

# Pruning vineyards

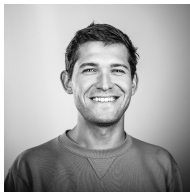
Updating barcodes and representative cycles by removing (and inserting) simplices



arxiv:2312.03925  
accompanying paper



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SUNY Albany



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bitbucket:phat-sirup  
accompanying code

October 14, 2025 / ComPer 2025: Workshop on Computational Topology  
Slides online at: [jlazovskis.com/talks](https://jlazovskis.com/talks)

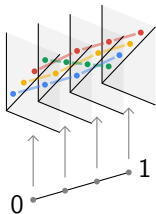
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Funding for JL: Activity 1.1.1.9 "Post-doctoral Research" of the Specific Objective 1.1.1 "Strengthening research and innovative capacities and introduction of advanced technologies in the common R&D system" of the European Union's Cohesion Policy Programme for 2021-2027 research application No 1.1.1.9/LZP/1/24/125 "Efficient topological signatures for representation learning in medical imaging"

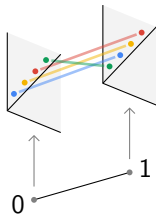


# Vineyards

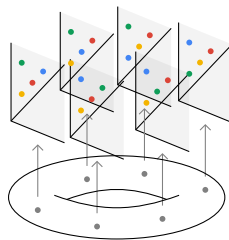
**Vines** and **vineyards** were introduced to capture a change in the persistence diagram coming from a change in the underlying data. Has since been generalized in different ways.



*vines and vineyards:*  
Cohen–Steiner, Edelsbrunner,  
Morozov (2006)



*moves:* Busaryev, Dey,  
Wang (2010)  
*move schedules:*  
Piekenbrock, Perea (2024)



*persistence diagram bundles:*  
Hickok (2022)



*vineyard modules:*  
Turner (2023)

These approaches are about **swapping** the order of simplices in a filtration.  
What about other possible changes, like **inserting** or **removing** simplices?

# Dynamic persistence

## Motivation:

- ▶ Potential to reuse costly computations
- ▶ Suited to non-static data (changing resolution, zooming in, moving around)
- ▶ Get a better understanding of class representatives

**Computational reality:** Given a simplicial complex with  $n$  simplices:

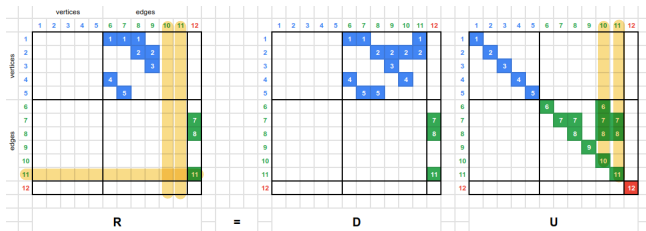
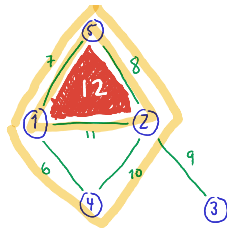
- ▶  $O(n^3)$  complexity to compute PH with standard algorithm (mat.mul. by Morozov, Skraba)
- ▶ Massive speedups with clever code, homological observations, dimensional restrictions
- ▶ Most **dynamic data** is read into / out of the computer as a sequence of static data points (changes inferred, not given directly). Makes sense to rerun everything with new data.

**But still:** Persistence cycles (and so persistence diagrams) are hard to track this way.

**Would be great if:** There was a PH implementation providing an **efficient, user-friendly** interface with **access to PD** at intermediate steps.

# Example

Consider a simplex-wise filtration on a simplicial complex. Consider 1-cycles and their reps.



**Inserting:** How are 1-classes affected if edge  $\{4, 3\}$  is inserted at position 10?

- One non-trivial 1-class appears. Is representative  $\{4, 2, 3\}$  or  $\{4, 1, 5, 2, 3\}$ ?

**Removing:** How are 1-classes affected if edge 8 is removed?

- One trivial 1-class disappears. Is other non-trivial representative different or the same?

# Inserting simplices: easy

Similarly to the **standard barcode algorithm**, add matching pivot columns left-to-right.

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**Algorithm 1: SBA - STANDARD BARCODE ALGORITHM** [ELZ00]

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**Input:** Boundary matrix  $D$ , identity matrix  $I$

**Output:** Reduced matrix  $R$ , operations matrix  $U$

```
1  $R \leftarrow D$ 
2  $U \leftarrow I$ 
3 for  $j = 1, \dots, n$ 
4   while there exists  $j' < j$  for which  $\text{piv}(R[j']) = \text{piv}(R[j]) \neq 0$ 
5     add  $R[j']$  to  $R[j]$ 
6     add  $U[j']$  to  $U[j]$ 
7 return  $R, U$ 
```

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**Complexity:** for one insertion is  $O(n^2)$ , vs  $O((n+1)^3)$  for computing from scratch.

