Cross-Layer Resource Allocation with elastic service scaling in Cloud Radio access network

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Assumption

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$$\mu_i, c_i > \lambda_i$$

$$\varphi_i(\mu_i) \geq 0, \forall \mu_i$$

$\varphi_i(\mu_i)$ is a convex and increasing function of μ_i

- μ_i service rate
- c_i transmission rate
- λ_i arrival rate
- $\varphi_i \mu_i = k_i \mu_i^{a_i}$
- $k_i > 0$
- $a_i > 1$

Aim

Minimizing system power consumtion in C-RAN, which consists of three components:

- Power consumption in BBU pool
- Power consumption in fibre links
- Power consumption in RRHs

Delay equation

Delay

$$d_i = \frac{1}{\mu_i - \lambda_i} + \frac{1}{c_i - \lambda_i}$$

- d_i total delay
- μ_i service rate
- λ_i arival rate
- c_i transmission rate

Equation

Recieved signal at UE i

$$\hat{x}_i = \sum_{j\in\mathcal{A}} h_{ij}^H w_{ij} x_i + \sum_{k\neq i}^{N} \sum_{j\in A} h_{ij}^H w_{kj} x_k + \delta_i$$

- $\mathcal{A}$ set of active RRHs
- x_i data symbol for ith user
- $w_{ij} \in \mathbb{C}^k$ transmit beamformer for UE i from RRH j
- $h_{ij}^H \in \mathbb{C}^k$ channel from RRH j to UE i
- $\delta_i \sim \mathcal{CN}(0, \sigma_i)$ additive white Gaussian noise(AWGN) at UE i
- $i \in \mathcal{N}$
- $j \in \mathcal{A}$

Equation

Signal-to-Interference-plus-Noise-Ratio(SINR) at UE i

$$\mathsf{SINR}_i(\mathcal{A}) = \frac{\left|\sum\limits_{j\in\mathcal{A}} h_{ij}^H w_{ij}\right|^2}{\sigma_i^2 + \sum\limits_{k\neq i}^{N}\left|\sum\limits_{j\in\mathcal{A}} h_{ij}^H w_{kj}\right|}$$

- σ_i
- $\mathcal{A}$ set of active RRHs
- $w_{ij} \in \mathbb{C}^k$ transmit beamformer for UE i from RRH j
- $h_{ij}^H \in \mathbb{C}^k$ channel from RRH j to UE i
- $i \in \mathcal{N}$
- $j \in \mathcal{A}$

Achieveable rate c_i

The achieveable rate c_i of UE i should satisfy

$$c_i \leq B_i \log(1 + \mathsf{SINR}_i(\mathcal{A}))$$

- B_i bandwidth for UE i
- SINR Signal-to-Interference-plus-Noise-Ratio
- $\mathcal{A}$ set of active RRHs

Transmitting Power

Each RRH j has maximum transmitting power constraint

$$\sum_{i=1}^{N} w_{ij}^{H} w_{ij} = \sum_{i=1}^{N} \|w_{ij}\|^2 \leq E_j$$

- $w_{ij} \in \mathbb{C}^k$ transmit beamformer for UE i from RRH j
- $i \in \mathcal{N}$
- $j \in \mathcal{L}$

Optimization Problem

$$\min_{\mu_i, c_i, w_{ij}, \mathcal{A}} \sum_{i}^{N} \varphi_i(\mu_i) + |\mathcal{A}| P_f + \frac{1}{\eta} \sum_{i=1}^{N} \sum_{j \in A} w_{ij}^H w_{ij}$$

subject to

$$\frac{1}{\mu_i - \lambda_i} + \frac{1}{c_i - \lambda_i} \leq \tau_i$$

$$\lambda_i < \mu_i, \lambda_i < c_i$$

$$c_i i \leq B_i \log(1 + SINR_i(\mathcal{A}))$$

$$\sum_{i=1}^{N} w_{ij}^H w_{ij} \leq E_i, \quad \forall i \in \mathcal{N}, \quad \forall j \in \mathcal{L}$$

- $\varphi_i(\mu_i)$ VM i's power consumption
- μ_i VM i's computation capacity/service rate/processing rate
- P_f power consumption active fibre links
- $\eta \in (0, 1)$ ineffficiency coefficient amplifier in RRH
- $\mathcal{A}$ set of active RRHs
- $w_{ij} \in \mathbb{C}^k$ transmit beamformer for UE i from RRH j
- $i \in \mathcal{N}$
- $j \in \mathcal{A}$

BnB algorithm for QWSRM problem

```
Input: $\mathcal{Q}_{\text{init}}, \mathcal{A},$ and{(c).
Initialize: Obtain $c_i$ by solving $\frac{\delta f(c)}{\delta c_i} = 0,$ for $i \in \mathcal{N}$. Set $k = 1, \mathcal{B} = \mathcal{Q}_{init}, u_1 =$
  $\gamma_{ub}(\mathcal{Q}_{init})$ and $l_1 = \gamma_{lb}(\mathcal{Q}_{init})$.
  Check the feasibility of problem (17) with given
  if feasible then
      $c_0 = \widetilde{c}$;
  else
      while $u_k - l_k > \epsilon$ do
          Branching:
            • Set $\mathcal{Q}_k = \mathcal{Q},$ where $\mathcal{Q}$ satisfies $\gamma_{lb}(\mathcal{Q}) = l_k$.
            • Split $\mathcal{Q}$ into $\mathcal{Q}_{\mathrm{I}}$ and $\mathcal{Q}_{\mathrm{II}},$ along one of its longest edges.
            • Update $\mathcal{B}_{k+1} = (\mathcal{B}_k \setminus \{\mathcal{Q}_k\}) \bigcup (\mathcal{Q}_{\mathrm{I}}, \mathcal{Q}_{\mathrm{II}})$.

          Bounding:
            • Update $u_{k+1} = \min_{\mathcal{Q} \in \mathcal{B}_{k+1}} \{\gamma_{ub}(\mathcal{Q})\}$
            • Update $l_{k+1} = \min_{\mathcal{Q} \in \mathcal{B}_{k+1}} \{\gamma_{lb}(\mathcal{Q})\}$

      end while
  Set $c_0 = c_{min}$;
  end if
Output: $c_0$.
```

Thank You