#### INTRODUCCIÓN AL MODELO DE DATOS OMOP-CDM

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More than 20 years teaching

open data - open source – open science

#### I am Alberto

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- European Bioinformatics Institute (Senior Data Scientist)
  - Led the Service Oriented Architecture
  - Genomics databases, web services, semantic web
- Navarrabiomed (Head of bioinformatics)
  - Led the genome sequencing bioinformatics infrastructure (NAGEN, PharmaNAGEN)
  - Multiomics data integration
- IOMED (Head of data)
  - Led data normalization efforts in more than 20 hospitals (OMOP-CDM, NLP)
  - 10M patients, 200million health records
- Barcelona Supercomputing Center (Head of Biomedical Data Hub)
  - Lead the Biomedical Data Hub
  - Open science, FAIR data

## Descripción del curso

El curso tiene como objetivo proporcionar una comprensión profunda del modelo de datos OMOP-CDM de la OHDSI, sus vocabularios y cómo usar la aplicación web Atlas para crear cohortes y extraer datos para estudios observacionales. El curso cubrirá los conceptos y principios básicos del modelo de datos OMOP-CDM, incluida su estructura, tablas clave y vocabularios. Además, el curso cubrirá los aspectos prácticos de trabajar con Atlas, incluido cómo crear y ejecutar consultas de **cohortes**.

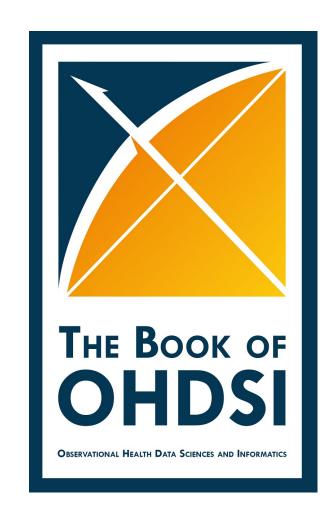
#### Esquema

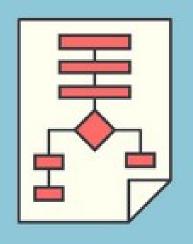
- Introducción al modelo de datos OMOP-CDM
  - Descripción general del modelo de datos OMOP-CDM
  - Conceptos y principios clave del modelo de datos
     OMOP-CDM
  - Comprensión de la estructura y tablas clave del modelo de datos OMOP-CDM
- Vocabularios OMOP-CDM

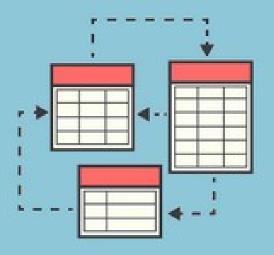


#### Recursos

- Sitio web: https://github.com/alabarga/omop-cdm-course
- Diapositivas
- El Libro de OHDSI
- Entorno local (docker)
  - Base de datos OMOP-CDM
  - Consultas de ejemplo
  - Atlas









## Databases

## Bases de datos

- Una base de datos es un conjunto de datos pertenecientes a un mismo contexto y almacenados sistemáticamente para su posterior uso.
- Un sistema gestor de bases de datos (SGBD) es un software que gestiona una o más bases de datos y nos permite explotar los datos almacenados en ellas de forma relativamente simple mediante SQL

## Modelo relacional

Edgar Frank Codd, en su artículo"A Relational Model of Data for Large Shared Data Banks" en 1970, definió el modelo relacional y publicó una serie de reglas para la evaluación de administradores de sistemas de datos relacionales y así nacieron las bases de datos relacionales.

Information Retrieval

P. BAXENDALE, Editor

#### A Relational Model of Data for Large Shared Data Banks

E. F. Codd IBM Research Laboratory, San Jose, California

Future users of large data banks must be protected from having to know how the data is organized in the machine (the internal representation). A prompting service which supplies such information is not a satisfactory solution. Activities of users at terminals and most application programs should remain unaffected when the internal representation of data is changed and even when some aspects of the external representation are changed. Changes in data representation will often be needed as a result of changes in query, update, and report traffic and natural growth in the types of stored information.