**Careful:** The representative cycles do not match up with computation from scratch.

This is still work in progress.

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**Algorithm 1: INSERTING A SIMPLEX**

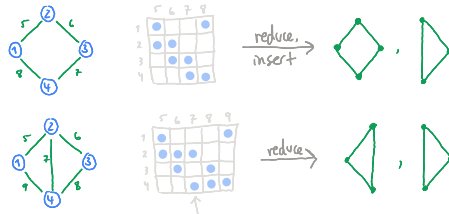
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**Input:** Reduced matrix  $R$ , operations matrix  $U$ , list of simplices to insert  $L$

**Output:** Updated reduced matrix  $P$ , updated operations matrix  $V$

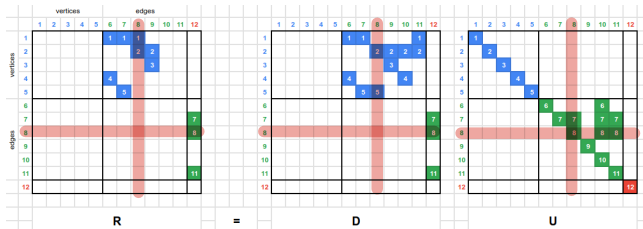
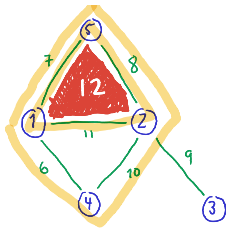
```
1  $P \leftarrow R$ 
2  $V \leftarrow U$ 
3 for  $\sigma_j \in L$ 
4   insert row and column at index  $j$  in  $R$ , with nonzero values at  $R[i, j]$ ,  $\sigma_i \in \partial \sigma_j$ 
5   insert row and column at index  $j$  in  $U$ , with one nonzero value at  $U[j, j]$ 
6    $r \leftarrow \text{piv}(\sigma_j)$ 
7   while there exists  $k \neq j$  with  $R[r, j] = R[r, k] \neq 0$ 
8     if  $k < j$ 
9       add column  $k$  to column  $j$  in  $P$  and in  $V$ 
10       $r \leftarrow \text{piv}(P[j])$ 
11     else if  $j < k$ 
12       add column  $j$  to column  $k$  in  $P$  and in  $V$ 
13        $r \leftarrow \text{piv}(P[k])$ 
14 return  $P, V$ 
```

---



# Removing simplices: motivating example

Back to example: To remove simplex 8 (after 12), “undo” the column ops it was involved in.



1. Check row 8 in matrix  $U$  to see in which representative cycles it appears.

► Add column 8 to columns 10, 11 in  $R$ ,  $U$ .

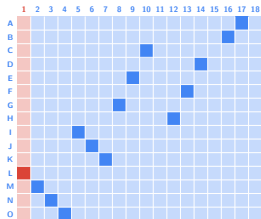
2. Ensure that  $R$  is reduced.

► Add column 10 to column 11 in  $R$ ,  $U$ .

**Observation:** Column 8 was “added twice” to column 11, cancelling itself out in the end. Only one column operation (8 to 10) was necessary.

# Removing simplices: in general

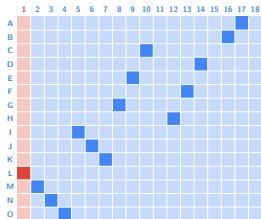
- ▶ Each “affected column” needs only one column addition to get it to where it should be.
- ▶ The column which should be added to it is determined by the left-most column with a pivot above its pivot (or below, if it itself is a left-most pivot column).



Consider the submatrix of  $R$  that contains only the columns “affected” by the removal of column 1.

# Removing simplices: in general

- ▶ Each “affected column” needs only one column addition to get it to where it should be.
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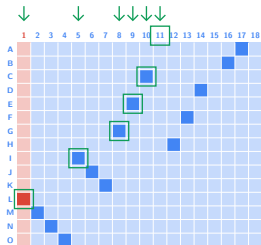
Consider the submatrix of  $R$  that contains only the columns “affected” by the removal of column 1.

**Step 1:** Identify the “left-most highest” pivots.



# Removing simplices: in general

- ▶ Each “affected column” needs only one column addition to get it to where it should be.
- ▶ The column which should be added to it is determined by the left-most column with a pivot above its pivot (or below, if it itself is a left-most pivot column).

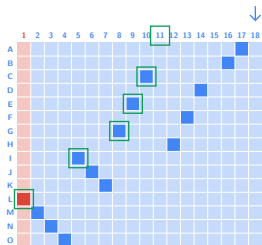


Consider the submatrix of  $R$  that contains only the columns “affected” by the removal of column 1.

