

Solutions Submissions Editorial

← All Solutions

Description

Java

```
public String alienOrder(String[] words) {
   Map<Character, Set<Character>> map=new HashMap<Character, Set<Character>>();
    Map<Character, Integer> degree=new HashMap<Character, Integer>();
    String result="";
    if(words==null || words.length==0) return result;
    for(String s: words){
        for(char c: s.toCharArray()){
            degree.put(c,0);
   }
    for(int i=0; i<words.length-1; i++){</pre>
        String cur=words[i];
        String next=words[i+1];
        int length=Math.min(cur.length(), next.length());
        for(int j=0; j<length; j++){</pre>
            char c1=cur.charAt(j);
            char c2=next.charAt(j);
            if(c1!=c2){
                Set<Character> set=new HashSet<Character>();
                if(map.containsKey(c1)) set=map.get(c1);
                if(!set.contains(c2)){
                    set.add(c2);
                    map.put(c1, set);
                    degree.put(c2, degree.get(c2)+1);
                }
                break;
            }
        }
    Queue<Character> q=new LinkedList<Character>();
    for(char c: degree.keySet()){
        if(degree.get(c)==0) q.add(c);
   }
   while(!q.isEmpty()){
        char c=q.remove();
        result+=c;
        if(map.containsKey(c)){
            for(char c2: map.get(c)){
                degree.put(c2,degree.get(c2)-1);
                if(degree.get(c2)==0) q.add(c2);
```

Comments (103)

Sort by: Best









6



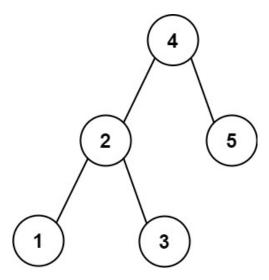
**■ Description** ■ Solutions Submissions □ Editorial

# 270. Closest Binary Search Tree Value Premium



Given the root of a binary search tree and a target value, return the value in the BST that is closest to the target. If there are multiple an smallest.

# Example 1:



**Input:** root = [4,2,5,1,3], target = 3.714286

Output: 4

#### Example 2:

**Input:** root = [1], target = 4.428571

Output: 1

# **Constraints:**

- The number of nodes in the tree is in the range [1, 10<sup>4</sup>].
- 0 <= Node.val <= 109
- $-10^9 <= target <= 10^9$

