Symmetric Monoidal ∞-Categories

- Note: we will call $(\infty, 1)$-categories ∞-categories.
- Main references: [4], [5], [2], [1], [3], [6]. Referencing Kerodon test: [6, Example 01UB].

1 Pseudofunctors, opfibrations, and symmetric monoidal categories

1.1 Pseudofunctors to Cat are Grothendieck opfibrations

- Grothendieck op-fibrations correspond to pseudofunctors to **Cat**.

The below gives an example to complement [2].

Example of a pseudofunctor to **Cat**: Let $\mathcal{C}$ be a category with pullbacks. Recall that for a map $f : C \to D$ in $\mathcal{C}$, we define a pullback functor

$$\begin{array}{rcl} f^* : \quad \mathcal{C}_{/D} & \to & \mathcal{C}_{/C}, \\ (h : X \to D) & \mapsto & (f^*h : P \to C), \end{array}$$

where we have formed a pullback

$$\begin{tikzcd} P \arrow[r, "h^*f"] \arrow[d, "f^*h"'] & X \arrow[d, "h"] \\ C \arrow[r, "f"'] & D \end{tikzcd}$$

in $\mathscr{C}$. For any map

$$\begin{tikzcd} X \arrow[rr, "\phi"] \arrow[dr, "h"'] & & X' \arrow[dl, "h'"] \\ & D & \end{tikzcd}$$

from h to h' in $\mathcal{C}_{/D}$, we define $f^*\phi$ to be the unique map making the diagram below commute.

$$\begin{tikzcd} X \arrow[rr, "\phi"] \arrow[dr, "h"'] & & X' \arrow[dl, "h'"] \\ & D & \\ & & \\ P \arrow[uuu, "h^*f"] \arrow[rr, dashed, "f^*\phi"] \arrow[dr, "f^*h"'] & & P' \arrow[uuu, "h'^*f"'] \arrow[dl, "f^*h'"] \\ & C \arrow[uuu, "f"] & \end{tikzcd}$$

Now, we may wish to define a functor

$$F : \mathcal{C}^{\mathrm{op}} \to \mathbf{Cat},$$
$$C \mapsto \mathcal{C}_{/C},$$

which sends a map $f : C \to D$ in $\mathcal{C}$ to a pullback functor $f^* : \mathcal{C}_{/D} \to \mathcal{C}_{/C}$. However a problem arises when we check that F respects composition: suppose $f : C \to D$, $g : D \to E$ are maps in $\mathcal{C}$. Then

$$F(g \circ f)(h : X \to E) = (g \circ f)^* h : P \to C,$$

corresponding to the pullback

$$\begin{tikzcd} P \arrow[r, "h^*(g\circ f)"] \arrow[d, "(g\circ f)^*h"'] & X \arrow[d, "h"] \\ C \arrow[r, "g\circ f"'] & E \end{tikzcd}$$

in $\mathcal{C}$. On the other hand,

$$(F(g) \circ F(f))(h : X \to E) = f^*(g^* h) : P'' \to C,$$

which corresponds to the diagram below.

$$\begin{tikzcd} P'' \arrow[r, "(g^*h)^*f"] \arrow[d, "f^*(g^*h)"'] & P' \arrow[r, "h^*g"] \arrow[d, "g^*h"] & X \arrow[d, "h"] \\ C \arrow[r, "f"'] & D \arrow[r, "g"'] & E \end{tikzcd}$$

The outer square is indeed a pullback square, since the inner two squares are, so we have a unique isomorphism $P \cong P''$. However, we do not in general have equality. This is because pullbacks are only unique up to unique isomorphism, and in defining a pullback functor we made arbitrary (and not necessarily compatible) choices of P, P' and P''. Thus, we have not defined a functor $\mathcal{C}^{\mathrm{op}} \to \mathbf{Cat}$, rather, we have defined what is known as a *pseudofunctor*.