**Step 1:** Identify the “left-most highest” pivots.

# Removing simplices: in general

- ▶ Each “affected column” needs only one column addition to get it to where it should be.
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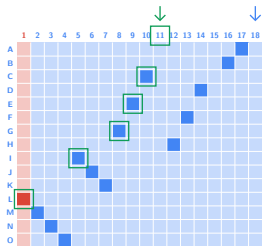
Consider the submatrix of  $R$  that contains only the columns “affected” by the removal of column 1.

**Step 1:** Identify the “left-most highest” pivots.

**Step 2:** Ensure column 18 is reduced:

# Removing simplices: in general

- ▶ Each “affected column” needs only one column addition to get it to where it should be.
- ▶ The column which should be added to it is determined by the left-most column with a pivot above its pivot (or below, if it itself is a left-most pivot column).



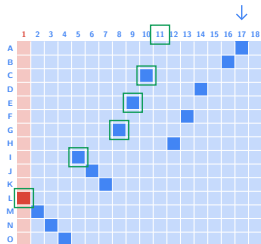
Consider the submatrix of  $R$  that contains only the columns “affected” by the removal of column 1.

**Step 1:** Identify the “left-most highest” pivots.

**Step 2:** Ensure column 18 is reduced: add column 11

# Removing simplices: in general

- ▶ Each “affected column” needs only one column addition to get it to where it should be.
- ▶ The column which should be added to it is determined by the left-most column with a pivot above its pivot (or below, if it itself is a left-most pivot column).



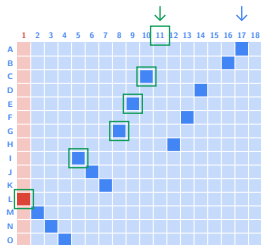
Consider the submatrix of  $R$  that contains only the columns “affected” by the removal of column 1.

**Step 1:** Identify the “left-most highest” pivots.

**Step 2:** Ensure column 17 is reduced:

# Removing simplices: in general

- ▶ Each “affected column” needs only one column addition to get it to where it should be.
- ▶ The column which should be added to it is determined by the left-most column with a pivot above its pivot (or below, if it itself is a left-most pivot column).



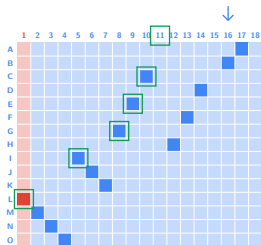
Consider the submatrix of  $R$  that contains only the columns “affected” by the removal of column 1.

**Step 1:** Identify the “left-most highest” pivots.

**Step 2:** Ensure column 17 is reduced: add column 11

# Removing simplices: in general

- ▶ Each “affected column” needs only one column addition to get it to where it should be.
- ▶ The column which should be added to it is determined by the left-most column with a pivot above its pivot (or below, if it itself is a left-most pivot column).



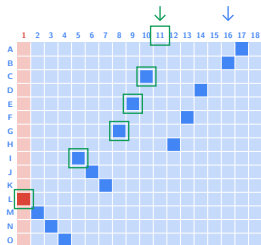
Consider the submatrix of  $R$  that contains only the columns “affected” by the removal of column 1.

**Step 1:** Identify the “left-most highest” pivots.

**Step 2:** Ensure column 16 is reduced:

# Removing simplices: in general

- ▶ Each “affected column” needs only one column addition to get it to where it should be.
- ▶ The column which should be added to it is determined by the left-most column with a pivot above its pivot (or below, if it itself is a left-most pivot column).



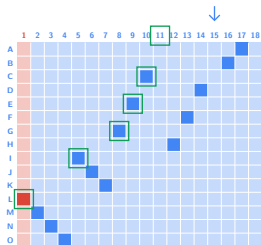
Consider the submatrix of  $R$  that contains only the columns “affected” by the removal of column 1.

**Step 1:** Identify the “left-most highest” pivots.

**Step 2:** Ensure column 16 is reduced: add column 11

# Removing simplices: in general

- ▶ Each “affected column” needs only one column addition to get it to where it should be.
- ▶ The column which should be added to it is determined by the left-most column with a pivot above its pivot (or below, if it itself is a left-most pivot column).



Consider the submatrix of  $R$  that contains only the columns “affected” by the removal of column 1.

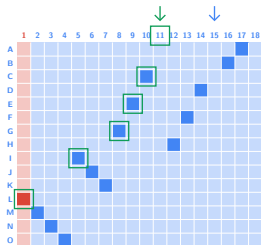
**Step 1:** Identify the “left-most highest” pivots.

**Step 2:** Ensure column 15 is reduced:



# Removing simplices: in general

- ▶ Each “affected column” needs only one column addition to get it to where it should be.
- ▶ The column which should be added to it is determined by the left-most column with a pivot above its pivot (or below, if it itself is a left-most pivot column).



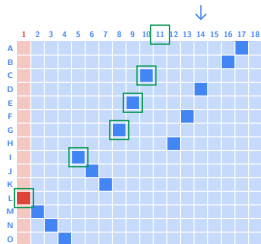
Consider the submatrix of  $R$  that contains only the columns “affected” by the removal of column 1.

