

# Test 1

1. a.  $L \Rightarrow R$   
 $[(C \Rightarrow A) \Rightarrow (A \wedge B \wedge \neg C \wedge D)] \Rightarrow [B \wedge D \wedge ((A \Rightarrow C) \Rightarrow \neg(C \Rightarrow A))]$

Assume  $L$  is True.

Case 1:  $C \Rightarrow A$  is False

Case 2:  $C \Rightarrow A$  is True and  $A \wedge B \wedge \neg C \wedge D$  is True

Case 1

$C$  must be True and  $A$  must be False so

$A \wedge B \wedge \neg C \wedge D$  is False because  $A$  is False.

On the right,  $((A \Rightarrow C) \Rightarrow \neg(C \Rightarrow A)) \Leftrightarrow (\text{True} \Rightarrow \text{True})$ .

$R$  becomes  $B \wedge D \wedge \text{True}$ . If either  $B$  or  $D$  are False, then  $L \Rightarrow R$  is False  $\blacksquare$

b.  $R \Rightarrow L$   
 $[B \wedge D \wedge ((A \Rightarrow C) \Rightarrow \neg(C \Rightarrow A))] \Rightarrow [(C \Rightarrow A) \Rightarrow (A \wedge B \wedge \neg C \wedge D)]$

Assume that  $R$  is True. This means the  $B$  and  $D$  are True

and that  $A$  and  $C$  are not equal ( $A$  True,  $C$  False or

$A$  False,  $C$  True). Then,  $L$  is False when  $C \Rightarrow A$  is

True and  $A \wedge B \wedge \neg C \wedge D$  is False. However, in order for

$C \Rightarrow A$  to be True,  $C$  must be False and  $A$  must be True,

so  $A \wedge B \wedge \neg C \wedge D$  is True. Therefore,  $R \Rightarrow L$  is True,  $\blacksquare$

2. L: There is a unique element in  $S$  with property P.  
 $R: \exists(x \in S): (P(x) \wedge \forall(y \in S): P(y) \Rightarrow x=y))$

$$L \Leftrightarrow R$$

Prove  $L \Rightarrow R$ :

First, assume  $L$  is True.

Let  $x$  be a unique element in  $S$  such that  $P(x)$ .

Since  $x$  is unique, there is only one of  $x$  in  $S$ .

Assume  $P(y)$  is True. Since there is only one element with property  $P$ ,  $y$  is just an alias of  $x$  so  $y=x$  and  $R$  has to be True.  $\blacksquare$

Prove  $R \Rightarrow L$ :

Assume  $\neg L$  so there isn't a unique element in  $S$  with property  $P$ . There are two cases here: either there doesn't exist an element with property  $P$  or there exists more than one element with property  $P$ . This can be written as

$\forall (x \in S) : (\neg P(x)) \vee \exists (y \in S) : P(y) \wedge x \neq y$ , which is just  $\neg R$  so  $\neg R$  is True.  $\blacksquare$

3. Property:  $\forall (B \in \mathbb{R}) \exists (c \in \mathbb{R}) \forall (x \geq c) : f(x) \geq B$

If  $f(x) = -x^2$ , it doesn't satisfy this property.

Because  $f(x) = -x^2$ ,  $f(x)$  is always  $\leq 0$ . If a  $B$

is chosen that is greater than 0, there is no  $x$  that makes  $f(x) \geq B$ .  $\blacksquare$

4.  $\mathbb{Z}_n = \{0, 1, \dots, n-1\}$   $f_b(x) = bx + 1 \pmod{n}$

$f_b$  is bijective  $\Leftrightarrow \gcd(b, n) = 1$

Prove:  $f_b$  is bijective  $\Rightarrow \gcd(b, n) = 1$

Assume  $\gcd(b, n) \neq 1$ , so  $\gcd(b, n) > 1$ .

Let  $\gcd(b, n)$  be  $k$ . Let  $b = ki$  and

$n = kj$ .  $f_b(x) = kix + 1 \pmod{kj}$ . Since

$b$  and  $n$  are not relatively prime, there exists  $x_1$  and  $x_2$  where  $x_1 \neq x_2$  and

$f_b(x_1) = f_b(x_2)$ . This means that

$f_b$  is not injective.



Prove:  $\gcd(b, n) = 1 \Rightarrow f_b$  is bijective

Assume that  $f_b$  is not bijective, so  $f_b$  is either not surjective or not injective.

When  $f_b$  is not injective, this means there exists  $x_1$  and  $x_2$  where  $x_1 \neq x_2$  such that  $f(x_1) = f(x_2)$ .

$f_b(x_1) = bx_1 + 1 \pmod{n} = f_b(x_2) = bx_2 + 1 \pmod{n}$ .

$bx_1 + 1 \equiv bx_2 + 1 \pmod{n}$

$(bx_1 + 1) - (bx_2 + 1) = kn$  for some  $k \in \mathbb{Z} \Rightarrow b(x_1 - x_2) = kn$

This means there is some common factor between  $b$  and  $n$  greater than 1 so  $\gcd(b, n) \neq 1$ .



EC. The kingfisher travels in a combination of vertical and horizontal movements. Going diagonally compared to going linearly doesn't matter as the perimeter is equivalent.

Prove: move to  $(m, n)$  in  $l$  hops iff

$$l \geq |m| + |n| \text{ and } l \equiv m+n \pmod{2}$$

Assume it can move to  $(m, n)$  in  $l$  hops

and  $l < |m| + |n|$  or  $l \not\equiv m+n \pmod{2}$ ,

$l < |m| + |n|$  because  $|m| + |n|$  is its shortest Manhattan distance.

Base Case:  $(m, n)$  is  $(0, 0)$

$l$  can't be  $< 0$  because can't be negative hops  
so  $l < |m| + |n|$  is always False. Let  $l = 0$   
since that is the min. number of hops.

$0 \equiv (0+0) \pmod{2}$ . Contradiction.  $\square$

Assume  $l \geq |m| + |n|$  and  $l \equiv m+n \pmod{2}$ .

One way to get to  $(m, n)$  is to move  
 $m$  times in the  $X$  direction and move  
 $n$  times in the  $Y$  direction. Therefore,  
it can move to  $(m, n)$  in  $l$  hops.  $\square$