(definition of a pseudofunctor)

The above example is one way in which pseudofunctors into **Cat** naturally arise; another common example is the pseudofunctor

$$\mathbf{CRing} \to \mathbf{Cat}$$
$$R \mapsto R\text{-}\mathbf{Mod},$$

which sends a ring homomorphism $\phi : R \to S$ to the functor $- \otimes_R S : R\text{-}\mathbf{Mod} \to S\text{-}\mathbf{Mod}$ (extension of scalars). However, to give the data of a pseudofunctor $F : \mathcal{C} \to \mathbf{Cat}$, we must specify not only the functions $\mathrm{ob}(\mathcal{C}) \to \mathrm{ob}(\mathbf{Cat})$ and $\mathrm{Hom}_{\mathcal{C}}(X, Y) \to \mathrm{Hom}_{\mathbf{Cat}}(F(X), F(Y))$ for each $X, Y \in \mathcal{C}$, but also natural isomorphisms

$$F(\mathrm{id}_X) \cong \mathrm{id}_{F(X)}, \quad F(g \circ f) \cong F(g) \circ F(f).$$

That's a pain, let's use Grothendieck opfibrations instead.

Definition 1.1.1. Let $p : X \to \mathcal{C}$ be a functor, and let $f : c \to c'$ be a morphism in $\mathcal{C}$. A morphism $\phi : x \to x'$ in X lying over f is *p-cocartesian* if for any other morphism $\psi : x \to x''$ in X, and for any morphism $g : c' \to p(x'')$ in $\mathcal{C}$ satisfying $g \circ f = p(\psi)$, there exists a unique morphism $\gamma : x' \to x''$ such that $p(\gamma) = g$ and $\psi = \gamma \circ \phi$.

$$
\begin{tikzcd}
& & x'' \arrow[ddd, mapsto] & \\
x \arrow[r, "\phi"'] \arrow[urr, "\forall\psi"] \arrow[d, mapsto] & x' \arrow[ur, dashed, "\exists!\gamma"'] \arrow[d, mapsto] & & X \arrow[d, "p"'] \\
c \arrow[r, "f"] \arrow[drr, "p(\psi)"'] & c' \arrow[dr, "\forall g"] & & \mathcal{C} \\
& & p(x'') &
\end{tikzcd}
$$

DEFINITION 1.1.2. Let $p : X \to \mathcal{C}$ be a functor. Then p is a *Grothendieck opfibration* if for any morphism of $\mathcal{C}$ and any lift of its source, there is a p-cocartesian morphism with that source lying over it.

1.2 Symmetric monoidal categories are special pseudofunctors to Cat

- Now that we know how to move between pseudofunctors to **Cat** and opfibrations, let's write the data of a symmetric monoidal category as a pseudofunctor to **Cat**.

1.3 ...which are special opfibrations

- The above implies there's some category $\mathcal{D}$ such that opfibrations $\mathcal{D} \to \mathrm{Fin}_*$ are the same as symmetric monoidal categories. Let's see what $\mathcal{D}$ is.
- (Usual) definition of symmetric monoidal category and translation into the language of op-fibrations into Fin_*, running example of $\mathbf{Vect}_k$ with $\otimes$ or $\times$.
- Possibly mention swapping out $\mathbf{Fin}_*$ for Δ^{op} gives a monoidal category rather than a symmetric monoidal category. (How do we get a braided monoidal category? Apparently there is no base 1-category we can look at opfibrations into, because the correct formulation is with E_2, which has higher homotopy groups on the mapping spaces.)
- Correspondence of symmetric monoidal functors with morphisms of opfibrations.

The construction below is in [4].

Let $(\mathcal{C}, \otimes)$ be a symmetric monoidal category. We define a new category $\mathcal{C}^{\otimes}$, whose objects are finite (possibly empty) sequences of objects of $\mathcal{C}$, denoted by $[C_1, ..., C_n]$. A morphism

$$[C_1, ..., C_n] \to [C'_1, ..., C'_m]$$

consists of a subset $S \subseteq \{1, ..., n\}$, a map of finite sets $\alpha : S \to \{1, ..., m\}$, and a collection of morphisms $\{f_j : \bigotimes_{i \in \alpha^{-1}\{j\}} C_i \to C'_j\}_{1 \le j \le m}$ in $\mathcal{C}$.

