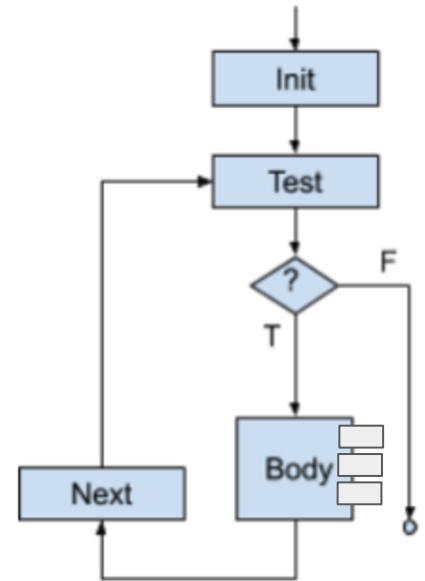


Control Flow, Call Graphs and Subroutine Construction

Control Flow Graph

- A graphic representation of the representation between basic blocks
- A basic block:
 - a list of instructions with
 - a single entry point (starting point)
 - a single exit point (last instruction)
- Such representations model the behavior of our code
- Recall the while loop, and other control structures
- What about subroutines calls
(subroutine: general term for ...
methods, functions, procedures, etc.)

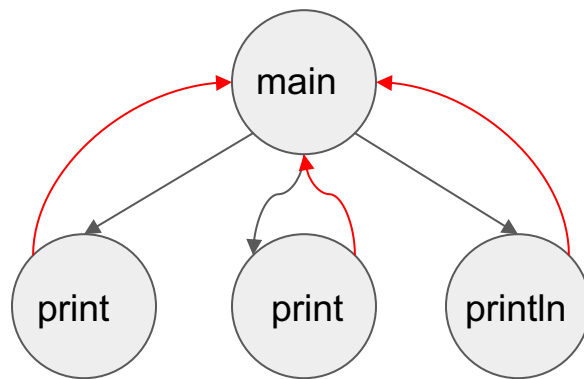




While Loop

Call Graph

- a control flow graph depicting the relationships between subroutines
- Call Graph for the "Hello World" program

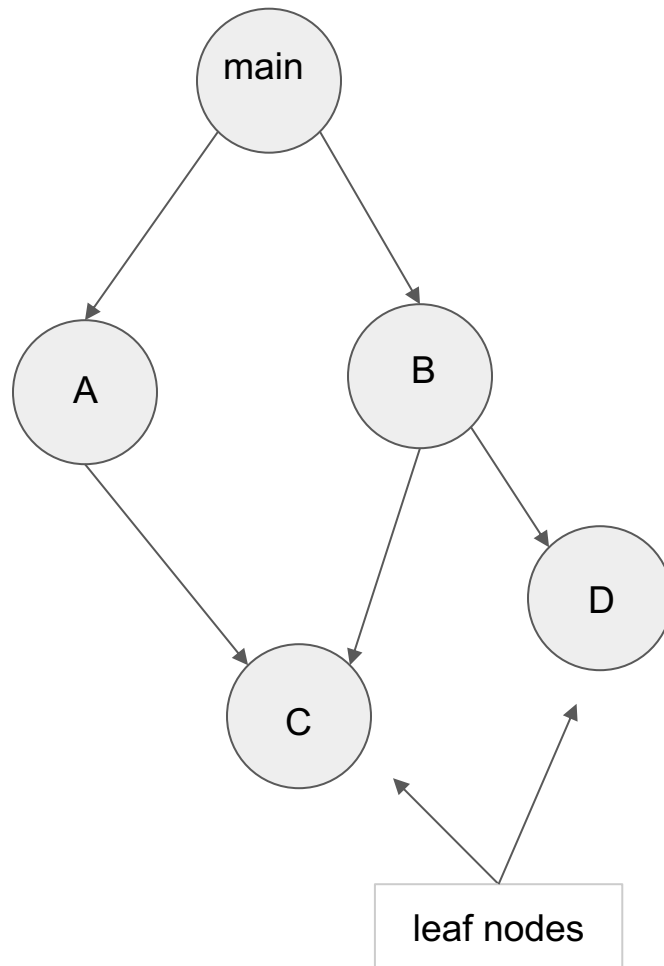
```
class HelloWorld
{
    public static void main(String args[])
    {
        System.out.print("Hello ");
        System.out.print("World");
        System.out.println("");
    }
}
```



call: 
return: 

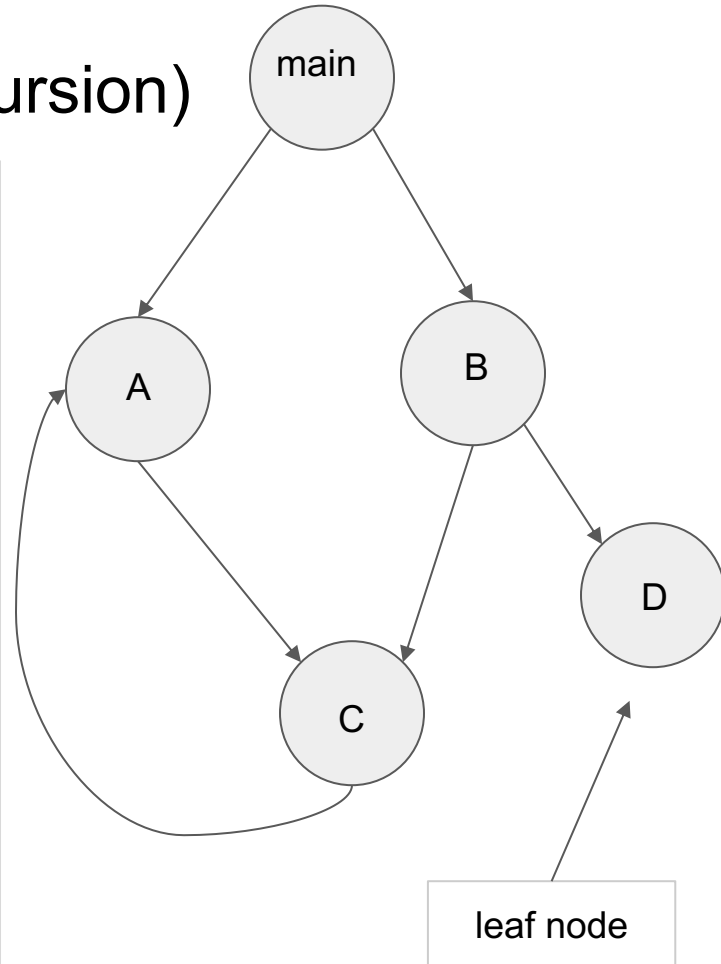
Call Graph II

```
public static void A(void) {  
    int x = 5;  
    C();  
}  
public static void B(void) {  
    C();  
    D();  
}  
public static void C(void) {  
    ;  
}  
public static void D(void) {  
    ;  
}  
  
public static void main(String args[])  
    {  
        A();  
        B();  
    }  
}
```



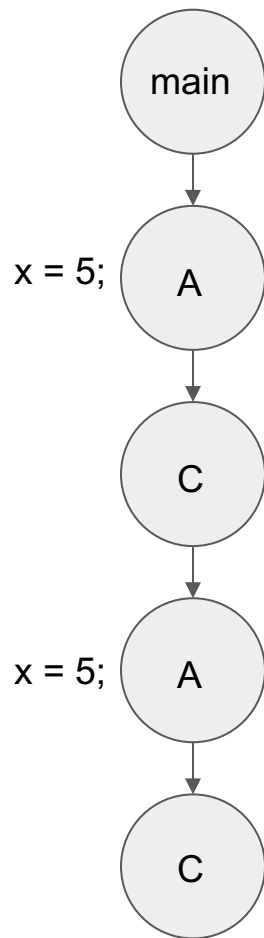
Call Graph with a Loop (Recursion)

```
public static void A(void) {  
    int x = 5;  
    C();  
}  
public static void B(void) {  
    C();  
    D();  
}  
public static void C(void) {  
    A();  
}  
public static void D(void) {  
    ;  
}  
  
public static void main(String args[])  
    {  
        A();  
        B();  
    }  
}
```



Dynamic Call Graph (Runtime)

```
public static void A(void) {  
    static int x = 5;  
    C();  
}  
public static void B(void) {  
    C();  
    D();  
}  
public static void C(void) {  
    A();  
}  
public static void D(void) {  
    ;  
}  
  
public static void main(String args[])  
    {  
        A();  
        B();  
    }  
}
```



Memory Organization (Java program)

```
class Main {  
  
    public static int x = 5;  
    int y = 7;  
  
    public int addNumbers(int a, int b) {  
        int sum = a + b;  
        return sum;  
    }  
  
    public static void main(String[] args) {  
        int num1 = 25;  
        int num2 = 15;  
  
        // create an object of Main  
        Main obj = new Main();  
        int result = obj.addNumbers(num1, num2);  
        System.out.println("Sum is: " + result);  
    }  
}
```

STACK
int a; int b;

HEAP

.data

.text
(INSTRUCTIONS)

Dynamic:

Locations
defined at
runtime.

