



Atomic Boolean

Get a clear understanding and use cases of the AtomicBoolean class and its differences with volatile boolean variables.

We'll cover the following



- Explanation
- Difference between Volatile boolean and AtomicBoolean

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Explanation#

The AtomicBoolean class belonging to Java's

java.util.concurrency.atomic package represents a boolean value that can be updated and modified atomically. The Atomic* family of classes extend the notion of volatile variables that are designed to be operated upon without locking using machine-level atomic instructions available on modern processors. However, on other platforms some form of internal locking may be used to serialize thread access.

Atomic* classes including AtomicBoolean offer a method compareAndSet(expectedValue, updatedValue) to conditionally update the value of the variable to updatedValue if it is set to expectedValue in one go, i.e. atomically. All read-and-update methods except for lazySet() and weakCompareAndSet() have memory effects equivalent of both reading and writing volatile variables.





The read and write methods i.e. get() and set() on instances of this class are similar in behavior to volatile variables i.e get() has the memory effect of reading a volatile variable and set() has the memory effect of writing a volatile variable.

Difference between Volatile boolean and AtomicBoolean#

It may be tempting to confuse and equate volatile boolean with AtomicBoolean, however, the two differ in several respects. Though it is fair to say that volatile boolean exhibits a subset of the functionality of AtomicBoolean. Recall, from our discussion on the volatile keyword that Java's memory model and compiler optimization can result in a situation where one thread doesn't see the latest value of a shared variable in main memory since it is using a cached stale value. Marking such a shared variable as volatile ensures that all threads see the latest state of the variable at all times. The quintessential example of this scenario found in online literature is as follows:









```
class SomeClass {
  boolean quit = false;

void thread1() {
    while (!quit) {
        // ... Take some action
    }
}

void thread2() {
    // thread 1 may never see the updated value of quit!
    quit = true;
}
```

Without volatile the thread executing the method thread1 in the above program can spin in an infinite loop and never quit. Apart from delivering a consistent view of the memory, the volatile keyword doesn't promise much in synchronization guarantees. Specifically, multiple threads accessing a volatile keyword don't do so in a serialized manner. The onus of making a volatile variable's accesses synchronized and thread-safe is on the developer. This is where the AtomicBoolean classes come in. For instance, the methods compareAndSet() and getAndSet() exposed by the AtomicBoolean class represent a series of operations executed atomically, which otherwise would require synchronization on the part of the developer and are unachievable as an atomic transaction using volatile variables.

To make the distinction clearer let's work with an example. Say, we instantiate several threads and have them race to change a shared boolean variable's won value to true. The first one to do so is considered to have won the race. The rest print a message to indicate they lost the race. We mark the variable won as volatile and run the program below:

```
€€}
    import java.util.concurrent.ExecutorService;
1
    import java.util.concurrent.Executors;
2
    import java.util.concurrent.Future;
4
5
   class Demonstration {
6
7
        static volatile boolean won = false;
8
9
        public static void main( String args[] ) throws Exception {
            ExecutorService es = Executors.newFixedThreadPool(25);
10
11
            try {
                int numThreads = 15;
12
13
                Runnable[] racers = new Runnable[numThreads];
14
                Future[] futures = new Future[numThreads];
15
16
                // create thread tasks
17
                for (int i = 0; i < numThreads; i++) {
                     racers[i] = new Runnable() {
18
19
                         @Override
                         public void run() {
20
21
                             race():
22
                         }
23
                    };
24
                }
25
26
                // submit threads for execution
27
                for (int i = 0; i < numThreads; i++) {
                     futures[i] = es.submit(racers[i]);
28
\triangleright
```

If you run the above program enough times especially on a machine with multiple processes you'll observe multiple threads printing the winning statement. The widget above may not show the faulty run since the VM executing the code may have. asingle vCPU. The volatile variable doesn't prevent multiple threads from accessing and then updating the variable won . A possible sequence resulting in multiple threads declaring themselves as winners can be:

- 1. Thread A reads the value of won which is false.
- 2 Thread B reads the value of won which is false

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- 3. Thread A changes the value of won to true and the variable is updated in main memory for all threads to see.
- 4. Thread B doesn't know Thread A had read the value of the variable won before Thread B accessed it. Thread B too updates the value of won to true.
- 5. Since won is a volatile variable, it always reflects its latest value to the next thread reading it but that's about it.