For two morphisms $f : [C_1, ..., C_n] \to [C'_1, ..., C'_m]$ and $g : [C'_1, ..., C'_m] \to [C''_1, ..., C''_l]$, determining two subsets $S \subseteq \{1, ..., n\}$ and $T \subseteq \{1, ..., m\}$ and maps $\alpha : S \to \{1, ..., m\}$, $\beta : T \to \{1, ..., l\}$, the composition $g \circ f$ is given by the subset $U = \alpha^{-1}T \subseteq \{1, ..., n\}$, the map $\beta \circ \alpha : U \to \{1, ..., l\}$ and the maps

$$\left\{ \bigotimes_{i \in (\beta\circ\alpha)^{-1}\{k\}} C_i \cong \bigotimes_{j \in \beta^{-1}\{k\}} \bigotimes_{i \in \alpha^{-1}\{j\}} C_i \to \bigotimes_{j \in \beta^{-1}\{k\}} C'_j \to C''_k \right\}_{1 \le k \le l}.$$

For example, let

$$f : [C_1, C_2, C_3, C_4] \to [C'_1, C'_2, C'_3]$$

be a morphism in $\mathcal{C}^{\otimes}$ consisting of the subset $\{1, 2, 3\} \subseteq \{1, 2, 3, 4\}$, the map

$$\begin{aligned} \alpha : \{1, 2, 3\} &\to \{1, 2, 3\}, \\ 1 &\mapsto 1, \\ 2 &\mapsto 2, \\ 3 &\mapsto 3, \end{aligned}$$

and morphisms

$$f_1 : C_1 \to C'_1, \quad f_2 : C_2 \otimes C_3 \to C'_2, \quad f_3 : \mathbf{1} \to C'_3,$$

and let

$$g : [C'_1, C'_2, C'_3] \to [C''_1, C''_2, C''_3]$$

be a morphism in $\mathcal{C}^{\otimes}$ consisting of the subset $\{1, 2, 3\} \subseteq \{1, 2, 3\}$, the map

$$\begin{aligned} \beta : \{1, 2, 3\} &\to \{1, 2, 3\}, \\ 1, 2, 3 &\mapsto 3, \end{aligned}$$

and morphisms

$$g_1 : \mathbf{1} \to C''_1, \quad g_2 : \mathbf{1} \to C''_2, \quad g_3 : C'_1 \otimes C'_2 \otimes C'_3 \to C''_3.$$

Then the composition $g \circ f$ consists of the subset $\alpha^{-1}\{1, 2, 3\} = \{1, 2, 3\} \subseteq \{1, 2, 3, 4\}$, the map

$$\begin{aligned} \beta \circ \alpha : \{1, 2, 3\} &\to \{1, 2, 3\}, \\ 1, 2, 3 &\mapsto 3, \end{aligned}$$

and the morphisms

$$(g \circ f)_1 = g_1, \quad (g \circ f)_2 = g_2, \quad (g \circ f)_3 = g_3 \circ (f_1 \otimes f_2 \otimes f_3).$$

(really?)

(some intuition on this, tensor along the fibres, etc)

Claim: the forgetful functor

$$\begin{aligned} p : \quad \mathcal{C}^{\otimes} &\to \mathrm{Fin}_*, \\ [C_1, ..., C_n] &\mapsto \langle n \rangle_* \end{aligned}$$

is an opfibration. (It almost tautologically is).

1.4 Generalisation to ∞-categories

- Translation of the above into ∞-categorical language.
- A functor $p : D \to C$ between ordinary categories is a Grothendieck opfibration if and only if the induced functor $N(p) : N(D) \to N(C)$ on nerves is a cocartesian fibration – I *think* I have finally managed to prove this!
- Some examples (nerve of an ordinary symmetric monoidal category, currently trying to find more examples – many people talk about **Sp**, but it seems like I'd need a lot of background to understand this).

- If an ∞-category has finite (co)products, there is a (co)cartesian monoidal structure on $\mathcal{C}$. And we would have hoped so, because it's definitely true for 1-categories!
- Algebra objects in monoidal (∞-)categories
- Possibly generalisation to ∞-operads, depending on how much the above comes to or if I find anything fun to do with symmetric monoidal ∞-categories.
- Might be cool to try to look at E_k algebras, to resolve the earlier mystery of how to write braided monoidal categories.

DEFINITION 1.4.1 ([1], Def 2.1). A functor $p : X \to Y$ between simplicial sets is an *inner fibration* if for all $n \geq 2$, all $0 < k < n$, and any solid arrow commutative square

$$\begin{tikzcd} \Lambda^n_k \arrow[r] \arrow[d, hook] & X \arrow[d, "p"] \\ \Delta^n \arrow[r] \arrow[ur, dashed] & Y \end{tikzcd}$$

there exists a dotted lift.