**Step 1:** Identify the “left-most highest” pivots.

**Step 2:** Ensure column 15 is reduced: add column 11

# Removing simplices: in general

- ▶ Each “affected column” needs only one column addition to get it to where it should be.
- ▶ The column which should be added to it is determined by the left-most column with a pivot above its pivot (or below, if it itself is a left-most pivot column).



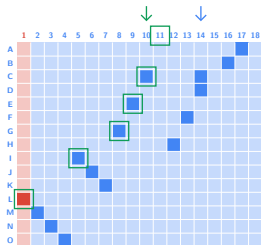
Consider the submatrix of  $R$  that contains only the columns “affected” by the removal of column 1.

**Step 1:** Identify the “left-most highest” pivots.

**Step 2:** Ensure column 14 is reduced:

# Removing simplices: in general

- ▶ Each “affected column” needs only one column addition to get it to where it should be.
- ▶ The column which should be added to it is determined by the left-most column with a pivot above its pivot (or below, if it itself is a left-most pivot column).



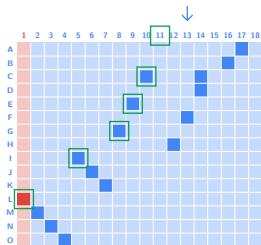
Consider the submatrix of  $R$  that contains only the columns “affected” by the removal of column 1.

**Step 1:** Identify the “left-most highest” pivots.

**Step 2:** Ensure column 14 is reduced: add column 10

# Removing simplices: in general

- ▶ Each “affected column” needs only one column addition to get it to where it should be.
- ▶ The column which should be added to it is determined by the left-most column with a pivot above its pivot (or below, if it itself is a left-most pivot column).



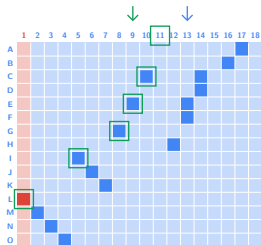
Consider the submatrix of  $R$  that contains only the columns “affected” by the removal of column 1.

**Step 1:** Identify the “left-most highest” pivots.

**Step 2:** Ensure column 13 is reduced:

# Removing simplices: in general

- ▶ Each “affected column” needs only one column addition to get it to where it should be.
- ▶ The column which should be added to it is determined by the left-most column with a pivot above its pivot (or below, if it itself is a left-most pivot column).



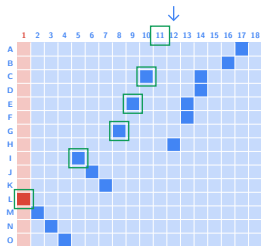
Consider the submatrix of  $R$  that contains only the columns “affected” by the removal of column 1.

**Step 1:** Identify the “left-most highest” pivots.

**Step 2:** Ensure column 13 is reduced: add column 9

# Removing simplices: in general

- ▶ Each “affected column” needs only one column addition to get it to where it should be.
- ▶ The column which should be added to it is determined by the left-most column with a pivot above its pivot (or below, if it itself is a left-most pivot column).



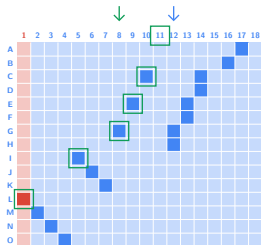
Consider the submatrix of  $R$  that contains only the columns “affected” by the removal of column 1.

**Step 1:** Identify the “left-most highest” pivots.

**Step 2:** Ensure column 12 is reduced:

# Removing simplices: in general

- ▶ Each “affected column” needs only one column addition to get it to where it should be.
- ▶ The column which should be added to it is determined by the left-most column with a pivot above its pivot (or below, if it itself is a left-most pivot column).



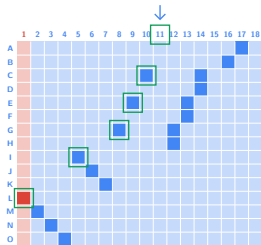
Consider the submatrix of  $R$  that contains only the columns “affected” by the removal of column 1.

**Step 1:** Identify the “left-most highest” pivots.

**Step 2:** Ensure column 12 is reduced: add column 8

# Removing simplices: in general

- ▶ Each “affected column” needs only one column addition to get it to where it should be.
- ▶ The column which should be added to it is determined by the left-most column with a pivot above its pivot (or below, if it itself is a left-most pivot column).



Consider the submatrix of  $R$  that contains only the columns “affected” by the removal of column 1.

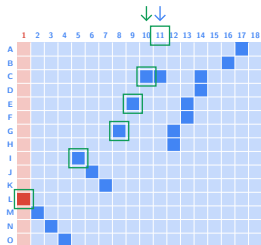
**Step 1:** Identify the “left-most highest” pivots.