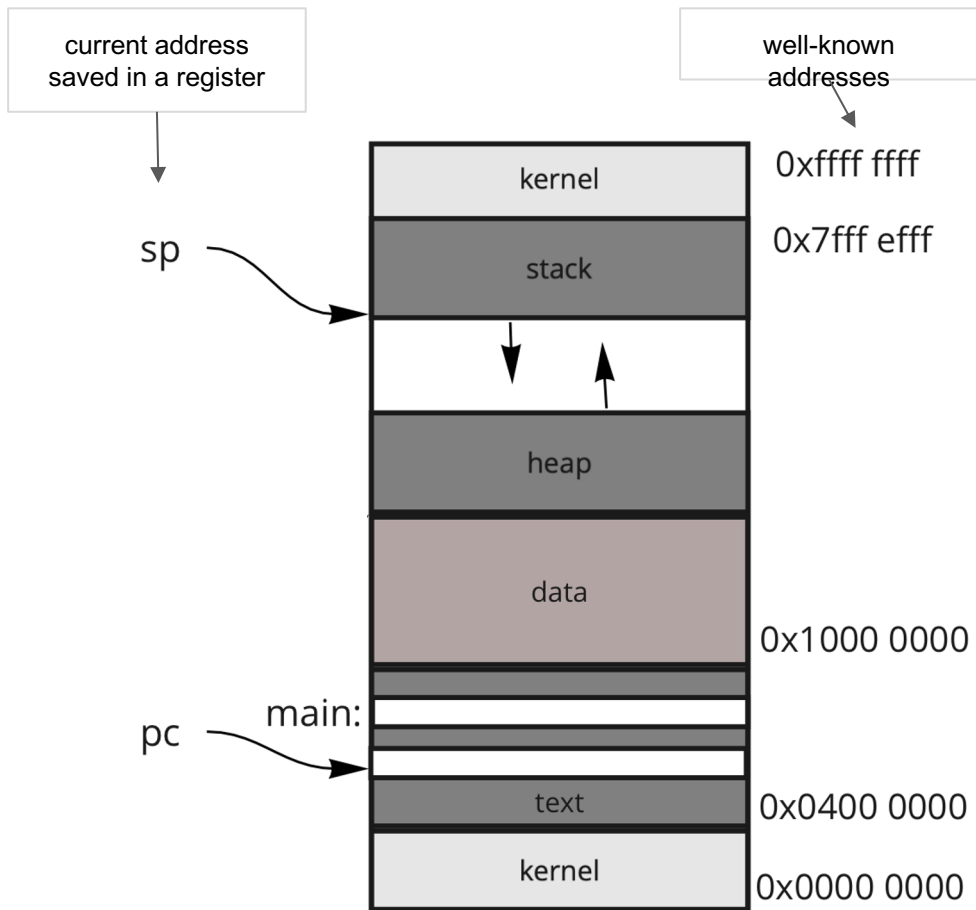
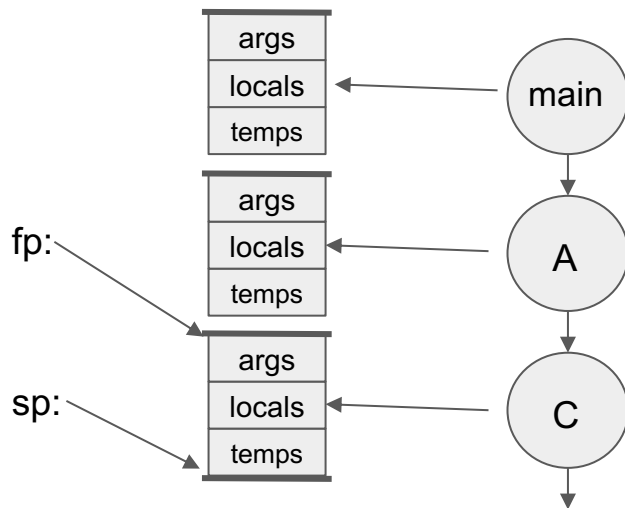
Static:

Locations are
defined when
the program
starts.

Frames

- Frame: a collection of variables:

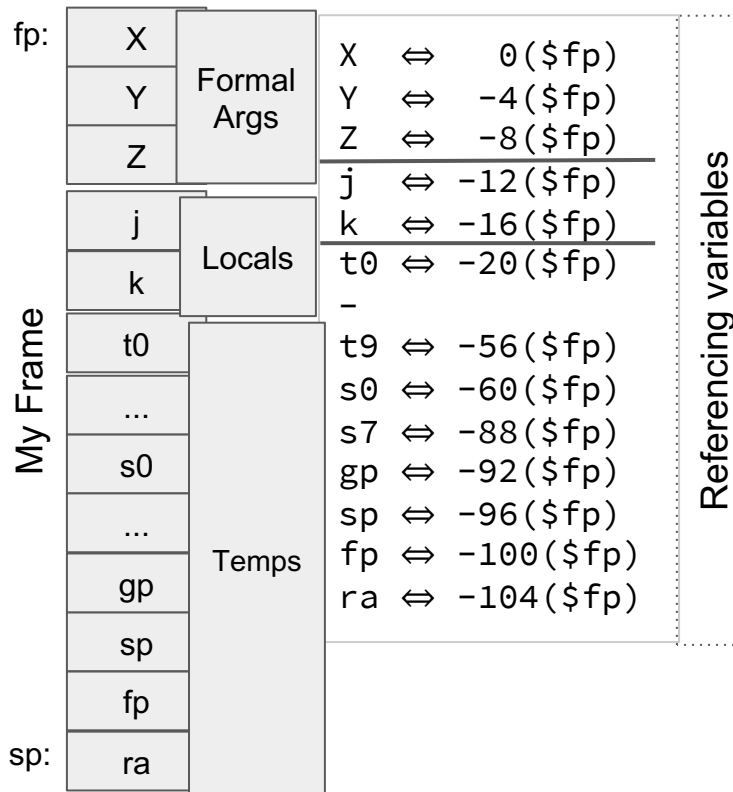
- Classes variables → "heap"
- Methods variables → "stack"
- Static variables → ".data"



Layout of the Frame

```
int my(int X, int Y, int Z) {  
    int j;  
    int k = Y + Z  
  
    j = sub(1, k, 3);  
    ;  
    return j;  
}
```

- When do we store values onto the frame?
 - in theory?
 - in practice?



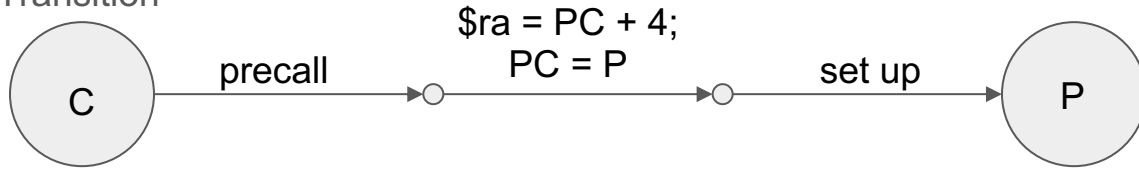
Subroutine Transition: Calling a Subroutine

1. The Client (C) needs to:

- Place actual args into the Frame
- Precall (preparation for the call)
- Transition

2. The Producer (P) needs to:

- Setup
- Do it's Thing



```
C:      nop
        ...
        nop      # precall
        jal P
        nop
        nop      # postcall
        ...
        jr $ra # return
```

```
P:      nop
        nop      # set up
        ...
        nop      # cleanup
        nop      # return value
        jr $ra
```