EXAMPLE 1.4.2. Let $\mathcal{C}, \mathcal{D}$ be categories, and $p : \mathcal{C} \to \mathcal{D}$ be a functor between them. Then $N(p) : N\mathcal{C} \to N\mathcal{D}$ is an inner fibration.

The following proposition is stated without proof in Section 2.3 of [5].

PROPOSITION 1.4.3. Let $p : X \to Y$ be an inner fibration, and suppose that the diagram below is a pullback square in **sSet**.

$$\begin{tikzcd} X' \arrow[r, "f"] \arrow[d, "p'"'] \arrow[dr, phantom, "\lrcorner", very near start] & X \arrow[d, "p"] \\ Y' \arrow[r, "g"'] & Y \end{tikzcd}$$

Then p' is also an inner fibration.

PROOF. Consider the (commutative) solid arrow diagram below.

$$\begin{tikzcd} \Lambda^n_k \arrow[r, "\lambda"] \arrow[d, hook, "\iota"'] & X' \arrow[r, "f"] \arrow[d, "p'"] \arrow[dr, phantom, "\lrcorner", very near start] & X \arrow[d, "p"] \\ \Delta^n \arrow[r, "\delta"'] \arrow[urr, dashed, "\phi"] & Y' \arrow[r, "g"'] & Y \end{tikzcd}$$

Since p is a fibration, there exists a dotted lift ϕ of $g\delta$; that is, $p\phi = g\delta$ and $\phi\iota = f\lambda$. Further, since the right square is a pullback diagram, there exists a unique map $\phi' : \Delta^n \to X'$ making the diagram below commute.

$$\begin{tikzcd} \Delta^n \arrow[dr, dashed, "\phi'"] \arrow[drr, bend left, "\phi"] \arrow[ddr, bend right, "\delta"'] & & \\ & X' \arrow[r, "f"] \arrow[d, "p'"'] \arrow[dr, phantom, "\lrcorner", very near start] & X \arrow[d, "p"] \\ & Y' \arrow[r, "g"'] & Y \end{tikzcd}$$

It remains to show that the triangle below commutes.

$$\begin{tikzcd} \Lambda^n_k \arrow[r, "\lambda"] \arrow[d, hook, "\iota"'] & X' \\ \Delta^n \arrow[ur, dashed, "\phi'"'] & \end{tikzcd}$$

Again, using the universal property of pullbacks, we see that there exist unique dotted maps such that the diagrams below commute.

$$
\begin{tikzcd}
\Lambda^n_k \arrow[dr, dashed] \arrow[drr, bend left, "f\lambda"] \arrow[ddr, bend right, "\delta\iota"'] & & \\
& X' \arrow[r, "f"] \arrow[d, "p'"'] \arrow[dr, phantom, "\lrcorner", very near start] & X \arrow[d, "p"] \\
& Y' \arrow[r, "g"'] & Y
\end{tikzcd}
\qquad
\begin{tikzcd}
\Lambda^n_k \arrow[dr, dashed] \arrow[drr, bend left, "f\phi'\iota"] \arrow[ddr, bend right, "\delta\iota"'] & & \\
& X' \arrow[r, "f"] \arrow[d, "p'"'] \arrow[dr, phantom, "\lrcorner", very near start] & X \arrow[d, "p"] \\
& Y' \arrow[r, "g"'] & Y
\end{tikzcd}
$$

The maps λ and $\phi'\iota$ make the left and right diagrams commute respectively. Further, we note that $f\phi' = \phi$ (by the second diagram) and $\phi\iota = f\lambda$ (since p is an inner fibration), so $f\phi'\iota = f\lambda$. Therefore, the above two diagrams are identical. Thus, by the uniqueness property of pullbacks, $\lambda = \phi'\iota$. $\square$

(Stupid note to self, very obvious but I forget it every now and again):

- If $X : \Delta^{\mathrm{op}} \to \mathbf{Set}$ is a simplicial set, and $\Delta^0 : \Delta^{\mathrm{op}} \to \mathbf{Set} := \mathrm{Hom}(-, [0])$, then a map $F : X \to \Delta^0$ is a natural transformation $(F_n : X_n \to *)_{n \in \mathbb{N}_0}$. That is, such a natural transformation is a family of maps down to a point. In other words, there's only really one natural transformation, so we really *can* view Δ^0 as a point.
- If Y is a simplicial set, and $y \in Y_0$ is a vertex of Y, we can view $\{y\}$ as a copy of Δ^0. Why is this? We can view $\{y\}$ as the constant simplicial set, sending everything to y. Then a natural isomorphism $\Delta^0 \cong \{y\}$ is a collection of isomorphisms $(* \to *)$, of which there is exactly one. Why is it natural? Well, there's only one map from a one-point set to another one-point set, so the square always commutes.

