

# Quantum Circuits with Classical and Quantum Control of Causal Orders

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joint work with  
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[arXiv:1807.10557 & arXiv:19??:soon™]

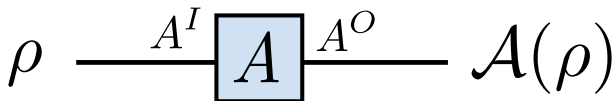


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DE GENÈVE**

# Quantum Channels

Quantum channels are most general physical map from quantum states to quantum states

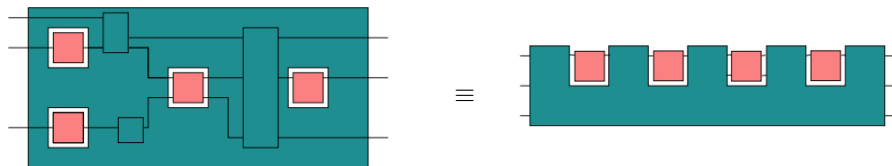
- *Completely*-positive trace preserving (CPTP) maps
- Extensively studied in quantum information



# Quantum Combs

How can channels themselves be transformed?

- Quantum combs introduced as a general type of physical transformation mapping quantum channels to a quantum channel



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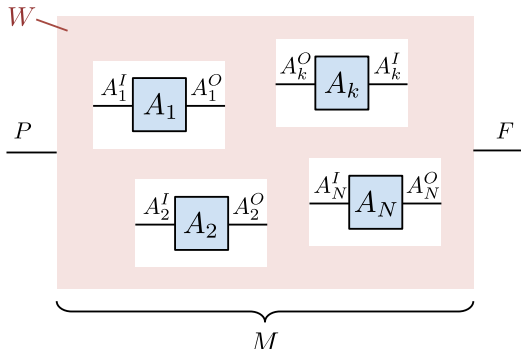
[Chiribella, D'Ariano and Perinotti, PRA 2018]

# Quantum Supermaps

Quantum combs not the most general transformation on channels

## Quantum Superchannel

A quantum superchannel is a transformation taking quantum channels into quantum channels even when applied to only part of channels' inputs. I.e., it is completely CPTP preserving.



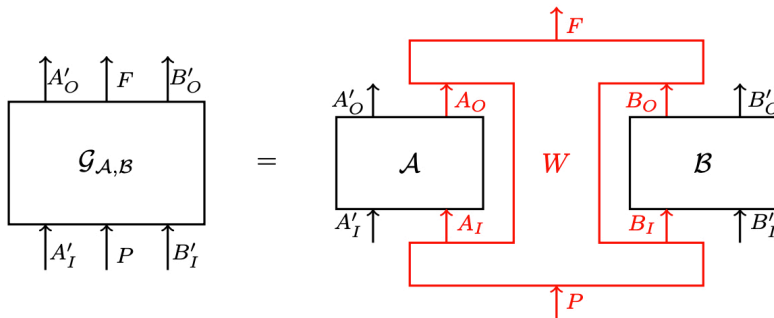
[Chiribella, D'Ariano and Perinotti, EurPL 2008], [Araújo, Feix, Navascués and Brukner, Quantum 2017], [Quintino, Dong, Shimbo, Soeda, Murao, arXiv 2018]

# Quantum Supermaps

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## Quantum Superchannel

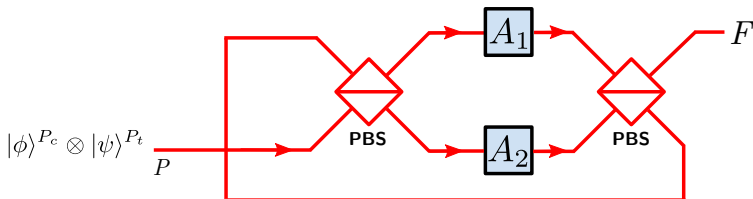
A quantum superchannel is a transformation taking quantum channels into quantum channels **even when applied to only part of channels' inputs**. I.e., it is completely CPTP preserving.



[Chiribella, D'Ariano and Perinotti, EurPL 2008], [Araújo, Feix, Navascués and Brukner, Quantum 2017], [Quintino, Dong, Shimbo, Soeda, Murao, arXiv 2018]

# Superchannels and Causal Order

- Superchannels also studied as **process matrices** for their ability to have **indefinite causal order**
- **Quantum switch** shows quantum mechanics allows interesting ways to compose channels beyond quantum combs



Do all superchannels have a physical implementation?

# Motivation

- Quantum combs are natural, physical transformations, but unnecessarily restrictive for some applications
- We know physical superchannels beyond combs exist, but have no good systematic approach to study them
  - Examples mostly *ad hoc*: the quantum switch and its generalisation, the  $N$ -switch

Goal: find natural physical classes of superchannels and understand their relation to causal structure

# Outline

## Quantum superchannels

- General setting and definition

## Quantum superchannels with a fixed causal structure

- Quantum circuits (quantum combs)

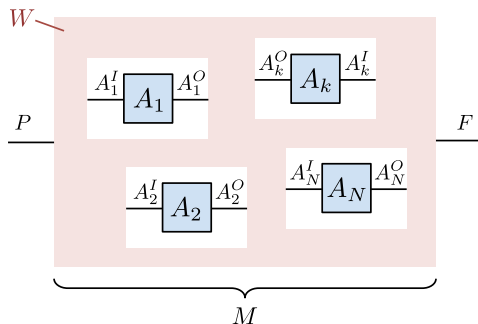
## Beyond superchannels with fixed causal structure

- Quantum circuits with classical control of causal order

- Quantum circuits with quantum (coherent) control of causal order



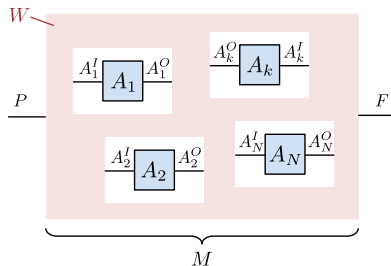
# Quantum Superchannels



Work in the Choi picture

- $|\mathbb{1}\rangle\rangle = \sum_i |i\rangle \otimes |i\rangle$  is the “pure Choi isomorphism” of an identity channel
- Pure Choi isomorphism: for an operator  $A$ ,  $|A\rangle\rangle = \mathbb{1} \otimes A |\mathbb{1}\rangle\rangle$
- Choi isomorphism: for a CP map  $\mathcal{M}$ ,  $M = \mathcal{I} \otimes \mathcal{M}(|\mathbb{1}\rangle\rangle\langle\langle\mathbb{1}|)$

# Quantum Supermaps



Superchannel represented by an operator  $W$  satisfying

$$W \geq 0, \quad \text{Tr } W = d_P d_{A_1^O} \dots d_{A_N^O}, \quad W \in \mathcal{L}^{\mathcal{N}}$$

$$\begin{aligned} M &= \text{Tr}_{A_1^O \dots A_N^O} [W(A_1^T \otimes \dots \otimes A_N^T \otimes \mathbb{1}^{PF})] \\ &= W * (A_1 \otimes \dots \otimes A_N) \in PF, \end{aligned}$$

■ Output state given by  $M * \rho = W * (\rho \otimes A_1 \otimes \dots \otimes A_N)$

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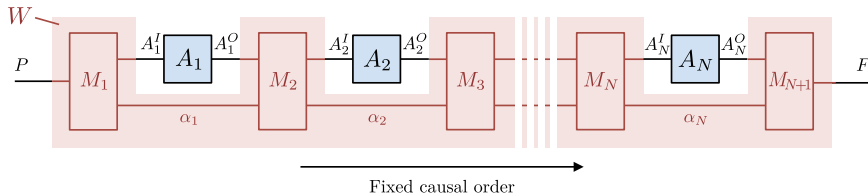
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# Quantum Combs as Fixed-Order Circuits

