

# Cross-Layer Resource Allocation with elastic service scaling in Cloud Radio access network

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February 24, 2018



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IEEE Transactions on Wireless Communications, vol 14, no. 9

September 2015



# Assumption

## Assumption

$$\mu_i, c_i > \lambda_i$$

$$\varphi_i(\mu_i) \geq 0, \forall \mu_i$$

$\varphi_i(\mu_i)$  is a convex and increasing function of  $\mu_i$

- $\mu_i$  service rate
- $c_i$  transmission rate
- $\lambda_i$  arrival rate
- $\varphi_i \mu_i = k_i \mu_i^{a_i}$
- $k_i > 0$
- $a_i > 1$



Minimizing system power consumption in C-RAN, which consists of three components:

- Power consumption in BBU pool
- Power consumption in fibre links
- Power consumption in RRHs



## Delay

$$d_i = \frac{1}{\mu_i - \lambda_i} + \frac{1}{c_i - \lambda_i}$$

- $d_i$  total delay
- $\mu_i$  service rate
- $\lambda_i$  arrival rate
- $c_i$  transmission rate



## Received signal at UE $i$

$$\hat{x}_i = \sum_{j \in \mathcal{A}} h_{ij}^H w_{ij} x_i + \sum_{k \neq i} \sum_{j \in \mathcal{A}} h_{ij}^H w_{kj} x_k + \delta_i$$

- $\mathcal{A}$  set of active RRHs
- $x_i$  data symbol for  $i$ th user
- $w_{ij} \in \mathbb{C}^k$  transmit beamformer for UE  $i$  from RRH  $j$
- $h_{ij}^H \in \mathbb{C}^k$  channel from RRH  $j$  to UE  $i$
- $\delta_i \sim \mathcal{CN}(0, \sigma_i)$  additive white Gaussian noise(AWGN) at UE  $i$
- $i \in \mathcal{N}$
- $j \in \mathcal{A}$



## Signal-to-Interference-plus-Noise-Ratio(SINR) at UE $i$

$$\text{SINR}_i(\mathcal{A}) = \frac{\left| \sum_{j \in \mathcal{A}} h_{ij}^H w_{ij} \right|^2}{\sigma_i^2 + \sum_{k \neq i} \left| \sum_{j \in \mathcal{A}} h_{ij}^H w_{kj} \right|^2}$$

- $\sigma_i$
- $\mathcal{A}$  set of active RRHs
- $w_{ij} \in \mathbb{C}^k$  transmit beamformer for UE  $i$  from RRH  $j$
- $h_{ij}^H \in \mathbb{C}^k$  channel from RRH  $j$  to UE  $i$
- $i \in \mathcal{N}$
- $j \in \mathcal{A}$



The achievable rate  $c_i$  of UE  $i$  should satisfy

$$c_i \leq B_i \log(1 + \text{SINR}_i(\mathcal{A}))$$

- $B_i$  bandwidth for UE  $i$
- SINR Signal-to-Interference-plus-Noise-Ratio
- $\mathcal{A}$  set of active RRHs





Each RRH  $j$  has maximum transmitting power constraint

$$\sum_{i=1}^N w_{ij}^H w_{ij} = \sum_{i=1}^N \|w_{ij}\|^2 \leq E_j$$

- $w_{ij} \in \mathbb{C}^k$  transmit beamformer for UE  $i$  from RRH  $j$
- $i \in \mathcal{N}$
- $j \in \mathcal{L}$



# Problem Formulation

$$\min_{\mu_i, c_i, w_{ij}, \mathcal{A}} \sum_i^N \varphi_i(\mu_i) + |\mathcal{A}| P_f + \frac{1}{\eta} \sum_{i=1}^N \sum_{j \in \mathcal{A}} w_{ij}^H w_{ij}$$

subject to

$$\frac{1}{\mu_i - \lambda_i} + \frac{1}{c_i - \lambda_i} \leq \tau_i$$

$$\lambda_i < \mu_i, \lambda_i < c_i$$

$$c_i i \leq B_i \log(1 + \text{SINR}_i(\mathcal{A}))$$

$$\sum_{i=1}^N w_{ij}^H w_{ij} \leq E_i, \quad \forall i \in \mathcal{N}, \quad \forall j \in \mathcal{L}$$