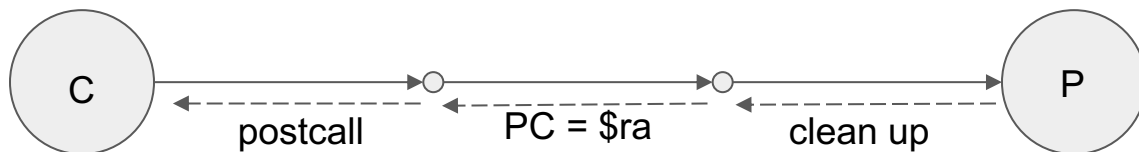
Subroutine Transition: Return from a Subroutine

2. The Client (C) needs to:

- Postcall
- Continue doing it's thing

1. The Producer (P) needs to:

- Clean up
- Position the return value
- Transition back



```
C:  nop
    ...
    nop    # precall
    jal P
    nop
    nop    # postcall
    ...
    jr $ra # return
```

```
P:  nop
    nop    # set up
    ...
    nop    # cleanup
    nop    # return value
    jr $ra
```

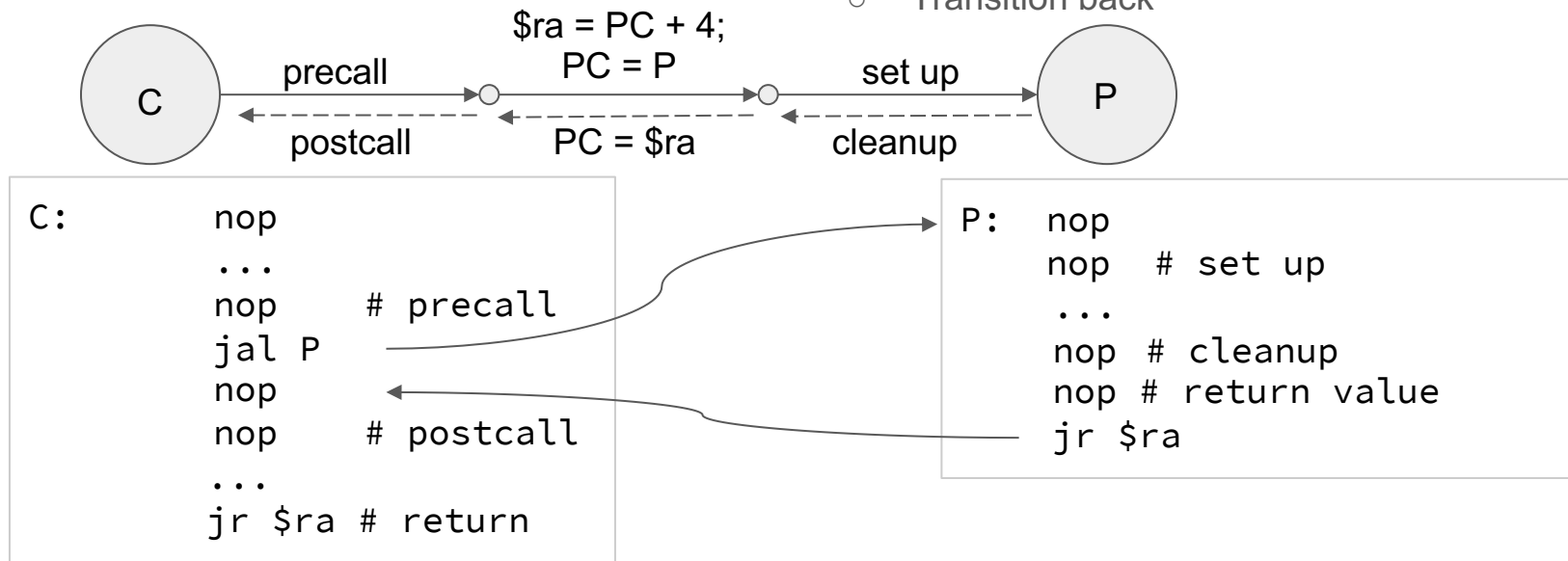
MIPS: Subroutine Process

2. The Client (C) needs to:

- Postcall ← Restore saved registers
- Do it's Thing

1. The Producer (P) needs to:

- Cleanup ← Restore S registers
- Position the return value
- Transition back



Subroutines

- Causes:

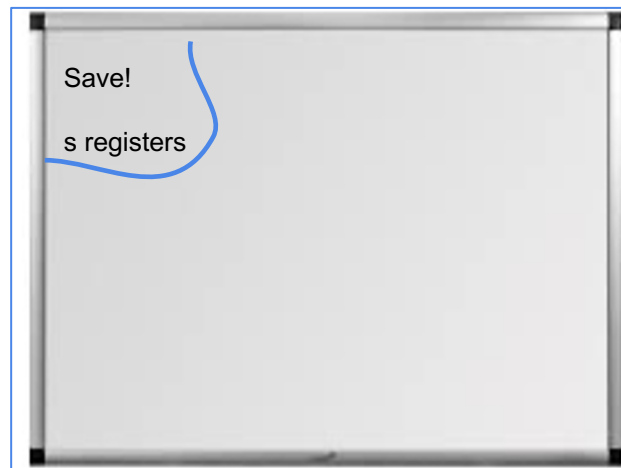
- a change in control-flow: `jal sub, jr $ra`
- a change in ownership of registers

- A Subroutine Calling Convention Exists

- pushing arguments onto the stack
 - MIPS Conventions (`$a0, $a1, $a2, $a3`) → `{ $v0, $v1 }`
- preserving registers (e.g., temps) onto the stack

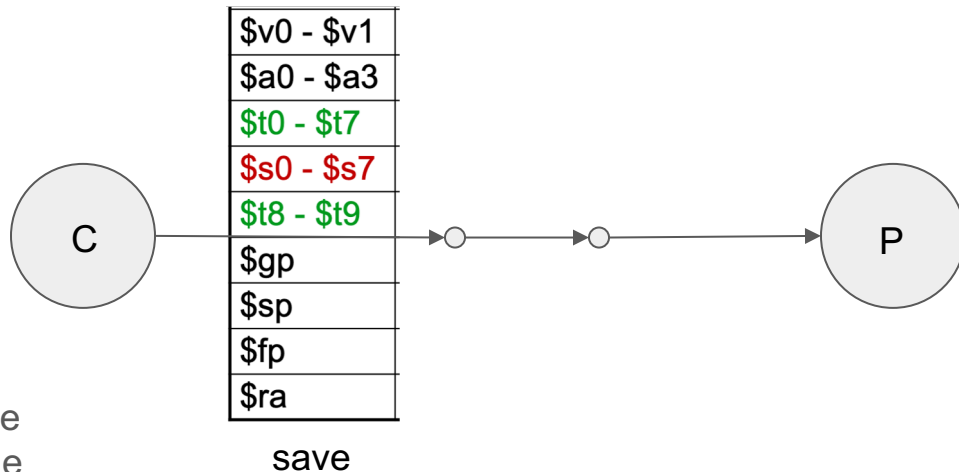
- Special cases (short circuit the MIPS Calling Convention)

- Main subroutine: the first subroutine in the dynamic call graph
 - No need to save the "s" registers upon entry
 - Give preference to "s" register utilization
- Leaf Subroutines: the last subroutine in the dynamic call graph
 - Give preference to "t" register utilization



Shared Resource: Registers

- You need to perform setup and cleanup routines for any shared resource!
- Precall:
 - Save what you need,
 - ~~Clear what you want private,~~
 - Leave alone what is passed along!
- Brute Force Approach:
 - ignore: \$zero, \$at, \$k1, \$k2
 - save all other registers
 - especially:
 - \$gp: might as well!
 - \$sp: this is the end of my frame
 - \$fp: this is the start of my frame
 - \$ra: this is my "return to" location

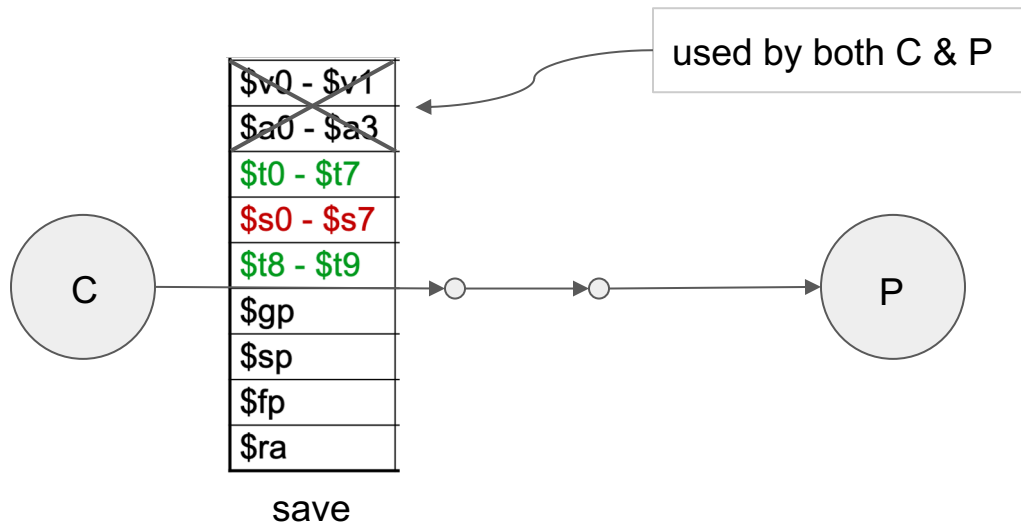


Shared Resource: Registers

- You need to perform setup and cleanup routines for any shared resource!

