

Digital Fundamentals

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Number Systems, Operations, and Codes

Chapter 2

Ch.2 Summary

The Decimal Number System

Decimal: 10 symbols (0 through 9)

If number > 9: need more digits

Weights: powers of ten

... 10^5 10^4 10^3 10^2 10^1 10^0 .

For fractional: weights are negative powers of ten

... 10^2 10^1 10^0 . 10^{-1} 10^{-2} 10^{-3} ...

Ch.2 Summary

The Decimal Number System

Decimal number: linear combination of powers

Ex) the number 9240

$$(9 \times 10^3) + (2 \times 10^2) + (4 \times 10^1) + (0 \times 10^0)$$

Ex) 480.52

$$(4 \times 10^2) + (8 \times 10^1) + (0 \times 10^0) + (5 \times 10^{-1}) + (2 \times 10^{-2})$$

Ch.2 Summary

The Binary Number System

Binary: 2 symbols (0 and 1)

If number > 1 : need more digits

Weights: powers of two

... 2^5 2^4 2^3 2^2 2^1 2^0 .

For fractional: weights are negative powers of two

... 2^2 2^1 2^0 . 2^{-1} 2^{-2} 2^{-3} 2^{-4} ...

Ch.2 Summary

The Binary Number System

Notice the pattern of zeros and ones in each column.

Digital counters frequently have this same pattern of digits:

TABLE 2-1				
DECIMAL NUMBER	BINARY NUMBER			
0	0	0	0	0
1	0	0	0	1
2	0	0	1	0
3	0	0	1	1
4	0	1	0	0
5	0	1	0	1
6	0	1	1	0
7	0	1	1	1
8	1	0	0	0
9	1	0	0	1
10	1	0	1	0
11	1	0	1	1
12	1	1	0	0
13	1	1	0	1
14	1	1	1	0
15	1	1	1	1

Ch.2 Summary

Binary Weights

The positional weights for binary numbers are assigned as shown below.

TABLE 2-2 • Binary weights.														
POSITIVE POWERS OF TWO (WHOLE NUMBERS)									NEGATIVE POWERS OF TWO (FRACTIONAL NUMBER)					
2^8	2^7	2^6	2^5	2^4	2^3	2^2	2^1	2^0	2^{-1}	2^{-2}	2^{-3}	2^{-4}	2^{-5}	2^{-6}
256	128	64	32	16	8	4	2	1	$\frac{1}{2}$ 0.5	$\frac{1}{4}$ 0.25	$\frac{1}{8}$ 0.125	$\frac{1}{16}$ 0.625	$\frac{1}{32}$ 0.03125	$\frac{1}{64}$ 0.015625

Ch.2 Summary

Binary-to-Decimal Conversion

Add all weights with 1 (ignoring 0)

Convert the binary number 100101.01 to decimal.

2^5	2^4	2^3	2^2	2^1	2^0	2^{-1}	2^{-2}
32	16	8	4	2	1	0.5	0.25
1	0	0	1	0	1	0	1
32			+4	+1		+0.25	= 37.25

Ch.2 Summary

Decimal-to-Binary Conversions - I

Iteratively find the maximum weight

Convert the decimal number 49 to binary.