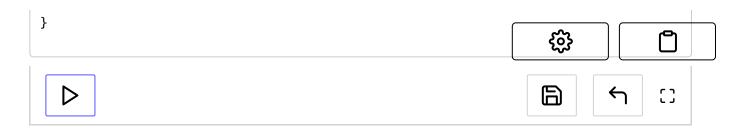
The above program can be fixed by using appropriate synchronization as shown in the widget below:

6





```
import java.util.concurrent.ExecutorService;
import java.util.concurrent.Executors;
import java.util.concurrent.Future;
class Demonstration {
    static Object syncObj = new Object();
    static boolean won = false;
    public static void main( String args[] ) throws Exception {
        ExecutorService es = Executors.newFixedThreadPool(25);
        try {
            int numThreads = 15;
            Runnable[] racers = new Runnable[numThreads];
            Future[] futures = new Future[numThreads];
            // create thread tasks
            for (int i = 0; i < numThreads; i++) {
                racers[i] = new Runnable() {
                    @Override
                    public void run() {
                        race();
                    }
                };
            }
            // submit threads for execution
            for (int i = 0; i < numThreads; i++) {
                futures[i] = es.submit(racers[i]);
            }
            // wait for threads to finish
            for (int i = 0; i < numThreads; i++) {
                futures[i].get();
        } finally {
            es.shutdown();
        }
    }
    static void race() {
        synchronized (sync0bj) {
            if (!won) {
                won = true;
                System.out.println(Thread.currentThread().getName() + " won the rac
            } else {
                System.out.println(Thread.currentThread().getName() + " lost.");
            }
        }
    }
```



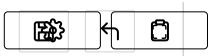
The above program synchronizes access to the won variable and only a single thread ever operates on the won variable at a time. Note that we have not marked won as volatile since the synchronized block establishes the *happens-before* relationship i.e. any thead reading the variables updated in the synchronized block see the latest value for those variables. We can rid of the synchronized block by using AtomicBoolean as shown in the following widget:





```
import java.util.concurrent.ExecutorService;
import java.util.concurrent.Executors;
import java.util.concurrent.Future;
import java.util.concurrent.atomic.AtomicBoolean;
class Demonstration {
    static AtomicBoolean won = new AtomicBoolean(false);
    public static void main( String args[] ) throws Exception {
        ExecutorService es = Executors.newFixedThreadPool(25);
        try {
            int numThreads = 15;
            Runnable[] racers = new Runnable[numThreads];
            Future[] futures = new Future[numThreads];
            // create thread tasks
            for (int i = 0; i < numThreads; i++) {
                racers[i] = new Runnable() {
                    @Override
                    public void run() {
                        race();
                    }
                };
            }
            // submit threads for execution
            for (int i = 0; i < numThreads; i++) {
                futures[i] = es.submit(racers[i]);
            }
            // wait for threads to finish
            for (int i = 0; i < numThreads; i++) {
                futures[i].get();
        } finally {
            es.shutdown();
        }
    }
    static void race() {
        if (won.compareAndSet(false, true)) {
            System.out.println(Thread.currentThread().getName() + " won the race.")
            System.out.println(Thread.currentThread().getName() + " lost.");
        }
    }
}
```





The method compareAndSet(expectedValue, updatedValue) atomically checks if the variable won is false and sets it to true and then returns true if the entire operation is successful. Since these compound steps occur atomically, we can be assured that only a single thread ever gets to set the variable won to true. Other threads all see the variable won as false and fall in the else clause. Using AtomicBoolean instead of synchronized block results in code that is more readable, concise and faster on machines with support for machine-level atomic operations.

In summary, you can think of the AtomicBoolean and its sibling classes as offering volatile-like capabilities and then some.

Consider the following statement and pick the correct statement:

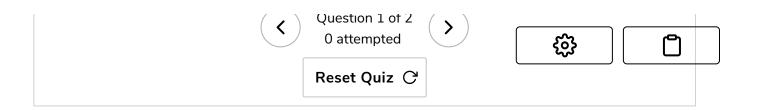
static volatile AtomicBoolean boolVar;

A) The statement won't compile as AtomicBoolean already mirrors the volatile functionality.

B) The volatile is superfluous and removing it doesn't break any functionality.

C) volatile is needed for synchronization purposes.

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