

## Monte Carlo

Goal: compute  $\alpha = EX$

Settings:

$$X_1, \dots, X_n$$

$$\alpha_n = \bar{X}_n$$

rate of convergence as a  
function of  $c$

$$\alpha(c) = \bar{X}_{N(c)}$$

$$c^{\frac{1}{2}} (\alpha(c) - \alpha) \Rightarrow \lambda^{-\frac{1}{2}} \mathcal{O}N(0,1)$$

$$\lambda^{-1} \sigma^2 = ET \cdot \text{Var} X$$

$$\cdot \tilde{S}_n = \sum_{i=1}^n \tilde{X}_i$$

$$\cdot \tilde{X}_i = X_i - EX$$

$$c^{\frac{1}{2}} \left( \frac{S_{N(c)}}{N(c)} - \alpha \right) = c^{\frac{1}{2}} \left( \frac{S_{N(c)} - N(c)\alpha}{N(c)} \right)$$

$$= c^{\frac{1}{2}} \frac{\tilde{S}_{N(c)}}{c} \left( \frac{c}{N(c)} \right)$$
$$= \frac{\tilde{S}_{N(c)}}{\sqrt{c}} \underbrace{\left( \frac{c}{N(c)} \right)}_{\substack{\text{a.s.} \\ \rightarrow ET}}$$

$$\frac{N(c)}{c} \xrightarrow{\text{a.s.}} \lambda$$

$$\frac{\tilde{S}_{N(c)}}{\sqrt{c}} \stackrel{?}{\approx} \frac{\tilde{S}_{\lambda c}}{\sqrt{c}}$$

$$\frac{\tilde{S}_{N(c)} - S_{\lambda c}}{\sqrt{c}} \xrightarrow{P} 0$$

Anscombe's Theorem

$$\text{On } \{ |N(c) - \lambda c| \leq \varepsilon c \}$$

↑  
prob converges to 1

$$\frac{|\tilde{S}_{N(c)} - S_{\lambda c}|}{\sqrt{c}} \leq \frac{\max_{1 \leq j \leq c} |\tilde{S}_j - \tilde{S}_{\lambda c}|}{\sqrt{c}}$$

Kolmogorov  
maximal inequality

## Uniform Random Generator

$u_1, u_2, \dots$  i.i.d. uniform on  $[0,1]$

↓

non-uniform r.v.'s

↓

model randomness needed

## Uniform R.N.G.

1) physical r.n.g.

e.g. coin flips

1, 0, 0, 1, ...

bits:  $\sum_{i=1}^n b_i 2^{-i}$

- expensive
- slow
- biased
- correlated

e.g. charged particles

follows a Poisson distribution

- need to estimate  $\lambda$
- long time

2) deterministic algorithm that produces unpredictable outcomes

• Midsquare Method

```
      x x x x
      x x x x
      -----
0 0 0 0 0 0 0 0
      ~~~~~
```

"trap state": 0000

- Full-Period Generator

$$X_n = \text{'large random integers'} \in \{0, 1, \dots, n-1\}$$

$$u_n = X_n / n$$

$$X_{n+1} = (a X_n + b) \bmod n$$

often 0
large

Full period (hit  $1 - n^{-1}$ )  $\leftarrow$  cyclic

$$\begin{cases} n & \text{prime} \\ a & \end{cases} \quad \text{e.g. } 2^{31} - 1 \quad (\sim 2 \text{ billion})$$

$$X \stackrel{D}{=} F^{-1}(u)$$

uniform  $\rightarrow$  non-uniform

e.g. Exponential

$$F(x) = 1 - e^{-\lambda x}$$

$$F^{-1}(u) = -\frac{1}{\lambda} \log(1-u) \stackrel{D}{=} \text{Exp}(\lambda)$$

smallest
 $\frac{1}{2 \text{ billion}}$ 
 $\Rightarrow \lambda$  largest 30

- Toussaint

Nonsense Twitter

$10^{120}$

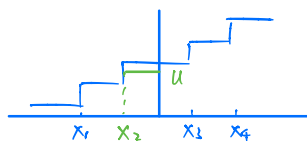
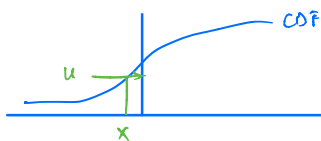
Statistical tests for random generators

### Non-uniform Generator

- Inversion

$X_1, X_2, \dots$  dist  $F$

$$X_i = F^{-1}(u_i)$$



• Acceptance - Rejection

$$\begin{array}{cc} P(\cdot) & Q(\cdot) \\ \uparrow & \uparrow \\ \text{want algo} & \text{good algo for} \\ & \text{generating samples} \end{array}$$

$$P(dw) = Q(dw) L(w)$$

$$f(x) = g(x) \left( \frac{f(x)}{g(x)} \right) \quad \sup_x \frac{f(x)}{g(x)} < c < \infty$$

Suppose  $L(w) \leq c < \infty, \forall w \in \Omega$

$$\begin{aligned} P(dw) &\propto Q(dw) \frac{L(w)}{c} \quad c \in [0, \infty] \\ &= Q(dw) P(u \leq \frac{L(w)}{c}) \end{aligned}$$

Steps:

1. Generate  $w$  under  $Q$
2. Generate on independent  $u$  and test

$$u \leq \frac{L(w)}{c}$$

- if yes, return  $w$  (Accept)
- if no, return to 1 (Reject)

{ Statistics  
 Machine Learning  
 Data Science

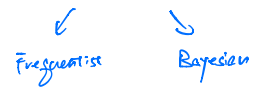
•  $X_1, X_2, \dots, X_n$  i.i.d.  $F$

sample from an underlying population



Model-free analysis "non-parametric"

- Parametric stat. modeling



- Non-parametric