- Semi-Optimal

- ignore: \$zero, \$at, \$k1, \$k2
- save only registers in local use
- but always save:
 - \$gp: might as well!
 - \$sp: this is the end of my frame
 - \$fp: this is the start of my frame
 - \$ra: this is my "return to" location

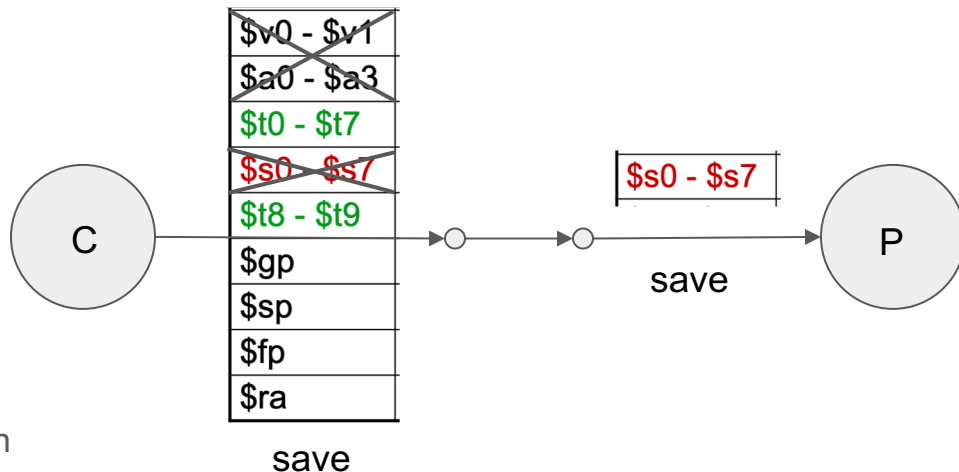


Shared Resource: Registers

- You need to perform setup and cleanup routines for any shared resource!

- MIPS Conventions

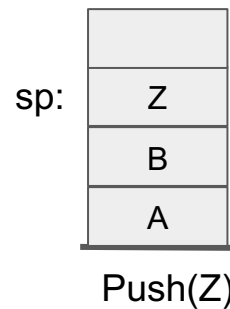
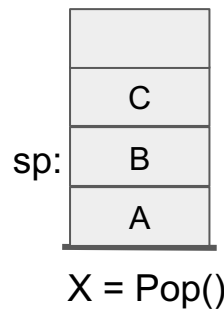
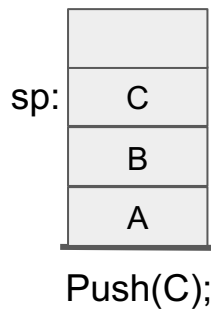
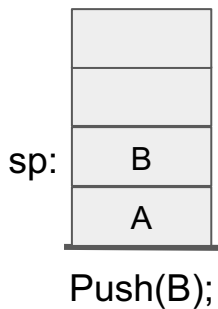
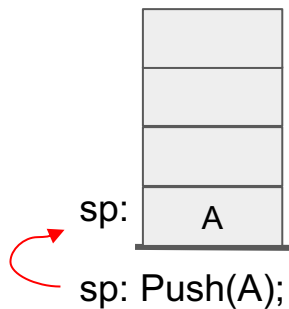
- Client (C): saves:
 - all T registers
 - \$gp: might as well!
 - \$sp: this is the end of my frame
 - \$fp: this is the start of my frame
 - \$ra: this is my "return to" location
- Provider (P): saves:
 - all S registers



Stack Operations

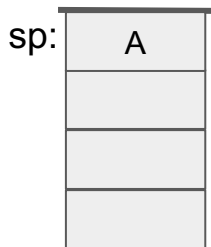
$\text{Push}(a) \Leftrightarrow$ $\text{sp} = \text{sp} + 1$ $\text{sp}[0] = a$	$x = \text{Pop}() \Leftrightarrow$ $x = \text{sp}[0]$ $\text{sp} = \text{sp} - 1$
---	---

- Stack is an abstract data structure
- The stack is an array of words
- Operations:
 - Push: $\text{Push}(A), \text{Push}(B), \text{Push}(C)$
 - Pop: $X = \text{Pop}();$
 - Push: $\text{Push}(Z);$

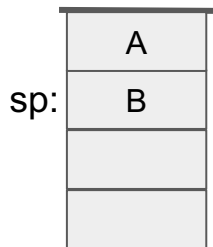


But the MIPS Way

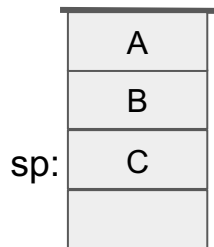
- Stack is an abstract data structure
- The stack is an array of words
- Operations:
 - Push: Push(A), Push(B), Push(C)
 - Pop: X = Pop();
- sp: points to the current top of stack



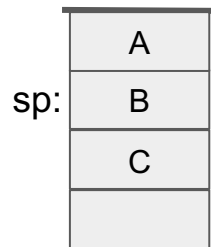
Push(A);



Push(B);



Push(C);



X = Pop();

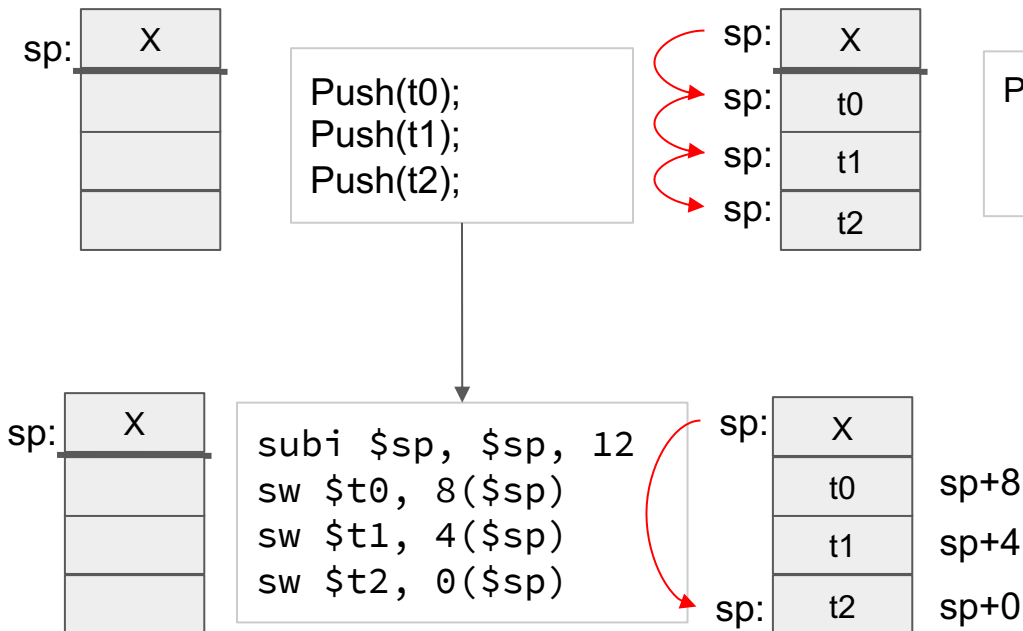
Push(a) \Leftrightarrow
sp = sp - 1
sp[0] = a

x = Pop() \Leftrightarrow
x = sp[0]
sp = sp + 1

Push(a) \Leftrightarrow
subi \$sp, \$sp, 4
sw \$a0, 0(\$sp)

x = Pop() \Leftrightarrow
lw \$v0, 0(\$sp)
addi \$sp, \$sp, 4

Multiple Pushes / Pops



Push(a) \Leftrightarrow
 $sp = sp - 1$
 $sp[0] = a$

$x = \text{Pop}() \Leftrightarrow$
 $x = sp[0]$
 $sp = sp + 1$

Push(a) \Leftrightarrow
`subi $sp, $sp, 4`
`sw $a0, 0($sp)`

$x = \text{Pop}() \Leftrightarrow$
`lw $v0, 0($sp)`
`addi $sp, $sp, 4`

`t0 = Pop();`
`t1 = Pop();`
`t2 = Pop();`

`lw $t0, 8($sp)`
`lw $t1, 4($sp)`
`lw $t2, 0($sp)`
`addi $sp, $sp, 12`

Frames in Detail

Client's Frame	fp:	X
		Y
		Z
		j
		k
		t0
		...
		s0
		...
		gp
		sp
	fp	fp
sp:		ra

