

10 September, 2015

ROBT305 – Embedded Systems

Lecture 7 – Real-Time Systems.

Basic definitions



## **Course Logistics**

#### **Reading Assignment:**

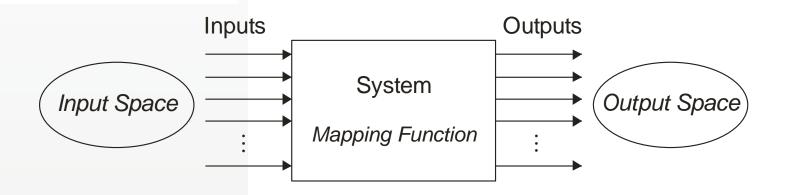
Chapter I of the Real-Time Systems Design and Analysis textbook

Homework Assignment #1 is out in Moodle and due to end of 13 September (Sunday)

Quiz #2 is on 17 September - Pthreads, Mutexes

### **Definition: System**

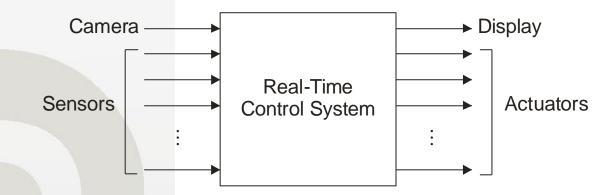
## A system is a mapping of a set of inputs into a set of outputs



- I. A system is an assembly of components connected together in an organized way
- 2. A system is fundamentally altered if a component joins or leaves it
- 3. It has a purpose
- 4. It has a degree of permanence
- 5. It has been defined as being of particular interest

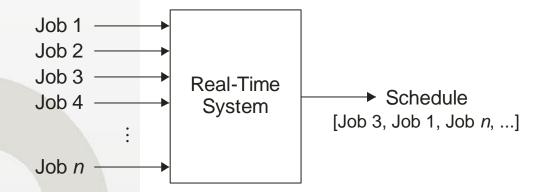
## **Example: A Real-Time Control System**

- Inputs are excitations and outputs are corresponding responses
- Inputs and outputs may be digital or analog
- Inputs are associated with sensors, cameras, etc.
- Outputs with actuators, displays, etc.



## Classic Representation

- A sequence of jobs to be scheduled and performance to be predicted
- Ignores the usual fact that the input sources and hardware under control may be highly complex



## **Definition: Response Time**

The time between the presentation of a set of inputs to a system and the realization of the required behavior, including the availability of all associated outputs, is called **the response time** of the system

- How fast and punctual does it need to be?
  - Depends on the specific real-time system
- But what is a real-time system?

## **Definitions: Real-Time System**

- A real-time system is a computer system that must satisfy bounded response-time constraints or risk severe consequences, including failure
- A real-time system is one whose logical correctness is based on both the correctness of the outputs and their timeliness

## **Definition: Failed System**

**A failed system** is a system that cannot satisfy one or more of the requirements stipulated in the system requirements specification

 Hence, rigorous specification of the system operating criteria, including timing constraints, is necessary

## **Definition: Embedded System**

An embedded system is a system containing one or more computers (or processors) having a central role in the functionality of the system, but the system is not explicitly called a computer

- A real-time system may be embedded or non-embedded
- But it is always reactive
  - Task scheduling is driven by ongoing interaction with the environment

## Degrees of "Real-Time"

- All practical systems are ultimately real-time systems
- Even a batch-oriented system—for example, grade processing at the end of a semester—is real-time
- Although the system may have response times of days, it must respond within a certain time
- Even a word-processing program should respond to commands within a reasonable amount of time
- Most of the literature refers to such systems as soft realtime systems

## Soft, Hard, and Firm "Real-Time"

#### Definition: Soft Real-Time System

A soft real-time system is one in which performance is degraded but not destroyed by failure to meet response-time constraints

#### Definition: Hard Real-Time System

A hard real-time system is one in which failure to meet even a single deadline may lead to complete or catastrophic system failure

#### Definition: Firm Real-Time System

A firm real-time system is one in which a few missed deadlines will not lead to total failure, but missing more than a few may lead to complete or catastrophic system failure

## **Example: Real-Time Classification**

System	Real-Time Classification	Explanation
Avionics weapons delivery system in which pressing a button launches an air-to-air missile	Hard	Missing the deadline to launch the missile within a specified time after pressing the button may cause the target to be missed, which will result in a catastrophe
Navigation controller for an autonomous weed-killer robot	Firm	Missing a few navigation deadlines causes the robot to veer out from a planned path and damage some crops
Console hockey game	Soft	Missing even several deadlines will only degrade performance

