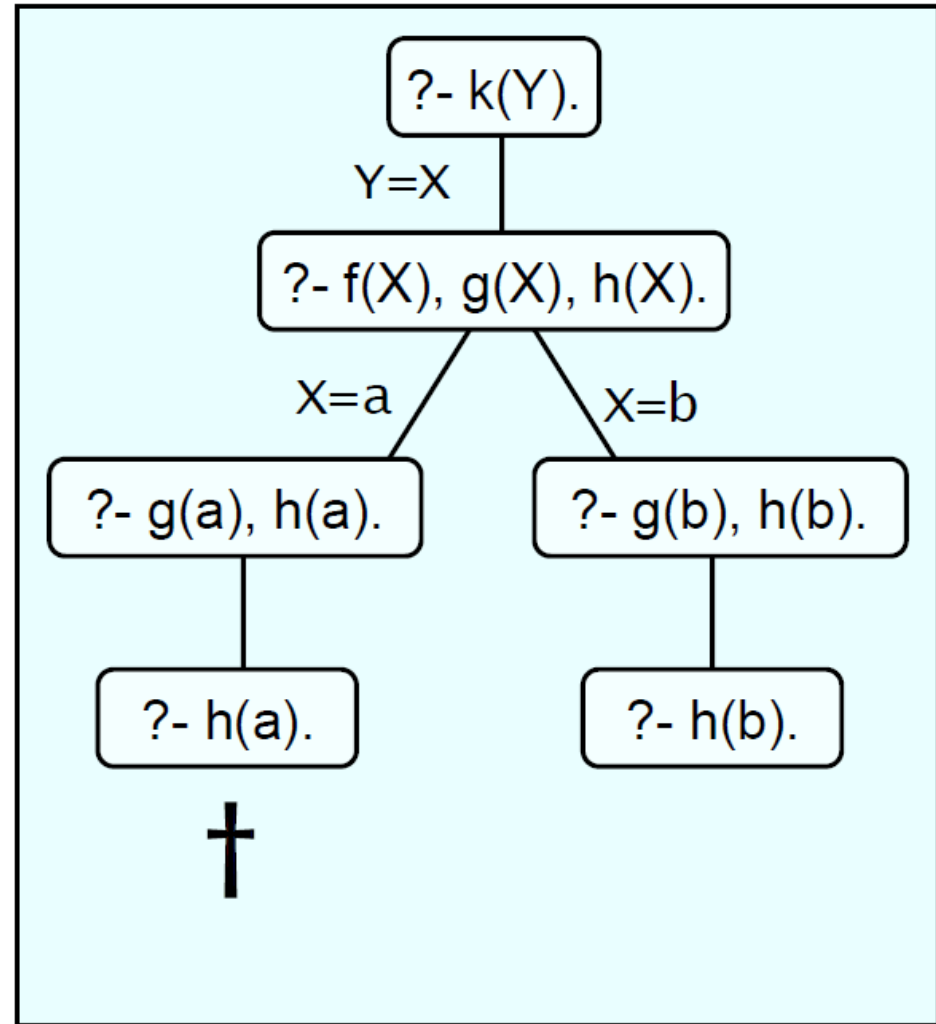


Example: Search Tree



f(a).
f(b).
g(a).
g(b).
h(b).
k(X):- f(X), g(X), h(X).

?- k(Y).