Existing noninferential, formatted data systems provide users with tree-structured files or slightly more general network models of the data. In Section 1, inadequacies of these models are discussed. A model based on n-ary relations, a normal form for data base relations, and the concept of a universal data sublanguage are introduced. In Section 2, certain operations on relations (other than logical inference) are discussed and applied to the problems of redundancy and consistency in the user's model.

KEY WORDS AND PHRASES: data bank, data base, data structure, data organization, hierarchies of data, networks of data, relations, derivability, redundancy, consistency, composition, join, retrieval language, predicate colculus, security, data integrity

CR CATEGORIES: 3.70, 3.73, 3.75, 4.20, 4.22, 4.29

#### 1. Relational Model and Normal Form

#### 1.1. Introduction

This paper is concerned with the application of elementary relation theory to systems which provide shared access to large banks of formatted data. Except for a paper by Childs [1], the principal application of relations to data systems has been to deductive question-answering systems. Levein and Maron [2] provide numerous references to work in this area.

In contrast, the problems treated here are those of data independence—the independence of application programs and terminal activities from growth in data types and changes in data representation—and certain kinds of data inconsistency which are expected to become troublesome even in nondeductive systems.

Volume 13 / Number 6 / June, 1970

The relational view (or model) of data described in Section 1 appears to be superior in several respects to the graph or network model [3, 4] presently in vogue for noninferential systems. It provides a means of describing data with its natural structure only—that is, without superimposing any additional structure for machine representation purposes. Accordingly, it provides a basis for a high level data language which will yield maximal independence between programs on the one hand and machine representation and organization of data on the other.

A further advantage of the relational view is that it forms a sound basis for treating derivability, redundancy, and consistency of relations—these are discussed in Section 2. The network model, on the other hand, has spawned a number of confusions, not the least of which is mistaking the derivation of connections for the derivation of relations (see remarks in Section 2 on the "connection trap").

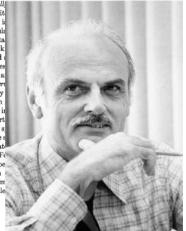
Finally, the relational view permits a clearer evaluation of the scope and logical limitations of present formatted data systems, and also the relative merits (from a logical standpoint) of competing representations of data within a single system. Examples of this clearer perspective are cited in various parts of this paper. Implementations of systems to support the relational model are not discussed.

#### 1.2. Data Dependencies in Present Systems

The provision of data description tables in recently developed information systems represents a major advance toward the goal of data independence [5, 6, 7]. Such tables facilitate changing certain characteristics of the data representation stored in a data bank. However, the variety of data representation characteristics which can be changed

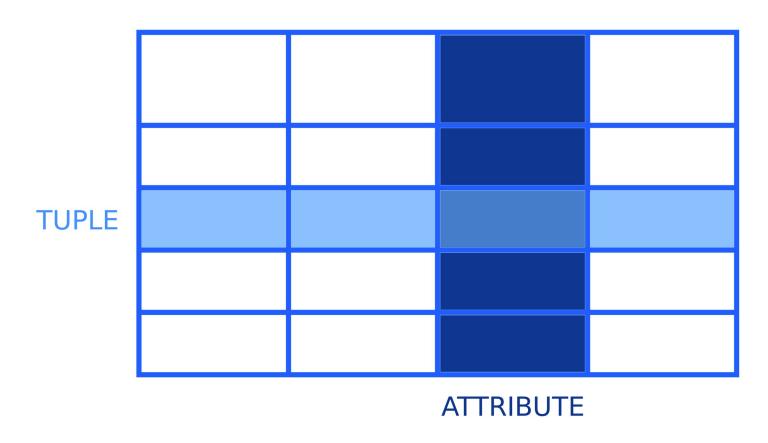
without logically still quite limit users interact is erties, particula lections of data the principal k to be removed; ence, and acces dependencies as

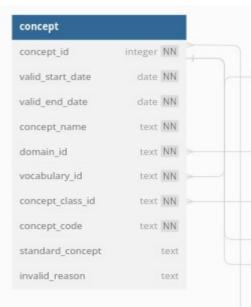
dependencies a
1.2.1. Order
data bank may
ing no concern
to participate ir
element to part
those existing s
elements to be s
closely associat
of addresses. Fe
parts might be
number. Such
grams to assum
from such a file