Example 1.4.4 ([1], Ex 2.2). Let $p : X \to \Delta^0$ be the canonical map, and suppose we have the diagram below, such that the outer square commutes.

$$
\begin{tikzcd}
\Lambda^n_k \arrow[r] \arrow[d, hook] & X \arrow[d, "p"] \\
\Delta^n \arrow[r] \arrow[ur, dashed] & \Delta^0
\end{tikzcd}
$$

The lower triangle commutes automatically, so the statement that p is an inner fibration is equivalent to the statement that for all $n \geq 2$, all $0 < k < n$, and any map $\Lambda^n_k \to X$, there exists a dotted lift.

$$
\begin{tikzcd}
\Lambda^n_k \arrow[r] \arrow[d, hook] & X \\
\Delta^n \arrow[ur, dashed] &
\end{tikzcd}
$$

That is, X is an ∞-category.

Now, combining the above argument with Proposition 1.4.3, we see that for any inner fibration $p : X \to Y$, each fibre $X \times_Y \{y\}$ is an ∞-category.

Definition 1.4.5 ([1], Def 3.1). Let $p : X \to Y$ be an inner fibration. An edge $f : \Delta^1 \to X$ of X is *p-cocartesian* if for all $n \geq 2$, any extension

$$
\begin{tikzcd}
\Delta^{\{0,1\}} \arrow[r, "f"] \arrow[d, hook] & X \\
\Lambda^n_0 \arrow[ur, "F"'] &
\end{tikzcd}
$$

and any solid arrow commutative diagram

$$\begin{tikzcd}
\Lambda^n_0 \arrow[r, "F"] \arrow[d, hook] & X \arrow[d, "p"] \\
\Delta^n \arrow[r] \arrow[ur, dashed] & Y
\end{tikzcd}$$

a dotted lift exists.

Definition 1.4.6. Let $p : X \to Y$ be an inner fibration. Then p is a cocartesian fibration if for any edge $\phi : y \to y'$ in Y_1, and for every $x \in X_0$ lying over y, there exists a p-cocartesian edge $f : x \to x'$ of X lying over ϕ.

The following proposition tells us that the above definition is a reasonable generalisation of Definition 1.1.2. It is also stated without proof in [5], which did not do wonders for my ego.

Proposition 1.4.7 ([5], Rmk 2.4.2.2). Let $\mathcal{C}$, $\mathcal{D}$ be categories, and let $p : \mathcal{C} \to \mathcal{D}$ be a functor between them. Then p is a Grothendieck opfibration if and only if the induced map $N(p) : N\mathcal{C} \to N\mathcal{D}$ is a cocartesian fibration of simplicial sets.

Proof. Let $f : d \to d'$ be a morphism of $\mathcal{D}$, and let c lie over d.

Suppose p is a Grothendieck opfibration, let $F : \Lambda^n_0 \to N\mathcal{C}$ be an extension of f, and let

$$\begin{tikzcd}
\Lambda^n_0 \arrow[r, "F"] \arrow[d, hook] & N\mathcal{C} \arrow[d, "N(p)"] \\
\Delta^n \arrow[r] & N\mathcal{D}
\end{tikzcd}$$

be a commutative diagram. If $n = 2$, it follows immediately from the fact that p is an opfibration that a dotted lift exists. Further, if $n > 3$, there is nothing to check, since an n-simplex in a category commutes if and only if all of its triangles commute, which is guaranteed for any extension $F : \Lambda^n_0 \to N\mathcal{C}$. We thus prove the proposition for $n = 3$.

