

ComPer 2025: Workshop on Computational Topology

October 10 - 14, University at Albany, State University of New York

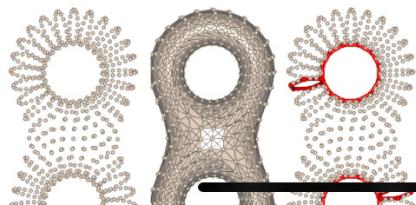


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Computing the zero-dimensional homology over rooted trees

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In some applications of topological data analysis one might consider only zero-dimensional homology. The zero-dimensional homology of a space is conceptually and computationally simpler than higher dimensional homology since it can be characterized as the linearization of the set of path-connected components of the space. In particular, the zero-dimensional homology modules of a filtration indexed by a poset P will in general form a proper subcategory of $\text{rep } P$ (to the contrary, $\text{add } \text{Im } H_{>i}$ is dense for i greater than zero). We show that when the indexing poset P is a rooted tree, only finitely many indecomposables can be obtained as zero-dimensional homology of a filtration. In fact, all these indecomposables are reduced rooted tree modules in the sense of Kinser. Moreover, we give an algorithm for the efficient decomposition of rooted tree modules, generalizing the so-called elder rule used in TDA for linear quivers. This is a report on joint work with Riju Bindua and Luis Scoccola.



1) Setting: k field

\mathcal{P} finite poset

$\text{rep } \mathcal{P} = \{M: \mathcal{P} \rightarrow \text{vec } k\}$

functor

TDA

$$X = (X_a)_{a \in \mathcal{P}}$$

i-th homology

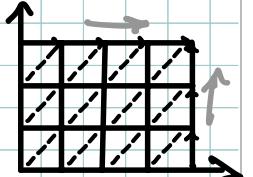


(persistence module) $M = H_i X \in \text{rep } \mathcal{P}$

Example:

\mathcal{P} 2d grid

(\Rightarrow multiparameter persistence)



\mathcal{P} -filtered top. space :

$$X_a \subseteq X_b \text{ for } a \leq b$$

$$H_i X_a \rightarrow H_i X_b \text{ linear map}$$

Theorem [Brodzki, Burfitt, Pirashvili: 2020]

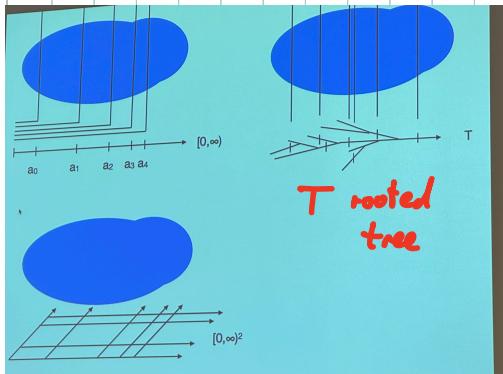
$\text{Im } H_i$ is dense in $\text{rep } \mathcal{P}$ for $i \geq 1$, k perfect field

2) Motivation:

Not true for H_0 : $\text{Im } H_0$ smaller than $\text{rep } P$

$\text{rep } P$ is wild in general

Question : Does $\text{Im } H_0$ become
manageable when restricting
to some tree $T \subseteq P$?



3) Result:

Theorem [Bridgeman, B., Scoccola 2024]

P rooted tree

$\Leftrightarrow \text{Hasse}(P)$ is tree

with unique

maximum

If $P = T$ is a rooted tree

then $| \text{ind}(\text{add}(\overline{T}_m H_0)) | < \infty$ ' H_0 is rep-finite'

Moreover, given a T -filtered graph G ,

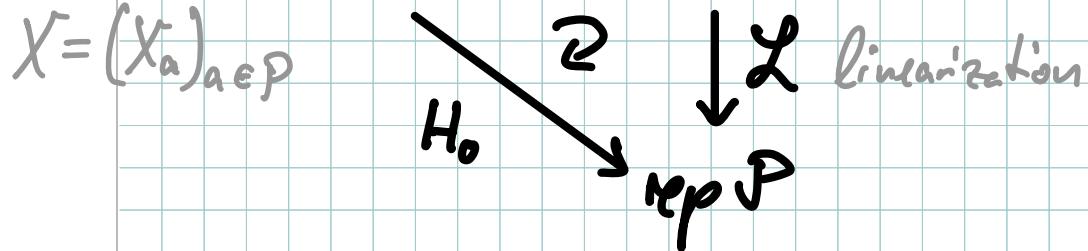
can compute $H_0(G) \cong U_1 \oplus \dots \oplus U_r$ in $O(|G|^2)$ time

4) Why rooted trees?

