# Cross-Layer Resource Allocation With Elastic Service Scaling in Cloud Radio Access Network

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#### Assumption of the Paper

#### Assumption

$$\mu_i, c_i > \lambda_i$$

$$\varphi_i(\mu_i) \ge 0, \forall \mu_i$$

 $\varphi_i(\mu_i)$  is a convex and increasing function of  $\mu_i$ 

 $\bullet$   $\mu_i$  service rate

ullet  $c_i$  transmission rate

 $\bullet$   $\lambda_i$  arrival rate

 $ullet \, arphi_i(\mu_i) = k_i \mu_i^{a_i} \, {\sf power \, consumption \, for \, } \mu_i$ 

•  $k_i > 0$  and  $a_i > 1$  are positive constants.





#### Aim of the Paper

Minimizing system power consumtion in C-RAN, which consists of three components:

- Power consumption in BBU pool
- Power consumption in fiber links
- Power consumption in RRHs





#### Delay

$$d_i = \frac{1}{\mu_i - \lambda_i} + \frac{1}{c_i - \lambda_i}$$
$$d_i \le \tau_i$$

#### where,

- $d_i$  total delay
- ullet  $\mu_i$  service rate
- ullet  $\lambda_i$  arrival rate
- ullet  $au_i$  is a predefined QoS requirement of UE i





#### Recieved signal at UE i

$$\hat{x}_i = \sum_{j \in \mathcal{A}} h_{ij}^H w_{ij} x_i + \sum_{k \neq i}^N \sum_{j \in \mathcal{A}} h_{ij}^H w_{kj} x_k + \delta_i, \quad i \in \mathcal{N}, j \in \mathcal{A}$$

#### where,

•  $h_{ij}^H \in \mathbb{C}^k$  channel from RRH j to UE i

•  $w_{ij} \in \mathbb{C}^k$  transmit beamformer for UE i from RRH j

ullet  $x_i$  data symbol for ith user

•  $\delta_i \sim \mathcal{CN}(0, \sigma_i)$  additive white Gaussian noise(AWGN) at UE i

ullet set of all UEs

set of active RRHs





#### Signal-to-Interference-plus-Noise-Ratio(SINR) at UE i

$$\mathsf{SINR}_i(\mathcal{A}) = \frac{\left|\sum\limits_{j \in \mathcal{A}} h^H_{ij} w_{ij}\right|^2}{\sigma^2_i + \sum\limits_{k \neq i}^{N} \left|\sum\limits_{j \in \mathcal{A}} h^H_{ij} w_{kj}\right|}, \quad i \in \mathcal{N}, j \in \mathcal{A}$$

where,

N

•  $h_{ij}^H \in \mathbb{C}^k$  channel from RRH j to UE i

 $\bullet \ w_{ij} \in \mathbb{C}^k \qquad \text{transmit beamformer for UE } i \text{ from RRH } j \\$ 

•  $\sigma_i$  noise(variance)

set of all UEs

set of active RRHs



#### The achieveable rate $c_i$ of UE i should satisfy

$$c_i \le B_i \log(1 + \mathsf{SINR}_i(\mathcal{A}))$$

where,

•  $B_i$  bandwidth for UE i

• SINR Signal-to-Interference-plus-Noise-Ratio

ullet set of active RRHs





#### Transmitting Power

#### Each RRH j has maximum transmitting power constraint

$$\sum_{i=1}^{N} w_{ij}^{H} w_{ij} = \sum_{i=1}^{N} ||w_{ij}||^{2} \le E_{j}, \quad i \in \mathcal{N}, j \in \mathcal{L}$$

where,

•  $w_{ij} \in \mathbb{C}^k$ 

transmit beamformer for UE i from RRH j

 $\bullet$   $E_j$ 

maximum power on a link





## Problem Formulation(Minimizing Sys Power Consumption)

$$\min_{\mu_{i}, c_{i}, w_{ij}, A} \sum_{i}^{N} \varphi_{i}(\mu_{i}) + |A| P_{f} + \frac{1}{\eta} \sum_{i=1}^{N} \sum_{j \in A} w_{ij}^{H} w_{ij}$$

#### subject to

$$\frac{1}{\mu_i - \lambda_i} + \frac{1}{c_i - \lambda_i} \le \tau_i$$
$$\mu_i, c_i > \lambda_i$$
$$c_i \le B_i \log(1 + \mathsf{SINR}_i(\mathcal{A}))$$

$$\sum_{i=1}^{N} w_{ij}^{H} w_{ij} \leq E_{i}, \quad \forall i \in \mathcal{N}, \quad \forall j \in \mathcal{L}$$