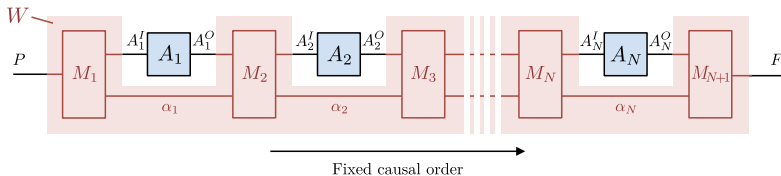
- Quantum combs can be seen as a circuit with a fixed order  $A_1 \prec \dots A_N$



Most general quantum circuit described by CPTP maps:

- $\mathcal{M}_1 : P \rightarrow A_1^I \alpha_1$ , where  $\alpha_1$  is an ancillary system
- $\mathcal{M}_{n+1} : A_n^O \alpha_n \rightarrow A_{n+1}^I \alpha_{n+1}$  for  $1 \leq n \leq N - 1$
- $\mathcal{M}_{N+1} : A_N^O \alpha_N \rightarrow F$

# Quantum Circuits with Fixed Causal Order



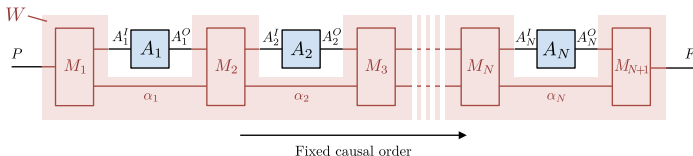
For input  $\rho$  output is

$$M_{N+1} * A_N * \cdots * M_2 * A_1 * M_1 * \rho = \underbrace{(M_1 * M_2 * \cdots * M_{N+1})}_W * (\rho \otimes A_1 \otimes \cdots \otimes A_N)$$

■  $W$  is defined uniquely by the maps  $M_n$  via the link product:

$$W = M_1 * M_2 * \cdots * M_{N+1} = \text{Tr}_{\alpha_1 \cdots \alpha_N} [M_1 \otimes M_2^{T_{\alpha_1}} \otimes \cdots \otimes M_{N+1}^{T_{\alpha_N}}]$$

# QC-FO Characterisation



The constraint that the  $M_n$  are CPTP maps and thus satisfy

$\text{Tr}_{A_{n+1}^I \alpha_{n+1}} M_{n+1} = \mathbb{1}^{A_n^O \alpha_n}$  allow the  $W$  of QC-FOs to be characterised

- They are precisely process matrices compatible with  $P \prec A_1 \prec \dots \prec A_N \prec F$

QC-FOs compatible with order  $P \prec A_1 \prec \dots \prec A_N \prec F$

$$\text{Tr}_F W = W_{(N)} \otimes \mathbb{1}^{A_N^O},$$

$$\text{Tr}_{A_{n+1}^I} W_{(n+1)} = W_{(n)} \otimes \mathbb{1}^{A_n^O} \quad \forall n = 1, \dots, N-1,$$

$$\text{and } \text{Tr}_{A_1^I} W_{(1)} = \mathbb{1}^P.$$

where  $W_{(n)} := \frac{1}{d_n^O d_{n+1}^O \dots d_N^O} \text{Tr}_{A_n^O A_{\{n+1, \dots, N\}}^I} W$  are reduced process matrices

# Example of a QC-FO



$$M_1 = |\mathbb{1}\rangle\langle\mathbb{1}|, \quad M_2 = |U\rangle\langle U|, \quad M_3 = |\mathbb{1}\rangle\langle\mathbb{1}|$$

We then have

$$\begin{aligned} W &= |\mathbb{1}\rangle\langle\mathbb{1}|^{PA^I} * |U\rangle\langle U|^{A^O B^I} * |\mathbb{1}\rangle\langle\mathbb{1}|^{B^O F} \\ &= |\mathbb{1}\rangle\langle\mathbb{1}|^{PA^I} \otimes |U\rangle\langle U|^{A^O B^I} \otimes |\mathbb{1}\rangle\langle\mathbb{1}|^{B^O F} \\ &= |w\rangle\langle w|, \end{aligned}$$

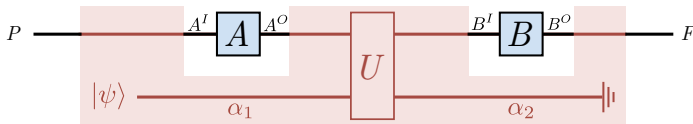
with  $|w\rangle = |\mathbb{1}\rangle \otimes |U\rangle \otimes |\mathbb{1}\rangle$

■ Conditions easily verified, e.g.

$$\mathrm{Tr}_F W = \underbrace{|\mathbb{1}\rangle\langle\mathbb{1}|^{PA^I} \otimes |U\rangle\langle U|^{A^O B^I}}_{W_{(2)}} \otimes \mathbb{1}^{B^O}, \text{ etc.}$$

# Example of a QC-FO

Channel with memory



$$M_1 = |\mathbb{1}\rangle\langle\mathbb{1}|^{PA^I} \otimes |\psi\rangle\langle\psi|^{\alpha_1}, \quad M_2 = |U\rangle\langle U|^{A^O\alpha_1 B^I\alpha_2}, \quad M_3 = |\mathbb{1}\rangle\langle\mathbb{1}|^{B^OF} \otimes \mathbb{1}^{\alpha_2}$$

This gives

$$W = \text{Tr}_{\alpha_2} \left[ |\mathbb{1}\rangle\langle\mathbb{1}|^{PA^I} \otimes (\langle\psi^*|U\rangle\langle U|\psi^*\rangle)^{A^OB^I\alpha_2} \otimes |\mathbb{1}\rangle\langle\mathbb{1}|^{B^OF} \right]$$



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# Classical Control of Causal Order

Mixing circuits with different orders is also physics

Bipartite Causal Separability [*Nat. Commun.* 2012]

A superchannel  $W$  is **causally separable** if it can be written

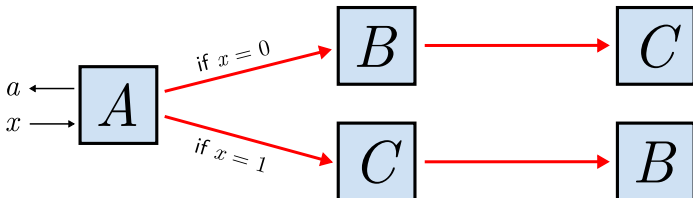
$$W = qW^{A \prec B} + (1 - q)W^{B \prec A}$$

for  $W^{A \prec B}, W^{B \prec A}$  QC-FOs. Otherwise it is **causally nonseparable**.

- The quantum switch can be proven to be causally nonseparable

# Dynamical Causal Order

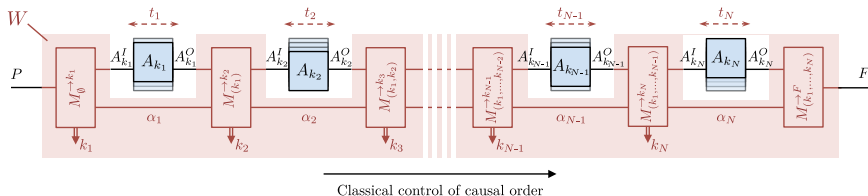
When considering more operations the situation is more subtle: one can have **dynamical causal order**



- Definition of causal separability and its characterisation is more subtle
- Can one give a circuit-type description of such superchannels?