Seen this question in a real interview before? 1/4

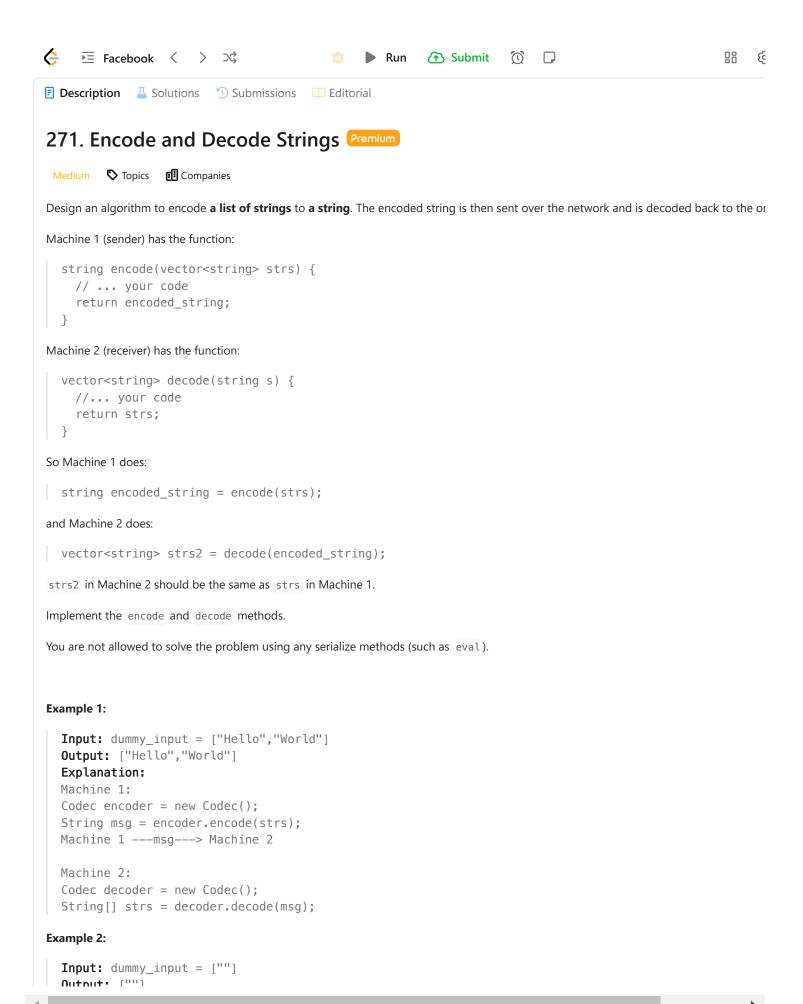
Yes No

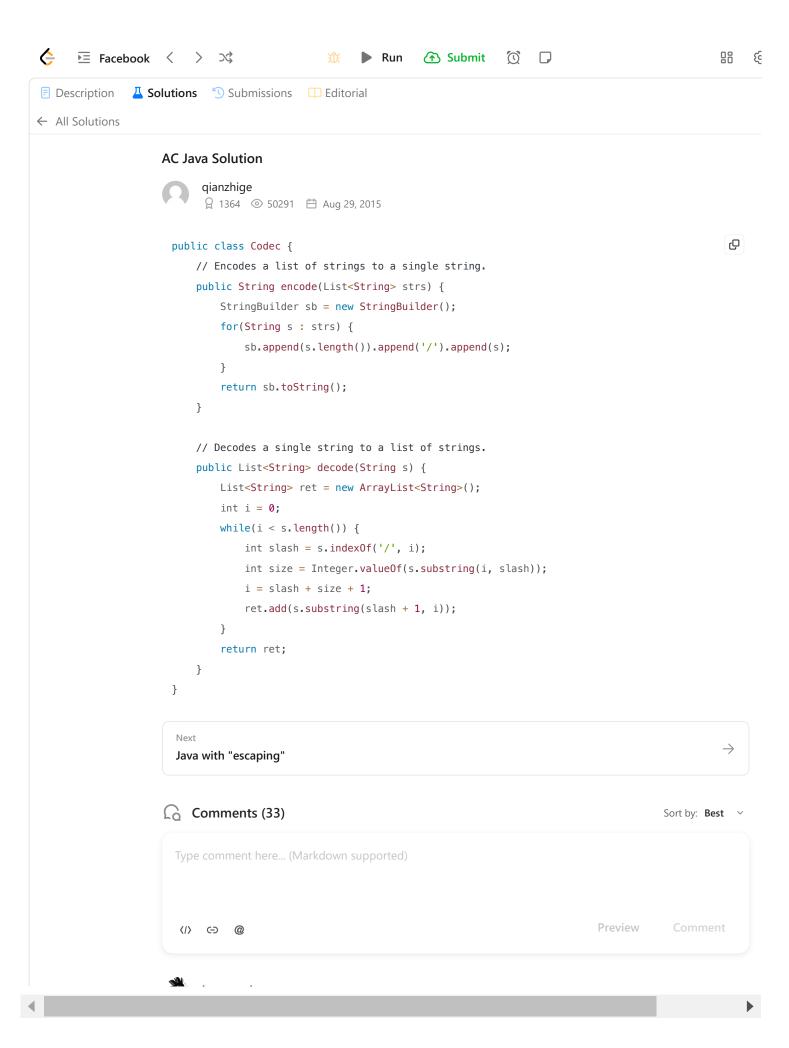
Submissions 623.1K Accepted 324.7K Acceptance Rate 52.1% ► Facebook 〈 > >☆

```
</>Code
Java ∨
           Auto
  1
       * Definition for a binary tree node.
  3
       * public class TreeNode {
             int val;
             TreeNode left;
  6
             TreeNode right;
             TreeNode() {}
  8
             TreeNode(int val) { this.val = val; }
  9
             TreeNode(int val, TreeNode left, TreeNode right) {
  10
                 this.val = val;
                 this.left = left;
  11
  12
                 this.right = right;
  13
             }
       * }
  14
       */
  15
      class Solution {
  16
          public int closestValue(TreeNode root, double target) {
  17
              int ret = root.val;
  19
               while(root != null){
  20
                   if(Math.abs(target - root.val) < Math.abs(target - ret)){</pre>
  21
                       ret = root.val;
  22
                   } else if (Math.abs(target - root.val) == Math.abs(target - ret)) {
                    if (root.val < ret) {</pre>
  23
  24
                       ret = root.val;
  25
                   }
  26
  27
                   root = root.val > target? root.left: root.right;
  28
  29
              return ret;
  30
          }
  31 }
○ Saved to cloud
```

Case 1 Case 2 + root = [4,2,5,1,3]

(4)





(1)

6

# 277. Find the Celebrity Premium

■ Description Solutions Submissions Editorial

Topics Companies Hint Medium

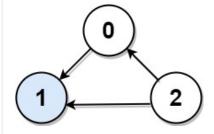
Suppose you are at a party with |n| people labeled from |0| to |n-1| and among them, there may exist one celebrity. The definition of a celebrate |n|other  $\begin{bmatrix} n - 1 \end{bmatrix}$  people know the celebrity, but the celebrity does not know any of them.

Now you want to find out who the celebrity is or verify that there is not one. You are only allowed to ask questions like: "Hi, A. Do you know information about whether A knows B. You need to find out the celebrity (or verify there is not one) by asking as few questions as possible ( sense).

You are given a helper function bool knows(a, b) that tells you whether a knows b. Implement a function int findCelebrity(n). Then celebrity if they are at the party.

Return the celebrity's label if there is a celebrity at the party. If there is no celebrity, return [-1].

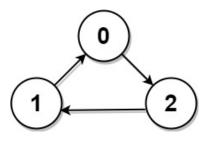
### Example 1:



Input: graph = [[1,1,0],[0,1,0],[1,1,1]]

Explanation: There are three persons labeled with 0, 1 and 2. graph[i][j] = 1 means person i knows p otherwise graph[i][j] = 0 means person i does not know person j. The celebrity is the person labelec both 0 and 2 know him but 1 does not know anybody.

