

Formal Languages and Compiler Design
Fifth laboratory
LL(1) Parser

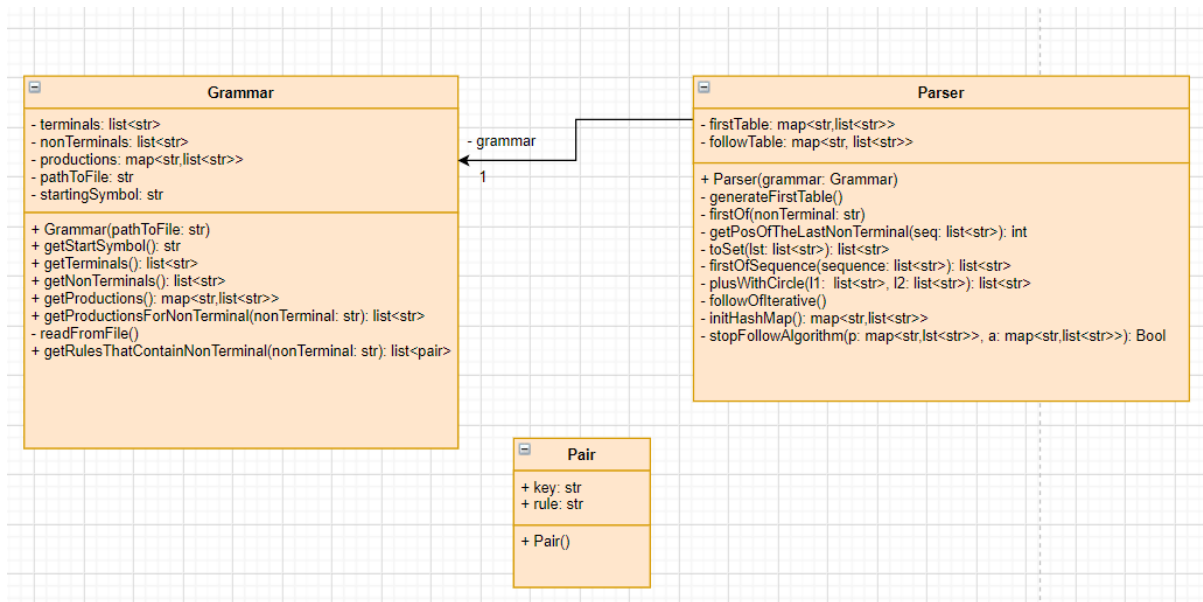
Part 1

Implemented classes:

1. *Pair*: simply represent a key, value pair
2. *Grammar*:
 - represents the grammar for a language described in a grammar.in file (does not require to have this name)
 - holds in “nonTerminals” the set of non-terminals from that grammar, also in the list “terminals” we keep the list for terminal symbols from that list
 - as for the productions, we keep them in a HashMap, so that we can map for each non-terminal the productions it is involved in.
 - the starting symbol is stored in a String field named “startingSymbol”
3. *Parser*:
 - represents an LL(1) parses
 - it contains a grammar, the one for which we build the parses
 - the First Table, containing the first values for the non-terminals
 - the Follow Table, containing the follow values for the non-terminals
 - the first values are computed using recursion and the follow values are computed in an iterative manner

Class Diagram:

- see next page



GRAMMAR.in file structure:

1. Description:

- first line, contains separated by “,” the set of non-terminals
- second line, contains separated by “,” the set of terminals
- third line, flag which marks the presence of “,” as a terminal
- last lines: productions

2. EBNF for grammar.in:

grammar ::= terminals “\n” nonTerminals “\n” flag “\n” productions

terminals ::= terminal | terminal “,” terminals

terminal ::= word

nonTerminals ::= nonTerminal | nonTerminal “,” nonTerminals

nonTerminal ::= word

flag ::= “true” | “false”

productions ::= production | production “\n” productions

production::= nonTerminal "->" ruleList "\n"
ruleList::=rule | rule "|" ruleList
rule::= compoundRule | ϵ
compoundRule::= terminal "\"" | nonTerminal "\"" | terminal "\"" compoundRule | nonTerminal "\"" compoundRule

3. File Example:

```
1 S,A,B,C,D
2 a,(,),+,*
3 false
4 S -> B A
5 A -> + B A |  $\epsilon$ 
6 B -> D C
7 C -> * D C |  $\epsilon$ 
8 D -> ( S ) | a
```