**Step 2:** Ensure column 11 is reduced:



# Removing simplices: in general

- ▶ Each “affected column” needs only one column addition to get it to where it should be.
- ▶ The column which should be added to it is determined by the left-most column with a pivot above its pivot (or below, if it itself is a left-most pivot column).



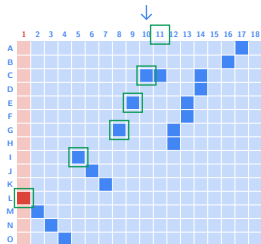
Consider the submatrix of  $R$  that contains only the columns “affected” by the removal of column 1.

**Step 1:** Identify the “left-most highest” pivots.

**Step 2:** Ensure column 11 is reduced: add column 10

# Removing simplices: in general

- ▶ Each “affected column” needs only one column addition to get it to where it should be.
- ▶ The column which should be added to it is determined by the left-most column with a pivot above its pivot (or below, if it itself is a left-most pivot column).



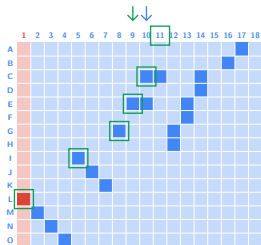
Consider the submatrix of  $R$  that contains only the columns “affected” by the removal of column 1.

**Step 1:** Identify the “left-most highest” pivots.

**Step 2:** Ensure column 10 is reduced:

# Removing simplices: in general

- ▶ Each “affected column” needs only one column addition to get it to where it should be.
- ▶ The column which should be added to it is determined by the left-most column with a pivot above its pivot (or below, if it itself is a left-most pivot column).



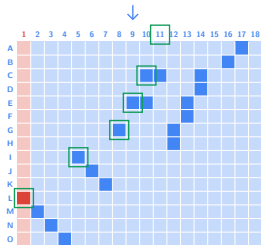
Consider the submatrix of  $R$  that contains only the columns “affected” by the removal of column 1.

**Step 1:** Identify the “left-most highest” pivots.

**Step 2:** Ensure column 10 is reduced: add column 9

# Removing simplices: in general

- ▶ Each “affected column” needs only one column addition to get it to where it should be.
- ▶ The column which should be added to it is determined by the left-most column with a pivot above its pivot (or below, if it itself is a left-most pivot column).



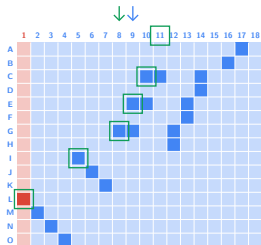
Consider the submatrix of  $R$  that contains only the columns “affected” by the removal of column 1.

**Step 1:** Identify the “left-most highest” pivots.

**Step 2:** Ensure column 9 is reduced:

# Removing simplices: in general

- ▶ Each “affected column” needs only one column addition to get it to where it should be.
- ▶ The column which should be added to it is determined by the left-most column with a pivot above its pivot (or below, if it itself is a left-most pivot column).



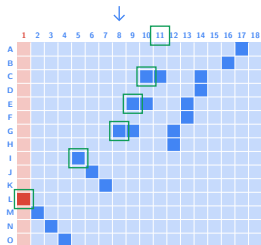
Consider the submatrix of  $R$  that contains only the columns “affected” by the removal of column 1.

**Step 1:** Identify the “left-most highest” pivots.

**Step 2:** Ensure column 9 is reduced: add column 8

# Removing simplices: in general

- ▶ Each “affected column” needs only one column addition to get it to where it should be.
- ▶ The column which should be added to it is determined by the left-most column with a pivot above its pivot (or below, if it itself is a left-most pivot column).



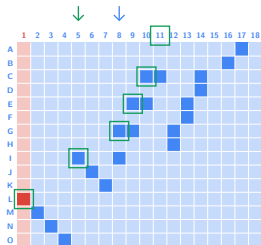
Consider the submatrix of  $R$  that contains only the columns “affected” by the removal of column 1.

**Step 1:** Identify the “left-most highest” pivots.

**Step 2:** Ensure column 8 is reduced:

# Removing simplices: in general

- ▶ Each “affected column” needs only one column addition to get it to where it should be.
- ▶ The column which should be added to it is determined by the left-most column with a pivot above its pivot (or below, if it itself is a left-most pivot column).



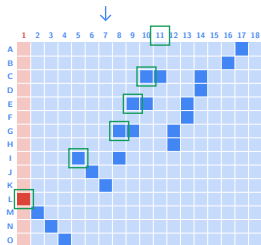
Consider the submatrix of  $R$  that contains only the columns “affected” by the removal of column 1.

**Step 1:** Identify the “left-most highest” pivots.

**Step 2:** Ensure column 8 is reduced: add column 5

# Removing simplices: in general

- ▶ Each “affected column” needs only one column addition to get it to where it should be.
- ▶ The column which should be added to it is determined by the left-most column with a pivot above its pivot (or below, if it itself is a left-most pivot column).