# Relation = Table

#### **RELATION**





#### concept\_synonym

concept\_id integer NN
language\_concept\_id integer NN
concept\_synonym\_name text NN

source_to_concept_map		
source_concept_id	integer	NN
target_concept_id	integer	NN
valid_start_date	date	NN
valid_end_date	date	NN
source_code	text	NN
source_vocabulary_id	text	NN
target_vocabulary_id	text	NN
source_code_description		text
invalid_reason		text

# vocabulary vocabulary\_concept\_id integer NN vocabulary\_id text NN vocabulary\_name text NN vocabulary\_reference text vocabulary\_version text

concept_id_1	ınteger	NN	
concept_id_2	integer	NN	
valid_start_date	date	NN	
valid_end_date	date	NN	
relationship_id	text	NN	
invalid_reason		text	

concept\_relationship

drug_strength	
drug_concept_id	integer NN
ingredient_concept_id	integer NN
valid_start_date	date NN
valid_end_date	date NN
amount_unit_concept_id	integer
numerator_unit_concept_id	integer
denominator_unit_concept_id	integer
box_size	integer
amount_value	numeric
numerator_value	numeric
denominator_value	numeric
invalid_reason	text

domain	
domain_concept_id	integer
domain_id	text NN
domain_name	text NN

relationship		
relationship_concept_id	integer	NN
relationship_id	text	NN
relationship_name	text	NN
is_hierarchical	text	NN
defines_ancestry	text	NN
reverse_relationship_id	text	NN

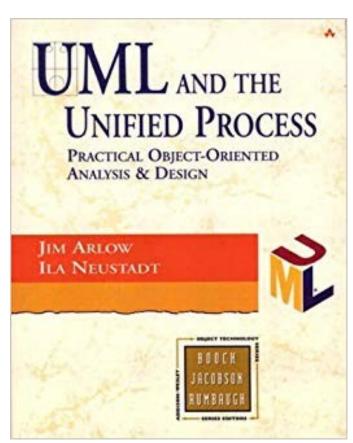
concept_class		
concept_class_concept_id	integer NN	
concept_class_id	text NN	
concept_class_name	text NN	

integer NN
integer NN
integer NN
integer NN

source_to_standard_vocab_r	nap
source_concept_id	integer
target_concept_id	integer
source_valid_start_date	date
source_valid_end_date	date
source_code	text
source_code_description	text
source_vocabulary_id	text
source_domain_id	text
source_concept_class_id	text
source_invalid_reason	text
target_concept_name	text
target_vocabulary_id	text
target_domain_id	text
target_concept_class_id	text

# Unified Modeling Language

- UML (Unified Modeling Language) es un lenguaje que permite modelar, construir y documentar los elementos que forman un sistema software orientado a objetos.
- El estándar UML no define un proceso de desarrollo específico, tan solo se trata de una notación.



# SQL

- El lenguaje principal de una base de datos relacional es SQL (Structured Query Language).
- Ofrece funcionalidades para:
  - Definir (DDL alteramos estructuras de las tablas)
  - Manipular (DML trabajamos con los datos)

# SQL components

- Data Definition Language (DDL)
  - Deals with structural aspect of the database: creation, modification, deletion of tables
- Data Manipulation Language (DML)
  - This allows modification of the data contained in the tables: insertion, deletion, selection, changing (even aggregation i.e count, sum, average)
- Data Control Language (DCL)
  - This deals with maintaining the security of the database using permissions and access control
- Transaction Control Language (TCL)
  - This deals with maintaining the integrity of the database using permissions, transactions

#### CREATE

- CREATE DATABASE my\_database
- CREATE SCHEMA my\_schema
- CREATE TABLE my\_table

#### **OMOP-CDM DB schemas**

- raw: datos originales
- cdm: tablas clínicas
- vocabularies: vocabularios
- nlp: resultados de nlp
- results: cohortes, calidad, etc
- webapi: Atlas
- tmp: tablas temporales (Oracle)

#### CREATE TABLE

```
CREATE TABLE [IF NOT EXISTS] [schema_name].table_name (
column_1 data_type PRIMARY KEY,
column_2 data_type NOT NULL UNIQUE,
column_3 data_type DEFAULT 0,
table_constraints
)
```

# Data types

- NULL. The value is a NULL value.
- INTEGER. The value is a signed integer, stored in 1, 2, 3, 4, 6, or 8 bytes depending on the magnitude of the value.
- REAL. The value is a floating point value, stored as an 8byte IEEE floating point number.
- TEXT. The value is a text string, stored using the database encoding (UTF-8, UTF-16BE or UTF-16LE).
- BLOB. The value is a blob of data, stored exactly as it was input.