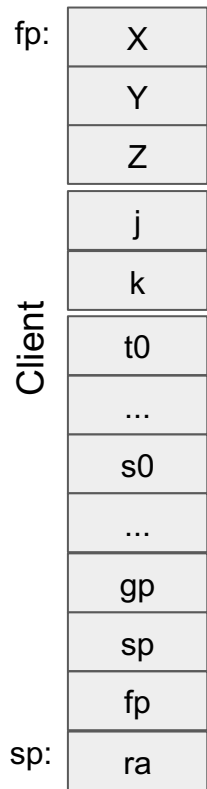
- Precall steps before "sub"
 - push args
 - save registers
 - jal sub # jump and link
- Steps to set up
 - build the frame
 - fp = sp + arg_size
 - sp = fp - frame_size
 - save S registers
- Steps to clean up
 - restore S registers
 - delete the frame (no need to!)
 - but sp = fp + 1
 - position the return value: (\$sp), \$v0
 - jr \$ra # jump register
- Postcall steps after "sub"
 - restore registers
 - move return value?
 - -4(\$sp), (\$fp), or \$v0

```
int sub(int X, int Y, int Z) {
    int j;
    int k = Y + Z;

    j = sub(1, k, 3);
    ;
    return j;
}
```

X	⇔	0(\$fp)
Y	⇔	-4(\$fp)
Z	⇔	-8(\$fp)
j	⇔	-12(\$fp)
k	⇔	-16(\$fp)
t0	⇔	-20(\$fp)
-		
t9	⇔	-56(\$fp)
s0	⇔	-60(\$fp)
s7	⇔	-88(\$fp)
gp	⇔	-92(\$fp)
sp	⇔	-96(\$fp)
fp	⇔	-100(\$fp)
ra	⇔	-104(\$fp)

Calling "sub"



➡ Precall steps before "sub"

- push args
- save registers
- jal sub # jump and link

● Steps to set up

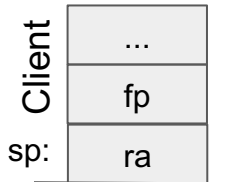
- build the frame
 - $fp = sp + arg_size$
 - $sp = fp - frame_size$
- save S registers

● Steps to clean up

- restore S registers
- delete the frame (no need to!)
 - but $sp = fp + 1$
- position the return value: (\$sp), \$v0
- jr \$ra # jump register

● Postcall steps after "sub"

- restore registers
- move return value?
 - $-4(\$sp)$, (\$fp), or \$v0



```
int sub(int X, int Y, int Z) {
    int j;
    int k = Y + Z

    j = sub(1, k, 3);
    ;
    return j;
}
```

X ⇔ 0(\$fp)

Y ⇔ -4(\$fp)

Z ⇔ -8(\$fp)

j ⇔ -12(\$fp)

k ⇔ -16(\$fp)

t0 ⇔ -20(\$fp)

...

t9 ⇔ -56(\$fp)

s0 ⇔ -60(\$fp)

s7 ⇔ -88(\$fp)

gp ⇔ -92(\$fp)

sp ⇔ -96(\$fp)

fp ⇔ -100(\$fp)

ra ⇔ -104(\$fp)

Calling "sub"

fp:	X
	Y
	Z
Client	j
	k
	t0
	...
	s0
	...
	gp
	sp
	fp
	ra

- Precall steps before "sub"

- push args
- save registers
- jal sub # jump and link

- Steps to set up

- build the frame
 - $fp = sp + arg_size$
 - $sp = fp - frame_size$
- save S registers

- Steps to clean up

- restore S registers
- delete the frame (no need to!)
 - but $sp = fp + 1$
- position the return value: (\$sp), \$v0
- jr \$ra # jump register

- Postcall steps after "sub"

- restore registers
- move return value?
 - $-4(\$sp)$, (\$fp), or \$v0

Client	...
	fp
	ra
sp:	1
	k
sp:	3
	args

```
int sub(int X, int Y, int Z) {
    int j;
    int k = Y + Z

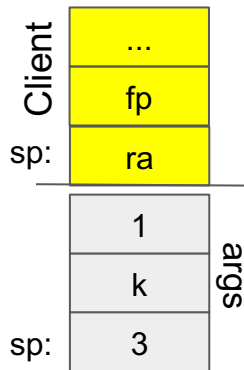
    j = sub(1, k, 3);
    ;
    return j;
}
```

```
X ⇔ 0($fp)
Y ⇔ -4($fp)
Z ⇔ -8($fp)
j ⇔ -12($fp)
k ⇔ -16($fp)
t0 ⇔ -20($fp)
...
t9 ⇔ -56($fp)
s0 ⇔ -60($fp)
s7 ⇔ -88($fp)
gp ⇔ -92($fp)
sp ⇔ -96($fp)
fp ⇔ -100($fp)
ra ⇔ -104($fp)
```

Calling "sub"



- Precall steps before "sub"
 - push args
 - save registers
 - jal sub # jump and link
- Steps to set up
 - build the frame
 - $fp = sp + arg_size$
 - $sp = fp - frame_size$
 - save S registers
- Steps to clean up
 - restore S registers
 - delete the frame (no need to!)
 - but $sp = fp + 1$
 - position the return value: (\$sp), \$v0
 - jr \$ra # jump register
- Postcall steps after "sub"
 - restore registers
 - move return value?
 - -4(\$sp), (\$fp), or \$v0



```
int sub(int X, int Y, int Z) {
    int j;
    int k = Y + Z

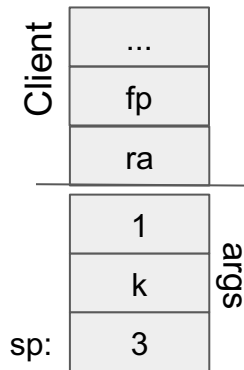
    j = sub(1, k, 3);
    ;
    return j;
}
```

X ⇔ 0(\$fp)
Y ⇔ -4(\$fp)
Z ⇔ -8(\$fp)
j ⇔ -12(\$fp)
k ⇔ -16(\$fp)
t0 ⇔ -20(\$fp)
...
t9 ⇔ -56(\$fp)
s0 ⇔ -60(\$fp)
s7 ⇔ -88(\$fp)
gp ⇔ -92(\$fp)
sp ⇔ -96(\$fp)
fp ⇔ -100(\$fp)
ra ⇔ -104(\$fp)

Transition to "sub"



- Precall steps before "sub"
 - push args
 - save registers
 - jal sub # jump and link
- Steps to set up
 - build the frame
 - $fp = sp + arg_size$
 - $sp = fp - frame_size$
 - save S registers
- Steps to clean up
 - restore S registers
 - delete the frame (no need to!)
 - but $sp = fp + 1$
 - position the return value: (\$sp), \$v0
 - jr \$ra # jump register
- Postcall steps after "sub"
 - restore registers
 - move return value?
 - $-4(\$sp)$, (\$fp), or \$v0



```
int sub(int X, int Y, int Z) {
    int j;
    int k = Y + Z

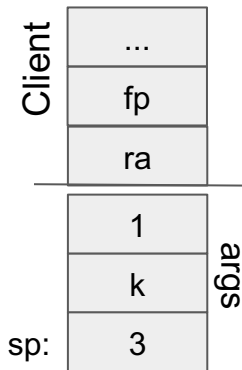
    j = sub(1, k, 3);
    ;
    return j;
}
```