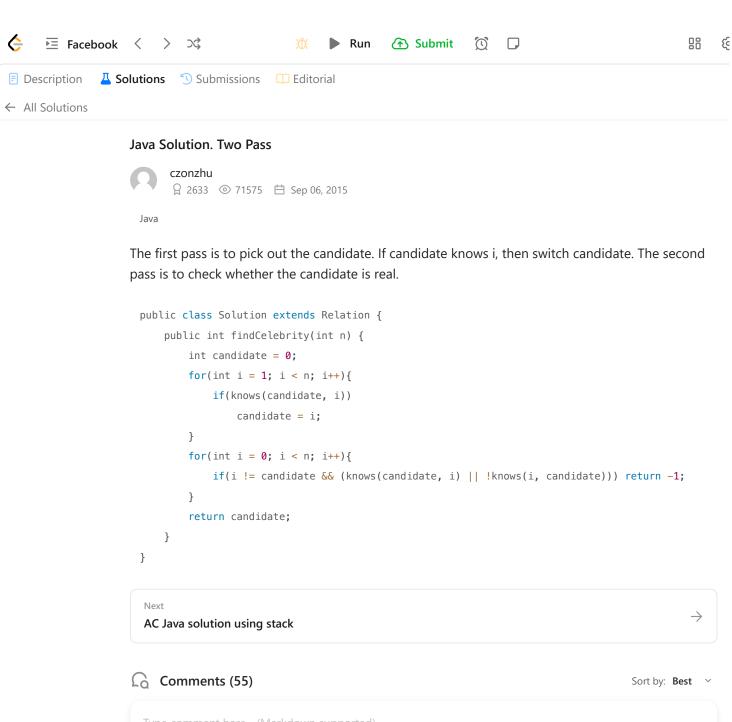
#### Example 2:



Input: graph = [[1,0,1],[1,1,0],[0,1,1]]

Output: -1

Explanation: There is no celebrity.



Comments (55)

Type comment here... (Markdown supported)

(I) © @ Preview Comment

ElementNotFoundExcepti...

Dec 12, 2015

હ

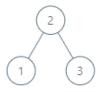
# 285. Inorder Successor in BST Premium

Medium Topics Companies

Given the root of a binary search tree and a node p in it, return the in-order successor of that node in the BST. If the given node has no in-c the tree, return null.

The successor of a node p is the node with the smallest key greater than p.val.

#### Example 1:

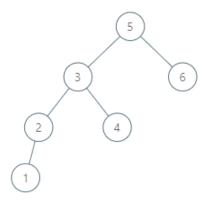


**Input:** root = [2,1,3], p = 1

Output: 2

Explanation: 1's in-order successor node is 2. Note that both p and the return value is of TreeNode

### Example 2:

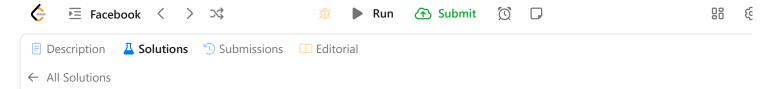


Input: root = [5,3,6,2,4,null,null,1], p = 6

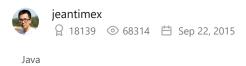
Output: null

Explanation: There is no in-order successor of the current node, so the answer is null.

- The number of nodes in the tree is in the range [1, 10<sup>4</sup>].
- $-10^5 <= Node.val <= 10^5$
- All Nodes will have unique values.



# Share my Java recursive solution



Just want to share my recursive solution for both getting the successor and predecessor for a given node in BST.

#### **Successor**

```
public TreeNode successor(TreeNode root, TreeNode p) {
  if (root == null)
    return null;

  if (root.val <= p.val) {
    return successor(root.right, p);
  } else {
    TreeNode left = successor(root.left, p);
    return (left != null) ? left : root;
  }
}</pre>
```

# **Predecessor**

```
public TreeNode predecessor(TreeNode root, TreeNode p) {
   if (root == null)
      return null;

   if (root.val >= p.val) {
      return predecessor(root.left, p);
   } else {
      TreeNode right = predecessor(root.right, p);
      return (right != null) ? right : root;
   }
}
```

Next

Java/Python solution, O(h) time and O(1) space, iterative

Comments (56)

Sort by: **Best**  $\vee$ 

 $\rightarrow$ 

હ

**■ Description** ■ Solutions Submissions □ Editorial

# 286. Walls and Gates Premium

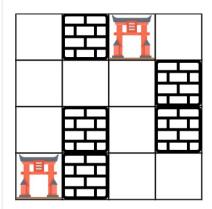
Medium Topics Companies

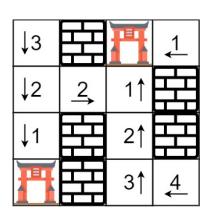
You are given an  $m \times n$  grid rooms initialized with these three possible values.

- -1 A wall or an obstacle.
- 0 A gate.
- INF Infinity means an empty room. We use the value  $2^{31} 1 = 2147483647$  to represent INF as you may assume that the distance to 2147483647.

Fill each empty room with the distance to its nearest gate. If it is impossible to reach a gate, it should be filled with INF.