Thus [classic]

$$\pi_0(X) = \{\text{connected comp}\}$$

P -filt. top $\xrightarrow{\pi_0}$ set P



H_0 factors through sets

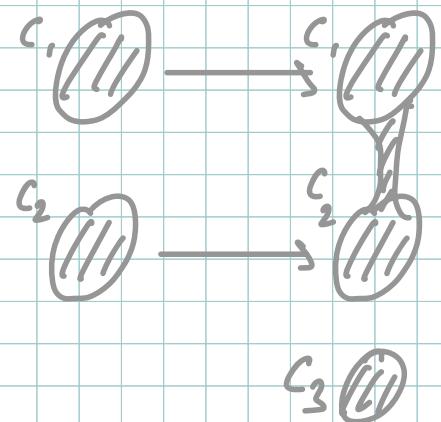
$$a \leq b \text{ in } P$$

$$\pi_0 X_a \longrightarrow \pi_0 X_b$$

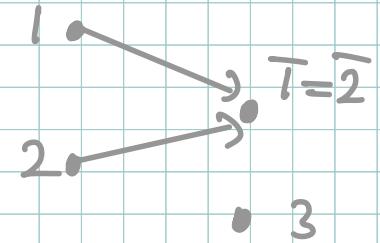
function

between sets

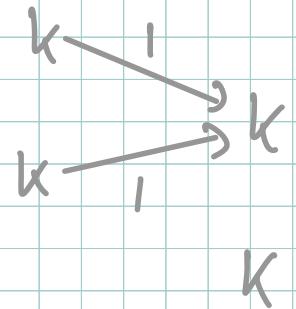
$$X_a \subseteq X_b$$



$$\pi_0(X)$$



$$H_0(X)$$



4) Why rooted trees?

Question: Which indecomposable representations do we get?

$$P = a \rightarrow b$$

$$C \begin{smallmatrix} \textcircled{1} \\ \textcircled{2} \end{smallmatrix} \longrightarrow C \begin{smallmatrix} \textcircled{1} \\ \textcircled{2} \end{smallmatrix}$$

$$C \begin{smallmatrix} \textcircled{1} \\ \textcircled{2} \end{smallmatrix}$$

$$\left. \begin{array}{l} K \xrightarrow{1} K \\ 0 \xrightarrow{0} K \end{array} \right\} \in \text{Im } H_0$$

$$\emptyset \neq X_a \subseteq X_b, \quad X_b \neq \emptyset : \quad K \xrightarrow{0} 0 \notin \text{Im } H_0$$

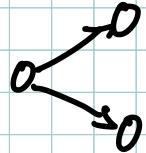
However, in $\text{Im } H_0$ we have

$$L \left(\begin{smallmatrix} 0 & 1 \\ 0 & 0 \end{smallmatrix} \right) = K^2 \xrightarrow{\left[\begin{smallmatrix} 1 & 1 \\ 0 & 1 \end{smallmatrix} \right]} K \cong K \xrightarrow{1} K \oplus K \xrightarrow{0} 0$$

So $K \xrightarrow{0} 0 \in \text{add Im } H_0$

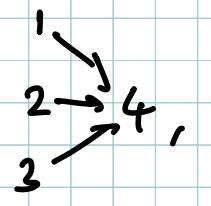
4) Why rooted trees? Which indecomposables do we get?

Observation:



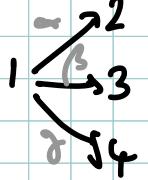
is not in Set^P ,

hence $K \xrightarrow{\begin{bmatrix} 1 \\ 1 \end{bmatrix}} K^2 \notin \text{Im } H_0$.