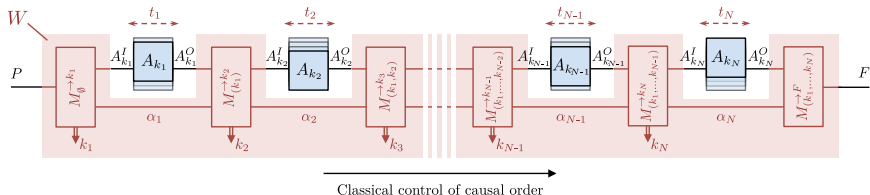
# Classically Controlled Circuits

- Oreshkov & Giarmatzi [NJP, 2016] suggested causal separability corresponds to **quantum circuits with classical control of causal order** (QC-CCs): “**classically controlled quantum circuits**”



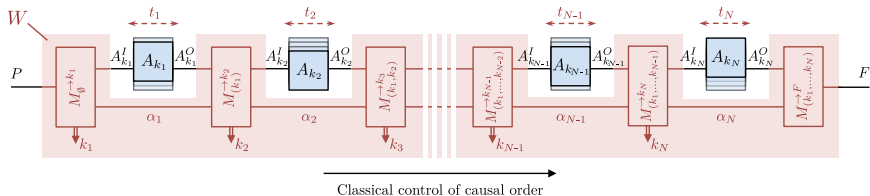
- At each time slot  $t_n$  exactly one operation  $A_{k_n}$  is applied
- Crucial requirement: **each operation applied once and only once**, irrespective of the operations themselves
  - Needed to ensure  $W$  gives a valid superchannel

# Classically Controlled Circuits



- Outcome of instrument  $\{M_{(k_1, \dots, k_n)}^{\rightarrow k_{n+1}}\}_{k_{n+1}}$  determines the  $(n+1)$ th operation to apply
- Technicality: the  $M_{(k_1, \dots, k_n)}^{\rightarrow k_{n+1}} \in A_{k_n}^O \alpha_n A_{k_{n+1}}^I \alpha_{n+1}$  belong to different spaces
  - Can solve by embedding in common direct-sum output space

# Process Matrix of a QC-CC



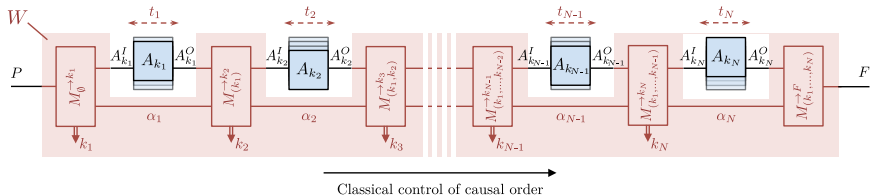
For input  $\rho$ , when operations applied in order  $k_1, \dots, k_N$ , output is

$$\begin{aligned}
 & M_{(k_1, \dots, k_N)}^{\rightarrow F} * A_{k_N} * M_{(k_1, \dots, k_{N-1})}^{\rightarrow k_N} * \dots * M_{(k_1, k_2)}^{\rightarrow k_3} * A_{k_2} * M_{(k_1)}^{\rightarrow k_2} * A_{k_1} * M_{\emptyset}^{\rightarrow k_1} * \rho \\
 &= \underbrace{M_{\emptyset}^{\rightarrow k_1} * M_{(k_1)}^{\rightarrow k_2} * M_{(k_1, k_2)}^{\rightarrow k_3} * \dots * M_{(k_1, \dots, k_{N-1})}^{\rightarrow k_N}}_{\widetilde{W}_{(k_1, \dots, k_N, F)}} * M_{(k_1, \dots, k_N)}^{\rightarrow F} * (\rho \otimes A_1 \otimes \dots \otimes A_N)
 \end{aligned}$$

Process matrix of a QC-CC

$$W = \sum_{(k_1, \dots, k_N)} \widetilde{W}_{(k_1, \dots, k_N, F)}$$

# QC-CC Characterisation



## Characterisation of circuits with classically controlled order

$W$  is the process matrix of a QC-CC iff  $\exists$  PSD matrices  $W_{(k_1, \dots, k_N, F)}, W_{(k_1, \dots, k_n)}$  for all  $1 \leq n \leq N$  with:

- $W = \sum_{(k_1, \dots, k_N)} W_{(k_1, \dots, k_N, F)}$
- $\forall (k_1, \dots, k_N), \text{Tr}_F W_{(k_1, \dots, k_N, F)} = W_{(k_1, \dots, k_N)} \otimes \mathbb{1}_{A_{k_N}^O}$
- $\forall n = 1, \dots, N-1, \forall (k_1, \dots, k_n)$   
 $\sum_{k_{n+1}} \text{Tr}_{A_{k_{n+1}}^I} W_{(k_1, \dots, k_n, k_{n+1})} = W_{(k_1, \dots, k_n)} \otimes \mathbb{1}_{A_{k_n}^O}$
- $\sum_{k_1} \text{Tr}_{A_{k_1}^I} W_{(k_1)} = \mathbb{1}^P$

# QC-CC Characterisation

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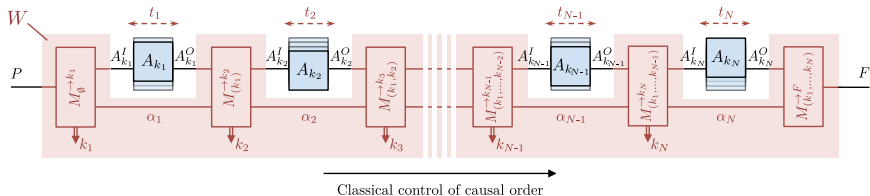
$W_{(k_1, \dots, k_N, F)}, W_{(k_1, \dots, k_n)}$  for all  $1 \leq n \leq N$  with:

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- $\forall n = 1, \dots, N-1, \forall (k_1, \dots, k_n)$   
 $\sum_{k_{n+1}} \text{Tr}_{A_{k_{n+1}}^I} W_{(k_1, \dots, k_n, k_{n+1})} = W_{(k_1, \dots, k_n)} \otimes \mathbb{1}^{A_{k_n}^O}$
- $\sum_{k_1} \text{Tr}_{A_{k_1}^I} W_{(k_1)} = \mathbb{1}^P$

- **Coincides with a sufficient (and possible necessary) condition for general causal separability** from Wechs, AA and Branciard [NJP 2019]
  - An operational interpretation for causal separability
- Can be checked with semidefinite programming techniques



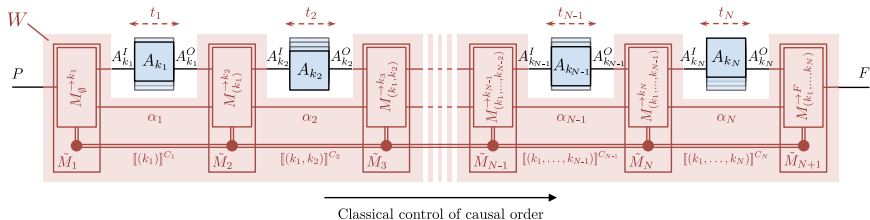
## Example of a QC-CC



### Classical Switch:

# Alternative Descriptions of QC-CCs

- Conditioning can be included in operations by introducing **(classical) control system**  $\llbracket (k_1, \dots, k_n) \rrbracket^{C_n} := |(k_1, \dots, k_n)\rangle\langle (k_1, \dots, k_n)|^{C_n}$

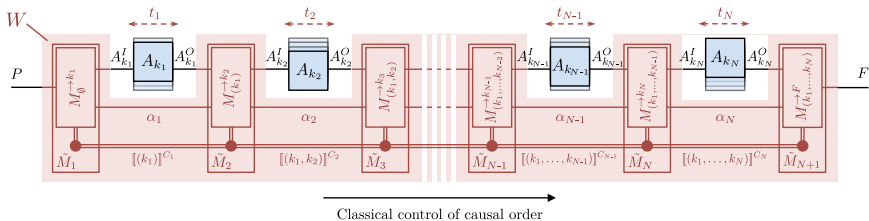


- Operations now given by the CPTP maps

$$\tilde{M}_{n+1} := \sum_{k_1, \dots, k_n, k_{n+1}} \tilde{M}_{(k_1, \dots, k_n)}^{\rightarrow k_{n+1}} \otimes \llbracket (k_1, \dots, k_n) \rrbracket^{C_n} \otimes \llbracket (k_1, \dots, k_n, k_{n+1}) \rrbracket^{C_{n+1}},$$

$$\tilde{M}_1 := \sum_{k_1} \tilde{M}_\emptyset^{\rightarrow k_1} \otimes \llbracket (k_1) \rrbracket^{C_1}, \quad M_{N+1} := \sum_{k_1, \dots, k_N} M_{(k_1, \dots, k_N)}^{\rightarrow F} \otimes \llbracket (k_1, \dots, k_N) \rrbracket^{C_N}$$