#### Example 1:





Input: rooms = [[2147483647,-1,0,2147483647],[2147483647,2147483647,2147483647,-1],[2147483647,-1,21

[0,-1,2147483647,2147483647]]

**Output:** [[3,-1,0,1],[2,2,1,-1],[1,-1,2,-1],[0,-1,3,4]]

#### Example 2:

**Input:** rooms = [[-1]]

Output: [[-1]]

- m == rooms.length
- n == rooms[i].length
- 1 <= m, n <= 250
- rooms[i][j] is -1, 0, or  $2^{31} 1$ .

Java Breadth-First Search

Push all gates into queue first. Then for each gate update its neighbor cells and push them to the queue.

Repeating above steps until there is nothing left in the queue.

```
public class Solution {
    public void wallsAndGates(int[][] rooms) {
        if (rooms.length == 0 || rooms[0].length == 0) return;
        Queue<int[]> queue = new LinkedList<>();
        for (int i = 0; i < rooms.length; i++) {</pre>
            for (int j = 0; j < rooms[0].length; <math>j++) {
                if (rooms[i][j] == 0) queue.add(new int[]{i, j});
            }
        while (!queue.isEmpty()) {
            int[] top = queue.remove();
            int row = top[0], col = top[1];
            if (row > 0 \&\& rooms[row - 1][col] == Integer.MAX_VALUE) {
                rooms[row - 1][col] = rooms[row][col] + 1;
                queue.add(new int[]{row - 1, col});
            }
            if (row < rooms.length - 1 && rooms[row + 1][col] == Integer.MAX_VALUE) {
                rooms[row + 1][col] = rooms[row][col] + 1;
                queue.add(new int[]{row + 1, col});
            }
            if (col > 0 && rooms[row][col - 1] == Integer.MAX_VALUE) {
                rooms[row][col - 1] = rooms[row][col] + 1;
                queue.add(new int[]{row, col - 1});
            }
            if (col < rooms[0].length − 1 && rooms[row][col + 1] == Integer.MAX_VALUE) {
                rooms[row][col + 1] = rooms[row][col] + 1;
                queue.add(new int[]{row, col + 1});
            }
    }
}
```

Previous

Benchmarks of DFS and BFS

 $\leftarrow$ 

My short java solution, very easy to under...

હ

■ Description Solutions Submissions Editorial

# 296. Best Meeting Point Premium

Topics Companies Hint

Given an m x n binary grid grid where each 1 marks the home of one friend, return the minimal total travel distance.

The **total travel distance** is the sum of the distances between the houses of the friends and the meeting point.

The distance is calculated using Manhattan Distance, where distance(p1, p2) = |p2.x - p1.x| + |p2.y - p1.y|.

### Example 1:

1	0	0	0	1
0	0	0	0	0
0	0	1	0	0

Input: grid = [[1,0,0,0,0,1],[0,0,0,0,0],[0,0,1,0,0]]

Output: 6

**Explanation:** Given three friends living at (0,0), (0,4), and (2,2).

The point (0,2) is an ideal meeting point, as the total travel distance of 2 + 2 + 2 = 6 is minimal. So return 6.

#### Example 2:

**Input:** grid = [[1,1]]

Output: 1

- m == grid.length
- n == grid[i].length
- 1 <= m, n <= 200
- grid[i][j] is either 0 or 1.
- There will be at least two friends in the grid.



# 14ms java solution

← All Solutions

```
larrywang2014
     public int minTotalDistance(int[][] grid) {
   int m = grid.length;
   int n = grid[0].length;
   List<Integer> I = new ArrayList<>(m);
   List<Integer> J = new ArrayList<>(n);
    for(int i = 0; i < m; i++){</pre>
       for(int j = 0; j < n; j++){
           if(grid[i][j] == 1){
               I.add(i);
               J.add(j);
           }
       }
   }
   return getMin(I) + getMin(J);
}
private int getMin(List<Integer> list){
   int ret = 0;
   Collections.sort(list);
   int i = 0;
   int j = list.size() - 1;
   while(i < j){</pre>
       ret += list.get(j--) - list.get(i++);
   }
   return ret;
}
```

Next

Java 2ms/Python 40ms two pointers solution no median no sort with explanation

 $\rightarrow$ 





# 298. Binary Tree Longest Consecutive Sequence Premium

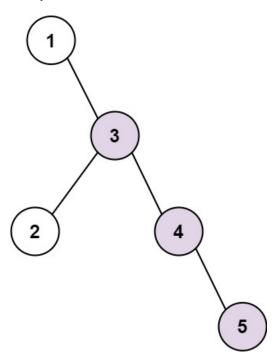
Medium Topics Companies

Given the [root] of a binary tree, return the length of the longest **consecutive sequence path**.

A consecutive sequence path is a path where the values increase by one along the path.

Note that the path can start at any node in the tree, and you cannot go from a node to its parent in the path.