#### **FOREIGN KEY**

FOREIGN KEY (foreign\_key\_columns)
 REFERENCES parent\_table(parent\_key\_columns)
 ON UPDATE action
 ON DELETE action;

 ALTER TABLE child ADD CONSTRAINT fk\_child\_parent FOREIGN KEY (parent\_id)
 REFERENCES parent(id);

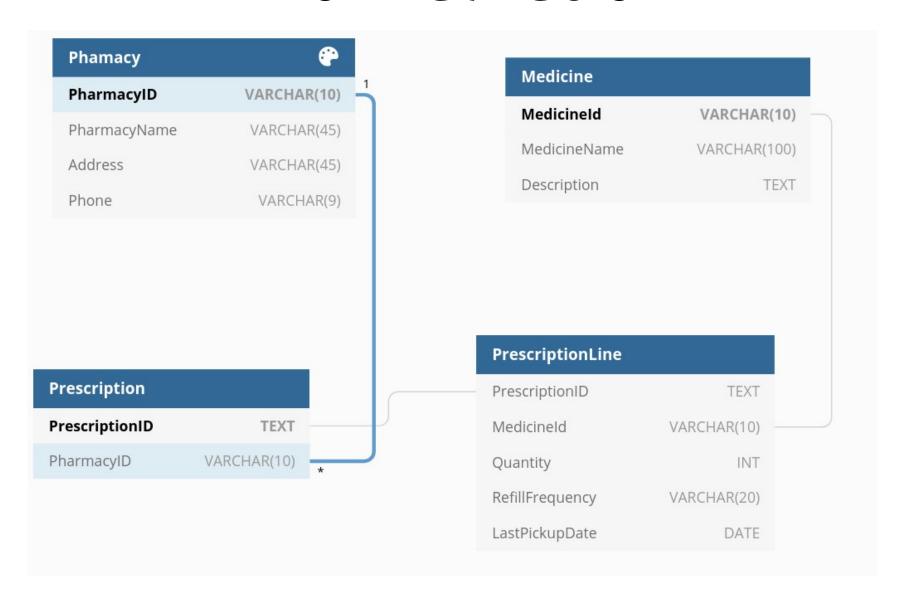
#### **DROP**

DROP TABLE table\_name

ALTER TABLE table\_name
 DROP COLUMN column\_name;

#### PrescriptionFilling

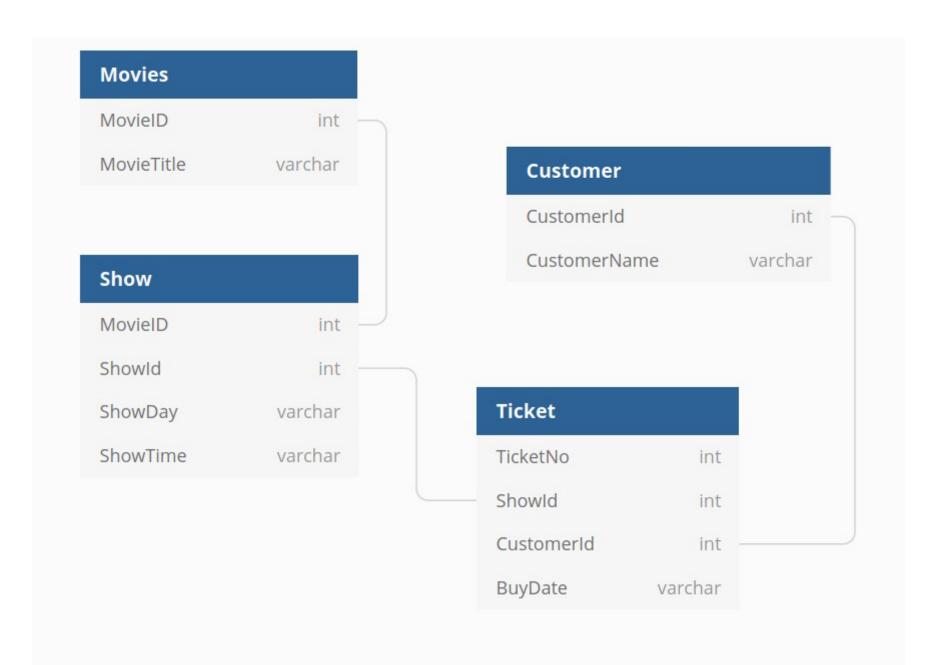
PharmacyID	PharmacyName	Address	Phone	PrescriptionID	MedicineID	MedicineName	Description	Quantity	LastPickupDate	RefillFrequency
PH1	PA	PAA	P111	PR1	M1	MA	MAAA	30	1/1/2013	Monthly
PH1	PA	PAA	P111	PR1	M2	МВ	MBBB	20	1/3/2013	Quarterly
PH1	PA	PAA	P111	PR2	M3	МС	мссс	50	1/1/2013	Monthly
PH2	РВ	PBB	P222	PR3	M1	MA	MAAA	60	1/30/2013	Semi-annually