Consider the submatrix of  $R$  that contains only the columns “affected” by the removal of column 1.

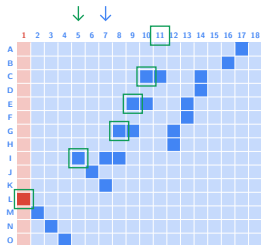
**Step 1:** Identify the “left-most highest” pivots.

**Step 2:** Ensure column 7 is reduced:



# Removing simplices: in general

- ▶ Each “affected column” needs only one column addition to get it to where it should be.
- ▶ The column which should be added to it is determined by the left-most column with a pivot above its pivot (or below, if it itself is a left-most pivot column).



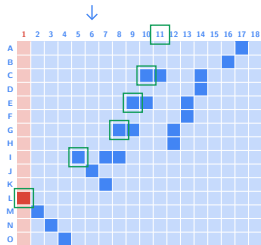
Consider the submatrix of  $R$  that contains only the columns “affected” by the removal of column 1.

**Step 1:** Identify the “left-most highest” pivots.

**Step 2:** Ensure column 7 is reduced: add column 5

# Removing simplices: in general

- ▶ Each “affected column” needs only one column addition to get it to where it should be.
- ▶ The column which should be added to it is determined by the left-most column with a pivot above its pivot (or below, if it itself is a left-most pivot column).



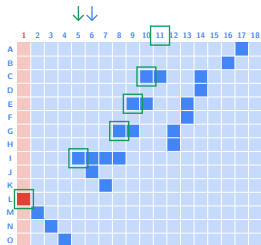
Consider the submatrix of  $R$  that contains only the columns “affected” by the removal of column 1.

**Step 1:** Identify the “left-most highest” pivots.

**Step 2:** Ensure column 6 is reduced:

# Removing simplices: in general

- ▶ Each “affected column” needs only one column addition to get it to where it should be.
- ▶ The column which should be added to it is determined by the left-most column with a pivot above its pivot (or below, if it itself is a left-most pivot column).



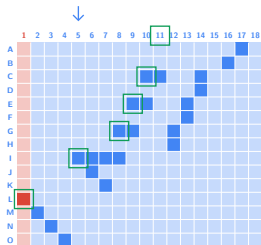
Consider the submatrix of  $R$  that contains only the columns “affected” by the removal of column 1.

**Step 1:** Identify the “left-most highest” pivots.

**Step 2:** Ensure column 6 is reduced: add column 5

# Removing simplices: in general

- ▶ Each “affected column” needs only one column addition to get it to where it should be.
- ▶ The column which should be added to it is determined by the left-most column with a pivot above its pivot (or below, if it itself is a left-most pivot column).



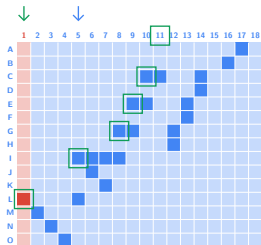
Consider the submatrix of  $R$  that contains only the columns “affected” by the removal of column 1.

**Step 1:** Identify the “left-most highest” pivots.

**Step 2:** Ensure column 5 is reduced:

## Removing simplices: in general

- ▶ Each “affected column” needs only one column addition to get it to where it should be.
- ▶ The column which should be added to it is determined by the left-most column with a pivot above its pivot (or below, if it itself is a left-most pivot column).



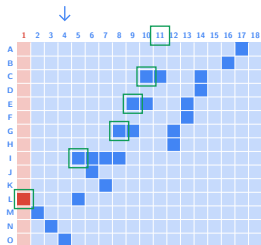
Consider the submatrix of  $R$  that contains only the columns “affected” by the removal of column 1.

**Step 1:** Identify the “left-most highest” pivots.

**Step 2:** Ensure column 5 is reduced: add column 1

# Removing simplices: in general

- ▶ Each “affected column” needs only one column addition to get it to where it should be.
- ▶ The column which should be added to it is determined by the left-most column with a pivot above its pivot (or below, if it itself is a left-most pivot column).



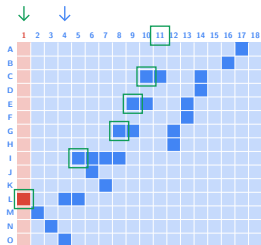
Consider the submatrix of  $R$  that contains only the columns “affected” by the removal of column 1.

**Step 1:** Identify the “left-most highest” pivots.

**Step 2:** Ensure column 4 is reduced:

## Removing simplices: in general

- ▶ Each “affected column” needs only one column addition to get it to where it should be.
- ▶ The column which should be added to it is determined by the left-most column with a pivot above its pivot (or below, if it itself is a left-most pivot column).



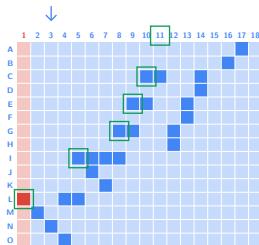
Consider the submatrix of  $R$  that contains only the columns “affected” by the removal of column 1.