## Example 1:

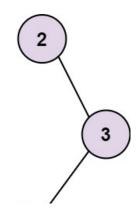


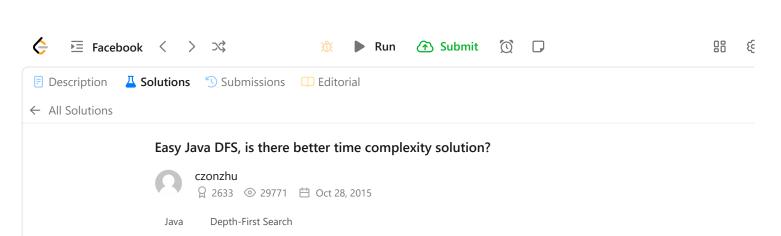
Input: root = [1,null,3,2,4,null,null,null,5]

Output: 3

Explanation: Longest consecutive sequence path is 3-4-5, so return 3.

## Example 2:

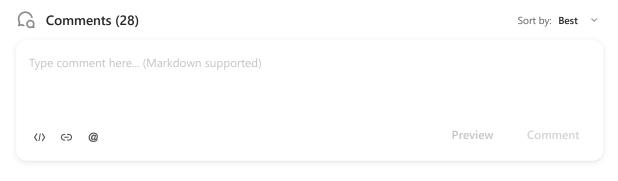




Just very intuitive depth-first search, send cur node value to the next level and compare it with the next level node.

```
public class Solution {
    private int max = 0;
    public int longestConsecutive(TreeNode root) {
        if(root == null) return 0;
        helper(root, 0, root.val);
        return max;
    }
    public void helper(TreeNode root, int cur, int target){
        if(root == null) return;
        if(root.val == target) cur++;
        else cur = 1;
        max = Math.max(cur, max);
        helper(root.left, cur, root.val + 1);
        helper(root.right, cur, root.val + 1);
    }
}
```

 $$\operatorname{\textsc{Next}}$$  Don't understand what is consecutive sequence



mmao3

Mar 19, 2019

6

# 311. Sparse Matrix Multiplication Premium

**Companies** 

Given two sparse matrices mat1 of size m x k and mat2 of size k x n, return the result of mat1 x mat2. You may assume that multiplica possible.

# Example 1:

**Input:** mat1 = [[1,0,0],[-1,0,3]], mat2 = [[7,0,0],[0,0,0],[0,0,1]]Output: [[7,0,0],[-7,0,3]]

# Example 2:

Input: mat1 = [[0]], mat2 = [[0]]

**Output:** [[0]]

#### **Constraints:**

- m == mat1.length
- k == mat1[i].length == mat2.length
- n == mat2[i].length
- 1 <= m, n, k <= 100
- -100 <= mat1[i][j], mat2[i][j] <= 100

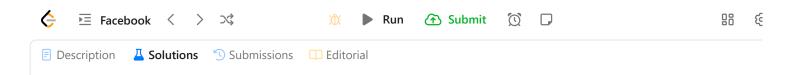
Seen this question in a real interview before? 1/4

Yes No

Accepted 187.3K Submissions 275.6K Acceptance Rate 68.0%

Topics

**Companies** 



#### **Easiest JAVA solution**

← All Solutions



UPDATE: Thanks to @stpeterh we have this 70ms concise solution:

The followings is the original 75ms solution:

The idea is derived from a CMU lecture.

A sparse matrix can be represented as a sequence of rows, each of which is a sequence of (column-number, value) pairs of the nonzero values in the row.

So let's create a non-zero array for A, and do multiplication on B.

Hope it helps!

```
public int[][] multiply(int[][] A, int[][] B) {
  int m = A.length, n = A[0].length, nB = B[0].length;
```

Solutions Submissions Editorial

The followings is the original 75ms solution:

The idea is derived from a CMU lecture.

A sparse matrix can be represented as a sequence of rows, each of which is a sequence of (column-number, value) pairs of the nonzero values in the row.

So let's create a non-zero array for A, and do multiplication on B.

Hope it helps!

```
public int[][] multiply(int[][] A, int[][] B) {
    int m = A.length, n = A[0].length, nB = B[0].length;
    int[][] result = new int[m][nB];
    List[] indexA = new List[m];
    for(int i = 0; i < m; i++) {</pre>
        List<Integer> numsA = new ArrayList<>();
        for(int j = 0; j < n; j++) {
            if(A[i][j] != 0){
                 numsA.add(j);
                 numsA.add(A[i][j]);
            }
        }
        indexA[i] = numsA;
    }
    for(int i = 0; i < m; i++) {</pre>
        List<Integer> numsA = indexA[i];
        for(int p = 0; p < numsA.size() - 1; <math>p += 2) {
            int colA = numsA.get(p);
            int valA = numsA.get(p + 1);
            for(int j = 0; j < nB; j ++) {
                 int valB = B[colA][j];
                 result[i][j] += valA * valB;
            }
    }
    return result;
}
```





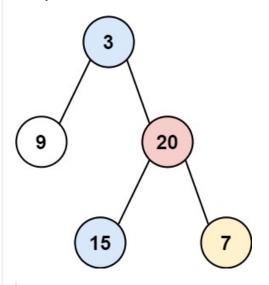
# 314. Binary Tree Vertical Order Traversal Premium

Medium Topics Companies Thint

Given the root of a binary tree, return the vertical order traversal of its nodes' values. (i.e., from top to bottom, column by column).

If two nodes are in the same row and column, the order should be from left to right.