https://dbdiagram.io/d/5f33ba32e1246d54aa2d0f95

#### MovieTicket

MovieID	MovieTitle	ShowID	ShowDay	ShowTime	CustomerID	CustomerName	<u>TicketNo</u>	BuyDate
1	AAA	1	1/1/2013	12 PM	1	X	1	12/31/2012
1	AAA	1	1/1/2013	12 PM	1	X	2	12/31/2012
1	AAA	1	1/1/2013	12 PM	2	Y	3	1/1/2013
1	AAA	1	1/1/2013	12 PM	2	Y	4	1/1/2013
1	AAA	2	1/2/2013	2 PM	3	Z	1	1/1/2013
1	AAA	2	1/2/2013	2 PM	3	Z	2	1/2/2013
2	BBB	1	1/3/2013	6 PM	4	W	1	1/2/2013
2	BBB	1	1/3/2013	6 PM	4	W	2	1/3/2013



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  - Declares the fields to be returned by the query
  - '\*' returns all fields
- FROM clause
  - Defines the table being queried
  - Can contain JOIN clauses to query multiple tables

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  - Sorts the records that are returned

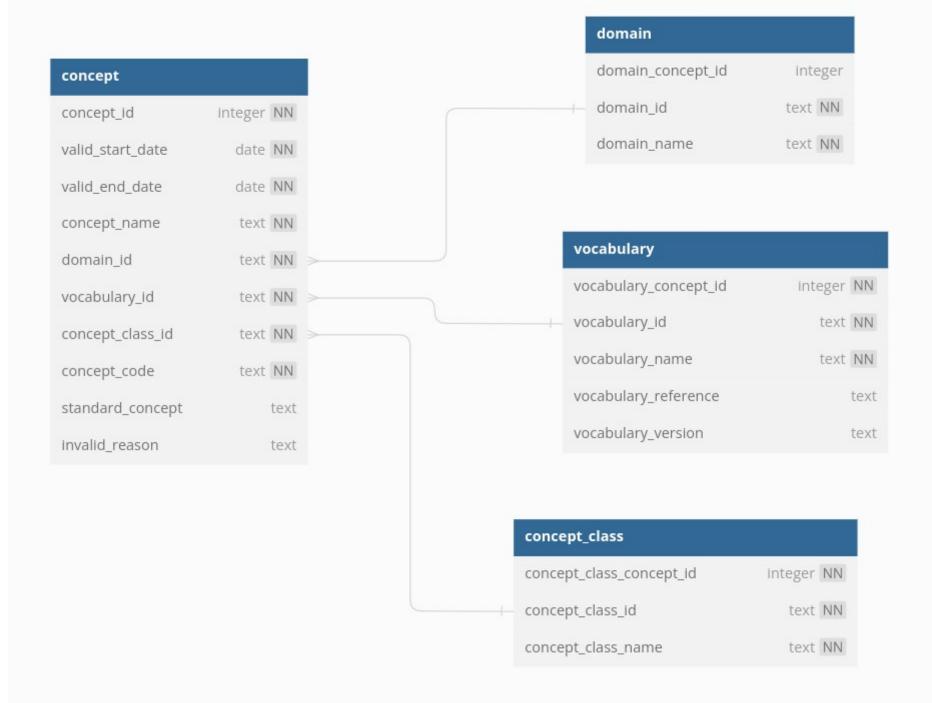
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## JOIN