X ⇔ 0(\$fp)
 Y ⇔ -4(\$fp)
 Z ⇔ -8(\$fp)
 j ⇔ -12(\$fp)
 k ⇔ -16(\$fp)
 t0 ⇔ -20(\$fp)
 ...
 t9 ⇔ -56(\$fp)
 s0 ⇔ -60(\$fp)
 s7 ⇔ -88(\$fp)
 gp ⇔ -92(\$fp)
 sp ⇔ -96(\$fp)
 fp ⇔ -100(\$fp)
 ra ⇔ -104(\$fp)

Producer: The set up



- Precall steps before "sub"
 - push args
 - save registers
 - jal sub # jump and link
- ➡ Steps to set up
 - build the frame
 - $fp = sp + arg_size$
 - $sp = fp - frame_size$
 - save S registers
- Steps to clean up
 - restore S registers
 - delete the frame (no need to!)
 - but $sp = fp + 1$
 - position the return value: (\$sp), \$v0
 - jr \$ra # jump register
- Postcall steps after "sub"
 - restore registers
 - move return value?
 - $-4(\$sp)$, (\$fp), or \$v0



```
int sub(int X, int Y, int Z) {
    ➡ int j;
      int k = Y + Z

      j = sub(1, k, 3);
      ;
      return j;
}
```

```
X  ⇔  0($fp)
Y  ⇔  -4($fp)
Z  ⇔  -8($fp)
j  ⇔  -12($fp)
k  ⇔  -16($fp)
t0 ⇔  -20($fp)
...
t9 ⇔  -56($fp)
s0 ⇔  -60($fp)
s7 ⇔  -88($fp)
gp ⇔  -92($fp)
sp ⇔  -96($fp)
fp ⇔  -100($fp)
ra ⇔  -104($fp)
```

Client

X
Y
Z
j
k
t0
...
s0
...
gp
sp
fp
ra

- Precall steps before "sub"
 - push args
 - save registers
 - jal sub # jump and link
- Steps to set up
 - ➡ ○ build the frame
 - $fp = sp + arg_size$
 - $sp = fp - frame_size$
 - save S registers
- Steps to clean up
 - restore S registers
 - delete the frame (no need to!)
 - but $sp = fp + 1$
 - position the return value: ($\$sp$), $\$v0$
 - jr $\$ra$ # jump register
- Postcall steps after "sub"
 - restore registers
 - move return value?
 - $-4(\$sp)$, ($\fp), or $\$v0$

The diagram illustrates a memory layout with two main sections: Client and Producer.

- Client Section:** Contains three pointers: `ra`, `fp`, and `...`.
- Producer Section:** Contains three main blocks:
 - args:** A block containing three values: `1`, `k`, and `3`. A red arrow points from the `fp` pointer in the Client section to the `k` value in this block.
 - locals:** A block containing three empty slots.
 - temps:** A block containing eight slots. The first slot is labeled `0`, and the last slot is labeled `sp:`.

```
int sub(int X, int Y, int Z) {  
    int j;  
    int k = Y + Z  
  
    j = sub(1, k, 3);  
    ;  
    return j;  
}
```

```
X  ⇔  0($fp)
Y  ⇔  -4($fp)
Z  ⇔  -8($fp)
j  ⇔  -12($fp)
k  ⇔  -16($fp)
t0 ⇔  -20($fp)
...
t9 ⇔  -56($fp)
s0 ⇔  -60($fp)
s7 ⇔  -88($fp)
gp ⇔  -92($fp)
sp ⇔  -96($fp)
fp ⇔  -100($fp)
ra ⇔  -104($fp)
```

Producer: The set up

Client	X
	Y
	Z
	j
	k
	t0
	...
	s0
	...
	gp
	sp
	fp
	ra

- Precall steps before "sub"
 - push args
 - save registers
 - jal sub # jump and link
- Steps to set up
 - build the frame
 - $fp = sp + arg_size$
 - $sp = fp - frame_size$
 - save S registers
- Steps to clean up
 - restore S registers
 - delete the frame (no need to!)
 - but $sp = fp + 1$
 - position the return value: (\$sp), \$v0
 - jr \$ra # jump register
- Postcall steps after "sub"
 - restore registers
 - move return value?
 - $-4(\$sp)$, $(\$fp)$, or $\$v0$

Client	...
	fp
	ra
<hr/>	
fp:	1
	k
	3
Producer	
sp:	

args

locals

temps

```
int sub(int X, int Y, int Z) {
    → int j;
      int k = Y + Z

      j = sub(1, k, 3);
      ;
      return j;
}
```

```
X ⇔ 0($fp)
Y ⇔ -4($fp)
Z ⇔ -8($fp)
j ⇔ -12($fp)
k ⇔ -16($fp)
t0 ⇔ -20($fp)
...
t9 ⇔ -56($fp)
s0 ⇔ -60($fp)
s7 ⇔ -88($fp)
gp ⇔ -92($fp)
sp ⇔ -96($fp)
fp ⇔ -100($fp)
ra ⇔ -104($fp)
```

Executing "sub"

Client

X
Y
Z
j
k
t0
...
s0
...
gp
sp
fp
ra

- Precall steps before "sub"
 - push args
 - save registers
 - jal sub # jump and link
- Steps to set up
 - build the frame
 - $fp = sp + arg_size$
 - $sp = fp - frame_size$
 - save S registers
- Steps to clean up
 - restore S registers
 - delete the frame (no need to!)
 - but $sp = fp + 1$
 - position the return value: (\$sp), \$v0
 - jr \$ra # jump register
- Postcall steps after "sub"
 - restore registers
 - move return value?
 - $-4(\$sp)$, $(\$fp)$, or $\$v0$

Client

...
fp
ra

fp:

1
k
3

args

Producer

j
k

locals

t0
...

s0
...

temps

gp
sp
fp

sp:

ra



```
int sub(int X, int Y, int Z) {
    int j;
    int k = Y + Z
    {
        j = sub(1, k, 3);
    }
    return j;
}
```

X \Leftrightarrow 0(\$fp)
 Y \Leftrightarrow -4(\$fp)
 Z \Leftrightarrow -8(\$fp)
 j \Leftrightarrow -12(\$fp)
 k \Leftrightarrow -16(\$fp)
 t0 \Leftrightarrow -20(\$fp)
 ...
 t9 \Leftrightarrow -56(\$fp)
 s0 \Leftrightarrow -60(\$fp)
 s7 \Leftrightarrow -88(\$fp)
 gp \Leftrightarrow -92(\$fp)
 sp \Leftrightarrow -96(\$fp)
 fp \Leftrightarrow -100(\$fp)
 ra \Leftrightarrow -104(\$fp)

Returning from "sub"

Client

X
Y
Z
j
k
t0
...
s0
...
gp
sp
fp
ra

- Precall steps before "sub"
 - push args
 - save registers
 - jal sub # jump and link
- Steps to set up
 - build the frame
 - $fp = sp + arg_size$
 - $sp = fp - frame_size$
 - save S registers
- Steps to clean up
 - restore S registers
 - delete the frame (no need to!)
 - but $sp = fp + 1$
 - position the return value: (\$sp), \$v0
 - jr \$ra # jump register
- Postcall steps after "sub"
 - restore registers
 - move return value?
 - $-4(\$sp)$, (\$fp), or \$v0

Client	...		
	fp		
	ra		
<hr/>			
Producer	fp:	1	args
		k	
		3	
		j	locals
		k	
		t0	temps
		...	
		s0	
		...	
		gp	
		sp	
		fp	
	sp:	ra	

```
int sub(int X, int Y, int Z) {
    int j;
    int k = Y + Z

    j = sub(1, k, 3);
    ;
    ➔ return j;
}
```

X ⇔ 0(\$fp)
 Y ⇔ -4(\$fp)
 Z ⇔ -8(\$fp)
 j ⇔ -12(\$fp)
 k ⇔ -16(\$fp)
 t0 ⇔ -20(\$fp)
 ...
 t9 ⇔ -56(\$fp)
 s0 ⇔ -60(\$fp)
 s7 ⇔ -88(\$fp)
 gp ⇔ -92(\$fp)
 sp ⇔ -96(\$fp)
 fp ⇔ -100(\$fp)
 ra ⇔ -104(\$fp)

Returning from "sub"