## Alternative Descriptions of QC-CCs



- Defining global controlled operations  $\tilde{A}_n := \bigoplus_{k_n \in \mathcal{N}} A_{k_n}$  we have

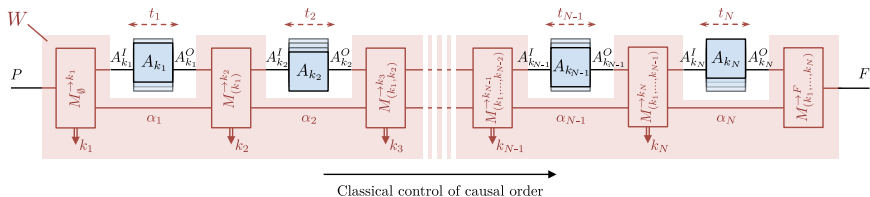
$$\tilde{M}_{N+1} * \tilde{A}_N * \tilde{M}_N * \cdots * \tilde{A}_1 * \tilde{M}_1 * \rho = \underbrace{\sum_{k_1, \dots, k_N} W_{(k_1, \dots, k_N, F)}}_W * (\rho \otimes A_1 \otimes \cdots \otimes A_N)$$

- Note that wlog we can take all operations to be purified isometries

$$M_{(k_1, \dots, k_{n-1})}^{\rightarrow k_n} = |V_{(k_1, \dots, k_{n-1})}^{\rightarrow k_n} \rangle \rangle \langle \langle V_{(k_1, \dots, k_{n-1})}^{\rightarrow k_n} |$$

- Suggests natural generalisation to quantum control of causal order

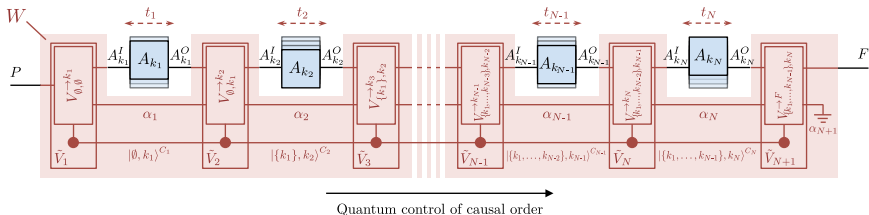
## Example Rehashed



Classical Switch again:

# From Classical to Coherent Control

- Relax the control state to store only *which* operations have been performed, **but not their order**:  $|\mathcal{K}_{n-1}, k_n\rangle^{C_n}$ 
  - Conditioning on  $\mathcal{K} = \{k_1, \dots, k_{n-1}\}$  allows **different orders** to **interfere**
  - Storing full history  $|(k_1, \dots, k_n)\rangle^{C_n}$  is more restrictive and included in this case by using ancillas



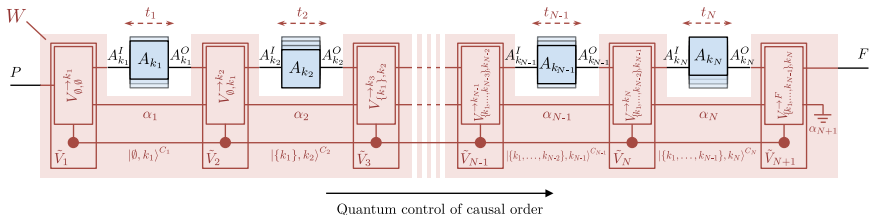
- Operations given by the isometries  $V_n$  with pure CJ representation

$$|\tilde{V}_{n+1}\rangle := \sum_{\substack{\mathcal{K}_{n-1} \\ k_n, k_{n+1}}} |\tilde{V}_{\mathcal{K}_{n-1}, k_n}^{\rightarrow k_{n+1}}\rangle \otimes |\mathcal{K}_{n-1}, k_n\rangle^{C_n} \otimes |\mathcal{K}_n, k_{n+1}\rangle^{C_{n+1}},$$

$$|\tilde{V}_1\rangle := \sum_{k_1} |\tilde{V}_{\emptyset, \emptyset}^{\rightarrow k_1}\rangle \otimes |\emptyset, k_1\rangle^{C_1}, \quad |\tilde{V}_{N+1}\rangle := \sum_{k_N} |\tilde{V}_{\mathcal{N} \setminus \{k_N\}, k_N}^{\rightarrow F}\rangle \otimes |\mathcal{N} \setminus \{k_N\}, k_N\rangle^{C_N}$$

# From Classical to Coherent Control

- Relax the control state to store only *which* operations have been performed, **but not their order**:  $|\mathcal{K}_{n-1}, k_n\rangle^{C_n}$ 
  - Conditioning on  $\mathcal{K} = \{k_1, \dots, k_{n-1}\}$  allows **different orders** to **interfere**
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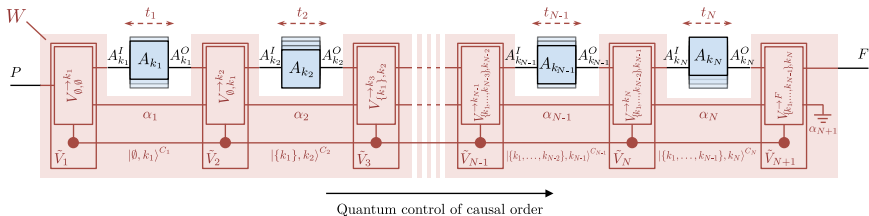


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$$|\tilde{V}_1\rangle := \sum_{k_1} |\tilde{V}_{\emptyset, \emptyset}^{\rightarrow k_1}\rangle \otimes |\emptyset, k_1\rangle^{C_1}, \quad |\tilde{V}_{N+1}\rangle := \sum_{k_N} |\tilde{V}_{\mathcal{N} \setminus \{k_N\}, k_N}^{\rightarrow F}\rangle \otimes |\mathcal{N} \setminus \{k_N\}, k_N\rangle^{C_N}$$

# Coherently (Quantum) Controlled Circuits



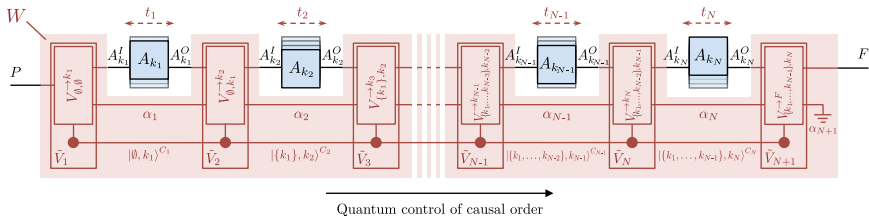
- Each  $V_{\mathcal{K}_{n-1}, k_n}^{I \rightarrow k_{n+1}} : \mathcal{H}^{A_{k_n}^O \alpha_n} \rightarrow \mathcal{H}^{A_{k_{n+1}}^I \alpha_{n+1}}$  embedded in larger space
- Control ensures that **each party applied once and only once**

For input  $|\psi\rangle$ , circuit applies transformation

$$|\tilde{V}_{N+1}\rangle\rangle * |\tilde{A}_N\rangle\rangle * |\tilde{V}_N\rangle\rangle * \cdots * |\tilde{V}_2\rangle\rangle * |\tilde{A}_1\rangle\rangle * |\tilde{V}_1\rangle\rangle * |\psi\rangle \in \mathcal{H}^{F\alpha_{N+1}}$$

with “pure link product”  $|a\rangle^A * |b\rangle^B := \langle\mathbb{1}|^{A \cap B} (|a\rangle \otimes |b\rangle) = \sum_i \langle i, i |^{(A \cap B)^{\otimes 2}} (|a\rangle \otimes |b\rangle)$