#### Where Do Deadlines Come from?

- Deadlines are based on the underlying physical phenomena of the system under control
- Punctuality is another measure related to response times
  - Particularly important in periodically sampled systems with high sampling rates (e.g., in audio and video signal processing)
- Rule of thumb for cost-effective real-time systems:
  - Process everything as slowly as possible and repeat tasks as seldom as possible

## **Definition: Real-Time Punctuality**

Real-time punctuality means that every response time has an average value,  $t_R$ , with upper and lower bounds of  $t_R + \varepsilon_U$  and  $t_R - \varepsilon_L$ , respectively, and  $\varepsilon_U, \varepsilon_L \to 0^+$ 

- In all practical systems, the values of  $\varepsilon_U$  and  $\varepsilon_L$  are nonzero, though they may be very small
- The nonzero values are due to cumulative latency and propagation-delay components (hardware/software)
- Such response times contain jitter within the interval  $t \in [-\varepsilon_{\rm L}, +\varepsilon_{\rm U}]$

# Example: Where a Response Time Comes from?

- An elevator door is automatically operated and it may have a capacitive safety edge for sensing possible passengers between the closing door blades
- Thus, the door blades can be quickly reopened before they touch the passenger and cause discomfort or even threaten the passenger's safety
- What is the required system response time from when it recognizes that a passenger is between the closing door blades and starting to reopen the door?



## **Door Reopening Example Cont'd**

This response time consists of five independent latency components:

Sensor:

Hardware:

System software:

Application software:

Door drive:

 $t_{\rm S min} = 5 \, \rm ms$ 

 $t_{\rm HW~min} = 1~\mu s$ 

 $t_{\rm SS~min} = 16~\mu s$ 

 $t_{\text{AS\_min}} = 0.5 \,\mu\text{s}$ 

 $t_{\rm DD\_min} = 300 \, \mathrm{ms}$ 

 $t_{\rm S\ max} = 15\ {\rm ms}$ 

 $t_{\text{HW}_{-}\text{max}} = 2 \,\mu\text{s}$ 

 $t_{\rm SS\_max} = 48 \, \mu \text{s}$ 

 $t_{\rm AS~max} = 0.5~\mu s$ 

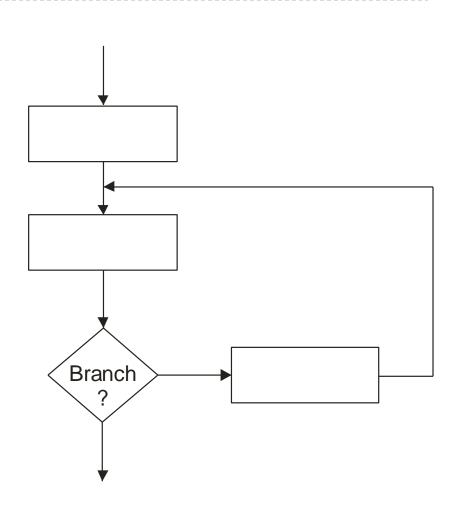
 $t_{\rm DD\_max} = 500 \, \mathrm{ms}$ 

Now, we can calculate the minimum and maximum values of the composite response time:  $t_{\min} \approx 305 \text{ ms}$ ,  $t_{\max} \approx 515 \text{ ms}$ 

The overall response time is dominated by the door drive's response time containing the deceleration time of the moving door blades.

## **Change in Flow of Control**

- In software systems a change in state results in a change in the flow-of-control
- The decision block suggests that the program flow can take alternative paths (see right)
- case, if-then, and while
  statements represent a
  possible change in flow-ofcontrol
- Invocation of procedures in C represent changes in flow-of-control



#### **Definitions: Event and Release Time**

#### Definition: **Event**

Any occurrence that causes the program counter to change non-sequentially is considered a change of flow-of-control, and thus an event

#### Definition: Release Time

The release time is the time at which an instance of a scheduled task is ready to run, and is generally associated with an interrupt

## **Taxonomy of Events**

- An event can be either synchronous or asynchronous
  - Synchronous events are those that occur at predictable times in the flow-of-control
  - Asynchronous events occur at unpredictable points in the flowof-control and are usually caused by external sources
- Moreover, events can be periodic, aperiodic or sporadic
  - A real-time clock that pulses regularly is a periodic event
  - Events that do not occur at regular periods are called aperiodic
  - Aperiodic events that tend to occur very infrequently are called sporadic

## **Example: Various Types of Events**

Туре	Periodic	Aperiodic	Sporadic
Synchronous	Cyclic code	Conditional branch	Divide-by-zero (trap) interrupt
Asynchronous	Clock interrupt	Regular, but not fixed-period interrupt	Power-loss alarm