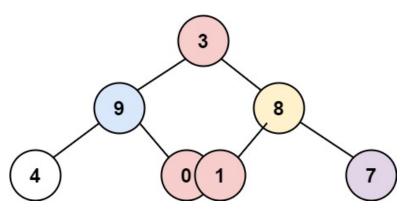
# Example 1:



Input: root = [3,9,20,null,null,15,7]

**Output:** [[9],[3,15],[20],[7]]

# Example 2:



Input: root = [3,9,8,4,0,1,7]

**Output:** [[4],[9],[3,0,1],[8],[7]]

# Example 3:

3

**E** Facebook <

> >3

```
</>Code
Java ∨
           Auto
  1
  2
       * Definition for a binary tree node.
  3
       * public class TreeNode {
             int val;
  5
             TreeNode left;
  6
             TreeNode right;
  7
             TreeNode() {}
  8
             TreeNode(int val) { this.val = val; }
  9
             TreeNode(int val, TreeNode left, TreeNode right) {
 10
                 this.val = val;
 11
                  this.left = left;
 12
                  this.right = right;
 13
             }
       * }
 14
       */
 15
      class Solution {
 16
 17
          public List<List<Integer>> verticalOrder(TreeNode root) {
 18
               List<List<Integer>> res = new ArrayList<>();
 19
               if (root == null) {
 20
                   return res;
 21
 22
 23
              Map<Integer, ArrayList<Integer>> map = new TreeMap<>();
 24
               Queue<TreeNode> q = new LinkedList<>();
 25
               Queue<Integer> cols = new LinkedList<>();
 26
 27
               q.add(root);
              cols.add(0);
 28
 29
              // int min = 0;
 30
 31
               // int max = 0;
 32
              while (!q.isEmpty()) {
 33
 34
                   TreeNode node = q.poll();
 35
                   int col = cols.poll();
 36
                   if (!map.containsKey(col)) {
 37
 38
                       map.put(col, new ArrayList<Integer>());
 39
 40
                   map.get(col).add(node.val);
 41
 42
                   if (node.left != null) {
 43
                       q.add(node.left);
 44
                       cols.add(col - 1);
 45
                       // min = Math.min(min, col - 1);
 46
 47
○ Saved to cloud
```

```
Case 1
             Case 2
                        Case 3
root =
  [3,9,20,null,null,15,7]
```

**E** Facebook <

> >3

```
</> Code
           Auto
Java ∨
 32
 33
               while (!q.isEmpty()) {
 34
                  TreeNode node = q.poll();
 35
                  int col = cols.poll();
 36
 37
                  if (!map.containsKey(col)) {
 38
                       map.put(col, new ArrayList<Integer>());
 39
 40
                  map.get(col).add(node.val);
 41
 42
                  if (node.left != null) {
 43
                       q.add(node.left);
 44
                       cols.add(col - 1);
 45
                       // min = Math.min(min, col - 1);
 46
 47
                  if (node.right != null) {
 48
 49
                       q.add(node.right);
 50
                       cols.add(col + 1);
 51
                       // max = Math.max(max, col + 1);
 52
                  }
 53
              }
 55
              // for (int i = min; i <= max; i++) {</pre>
 56
              //
                     res.add(map.get(i));
 57
              // }
 58
               for (var list : map.values()) {
 59
                res.add(list);
 60
 61
 62
              return res;
 63
          }
      }
 64
      //dfs below won't work because it requires top to buttom, but dfs stack can't follow that order
          // private void buildMap(TreeNode node, int index, Map<Integer, List<Integer>> indexMap) {
 67
               if (node == null)
 68
          //
 69
          //
                 return;
 70
          //
 71
          //
               indexMap.putIfAbsent(index, new ArrayList<Integer>());
 72
               indexMap.get(index).add(node.val);
 73
               buildMap(node.left, index - 1, indexMap);
 74
               buildMap(node.right, index + 1, indexMap);
          //
 75
          // }
 76
○ Saved to cloud
```

Case 1 Case 2 Case 3 root = [3,9,20,null,null,15,7]

6

# 317. Shortest Distance from All Buildings Premium

Topics Companies

You are given an  $m \times n$  grid grid of values [0, 1], or [2], where:

- each 0 marks an empty land that you can pass by freely,
- each 1 marks a building that you cannot pass through, and
- each 2 marks an obstacle that you cannot pass through.

You want to build a house on an empty land that reaches all buildings in the shortest total travel distance. You can only move up, down, le

Return the **shortest travel distance** for such a house. If it is not possible to build such a house according to the above rules, return [-1].

The total travel distance is the sum of the distances between the houses of the friends and the meeting point.

The distance is calculated using Manhattan Distance, where distance(p1, p2) = |p2.x - p1.x| + |p2.y - p1.y|.

#### Example 1:

1	0	2	0	1
0	0	0	0	0
0	0	1	0	0

Input: grid = [[1,0,2,0,1],[0,0,0,0,0],[0,0,1,0,0]]

**Explanation:** Given three buildings at (0,0), (0,4), (2,2), and an obstacle at (0,2).

The point (1,2) is an ideal empty land to build a house, as the total travel distance of 3+3+1=7 is So return 7.