# QCs with Quantum Control of Causal Order



- To identify the process matrix, note that for input  $|\psi\rangle$

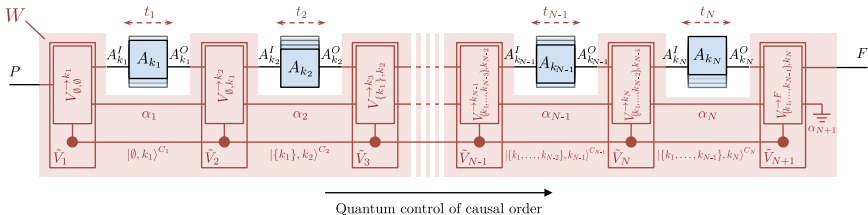
$$\begin{aligned}
& |\tilde{V}_{N+1}\rangle\rangle * |\tilde{A}_N\rangle\rangle * |\tilde{V}_N\rangle\rangle * \cdots * |\tilde{V}_2\rangle\rangle * |\tilde{A}_1\rangle\rangle * |\tilde{V}_1\rangle\rangle * |\psi\rangle \\
&= \sum_{k_1, \dots, k_N} \underbrace{|V_{\emptyset, \emptyset}^{\rightarrow k_1}\rangle\rangle * |V_{\emptyset, k_1}^{\rightarrow k_2}\rangle\rangle * |V_{\{k_1\}, k_2}^{\rightarrow k_3}\rangle\rangle * \cdots * |V_{\{k_1, \dots, k_{N-1}\}, k_N}^{\rightarrow F}\rangle\rangle}_{|w_{(k_1, \dots, k_N, F)}\rangle\rangle} * (|\psi\rangle \otimes |A_1\rangle\rangle \otimes \cdots \otimes |A_N\rangle\rangle)
\end{aligned}$$

## Process matrix of a QC-QC

$$W = \text{Tr}_{\alpha_{N+1}} |w\rangle\langle w|, \quad \text{with} \quad |w\rangle := \sum_{k_1, \dots, k_N} |w_{(k_1, \dots, k_N, F)}\rangle$$



# QC-QC Characterisation



- As for QC-CCs, can characterise such  $W$  with SDP constraints

## Characterisation of circuits with quantum control of causal order

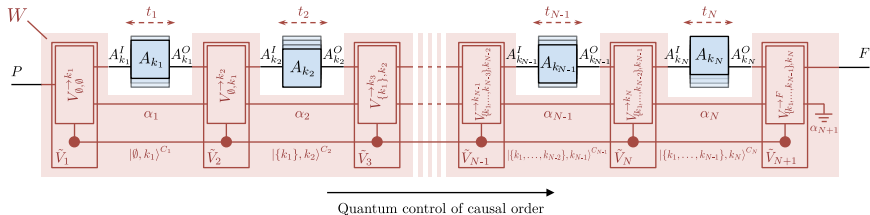
$W$  is the process matrix of a QC-QC iff  $\exists$  PSD matrices  $W_{(\mathcal{K}, \ell)} \in PA_{\mathcal{K}}^{IO} A_{\ell}^I$   
 $\forall \mathcal{K} \subsetneq \mathcal{N}, \ell \in \mathcal{N} \setminus \mathcal{K}$  satisfying

$$\mathrm{Tr}_F W = \sum_{k \in \mathcal{N}} \widetilde{W}_{(\mathcal{N} \setminus \{k\}, k)} \otimes \mathbb{1}^{A_k^O},$$

$$\forall \emptyset \subsetneq \mathcal{K} \subsetneq \mathcal{N}, \sum_{\ell \in \mathcal{N} \setminus \mathcal{K}} \text{Tr}_{A_\ell^I} \widetilde{W}_{(\mathcal{K}, \ell)} = \sum_{k \in \mathcal{K}} \widetilde{W}_{(\mathcal{K} \setminus \{k\}, k)} \otimes \mathbb{1}^{A_k^O},$$

and  $\sum_{\ell \in \mathcal{N}} \text{Tr}_{A_\ell^I} W_{(\emptyset, \ell)} = \mathbb{1}^P$ .

# QC-QC Characterisation

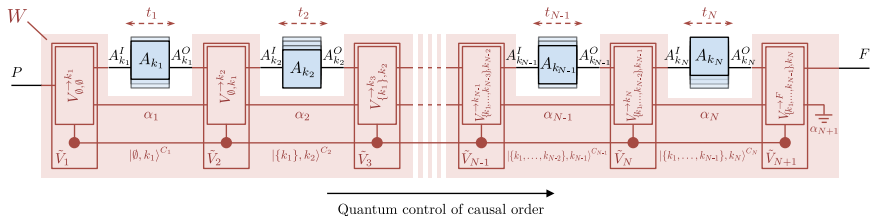


- As for QC-CCs, can characterise such  $W$  with SDP constraints

## Three-operation QC-QC Characterisation

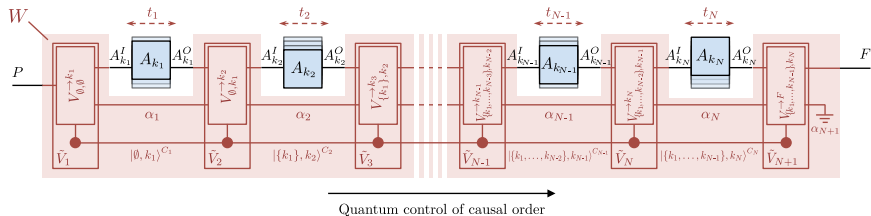
- $\text{Tr}_F W = \widetilde{W}_{(\{B,C\},A)} \otimes \mathbb{1}^{A^O} + \widetilde{W}_{(\{A,C\},B)} \otimes \mathbb{1}^{B^O} + \widetilde{W}_{(\{A,B\},C)} \otimes \mathbb{1}^{C^O}$
- $\text{Tr}_{C^I} \widetilde{W}_{(\{A,B\},C)} = \widetilde{W}_{(\{A\},B)} \otimes \mathbb{1}^{B^O} + \widetilde{W}_{(\{B\},A)} \otimes \mathbb{1}^{A^O}$ , etc.
- $\text{Tr}_{B^I} \widetilde{W}_{(\{A\},B)} + \text{Tr}_{C^I} \widetilde{W}_{(\{A\},C)} = W_{(\{\emptyset\},A)} \otimes \mathbb{1}^{A^O}$ , etc.
- $\text{Tr}_{A^I} W_{(\{\emptyset\},A)} + \text{Tr}_{B^I} W_{(\{\emptyset\},B)} + \text{Tr}_{C^I} W_{(\{\emptyset\},C)} = \mathbb{1}^P$

# QC-QC Summary



- QC-QCs are physically realisable, e.g., with a “quantum router”
- Realisation – in terms of the  $\tilde{V}_n$  – can be effectively obtained from the any  $W$  satisfying the characterisation
  - Can be checked and obtained via SDPs, or witnesses obtained
- Classically controlled circuits are recovered as a special case
  - But QC-QCs can be causally nonseparable in general

## QC-QC Summary

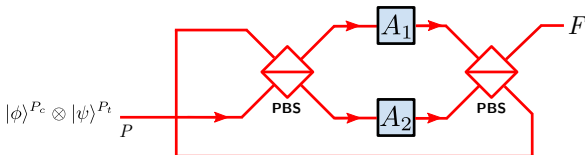


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  - But QC-QCs can be causally nonseparable in general

# Example: Quantum Switch

$N = 2$ ,  $d$ -dimensional target system and 2-dimensional “control”:

$$\mathcal{H}^P = \mathcal{H}^{P_t} \otimes \mathcal{H}^{P_c} \text{ and } \mathcal{H}^F = \mathcal{H}^{F_t} \otimes \mathcal{H}^{F_c}$$



The controlled operations are

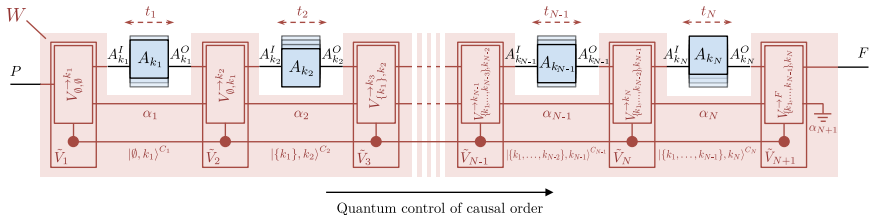
$$|V_{\emptyset, \emptyset}^{\rightarrow k_1}\rangle = |k_1\rangle^{P_c} |\mathbb{1}\rangle^{P_t A_1^I}, \quad |V_{\emptyset, k_1}^{\rightarrow k_2}\rangle = |\mathbb{1}\rangle^{A_{k_1}^O A_{k_2}^I}, \quad |V_{\{k_1\}, k_2}^{\rightarrow F}\rangle = |k_1\rangle^{F_c} |\mathbb{1}\rangle^{A_{k_2}^O F_t}$$