Suppose we have an extension $F : \Lambda^3_0 \to N\mathcal{C}$ of f; that is, a tetrahedron

$$\begin{tikzcd}
& c \arrow[rr, "\chi"] \arrow[dl, "\phi"'] \arrow[dr, "\psi"] & & c''' \\
c' \arrow[rr, "\gamma"'] \arrow[urrr, "\gamma'"] & & c'' \arrow[ur, "\gamma''"'] &
\end{tikzcd}$$

such that all faces containing the vertex c commute. Let

$$\begin{tikzcd}
& d \arrow[rr, "p(\chi)"] \arrow[dl, "f"'] \arrow[dr, "p(\psi)"] & & p(c''') \\
d' \arrow[rr, "p(\gamma)"'] \arrow[urrr, "p(\gamma')"] & & p(c'') \arrow[ur, "p(\gamma'')"'] &
\end{tikzcd}$$

be a commutative tetrahedron in $\mathcal{D}$. We claim that the tetrahedron in $\mathcal{C}$ commutes. First, note that $\gamma'' \circ \gamma$ is a lift of $p(\gamma')$, since $p(\gamma') = p(\gamma'') \circ p(\gamma) = p(\gamma'' \circ \gamma)$. Further,

$$\begin{aligned}
(\gamma \circ \gamma'') \circ \phi &= \gamma'' \circ \psi \\
&= \chi.
\end{aligned}$$

Thus, by the uniqueness in the universal property of ϕ, we have that $\gamma' = \gamma'' \circ \gamma$, as required.

Now, suppose $N(p)$ is a cocartesian fibration. Then there exists a lift $\phi : c \to c'$ of f, and, in particular, for any diagram

$$\begin{tikzcd} & c \arrow[dl, "\phi"'] \arrow[dr, "\psi"] & \\ c' & & c'' \end{tikzcd}$$

in $\mathcal{C}$, and any commutative diagram

$$\begin{tikzcd} & d \arrow[dl, "f"'] \arrow[dr, "p(\psi)"] & \\ d' \arrow[rr, "g"'] & & p(c'') \end{tikzcd}$$

in $\mathcal{D}$, there exists a map $\gamma : c' \to c''$ such that γ lies over g and $\gamma \circ \phi = \psi$. It remains to show that γ is unique.

Suppose that there were two maps $\gamma_1, \gamma_2 : c' \to c''$ lying over g and satisfying $\gamma_1 \circ \phi = \gamma_2 \circ \phi = \psi$. Then we would have a tetrahedron

$$\begin{tikzcd} & c \arrow[rr, "\psi"] \arrow[dl, "\phi"'] \arrow[dr, "\psi"] & & c'' \\ c' \arrow[rr, "\gamma_1"'] \arrow[urrr, "\gamma_2"] & & c'' \arrow[ur, "\mathrm{id}"'] & \end{tikzcd}$$

where all faces containing the vertex c commute. The image of this tetrahedron under p commutes in $\mathcal{D}$, so the original tetrahedron must commute in $\mathcal{C}$; that is, $\gamma_1 = \gamma_2$. □

1.5 A nontrivial example

Throughout this section, $\mathcal{A}$ is an abelian category, and $\mathcal{A}_{\text{proj}}$ is the full subcategory of $\mathcal{A}$ spanned by the projective objects.

Definition 1.5.1 ([4], Def 1.2.3.1). A *chain complex* with values in $\mathcal{A}$ is a composable sequence of morphisms

$$\cdots \to A_2 \xrightarrow{d(2)} A_1 \xrightarrow{d(1)} A_0 \xrightarrow{d(0)} A_{-1} \to \cdots$$

in $\mathcal{A}$ such that $d(n-1) \circ d(n) = 0$ for all $n \in \int$. The collection of chain complexes with values in $\mathcal{A}$ is an additive category, $\mathrm{Ch}(\mathcal{A})$.

Definition 1.5.2 ([4], Not 1.3.2.6). $\mathrm{Ch}^-(\mathcal{A})$ is the full subcategory of $\mathrm{Ch}(\mathcal{A})$ spanned by those chain complexes M_* such that $M_n \simeq 0$ for $n << 0$.

Definition 1.5.3 ([4], Def 1.3.2.7). Suppose $\mathcal{A}$ has enough projective objects. We let $\mathcal{D}^-(\mathcal{A})$ denote the ∞-category $N_{\text{dg}}(\mathrm{Ch}^-(\mathcal{A}_{\text{proj}}))$. We refer to $\mathcal{D}^-(\mathcal{A})$ as the *derived* ∞*-category of* $\mathcal{A}$.

References

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