## **Deterministic Systems**

- For any physical system, certain states exist under which the system is considered to be out of control
- The software controlling such a system must therefore avoid these states
- In embedded real-time systems, maintaining overall control is extremely important
- Software control of any real-time system and associated hardware is maintained when the next state of the system, given the current state and a set of inputs, is predictable

#### Definition: Deterministic System

A system is deterministic, if for each possible state and each set of inputs, a unique set of outputs and next state of the system can be determined

## **Deterministic Systems Cont'd**

- Event determinism means the next states and outputs of a system are known for each set of inputs that trigger events
- Thus, a system that is deterministic is also event deterministic
- However, event determinism may not imply determinism
- While it is a significant challenge to design systems that are completely event deterministic, it is possible to inadvertently end up with a system that is non-deterministic
- Finally, if in a deterministic system the response time for each set of outputs is known, then, the system also exhibits temporal determinism
- A side benefit of designing deterministic systems is that guarantees can be given that the system will be able to respond at any time, and in the case of temporally deterministic systems, when they will respond

#### **CPU Utilization or Time-Loading Factor**

- The final term to be defined is a critical measure of real-time system performance
- Because the CPU continues to execute instructions as long as power is applied, it will more or less frequently execute instructions that are not related to the fulfillment of a specific deadline
- The measure of the relative time spent doing non-idle processing indicates how much real-time processing is occurring

Definition: CPU Utilization Factor

The CPU utilization or time-loading factor, *U*, is a relative measure of the non-idle processing taking place

#### **CPU Utilization Zones**

- A system is said to be time-overloaded if U > 100%
- Systems that are too highly utilized are problematic
  - Additions, changes, or corrections cannot be made to the system without risk of time-overloading
- On the other hand, systems that are not sufficiently utilized are not necessarily cost-effective
  - The system was over-engineered and that costs could likely be reduced with less expensive hardware
- While a utilization of 50% is common for new products, 80% might be acceptable for systems that do not expect growth
- However, 70% as a target for *U* is one of the potentially useful results in the theory of real-time systems where tasks are periodic and independent

### **CPU Utilization Zones Cont'd**

<b>Utilization</b> %	<b>Z</b> one <b>T</b> ype	Typical Application
< 26	Unnecessarily safe	Various
26 – 50	Very safe	Various
51 – 68	Safe	Various
69	Theoretical limit	Embedded systems
70 – 82	Questionable	Embedded systems
83 – 99	Dangerous	Embedded systems
100	Critical	Marginally stressed system
> 100	Overloaded	Stressed system

#### Calculation of U

- $\triangleright$  U is calculated by summing the contribution of utilization factors for each task
- lacksquare Suppose a system has  $n\geq 1$  periodic tasks, each with an execution period of  $p_i$
- If task *i* is known to have a worst-case execution time of  $e_i$ , then the utilization factor,  $u_i$ , for task *i* is

$$u_i = e_i/p_i \tag{I.I}$$

Furthermore, the overall system utilization factor is

$$U = \sum_{i=1}^{n} u_i = \sum_{i=1}^{n} e_i / p_i$$
 (1.2)

- In practice, the determination of  $e_i$  can be difficult, in which case estimation or measuring must be used
- For aperiodic and sporadic tasks,  $u_i$  is calculated by assuming a worst-case execution period

## Example: Calculation of U

Suppose, an individual elevator controller in a bank of elevators has the following tasks with execution periods of  $p_i$  and worst-case execution times of  $e_i$ ,  $i \in [1,2,3,4]$ :

Task I: Communicate with the group dispatcher.

Task 2: Update the car position information and manage floor-to-floor runs as well as door control.

Task 3: Register and cancel car calls.

Task 4: Miscellaneous system supervisions.

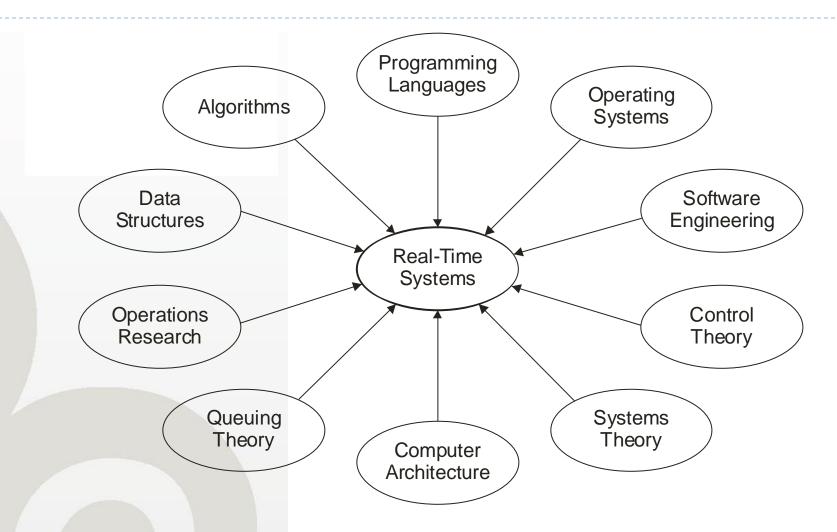
i	$e_i$	$p_i$
I	17 ms	500 ms
2	4 ms	25 ms
3	I ms	75 ms
4	20 ms	200 ms