- $\varphi_i(\mu_i)$  VM i's power consumption
- P<sub>f</sub> power consumption of active fibre links
- $\begin{array}{ll} \bullet & \eta \in (0,1) & \text{ineffficiency} \\ \text{coefficient of amplifier in} \\ \text{RRH} \end{array}$
- $w_{ij} \in \mathbb{C}^k$  transmit beamformer
- ullet  $au_i$  is a predefined QoS requirement
- ullet  ${\cal A}$  set of active RRHs
- $B_i$  bandwidth for UE i
- $\mu_i$  VM i's service rate
- ullet  $c_i$  transmission rate
- ullet  $\lambda_i$  arrival rate



#### **QWSRM**

Quasi Weighted Sum-Rate Maximization Problem





## Quasi Weighted Sum-Rate Maximization Problem

#### QWSRM

$$\min_{c_i, w_{ij}} \sum_{i=1}^{N} -\varepsilon_i c_i$$

subject to,

$$c_i \le B_i \log(1 + \mathsf{SINR}_i(\mathcal{A})), \quad \forall i \in \mathcal{N}$$

$$\sum_{i=1}^{N} w_{ij}^{H} w_{ij} \le E_i, \quad \forall j \in \mathcal{A}$$

 $\bullet$   $\varepsilon_i$ 

arbitrary nonnegative weight

 $c_i$   $B_i$ 

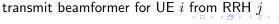
transmission rate bandwidth for UF *i* 

• A

set of active RRHs

 $\bullet$  E

- maximum power on a link
- lacksquare  $w_{ij} \in \mathbb{C}^k$





#### Generalization of WSRM

$$\min_{c_i, w_{ij}} f(\mathbf{c})$$

subject to,

$$\sum_{i=1}^{N} w_{ij}^{H} w_{ij} \le E_i, \quad \forall j \in \mathcal{A}$$

Where the objective function  $f(\mathbf{c})$ , for  $0 \le \mathbf{c} \le \bar{\mathbf{c}}$  has the following properties:

transmission rate

- $f(\mathbf{c})$
- $f(\mathbf{c}) < \infty$
- $f(\mathbf{c})$
- $w_{ij} \in \mathbb{C}^k$
- lacksquare  $c_i$
- $\bullet$   $E_j$
- lacktriangledown lacktriangledown
  - $\mathbf{c} = [c_1, \dots, c_N]^T$

is a function only of  ${\bf c}$ , and is continuously differentiable, and is convex in the feasible region transmit beamformer for UE i from RRH j

maximum power on a link set of active RRHs



#### Limits of BnB Algorithm

$$\gamma_{ub}(\mathcal{Q}) = \begin{cases}
f(\mathbf{c}_{min}), & \mathbf{c}_{min} \in \mathcal{F} \\
+\infty, & otherwise
\end{cases}$$

$$\gamma_{lb}(\mathcal{Q}) = \begin{cases}
f(\mathbf{c}_{max}), & \mathbf{c}_{max} \in \mathcal{F} \\
+\infty, & otherwise
\end{cases}$$

#### where,

$$\bullet \mathbf{c}_{min} = \left[c_{1,min}, \dots, c_{N,min}\right]^{T}$$

• 
$$\mathbf{c}_{max} = \begin{bmatrix} c_{1,max}, \dots, c_{N,max} \end{bmatrix}^T$$

- $\bullet$   $\gamma_{ub}(\mathcal{Q})$  upper bound
- ullet  $\gamma_{lb}(\mathcal{Q})$  lower bound





## BnB algorithm for QWSRM problem

```
Input: Q_{init}, A, and \{(c).
Initialize: Obtain c_i by solving \frac{\delta f(c)}{\delta c_i} = 0, for i \in \mathcal{N}. Set k = 1, \mathcal{B} = \mathcal{Q}_{init}, u_1 = 0
   \gamma_{ub}(\mathcal{Q}_{init}) and l_1 = \gamma_{lb}(\mathcal{Q}_{init}).
   Check the feasibility of problem (17) with given
   if feasible then
         c_0 = \tilde{c};
   else
         while u_k - l_k > \epsilon do
               Branching:
                   • Set Q_k = Q, where Q satisfies \gamma_{lb}(Q) = l_k.
                   • Split Q into Q_I and Q_{II}, along one of its longest edges.
                   • Update \mathcal{B}_{k+1} = (\mathcal{B}_k \setminus \{\mathcal{Q}_k\}) \bigcup (\mathcal{Q}_{\mathrm{I}}, \mathcal{Q}_{\mathrm{II}}).
```

#### Bounding:

- Update  $u_{k+1} = \min_{\mathcal{Q} \in \mathcal{B}_{k+1}} \{ \gamma_{ub}(\mathcal{Q}) \}$
- Update  $l_{k+1} = \min_{\mathcal{Q} \in \mathcal{B}_{k+1}} \{ \gamma_{lb}(\mathcal{Q}) \}$

#### end while

Set 
$$c_0 = c_{min}$$
; end if

Output:  $c_0$ .

## Thank You