Process vector is then

$$\begin{aligned} |w_s\rangle &:= |w_{(P, A_1, A_2, F)}\rangle + |w_{(P, A_2, A_1, F)}\rangle \\ &= |V_{\emptyset, \emptyset}^{\rightarrow A_1}\rangle * |V_{\emptyset, A_1}^{\rightarrow A_2}\rangle * |V_{\{A_1\}, A_2}^{\rightarrow F}\rangle + |V_{\emptyset, \emptyset}^{\rightarrow A_2}\rangle * |V_{\emptyset, A_2}^{\rightarrow A_1}\rangle * |V_{\{A_2\}, A_1}^{\rightarrow F}\rangle \\ &= |1\rangle^{P_c} |\mathbb{1}\rangle^{P_t A_1^I} |\mathbb{1}\rangle^{A_1^O A_2^I} |\mathbb{1}\rangle^{A_2^O F_t} |1\rangle^{F_c} + |2\rangle^{P_c} |\mathbb{1}\rangle^{P_t A_2^I} |\mathbb{1}\rangle^{A_2^O A_1^I} |\mathbb{1}\rangle^{A_1^O F_t} |2\rangle^{F_c} \end{aligned}$$

Standard four-partite switch recovered as  $W_{\text{switch}} = |w_s\rangle\langle w_s|$

## Example: Quantum $N$ -Switch



# Beyond the Quantum Switch?

- $N$ -partite generalisation of the quantum switch is a QC-QC
  - Essentially the extent of known “interesting” causally nonseparable processes
- Do QC-QCs offer something new, or are they all “equivalent” to the switch?
- Need a better understanding of causally nonseparable resources and free operations
  - Taddei, Nery and Aolita [arXiv:1903.06180]: local operations and controlled non-signalling operations of bipartite processes
  - Composition: Possible compositions severely restricted [Guérin et al., NJP 2019], but can, e.g., concatenate switches, or insert them inside other switches

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# Beyond the Quantum Switch?

However, recall characterisation of QC-QCs

$$\begin{aligned}\mathrm{Tr}_F W &= \sum_{k \in \mathcal{N}} \widetilde{W}_{(\mathcal{N} \setminus \{k\}, k)} \otimes \mathbb{1}^{A_k^O}, \\ \forall \emptyset \subsetneq \mathcal{K} \subsetneq \mathcal{N}, \quad \sum_{\ell \in \mathcal{N} \setminus \mathcal{K}} \mathrm{Tr}_{A_\ell^I} \widetilde{W}_{(\mathcal{K}, \ell)} &= \sum_{k \in \mathcal{K}} \widetilde{W}_{(\mathcal{K} \setminus \{k\}, k)} \otimes \mathbb{1}^{A_k^O}, \\ \text{and} \quad \sum_{\ell \in \mathcal{N}} \mathrm{Tr}_{A_\ell^I} W_{(\emptyset, \ell)} &= \mathbb{1}^P.\end{aligned}$$

- The  $\widetilde{W}_{(\mathcal{N} \setminus \{k\}, k)}$  need not, *a priori*, be valid process matrices
  - $\mathrm{Tr}_F W$  not necessarily a mixture of valid process matrices compatible with fixed last parties, i.e.

$$\mathrm{Tr}_F W \neq q_1 W^{\{A_1, A_2\} \prec A_3} + q_2 W^{\{A_1, A_3\} \prec A_2} + q_3 W^{\{A_2, A_3\} \prec A_1}$$

- Seems like such  $W$  can't be obtained by composing switches
  - Numerically, such processes seem to exist: further study needed to find (and interpret) nice examples
  - In particular, can one obtain new types of advantages with QC-QCs – which, by construction, are realisable

# QC-QCs and Causal Correlations

Can quantum circuits with quantum control of causal order violate causal inequalities?

QC-QC correlations are causal

Let  $W$  be a QC-QC with trivial spaces  $\mathcal{H}^P$  and  $\mathcal{H}^F$ . Then the correlations

$$P(a_1, \dots, a_N | x_1, \dots, x_N) = \text{Tr}[W \cdot (M_{a_1|x_1}^T \otimes \dots \otimes M_{a_N|x_N}^T)]$$

are causal for any instruments  $\{M_{a_i|x_i}\}_{x_i}$ .

- Can noncausal correlations be realised in nature?
  - Would require going beyond this type of generic, coherently controlled circuit
- QC-QCs nevertheless have potential for new advantages arising from indefinite causal order
  - New classes of physically realisable, causally nonseparable processes
  - Use as “quantum super-instruments”, generalising quantum testers, for transformation tasks

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# Summary & Outlook

- Circuit-based (physical) families of superchannels of increasing power
- Quantum circuits with classical control of causal order
  - Coincide with sufficient condition for causal separability
- Quantum circuits with quantum control of causal order
  - Potential new realisable, causally nonseparable, circuits beyond the quantum switch?
  - Do QC-QCs provide new information theoretical advantages?
  - Need for resource theoretical treatment for such processes
  - Are there other classes of physically realisable processes?

[arXiv:1807.10557 + new paper soon]

# Choi Isomorphism and Link Product

- $|\mathbb{1}\rangle\rangle = \sum_i |i\rangle \otimes |i\rangle$  is the “pure Choi isomorphism” of an identity channel
- Pure Choi isomorphism: for an operator  $A$ ,  $|A\rangle\rangle = \mathbb{1} \otimes A |\mathbb{1}\rangle\rangle$
- Mixed Choi isomorphism: for a CP map  $\mathcal{M}$ ,  $M = \mathcal{I} \otimes \mathcal{M}(|\mathbb{1}\rangle\rangle\langle\langle\mathbb{1}|)$
- Inverse Choi isomorphism given by the link product:  $\mathcal{M}(\rho) = M * \rho$ ;  
 $A |\psi\rangle = |A\rangle\rangle * |\psi\rangle$

# Constraints for Process Matrix Validity

Recall the notation:

$${}_X W := (\text{Tr}_X W) \otimes \frac{\mathbb{1}^X}{d_X}, \quad {}_1 W := W, \quad [\sum_X \alpha_X X] W := \sum_X \alpha_X {}_X W,$$

## Space of valid process matrices

$$W \in \mathcal{L}^{\mathcal{N}} \Leftrightarrow \forall \chi \subseteq \mathcal{N}, \chi \neq \emptyset, \prod_{i \in \chi} [1 - A_O^i] A_{IO}^{\mathcal{N} \setminus \chi} W = 0,$$

## Space of valid process compatible with $A$ first

$$W \in \mathcal{L}^{A_k \prec (\mathcal{N} \setminus A_k)}$$

$$\Leftrightarrow [1 - A_O^k] A_{IO}^{\mathcal{N} \setminus k} W = 0 \quad \text{and} \quad \forall \chi \subseteq \mathcal{N} \setminus k, \chi \neq \emptyset, \prod_{i \in \chi} [1 - A_O^i] A_{IO}^{\mathcal{N} \setminus k \setminus \chi} W = 0,$$