Client	X
	Y
	Z
	j
	k
	t0
	...
	s0
	...
	gp
	sp
	fp
	ra

- Precall steps before "sub"
 - push args
 - save registers
 - jal sub # jump and link
- Steps to set up
 - build the frame
 - $fp = sp + arg_size$
 - $sp = fp - frame_size$
 - save S registers
- Steps to clean up
 - restore S registers
 - delete the frame (no need to!)
 - but $sp = fp + 1$
 - position the return value: (\$sp), \$v0
 - jr \$ra # jump register
- Postcall steps after "sub"
 - restore registers
 - move return value?
 - $-4(\$sp)$, $(\$fp)$, or $\$v0$

Client	...
	fp
	ra
<hr/>	
fp:	1
	k
	3
Producer	j
	k
	t0
	...
	s0
	...
	gp
	sp
	fp
	ra
sp:	ra

args

locals

temps

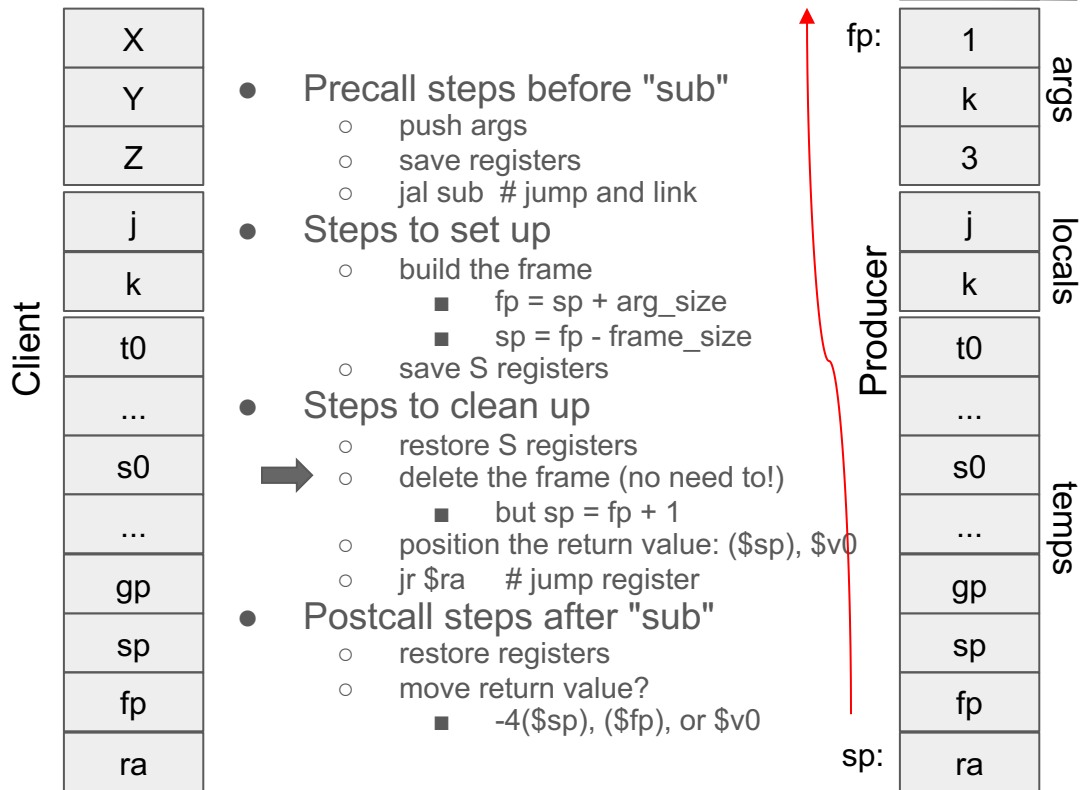
```
int sub(int X, int Y, int Z) {
    int j;
    int k = Y + Z

    j = sub(1, k, 3);
    ;
    ➔ return j;
}
```

X ⇔ 0(\$fp)
 Y ⇔ -4(\$fp)
 Z ⇔ -8(\$fp)
 j ⇔ -12(\$fp)
 k ⇔ -16(\$fp)
 t0 ⇔ -20(\$fp)
 ...
 t9 ⇔ -56(\$fp)
 s0 ⇔ -60(\$fp)
 s7 ⇔ -88(\$fp)
 gp ⇔ -92(\$fp)
 sp ⇔ -96(\$fp)
 fp ⇔ -100(\$fp)
 ra ⇔ -104(\$fp)

lw \$s0, -60(\$fp)
 lw \$s1, -64(\$fp)
 lw \$s2, -68(\$fp)
 ...
 lw \$s7, -88(\$fp)

Returning from "sub"



```
int sub(int X, int Y, int Z) {
    int j;
    int k = Y + Z

    j = sub(1, k, 3);
    ;
    return j;
}
```

X ⇔ 0(\$fp)
 Y ⇔ -4(\$fp)
 Z ⇔ -8(\$fp)
 j ⇔ -12(\$fp)
 k ⇔ -16(\$fp)
 t0 ⇔ -20(\$fp)
 ...
 t9 ⇔ -56(\$fp)
 s0 ⇔ -60(\$fp)
 s7 ⇔ -88(\$fp)
 gp ⇔ -92(\$fp)
 sp ⇔ -96(\$fp)
 fp ⇔ -100(\$fp)
 ra ⇔ -104(\$fp)

Client

X
Y
Z
j
k
t0
...
s0
...
gp
sp
fp
ra

- Precall steps before "sub"
 - push args
 - save registers
 - jal sub # jump and link
- Steps to set up
 - build the frame
 - fp = sp + arg_size
 - sp = fp - frame_size
 - save S registers
- Steps to clean up
 - restore S registers
 - delete the frame (no need to!)
 - but sp = fp + 1
 - ➡ ○ position the return value: (\$sp), \$v0
 - jr \$ra # jump register
- Postcall steps after "sub"
 - restore registers
 - move return value?
 - -4(\$sp), (\$fp), or \$v0

Client	...	
	fp	
	ra	
sp: fp:	j	args
	k	
	3	
Producer	j	locals
	k	
	t0	temps
	...	
	s0	
	...	
	gp	
	sp	
	fp	
	ra	

```
int sub(int X, int Y, int Z) {
    int j;
    int k = Y + Z

    j = sub(1, k, 3);
    ;
    ➡ return j;
}
```

```
X ⇔ 0($fp)
Y ⇔ -4($fp)
Z ⇔ -8($fp)
j ⇔ -12($fp)
k ⇔ -16($fp)
t0 ⇔ -20($fp)
...
t9 ⇔ -56($fp)
s0 ⇔ -60($fp)
s7 ⇔ -88($fp)
gp ⇔ -92($fp)
sp ⇔ -96($fp)
fp ⇔ -100($fp)
ra ⇔ -104($fp)
```


Transition back

Client	X
	Y
	Z
	j
	k
	t0
	...
	s0
	...
	gp
	sp
	fp
	ra

- Precall steps before "sub"
 - push args
 - save registers
 - jal sub # jump and link
- Steps to set up
 - build the frame
 - $fp = sp + arg_size$
 - $sp = fp - frame_size$
 - save S registers
- Steps to clean up
 - restore S registers
 - delete the frame (no need to!)
 - but $sp = fp + 1$
 - position the return value: (\$sp), \$v0
 - jr \$ra # jump register
- Postcall steps after "sub"
 - restore registers
 - move return value?
 - -4(\$sp), (\$fp), or \$v0

Client	...	
	fp	
	ra	
sp: fp:		args
Producer	i	
	k	
	3	
	j	locals
	k	
	t0	temps
	...	
	s0	
	...	
	gp	
	sp	
	fp	
	ra	

```
int sub(int X, int Y, int Z) {
    int j;
    int k = Y + Z

    j = sub(1, k, 3);
    ;
    → return j;
}
```

```
X ⇔ 0($fp)
Y ⇔ -4($fp)
Z ⇔ -8($fp)
j ⇔ -12($fp)
k ⇔ -16($fp)
t0 ⇔ -20($fp)
...
t9 ⇔ -56($fp)
s0 ⇔ -60($fp)
s7 ⇔ -88($fp)
gp ⇔ -92($fp)
sp ⇔ -96($fp)
fp ⇔ -100($fp)
ra ⇔ -104($fp)
```