**Step 1:** Identify the “left-most highest” pivots.

**Step 2:** Ensure column 4 is reduced: add column 1

# Removing simplices: in general

- ▶ Each “affected column” needs only one column addition to get it to where it should be.
- ▶ The column which should be added to it is determined by the left-most column with a pivot above its pivot (or below, if it itself is a left-most pivot column).



Consider the submatrix of  $R$  that contains only the columns “affected” by the removal of column 1.

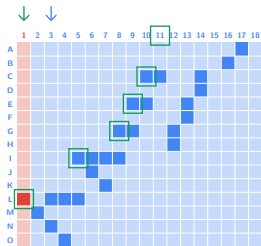
**Step 1:** Identify the “left-most highest” pivots.

**Step 2:** Ensure column 3 is reduced:



# Removing simplices: in general

- ▶ Each “affected column” needs only one column addition to get it to where it should be.
- ▶ The column which should be added to it is determined by the left-most column with a pivot above its pivot (or below, if it itself is a left-most pivot column).



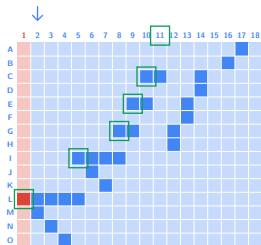
Consider the submatrix of  $R$  that contains only the columns “affected” by the removal of column 1.

**Step 1:** Identify the “left-most highest” pivots.

**Step 2:** Ensure column 3 is reduced: add column 1

# Removing simplices: in general

- ▶ Each “affected column” needs only one column addition to get it to where it should be.
- ▶ The column which should be added to it is determined by the left-most column with a pivot above its pivot (or below, if it itself is a left-most pivot column).



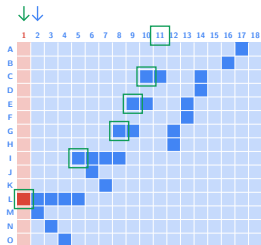
Consider the submatrix of  $R$  that contains only the columns “affected” by the removal of column 1.

**Step 1:** Identify the “left-most highest” pivots.

**Step 2:** Ensure column 2 is reduced:

## Removing simplices: in general

- ▶ Each “affected column” needs only one column addition to get it to where it should be.
- ▶ The column which should be added to it is determined by the left-most column with a pivot above its pivot (or below, if it itself is a left-most pivot column).



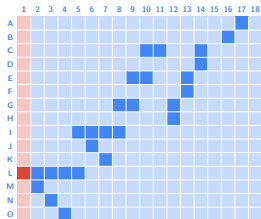
Consider the submatrix of  $R$  that contains only the columns “affected” by the removal of column 1.

**Step 1:** Identify the “left-most highest” pivots.

**Step 2:** Ensure column 2 is reduced: add column 1

# Removing simplices: in general

- ▶ Each “affected column” needs only one column addition to get it to where it should be.
- ▶ The column which should be added to it is determined by the left-most column with a pivot above its pivot (or below, if it itself is a left-most pivot column).



Consider the submatrix of  $R$  that contains only the columns “affected” by the removal of column 1.

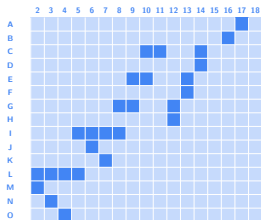
**Step 1:** Identify the “left-most highest” pivots.

**Step 2:** Ensure all columns are reduced

**Step 3:** Drop column 1

# Removing simplices: in general

- ▶ Each “affected column” needs only one column addition to get it to where it should be.
- ▶ The column which should be added to it is determined by the left-most column with a pivot above its pivot (or below, if it itself is a left-most pivot column).



Consider the submatrix of  $R$  that contains only the columns “affected” by the removal of column 1.

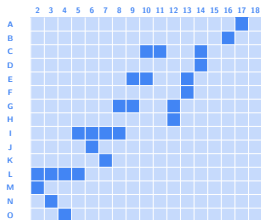
**Step 1:** Identify the “left-most highest” pivots.

**Step 2:** Ensure all columns are reduced

**Step 3:** Drop column 1

# Removing simplices: in general

- ▶ Each “affected column” needs only one column addition to get it to where it should be.
- ▶ The column which should be added to it is determined by the left-most column with a pivot above its pivot (or below, if it itself is a left-most pivot column).



Consider the submatrix of  $R$  that contains only the columns “affected” by the removal of column 1.

**Step 1:** Identify the “left-most highest” pivots.

**Step 2:** Ensure all columns are reduced

**Step 3:** Drop column 1

Each removal performs  $m$  column additions in  $R$ ,  $U$  (ensuring updated rep. cycles)

**Complexity:** for one removal is  $O(mn)$ , vs  $O((n-1)^3)$  for computing from scratch.

# SiRUP: Simplicial Removal Update Procedure

- ▶ Proof of correctness shows equivalence with “naive” method, on sets  $A, B$ .
- ▶ (Strong) induction on these “affected” simplices.