# Causal Separability: Necessary Conditions

- Explicit necessary conditions can be obtained by choosing specific CP maps and ancillas at each level of the recursive definition
- Ognyan and Giarmatzi showed how such a choice proves sufficient conditions also necessary in tripartite case
  1.  $\rho$ : maximally entangled state for each pair of parties
  2.  $M_{A_k}: |\Phi^+\rangle\langle\Phi^+|$  – M.E.S. between  $A_k^{IO}$  and half of ancilla between  $A_k$  and some  $A_{k'}$
- “Teleports”  $A_k$ ’s system on  $A_k^{IO}$  to  $A_{k'}^{I'}$

$$\underbrace{W_{(k)}^\rho}_{N\text{-partite, } A_k \text{ first}} \longrightarrow \underbrace{W_{(k)}^{A_k^{IO} \rightarrow A_{k'}^{I'}}}_{(N-1)\text{-partite, formally equivalent to } W_{(k)}} \otimes \rho' := (W_{(k)}^\rho)_{|M_{A_k}}$$

- Any constraints obeyed by  $W_{(k)}^{A_k^{IO} \rightarrow A_{k'}^{I'}}$  must be obeyed by  $W_{(k)}$  once Hilbert spaces relabelled
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An  $N$ -partite  $W^{\text{sep}} \in \mathcal{W}^{\text{sep}}$  must have a decomposition  $W = \sum_{k \in \mathcal{N}} W_{(k)}$  where:

1.  $W_{(k)}$  is a valid process compatible with  $A_k \prec (\mathcal{N} \setminus A_k)$
2. For each  $k' \neq k$ ,  $W_{(k)}^{A_k^{IO} \rightarrow A_{k'}^{I'}}$  is an  $(N - 1)$ -partite causally separable process
  - $\implies$  obeys the necessary conditions for  $(N - 1)$ -partite processes

- Coincides with separable condition for  $N = 3$  [Oreshkov & Giarmatzi, NJP 2016]
- Also reduced 4-partite scenario (no output for  $D$ , c.f. quantum switch)
- Note that decomposition may differ for each  $k'$ 
  - Satisfying these conditions with a unique decomposition would imply the sufficient conditions
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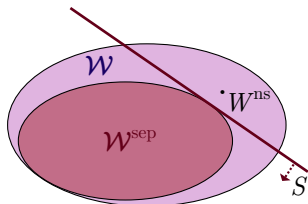
# Witnesses of Causal Nonseparability

## Causally separable process matrix

$$W^{\text{sep}} = q W^{A \prec B} + (1 - q) W^{B \prec A},$$

- Convex cone of (non-normalised) causally separable processes:

$$\mathcal{W}^{\text{sep}} = (\mathcal{P} \cap \mathcal{L}^{A \prec B}) + (\mathcal{P} \cap \mathcal{L}^{B \prec A})$$



## Witness of causal nonseparability

$$\forall W^{\text{ns}} \notin \mathcal{W}^{\text{sep}}, \exists S :$$

$$\text{Tr}[S^T \cdot W^{\text{ns}}] < 0, \text{ and}$$

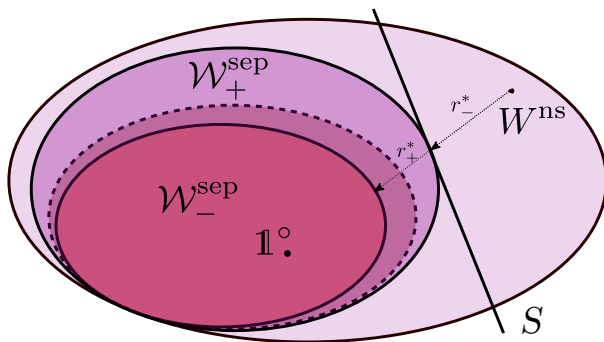
$$\text{Tr}[S^T \cdot W^{\text{sep}}] \geq 0 \quad \forall W^{\text{sep}} \in \mathcal{W}^{\text{sep}}$$

[Araújo et al., NJP 2015; Branciard, Sci. Rep. 2016]

- Witnesses can be efficiently constructed by semidefinite programming (SDP)
- Witnesses can be measured experimentally

# Witnessing Causal Nonseparability

- Both necessary and sufficient conditions define convex cones  $\mathcal{W}_+^{\text{sep}}$ ,  $\mathcal{W}_-^{\text{sep}}$  of (non-normalised) process matrices



- Membership can be tested with SDP
- Dual SDP from necessary condition gives causal witnesses
- So far no numerical evidence that  $\mathcal{W}_-^{\text{sep}} \neq \mathcal{W}_+^{\text{sep}}$ , but...

# Cones $\mathcal{W}^{\text{sep}}$ and $\mathcal{S}$ for tripartite case

Adopt the notation  $\mathcal{L}_X = \{W|_X W = 0\}$ .

$$\begin{aligned}\mathcal{W}^{\text{sep}} = & \mathcal{L}_{[1-A_O]B_{IO}C_{IO}} \cap (\mathcal{P} \cap \mathcal{L}_{[1-B_O]C_{IO}} \cap \mathcal{L}_{[1-C_O]} + \mathcal{P} \cap \mathcal{L}_{[1-C_O]B_{IO}} \cap \mathcal{L}_{[1-B_O]}) \\ & + \mathcal{L}_{[1-B_O]A_{IO}C_{IO}} \cap (\mathcal{P} \cap \mathcal{L}_{[1-A_O]C_{IO}} \cap \mathcal{L}_{[1-C_O]} + \mathcal{P} \cap \mathcal{L}_{[1-C_O]A_{IO}} \cap \mathcal{L}_{[1-A_O]}) \\ & + \mathcal{L}_{[1-C_O]A_{IO}B_{IO}} \cap (\mathcal{P} \cap \mathcal{L}_{[1-A_O]B_{IO}} \cap \mathcal{L}_{[1-B_O]} + \mathcal{P} \cap \mathcal{L}_{[1-B_O]A_{IO}} \cap \mathcal{L}_{[1-A_O]}),\end{aligned}$$

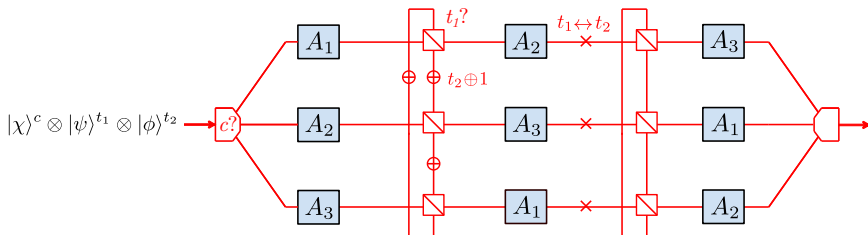
$$\begin{aligned}\mathcal{S} = & \left( \mathcal{L}_{[1-A_O]B_{IO}C_{IO}} + (\mathcal{P} + \mathcal{L}_{[1-B_O]C_{IO}} + \mathcal{L}_{[1-C_O]}) \cap (\mathcal{P} + \mathcal{L}_{[1-C_O]B_{IO}} + \mathcal{L}_{[1-B_O]}) \right) \\ & \cap \left( \mathcal{L}_{[1-B_O]A_{IO}C_{IO}} + (\mathcal{P} + \mathcal{L}_{[1-A_O]C_{IO}} + \mathcal{L}_{[1-C_O]}) \cap (\mathcal{P} + \mathcal{L}_{[1-C_O]A_{IO}} + \mathcal{L}_{[1-A_O]}) \right) \\ & \cap \left( \mathcal{L}_{[1-C_O]A_{IO}B_{IO}} + (\mathcal{P} + \mathcal{L}_{[1-A_O]B_{IO}} + \mathcal{L}_{[1-B_O]}) \cap (\mathcal{P} + \mathcal{L}_{[1-B_O]A_{IO}} + \mathcal{L}_{[1-A_O]}) \right).\end{aligned}$$

# Example: New type of 3-operation QC-QC

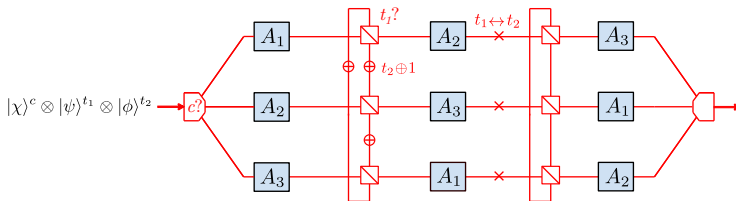
$N = 3$ , two qubit “targets” and 3-dimensional “control”