Client: The Postcall

Client	X
	Y
	Z
	j
	k
	t0
	...
	s0
	...
	gp
	sp
	fp
	ra

- Precall steps before "sub"
 - push args
 - save registers
 - jal sub # jump and link
- Steps to set up
 - build the frame
 - $fp = sp + arg_size$
 - $sp = fp - frame_size$
 - save S registers
- Steps to clean up
 - restore S registers
 - delete the frame (no need to!)
 - but $sp = fp + 1$
 - position the return value: (\$sp), \$v0
 - jr \$ra # jump register



Postcall steps after "sub"

- restore registers
- move return value?
 - $-4(\$sp)$, $(\$fp)$, or $\$v0$

Client	...
	fp
	ra
<hr/>	
sp: fp:	j
	k
	Z
Producer	j
	k
	t0
	...
	s0
	...
	gp
	sp
	fp
	ra
sp:	

args

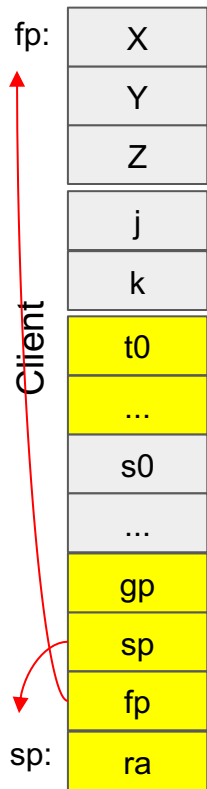
locals

temps

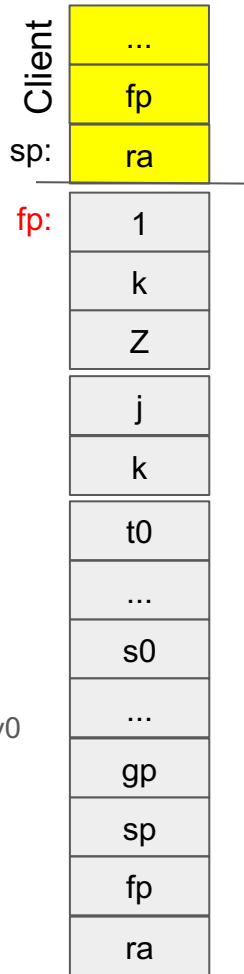
```
int sub(int X, int Y, int Z) {
    int j;
    int k = Y + Z
    j = sub(1, k, 3);
    return j;
}
```

X \Leftrightarrow 0(\$fp)
 Y \Leftrightarrow -4(\$fp)
 Z \Leftrightarrow -8(\$fp)
 j \Leftrightarrow -12(\$fp)
 k \Leftrightarrow -16(\$fp)
 t0 \Leftrightarrow -20(\$fp)
 ...
 t9 \Leftrightarrow -56(\$fp)
 s0 \Leftrightarrow -60(\$fp)
 s7 \Leftrightarrow -88(\$fp)
 gp \Leftrightarrow -92(\$fp)
 sp \Leftrightarrow -96(\$fp)
 fp \Leftrightarrow -100(\$fp)
 ra \Leftrightarrow -104(\$fp)

Client: The set up



- Precall steps before "sub"
 - push args
 - save registers
 - jal sub # jump and link
- Steps to set up
 - build the frame
 - $fp = sp + arg_size$
 - $sp = fp - frame_size$
 - save S registers
- Steps to clean up
 - restore S registers
 - delete the frame (no need to!)
 - but $sp = fp + 1$
 - position the return value: (\$fp), \$v0
 - jr \$ra # jump register
- Postcall steps after "sub"
 - restore registers
 - move return value?
 - $-4(\$sp)$, (\$fp), or \$v0



```
int sub(int X, int Y, int Z) {
    int j;
    int k = Y + Z
    j = sub(1, k, 3);
    ;
    return j;
}
```

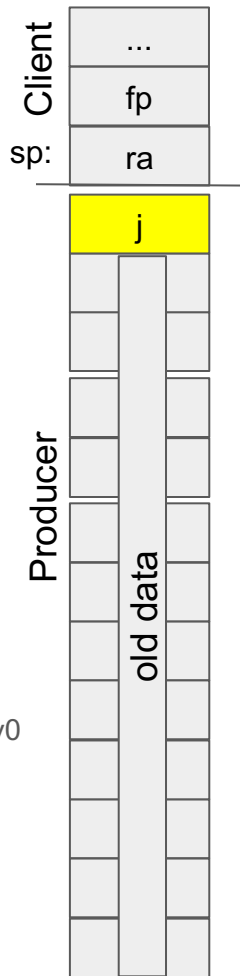
X \Leftrightarrow 0(\$fp)
 Y \Leftrightarrow -4(\$fp)
 Z \Leftrightarrow -8(\$fp)
 j \Leftrightarrow -12(\$fp)
 k \Leftrightarrow -16(\$fp)
 t0 \Leftrightarrow -20(\$fp)
 ...
 t9 \Leftrightarrow -56(\$fp)
 s0 \Leftrightarrow -60(\$fp)
 s7 \Leftrightarrow -88(\$fp)
 gp \Leftrightarrow -92(\$fp)
 sp \Leftrightarrow -96(\$fp)
 fp \Leftrightarrow -100(\$fp)
 ra \Leftrightarrow -104(\$fp)

```
lw $sp, 4($fp)
lw $fp, 4($sp)
-----
lw $t0, -20($fp)
lw $t1, -24($fp)
...
lw $ra, -104($fp)
```

Client: The Postcall



- Precall steps before "sub"
 - push args
 - save registers
 - jal sub # jump and link
- Steps to set up
 - build the frame
 - fp = sp + arg_size
 - sp = fp - frame_size
 - save S registers
- Steps to clean up
 - restore S registers
 - delete the frame (no need to!)
 - but sp = fp + 1
 - position the return value: (\$fp), \$v0
 - jr \$ra # jump register
- Postcall steps after "sub"
 - restore registers
 - move return value?
 - -4(\$sp), (\$fp), or \$v0



```
int sub(int X, int Y, int Z) {
    int j;
    int k = Y + Z
    j = sub(1, k, 3);
    ;
    return j;
}
```

```
X ⇔ 0($fp)
Y ⇔ -4($fp)
Z ⇔ -8($fp)
j ⇔ -12($fp)
k ⇔ -16($fp)
t0 ⇔ -20($fp)
...
t9 ⇔ -56($fp)
s0 ⇔ -60($fp)
s7 ⇔ -88($fp)
gp ⇔ -92($fp)
sp ⇔ -96($fp)
fp ⇔ -100($fp)
ra ⇔ -104($fp)
```

The Next Instruction:



- Precall steps before "sub"
 - push args
 - save registers
 - jal sub # jump and link
- Steps to set up
 - build the frame
 - fp = sp + arg_size
 - sp = fp - frame_size
 - save S registers
- Steps to clean up
 - restore S registers
 - delete the frame (no need to!)
 - but sp = fp + 1
 - position the return value: (\$fp), \$v0
 - jr \$ra # jump register
- Postcall steps after "sub"
 - restore registers
 - move return value?
 - -4(\$sp), (\$fp), or \$v0

```
int sub(int X, int Y, int Z) {
    int j;
    int k = Y + Z

    j = sub(1, k, 3);
    → ;
    return j;
}
```

```
X  ⇔  0($fp)
Y  ⇔  -4($fp)
Z  ⇔  -8($fp)
j  ⇔  -12($fp)
k  ⇔  -16($fp)
t0 ⇔  -20($fp)
...
t9 ⇔  -56($fp)
s0 ⇔  -60($fp)
s7 ⇔  -88($fp)
gp ⇔  -92($fp)
sp ⇔  -96($fp)
fp ⇔  -100($fp)
ra ⇔  -104($fp)
```

Client -- Producer Convention Caveats:

- Main Memory is slow:
 - first 4 arguments should not be passed via the stack but via: \$a0, \$a1, \$a2, \$a3
 - the 2 return values should not be passed via the stack but via: \$v0, \$v1
- Although there are 32 general purpose registers:
 - Can't use: \$zero, \$at, \$k1, \$k2
 - If you use: \$gp, \$sp, \$fp, \$ra
 - you must take steps to save--restore these registers at call boundaries
 - if you use: \$a0, \$a1, \$a2, \$a3, \$v0, \$v1
 - you must take steps to save--restore these registers at call boundaries
- A compiler **MUST** follow this convention,
 - but the assembly level programmer can "optimize" their code!