- SELECT
- FROM table1 t1
- INNER JOIN table 2 t2
- ON t1.pk = t2.fk



#### **INSERT INTO**

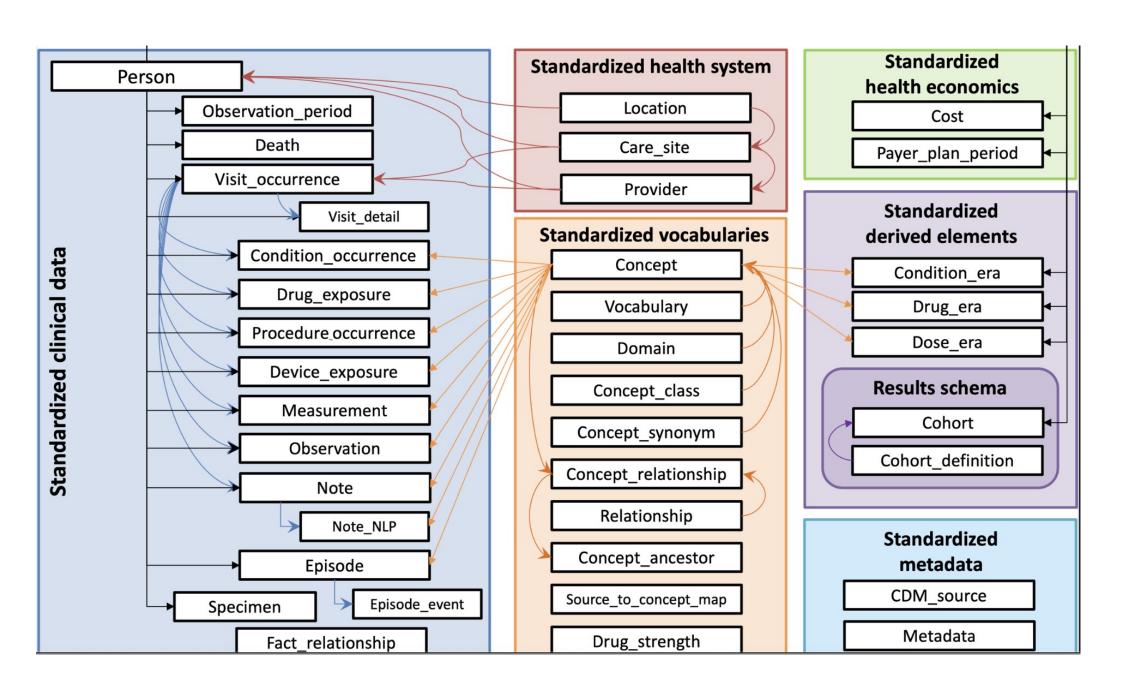
```
INSERT INTO table_2
SELECT field_1, field_2
FROM table1
[WHERE conditions];
```

#### UPDATE

```
UPDATE table
SET column_1 = new_value_1,
  column_2 = new_value_2
WHERE
 search condition
```

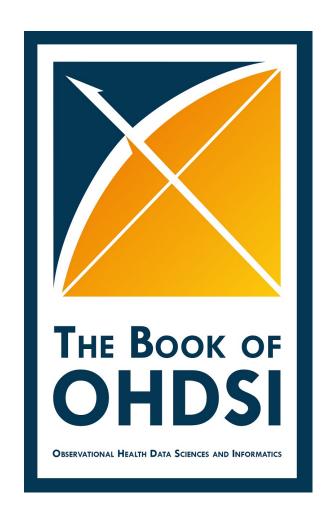
## DELETE

DELETE FROM table\_name
WHERE condition



#### Recursos

- Sitio web: https://github.com/alabarga/omop-cdm-course
- Diapositivas
- El Libro de OHDSI
- EHDEN academy
- OHDSI workshops
- Atlas tutorials
- Entorno local (docker)

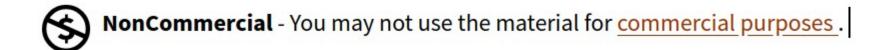


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