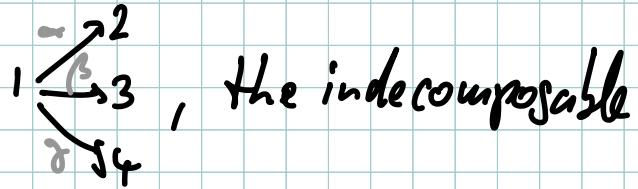
\Rightarrow For $T =$ 

rooted tree, the indecomposable

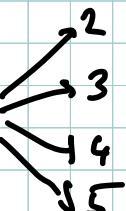
$\notin \text{add Im } H_0$

But For $T =$ 

not rooted tree



$\in \text{add Im } H_0$

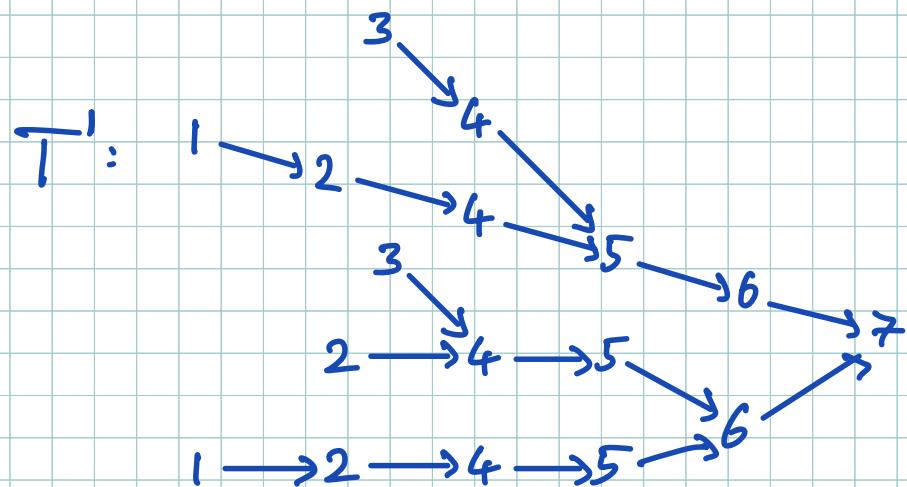
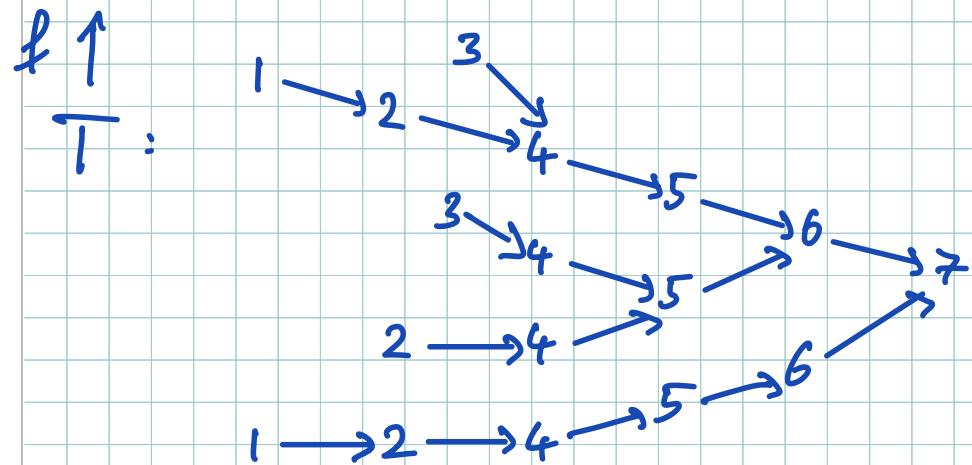
In fact, for $T =$ 

$|\text{ind}(\text{add Im } H_0)| = \infty$



5) Tree Modules kT :

$P:$ $1 \rightarrow 2 \rightarrow 4 \rightarrow 5 \rightarrow 6 \rightarrow 7$



Examples

$kT:$ k^2

$$k^2 \xrightarrow{\quad} k^3 \xrightarrow{\quad} k^4 \xrightarrow{\quad} k^3 \xrightarrow{\begin{bmatrix} 1 & 1 & 0 \\ 0 & 0 & 1 \end{bmatrix}} k^2 \xrightarrow{\quad} k$$

T is a rooted tree / P :

\exists graph maps $f: T \rightarrow P$

$kT := f^* L(T) \in \text{rep } P$

In fact $kT \in \sum_n H_0$

Test : Which of kT, kT'
is indecomposable?

6) Rooted Tree Modules

Prop [BBS] Let P be a rooted tree. Then

$$\text{add } \mathbb{T}_m H_0 = \text{add} \left(kT \mid T \xrightarrow{k} P \text{ rooted tree}/P \right)$$

\in rooted tree module

Question : When is kT indecomposable?