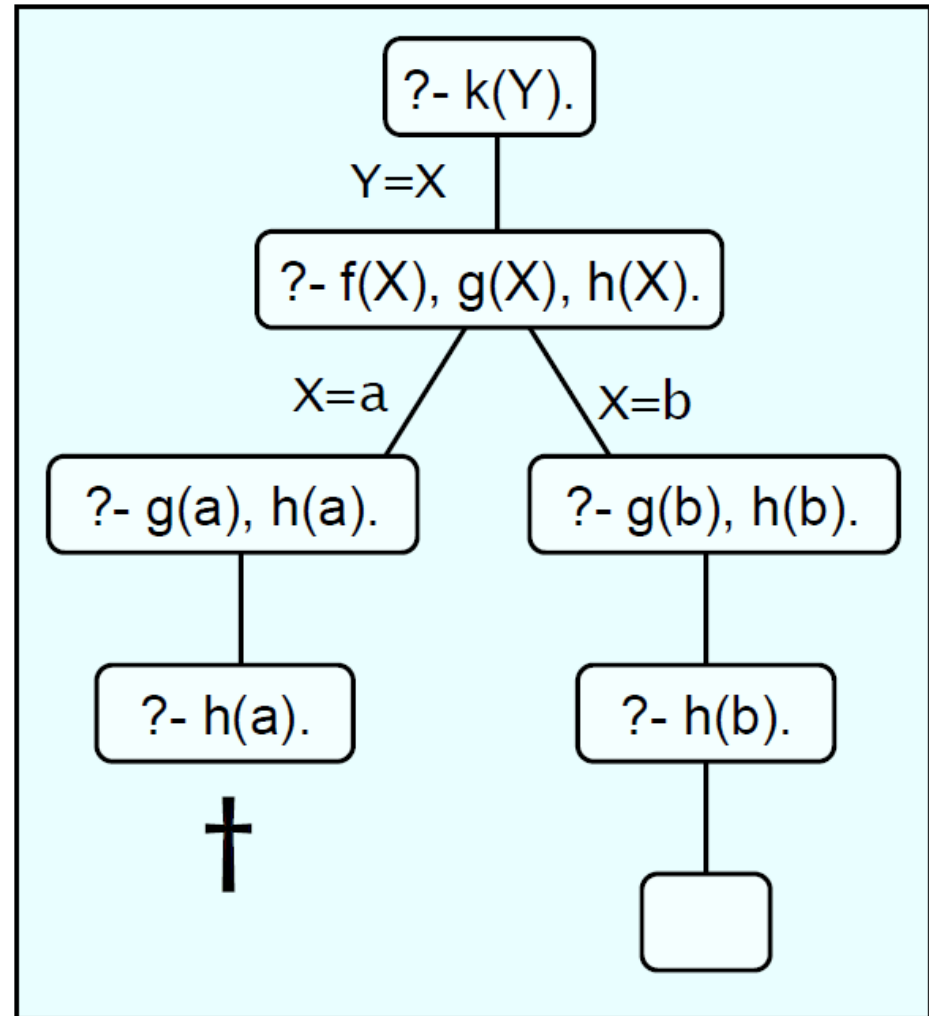


Example: Search Tree



f(a).
f(b).
g(a).
g(b).
h(b).
k(X):- f(X), g(X), h(X).

?- k(Y).
Y=b

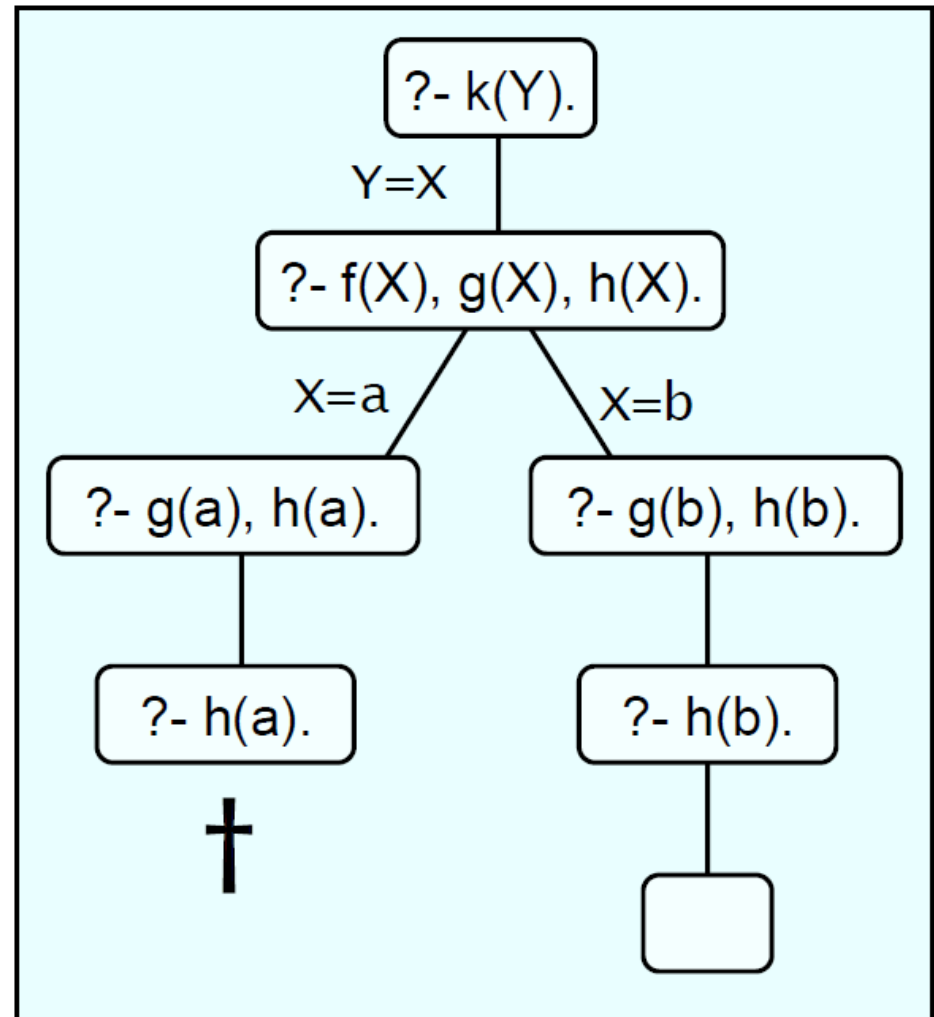


Example: Search Tree



f(a).
f(b).
g(a).
g(b).
h(b).
k(X):- f(X), g(X), h(X).

?- k(Y).
Y=b;
no
?-



Exercise



We are working with the following knowledge base:

house_elf(dobby).

witch(hermione).

witch('McGonagall').

witch(rita_skeeter).

magic(X):- house_elf(X).

magic(X):- wizard(X).

magic(X):- witch(X).

Draw the search tree for:

?- magic(Hermione).

Next...



Recursion
Lists
Arithmetic