- $\varphi_i(\mu_i)$  VM  $i$ 's power consumption
- $\mu_i$  VM  $i$ 's computation capacity/service rate/processing rate
- $P_f$  power consumption active fibre links
- $\eta \in (0, 1)$  inefficiency coefficient amplifier in RRH
- $\mathcal{A}$  set of active RRHs
- $w_{ij} \in \mathbb{C}^k$  transmit beamformer for UE  $i$  from RRH  $j$
- $i \in \mathcal{N}$
- $j \in \mathcal{A}$



## Quasi Weighted Sum-Rate Maximization Problem



# Quasi Weighted Sum-Rate Maximization Problem

## QWSRM

$$\min_{c_i, w_{ij}} \sum_{i=1}^N -\varepsilon_i c_i$$

subject to,

$$c_i \leq B_i \log(1 + \text{SINR}_i(\mathcal{A})), \quad \forall i \in \mathcal{N}$$

$$\sum_{i=1}^N w_{ij}^H w_{ij} \leq E_j, \quad \forall j \in \mathcal{A}$$

- $\varepsilon_i$  an arbitrary nonnegative weight
- $c_i$  transmission rate
- $B_i$  bandwidth for UE  $i$
- $\mathcal{A}$  set of active RRHs
- $E_j$  maximum power on a link
- $w_{ij} \in \mathbb{C}^k$  transmit beamformer for UE  $i$  from RRH  $j$



# Generalization of WSRM

$$\min_{\mathbf{c}_i, w_{ij}} f(\mathbf{c})$$

subject to,

$$\sum_{i=1}^N w_{ij}^H w_{ij} \leq E_i, \quad \forall j \in \mathcal{A}$$

Where the objective function  $f(\mathbf{c})$ , for  $0 \leq \mathbf{c} \leq \bar{\mathbf{c}}$  has the following properties:

- $f(\mathbf{c})$  is a function only of  $\mathbf{c}$ , and
- $f(\mathbf{c}) < \infty$  is continuously differentiable, and
- $f(\mathbf{c})$  is convex in the feasible region
- $\mathbf{c} = [c_1, \dots, c_N]^T$
- $w_{ij} \in \mathbb{C}^k$  transmit beamformer for UE  $i$  from RRH  $j$
- $c_i$  transmission rate
- $E_j$  maximum power on a link
- $\mathcal{A}$  set of active RRHs



$$\gamma_{ub}(\mathcal{Q}) = \begin{cases} f(\mathbf{c}_{min}), & \mathbf{c}_{min} \in \mathcal{F} \\ +\infty, & \text{otherwise} \end{cases}$$

$$\gamma_{lb}(\mathcal{Q}) = \begin{cases} f(\mathbf{c}_{max}), & \mathbf{c}_{max} \in \mathcal{F} \\ +\infty, & \text{otherwise} \end{cases}$$

- $\mathbf{c}_{min} = [c_{1,min}, \dots, c_{N,min}]^T$
- $\mathbf{c}_{max} = [c_{1,max}, \dots, c_{N,max}]^T$



# BnB algorithm for QWSRM problem

**Input:**  $Q_{init}$ ,  $\mathcal{A}$ , and  $\{(c)\}$ .

**Initialize:** Obtain  $c_i$  by solving  $\frac{\delta f(c)}{\delta c_i} = 0$ , for  $i \in \mathcal{N}$ . Set  $k = 1, \mathcal{B} = Q_{init}, u_1 = \gamma_{ub}(Q_{init})$  and  $l_1 = \gamma_{lb}(Q_{init})$ .

Check the feasibility of problem (17) with given

**if feasible then**

$$c_0 = \tilde{c};$$

**else**

**while**  $u_k - l_k > \epsilon$  **do**

Branching:

- Set  $Q_k = Q$ , where  $Q$  satisfies  $\gamma_{lb}(Q) = l_k$ .
- Split  $Q$  into  $Q_I$  and  $Q_{II}$ , along one of its longest edges.
- Update  $\mathcal{B}_{k+1} = (\mathcal{B}_k \setminus \{Q_k\}) \cup (Q_I, Q_{II})$ .

Bounding:

- Update  $u_{k+1} = \min_{Q \in \mathcal{B}_{k+1}} \{\gamma_{ub}(Q)\}$
- Update  $l_{k+1} = \min_{Q \in \mathcal{B}_{k+1}} \{\gamma_{lb}(Q)\}$

**end while**

Set  $c_0 = c_{min}$ ;

**end if**

**Output:**  $c_0$ .

# Thank You