Ringel : no answer (1995; study tree modules)

Crawley-Boevey : In case $T = P$, kT is always indecomposable

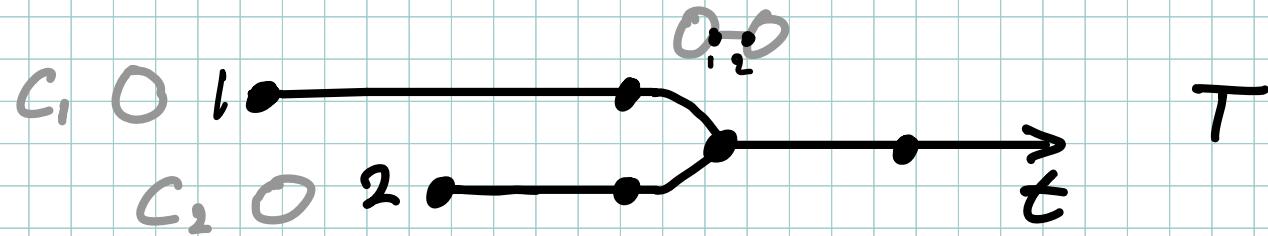
Kinser [2010] : kT indecomposable $\Leftrightarrow T$ reduced

Def: T reduced if ... (recursive def of \leq on reduced trees used to define
when T is reduced, recursively)

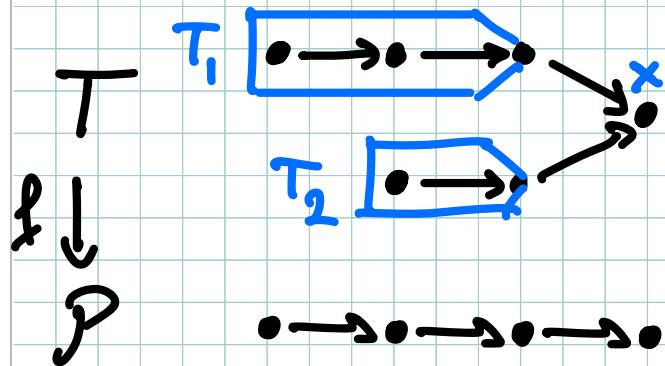
7) The elder rule

TDA

$$\mathcal{P} = \mathbb{R}_{\geq 0}$$



elder rule: When 2 components merge, kill the younger one!



T_1, T_2 subtrees merging at x
 $T = \text{merge}(T_1, T_2)$

\exists morphism $T_2 \rightarrow T_1$
(as trees over \mathcal{P})

$$kT \cong kT_2 \oplus K(T_1 \rightarrow x)$$

in rep \mathcal{P}

8) Generalized elder rule

Thm [BBS] Let P be a rooted tree,

and $\overline{T} = \text{merge}(T_1, \dots, T_e)$ a rooted tree over P .

(a) If \exists morphism $T_2 \rightarrow T_1$, then $K\overline{T} \cong K\overline{T}_2 \oplus K(\overline{T} \setminus \overline{T}_2)$
(as trees over P) in rep \mathcal{P}

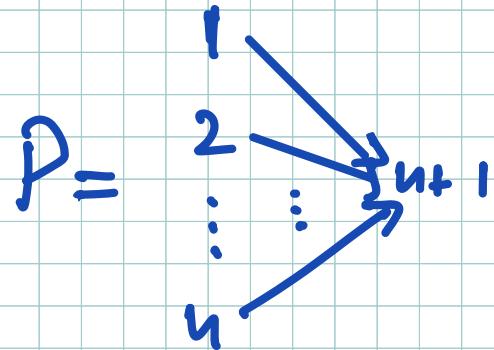
(b) $K\overline{T}$ indecomposable $\xrightleftharpoons[\text{Kinser}]{}$ \overline{T} reduced
 $\Leftrightarrow \text{Hom}_{\text{tree}/P}^{(T, \overline{T})} = \{\text{id}_{\overline{T}}\}$

(c) $\exists O(|\overline{T}|^2)$ -time algorithm decomposing $K\overline{T}$

~ examples over E_7 !

Examples :

$$S = \{s_1, \dots, s_j\} \subseteq \{1, \dots, n\}$$



Then $T_S = \begin{matrix} s_1 \\ \vdots \\ s_j \end{matrix} \xrightarrow{\quad} u+1$ is reduced,

$$\text{so } |\text{ind}(\text{add}(\overline{T_m} H_0))| \geq 2^n$$

$$P: 1 \xrightarrow{} 2 \xrightarrow{} \dots \xrightarrow{} n$$

\overline{T}_1 with root at $n-1$, or empty

$$\overline{T} = \text{merge}(\overline{T}_1)$$

$n-1$

+

|

trees with root n

g) Finiteness

If P is a rooted tree
then $| \text{ind}(\text{add}(\bar{T} \cap H_0)) | < \infty$ ' H_0 is rep-finite'

Proof: Let T be a reduced rooted tree over P .

Write $\bar{T} = \text{merge}(T_1, \dots, T_e)$. Then each T_i is reduced (otherwise, T_i would split off a summand).

By induction, there are only finitely many possibilities for each T_i , and Q must be bounded by the product of these numbers (otherwise there are repetitions, hence morphisms)

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