---

**Algorithm 2: SiRUP- SIMPLICIAL REMOVAL UPDATE PROCEDURE**

---

**Input:** Reduced matrix  $R$ , operations matrix  $U$ , list of simplices to be removed  $L$

**Output:** Updated reduced matrix  $P$ , updated operations matrix  $V$

```
1  $P \leftarrow R$ 
2  $V \leftarrow U$ 
3 for  $\sigma_j \in L$ 
4    $A = \{i : V[j, i] \neq 0, i \neq j\} = \{A_1, A_2, \dots\}$  in increasing order
5    $B = \{A_0 = j\} \cup \{A_t \in A : \text{piv}(P[A_i]) < \text{piv}(P[A_t]) \ \forall t < i\}$  in increasing order
6   for  $i \in A$  in decreasing order
7      $k \leftarrow$  the smallest index in  $B$  such that  $\text{piv}(P[B_k]) \leq \text{piv}(P[i])$ 
8     if  $i = B_k$ 
9        $k \leftarrow k - 1$ 
10    add column  $B_k$  to column  $i$  in  $P$  and in  $V$ 
11  remove  $P[j]$ ,  $V[j]$ , and row  $j$  in  $P$  and in  $V$ 
12 return  $P, V$ 
```

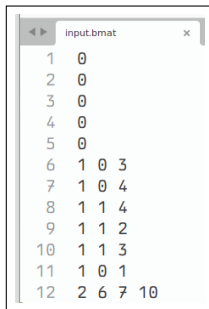
---

**Implementation:** On top of PHAT software, publicly available

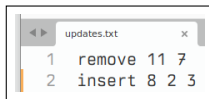
- ▶ Bottleneck is accessing rows to get “affected” simplices
- ▶ Tested and timed, see arxiv - time vs. space tradeoff (need matrix  $U$ )

# Usage

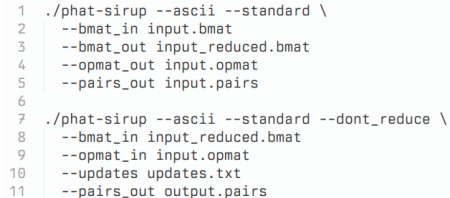
Input boundary matrix and output persistence pairs are in PHAT format:



```
input.bmat
1 0
2 0
3 0
4 0
5 0
6 1 0 3
7 1 0 4
8 1 1 4
9 1 1 2
10 1 1 3
11 1 0 1
12 2 6 7 10
```



```
updates.txt
1 remove 11 7
2 insert 8 2 3
```



```
1 ./phat-sirup --ascii --standard \
2   --bmat_in input.bmat
3   --bmat_out input_reduced.bmat
4   --opmat_out input.opmat
5   --pairs_out input.pairs
6
7 ./phat-sirup --ascii --standard --dont_reduce \
8   --bmat_in input_reduced.bmat
9   --opmat_in input.opmat
10  --updates updates.txt
11  --pairs_out output.pairs
```

**Goal:** Implement removal, insertion, swapping (see also Dionysus).

**Availability:** Main (remove) and dev (remove, insert) branches on BitBucket.

Work in progress!!



# Thank you for your attention

- ▶ Bauer, Kerber, Reininghaus, Wagner. *PHAT - Persistent Homology Algorithms Toolbox*. Journal of Symbolic Computation, 2016.
- ▶ Busaryev, Dey, Wang. *Tracking a Generator by Persistence*. Computing and Combinatorics, 2010.
- ▶ Cohen–Steiner, Edelsbrunner, Morozov. *Vines and vineyards by updating persistence in linear time*. Symposium of Computational Geometry, 2006.
- ▶ Edelsbrunner, Letscher, Zomorodian. *Topological persistence and simplification*. Symposium on Foundations of Computer Science, 2000.
- ▶ Giunti, Lazovskis. *Pruning vineyards: updating barcodes by removing simplices*. arxiv, 2025.
- ▶ Giunti, Lazovskis. *PHAT-SIRUP*. BitBucket, 2025.
- ▶ Hickok. *Persistence Diagram Bundles: A Multidimensional Generalization of Vineyards*. arxiv, 2022.
- ▶ Morozov, Skraba. *Persistent (Co)Homology in Matrix Multiplication Time*. arxiv, 2024.
- ▶ Morozov. *Dionysus, a C++ library for computing persistent homology*. 2007.
- ▶ Piekenbrock, Perea. *Move schedules: fast persistence computations in coarse dynamic settings*. Journal of Applied and Computational Topology, 2024.
- ▶ Turner. *Representing Vineyard Modules*. arxiv, 2023.