$$U = \sum_{i=1}^{4} e_i / p_i = 0.31$$

## **Usual Misconceptions**

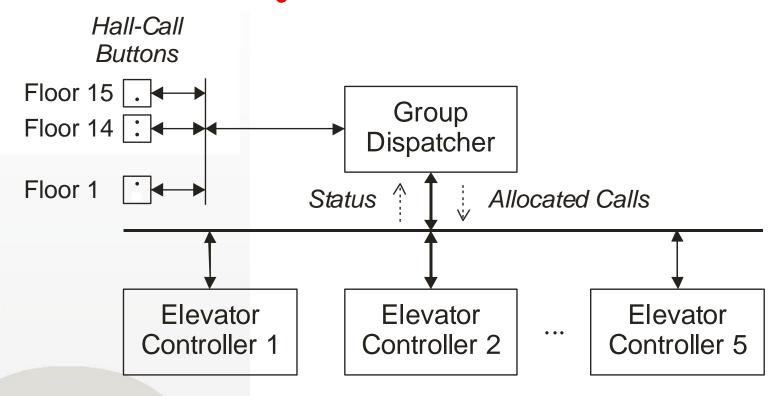
- ► Real-time systems are synonymous with "fast" systems
  - Many (but not all) hard real-time systems deal with deadlines in the tens of milliseconds
- There are universal, widely accepted methodologies for real-time systems specification and design
  - There is still no methodology available that answers all of the challenges of real-time specification and design all the time and for all applications
- There is no more a need to build a real-time operating system, because many commercial products exist
  - Commercial solutions have certainly their place, but choosing when to use an off-the-shelf solution and choosing the right one are continuing challenges

### Rich Variety of "Real-Time" Disciplines



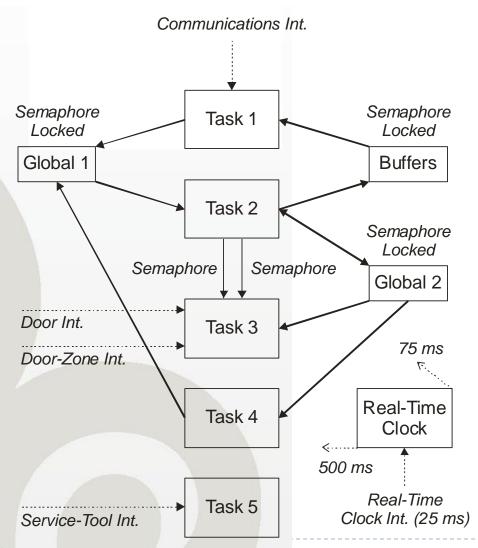
Real-time systems is a truly multi-dimensional subdiscipline of computer systems engineering

## Example: Architecture of an Elevator Bank Control System



- The group dispatcher register and cancel hall calls, and allocates registered calls dynamically to the most suitable elevators depending on their current status (e.g. occupancy, car position, running direction, etc.) with minimum waiting time
- The group dispatcher needs to periodically collect status information from each individual elevator

# **Example: Real-Time Structure of an Elevator Controller**



T_	
Purpose	Rate
Group dispatcher communication.	500 ms
Pass data to Task 2 via global	
variable Global 1.	
Task execution time -15 ms	
Update car position, register car	75 ms
calls, determine destination floors	
Task execution time -17 ms	
Create floor run schedule, control	Aperio
door opening and closing and	-dic
update displays.	
Variable execution time	
It is a Finite State Machine	
Backup services in case of	500 ms
communications failure	
Diagnostics and debugger, Runs	Backg-
1 2	round
	Pass data to Task 2 via global variable Global 1.  Task execution time -15 ms  Update car position, register car calls, determine destination floors Task execution time -17 ms  Create floor run schedule, control door opening and closing and update displays.  Variable execution time  It is a Finite State Machine  Backup services in case of

## **Any Questions?**