#### Example 2:

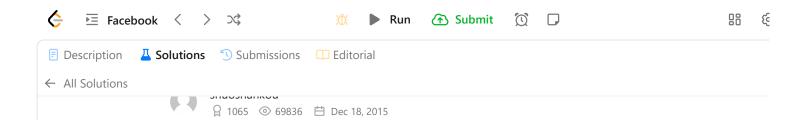
**Input:** grid = [[1,0]]

Output: 1

#### Example 3:

**Input:** grid = [[1]]

Output: -1



Inspired by previous solution.

Java

The main idea is the following:

Traverse the matrix. For each building, use BFS to compute the shortest distance from each '0' to this building. After we do this for all the buildings, we can get the sum of shortest distance from every '0' to all reachable buildings. This value is stored

in 'distance[][]'. For example, if grid[2][2] == 0, distance[2][2] is the sum of shortest distance from this block to all reachable buildings.

Time complexity: O(number of 1) $O(number of 0) \sim O(m^2n^2)$ 

We also count how many building each '0' can be reached. It is stored in reach[][]. This can be done during the BFS. We also need to count how many total buildings are there in the matrix, which is stored in 'buildingNum'.

Finally, we can traverse the distance[][] matrix to get the point having shortest distance to all buildings. O(m\*n)

The total time complexity will be  $O(m^2*n^2)$ , which is quite high!. Please let me know if I did the analysis wrong or you have better solution.

```
public class Solution {
    public int shortestDistance(int[][] grid) {
        if (grid == null || grid[0].length == 0) return 0;
        final int[] shift = new int[] \{0, 1, 0, -1, 0\};
        int row = grid.length, col = grid[0].length;
        int[][] distance = new int[row][col];
        int[][] reach = new int[row][col];
        int buildingNum = 0;
        for (int i = 0; i < row; i++) {
            for (int j =0; j < col; j++) {
                if (grid[i][j] == 1) {
                    buildingNum++;
                    Queue<int[]> myQueue = new LinkedList<int[]>();
                    myQueue.offer(new int[] {i,j});
                    boolean[][] isVisited = new boolean[row][col];
                    int level = 1;
```

```
public class Solution {
    public int shortestDistance(int[][] grid) {
        if (grid == null || grid[0].length == 0) return 0;
        final int[] shift = new int[] {0, 1, 0, -1, 0};
        int row = grid.length, col = grid[0].length;
        int[][] distance = new int[row][col];
        int[][] reach = new int[row][col];
        int buildingNum = 0;
        for (int i = 0; i < row; i++) {
            for (int j =0; j < col; j++) {
                if (grid[i][j] == 1) {
                    buildingNum++;
                    Queue<int[]> myQueue = new LinkedList<int[]>();
                    myQueue.offer(new int[] {i,j});
                    boolean[][] isVisited = new boolean[row][col];
                    int level = 1;
                    while (!myQueue.isEmpty()) {
                        int qSize = myQueue.size();
                        for (int q = 0; q < qSize; q++) {
                            int[] curr = myQueue.poll();
                            for (int k = 0; k < 4; k++) {
                                 int nextRow = curr[0] + shift[k];
                                 int nextCol = curr[1] + shift[k + 1];
                                if (nextRow >= 0 && nextRow < row && nextCol >= 0 &&
nextCol < col</pre>
                                     && grid[nextRow][nextCol] == 0 &&
!isVisited[nextRow][nextCol]) {
                                         //The shortest distance from
[nextRow][nextCol] to thic building
                                         // is 'level'.
                                         distance[nextRow][nextCol] += level;
                                         reach[nextRow][nextCol]++;
                                         isVisited[nextRow][nextCol] = true;
                                         myQueue.offer(new int[] {nextRow, nextCol});
                                     }
                            }
                        }
                        level++;
                    }
                }
            }
        }
        int shortest = Integer.MAX_VALUE;
```

```
for (int i = 0; i < row; i++) {
    for (int j = 0; j < col; j++) {
        if (grid[i][j] == 0 && reach[i][j] == buildingNum) {
            shortest = Math.min(shortest, distance[i][j]);
        }
    }
}
return shortest == Integer.MAX_VALUE ? -1 : shortest;
}
</pre>
```

6

🗏 **Description** 🚨 Solutions " Submissions 🕮 Editorial

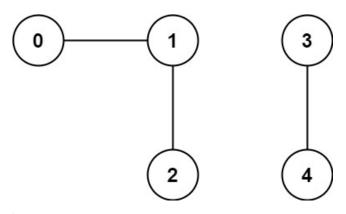
# 323. Number of Connected Components in an Undirected Graph Premium

Medium Topics Companies

You have a graph of n nodes. You are given an integer n and an array edges where edges [i] = [ai, bi] indicates that there is an edge  $b_i$  in the graph.

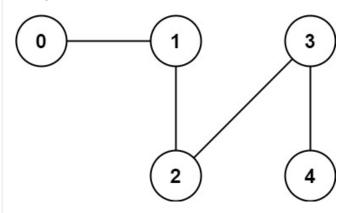
Return the number of connected components in the graph.