- Initial control state  $|k_1\rangle^{P_c}$  determines first party
- First party acts on first target qubit  $|\psi\rangle^{P_{t_1}}$
- Output of  $A_{k_1}$  determines (dynamically, coherently)  $k_2$  and conditions a flip on second target  $|\phi\rangle^{P_{t_2}}$ , which is swapped to become “active” target *after*  $A_{k_2}$
- $A_{k_3}$  then acts on this second target qubit

Can be represented in an “unravelled” form as:



# Example: New type of 3-operation QC-QC



Controlled operations can be written

$$|V_{\emptyset, \emptyset}^{\rightarrow k_1}\rangle\rangle = |k_1\rangle^{P_c} |\mathbb{1}\rangle\rangle^{P_{t_1} A_{k_1}^I} |\mathbb{1}\rangle\rangle^{P_{t_2} \alpha_1}$$

$$|V_{\emptyset, k_1}^{\rightarrow k_2}\rangle\rangle = |0\rangle^{A_{k_1}^O} |0\rangle^{A_{k_2}^I} |\mathbb{1}\rangle\rangle^{\alpha_1 \alpha_2}, \text{ and } |V_{\emptyset, k_1}^{\rightarrow k'_2}\rangle\rangle = |1\rangle^{A_{k_1}^O} |1\rangle^{A_{k'_2}^I} |X\rangle\rangle^{\alpha_1 \alpha_2} \text{ for } k_2 \neq k'_2$$

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$$|V_{\{k_1, k_2\}, k_3}^{\rightarrow F}\rangle\rangle = |k_3\rangle^{F_c} |\mathbb{1}\rangle\rangle^{A_{k_3}^O F_{t_2}} |\mathbb{1}\rangle\rangle^{\alpha_3 F_{t_1}}$$

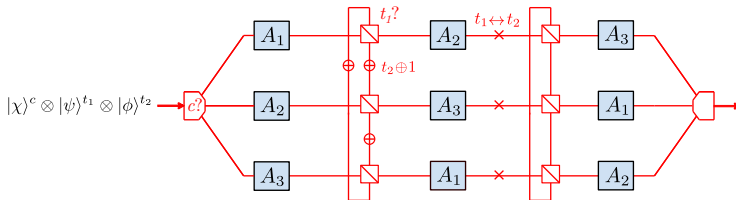
giving (with cyclic permutations for  $k_1 = 2, 3$ )

$$|w_{(P, A_1, A_2, A_3, F)}\rangle\rangle = |1\rangle^{P_c} |\mathbb{1}\rangle\rangle^{P_{t_1} A_1^I} |00\rangle^{A_1^O A_2^I} |\mathbb{1}\rangle\rangle^{A_2^O F_{t_1}} |\mathbb{1}\rangle\rangle^{P_{t_2} A_3^I} |\mathbb{1}\rangle\rangle^{A_3^O F_{t_2}} |3\rangle^{F_c}$$

$$|w_{(P, A_1, A_3, A_2, F)}\rangle\rangle = |1\rangle^{P_c} |\mathbb{1}\rangle\rangle^{P_{t_1} A_1^I} |11\rangle^{A_1^O A_3^I} |\mathbb{1}\rangle\rangle^{A_3^O F_{t_1}} |X\rangle\rangle^{P_{t_2} A_2^I} |\mathbb{1}\rangle\rangle^{A_2^O F_{t_2}} |2\rangle^{F_c}$$



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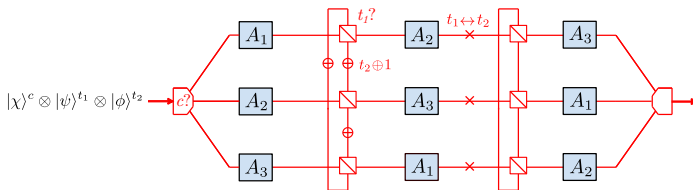
$$|V_{\{k_1, k_2\}, k_3}^{\rightarrow F}\rangle\rangle = |k_3\rangle^{F_c} |\mathbb{1}\rangle\rangle^{A_{k_3}^O F_{t_2}} |\mathbb{1}\rangle\rangle^{\alpha_3 F_{t_1}}$$

giving (with cyclic permutations for  $k_1 = 2, 3$ )

$$|w_{(P, A_1, A_2, A_3, F)}\rangle\rangle = |1\rangle^{P_c} |\mathbb{1}\rangle\rangle^{P_{t_1} A_1^I} |00\rangle^{A_1^O A_2^I} |\mathbb{1}\rangle\rangle^{A_2^O F_{t_1}} |\mathbb{1}\rangle\rangle^{P_{t_2} A_3^I} |\mathbb{1}\rangle\rangle^{A_3^O F_{t_2}} |3\rangle^{F_c}$$

$$|w_{(P, A_1, A_3, A_2, F)}\rangle\rangle = |1\rangle^{P_c} |\mathbb{1}\rangle\rangle^{P_{t_1} A_1^I} |11\rangle^{A_1^O A_3^I} |\mathbb{1}\rangle\rangle^{A_3^O F_{t_1}} |X\rangle\rangle^{P_{t_2} A_2^I} |\mathbb{1}\rangle\rangle^{A_2^O F_{t_2}} |2\rangle^{F_c}$$

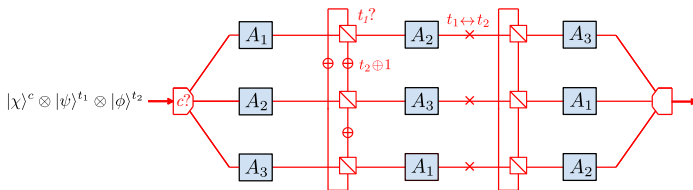
# Example: New type of 3-operation QC-QC



As before, define  $W = |w\rangle\rangle\langle\langle w|$  with  $|w\rangle\rangle = \sum_{k_1, k_2, k_3} |w_{(P, k_1, k_2, k_3, F)}\rangle\rangle$

- $W$  is easily seen to be causally nonseparable
  - Pure process, not compatible with any of  $A$ ,  $B$  or  $C$  being first (after  $P$ )
- Appears qualitatively different to the quantum switch
  - Two target qubits, one of which is also used to control the causal order
- But how to prove  $W$  is fundamentally inequivalent to quantum switch?
  - Could imagine composing switches, using control of one as target for another, etc.
  - Need a more complete resource theory, e.g. generalising Taddei, Nery and Aolita's proposal for bipartite processes [arXiv:1903.06180]

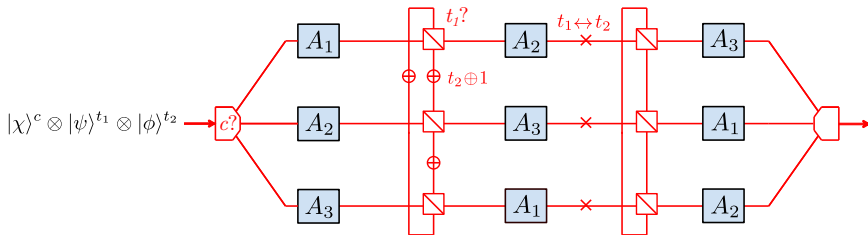
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# Example: New type of 3-operation QC-QC



- Crucial difference:  $\text{Tr}_F W$  is causally nonseparable and cannot be written as a mixture of valid process matrices with fixed last parties, i.e.

$$\text{Tr}_F W \neq q_1 W^{\{A_1, A_2\} \prec A_3} + q_2 W^{\{A_1, A_3\} \prec A_2} + q_3 W^{\{A_2, A_3\} \prec A_1}$$

- Recall characterisation:  $\text{Tr}_F W = \sum_{k \in \mathcal{N}} \widetilde{W}_{(\mathcal{N} \setminus \{k\}, k)} \otimes \mathbb{1}^{A_k^O}$ 
  - the  $\widetilde{W}_{(\mathcal{N} \setminus \{k\}, k)}$  need not be valid process matrices
- Seems like no composition of quantum switches could give this property!