## Example 1:



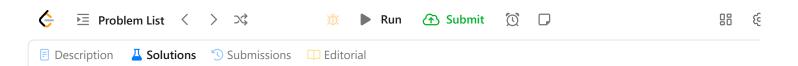
**Input:** n = 5, edges = [[0,1],[1,2],[3,4]] Output: 2

# Example 2:



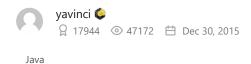
**Input:** n = 5, edges = [[0,1],[1,2],[2,3],[3,4]] **Output:** 1

- 1 <= n <= 2000
- 1 <= edges.length <= 5000



#### **Easiest 2ms Java Solution**

← All Solutions



This is 1D version of Number of Islands II. For more explanations, check out this 2D solution.

- 1. n points = n islands = n trees = n roots.
- 2. With each edge added, check which island is e[0] or e[1] belonging to.
- 3. If e[0] and e[1] are in same islands, do nothing.
- 4. Otherwise, **union** two islands, and reduce islands count by 1.
- 5. Bonus: path compression can reduce time by 50%.

Hope it helps!

```
public int countComponents(int n, int[][] edges) {
    int[] roots = new int[n];
    for(int i = 0; i < n; i++) roots[i] = i;</pre>
    for(int[] e : edges) {
        int root1 = find(roots, e[0]);
        int root2 = find(roots, e[1]);
        if(root1 != root2) {
            roots[root1] = root2; // union
            n--;
        }
   }
    return n;
}
public int find(int[] roots, int id) {
   while(roots[id] != id) {
        roots[id] = roots[roots[id]]; // optional: path compression
        id = roots[id];
   }
    return id;
}
```

Next

[Java] Union-Find, DFS, BFS Solutions - Complexity Explain - Clean code

E Similar Questions

```
public int maxSubArrayLen(int[] nums, int k) {
        int currSum = 0, maxLen = 0; // set initial values for cumulative sum and
max length sum to k
       HashMap<Integer, Integer> sumToIndexMap = new HashMap<Integer, Integer>();
// key: cumulative sum until index i, value: i
        for (int i = 0; i < nums.length; i++) {</pre>
            currSum = currSum + nums[i]; // update cumulative sum
            // two cases where we can update maxLen
            if (currSum == k) maxLen = i + 1; // case 1: cumulative sum is k, update
maxLen for sure
            else if (sumToIndexMap.containsKey(currSum - k)) maxLen =
Math.max(maxLen, i - sumToIndexMap.get(currSum - k)); // case 2: cumulative sum is
more than k, but we can truncate a prefix of the array
            // store cumulative sum in map, only if it is not seen
            // because only the earlier (thus shorter) subarray is valuable, when we
want to get the maxLen after truncation
            if (!sumToIndexMap.containsKey(currSum)) sumToIndexMap.put(currSum, i);
        return maxLen;
    }
```

હ

# 333. Largest BST Subtree Premium

Medium Topics **Companies** Hint

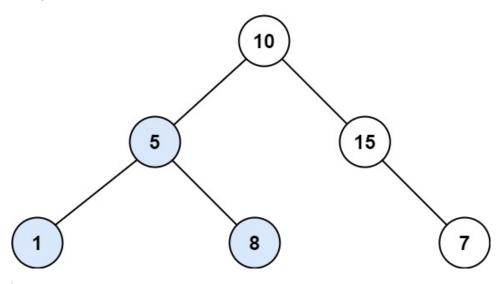
Given the root of a binary tree, find the largest subtree, which is also a Binary Search Tree (BST), where the largest means subtree has the lar nodes.

A Binary Search Tree (BST) is a tree in which all the nodes follow the below-mentioned properties:

- The left subtree values are less than the value of their parent (root) node's value.
- The right subtree values are greater than the value of their parent (root) node's value.

Note: A subtree must include all of its descendants.

#### Example 1:



Input: root = [10,5,15,1,8,null,7]

Output: 3

Explanation: The Largest BST Subtree in this case is the highlighted one. The return value is the su which is 3.

## Example 2:

Input: root = [4,2,7,2,3,5,null,2,null,null,null,null,null,1]

Output: 2

- The number of nodes in the tree is in the range [0, 10<sup>4</sup>].
- $-10^4 <= Node.val <= 10^4$

```
public class Solution {
    class Result { // (size, rangeLower, rangeUpper) -- size of current tree, range
of current tree [rangeLower, rangeUpper]
        int size;
        int lower;
        int upper;
        Result(int size, int lower, int upper) {
            this.size = size;
            this.lower = lower;
            this.upper = upper;
        }
    }
    int max = 0;
    public int largestBSTSubtree(TreeNode root) {
        if (root == null) { return 0; }
        traverse(root);
        return max;
    }
    private Result traverse(TreeNode root) {
        if (root == null) { return new Result(0, Integer.MAX_VALUE,
Integer.MIN_VALUE); }
        Result left = traverse(root.left);
        Result right = traverse(root.right);
        if (left.size == -1 || right.size == -1 || root.val <= left.upper ||
root.val >= right.lower) {
            return new Result(-1, 0, 0);
        int size = left.size + 1 + right.size;
        max = Math.max(size, max);
        return new Result(size, Math.min(left.lower, root.val),
Math.max(right.upper, root.val));
    }
}
```