49 is in $[32, 64)$: the maximum weight is 32

remainder is $49 - 32 = 17$

17 is in $[16, 32)$: the maximum weight is 16

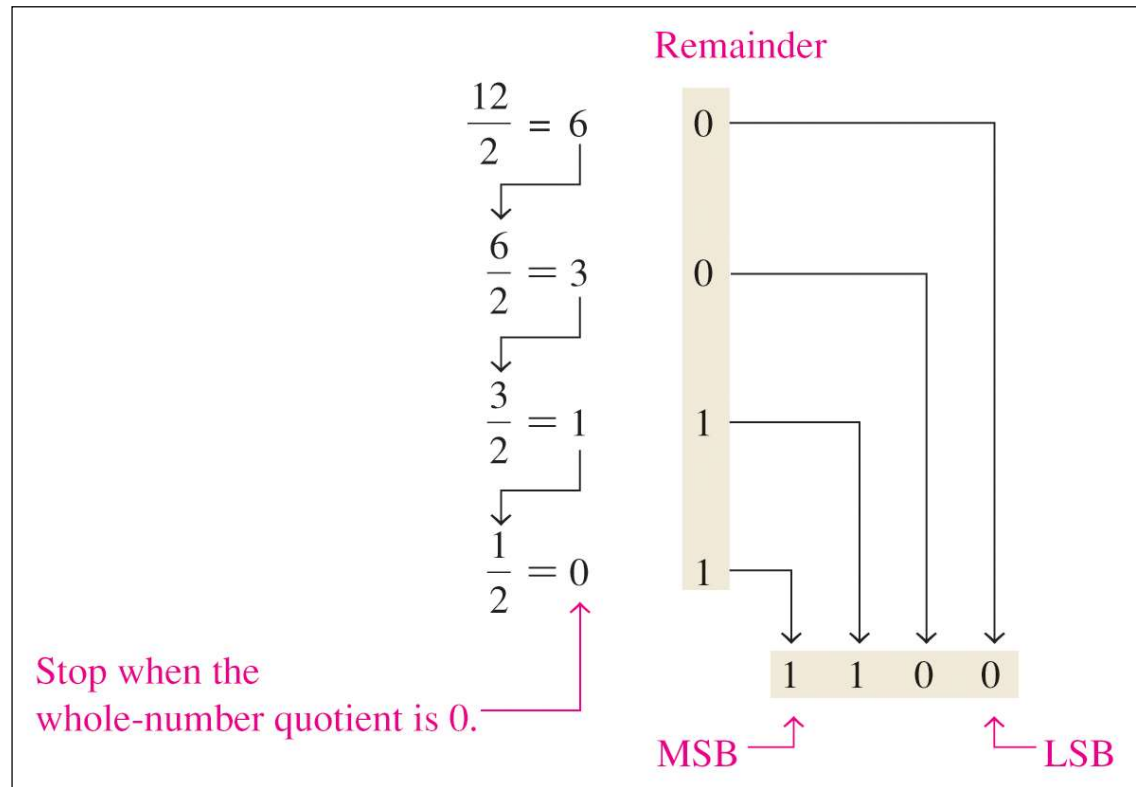
remainder is $17 - 16 = 1$, etc...

2^6	2^5	2^4	2^3	2^2	2^1	2^0
64	32	16	8	4	2	1
0	1	1	0	0	0	1

Ch.2 Summary

Decimal-to-Binary Conversions - II

Iteratively dividing by 2 and get remainders.

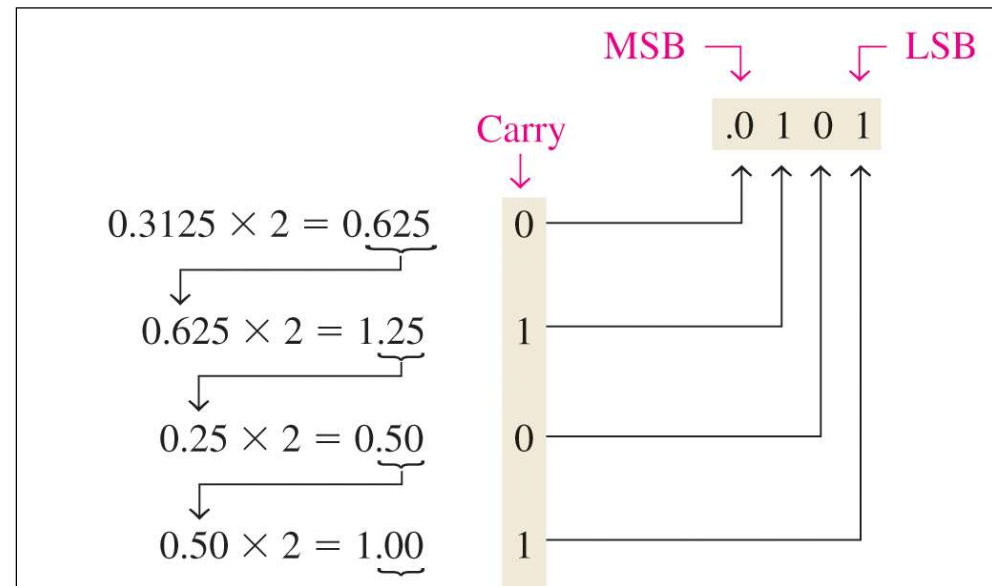


Ch.2 Summary

Decimal-to-Binary Conversions - II

For fraction,

Iteratively multiplying by 2
and get carries.



Ch.2 Summary

Binary Addition

The rules for binary addition are

$0 + 0 = 0$	Sum = 0, carry = 0
$0 + 1 = 0$	Sum = 1, carry = 0
$1 + 0 = 1$	Sum = 1, carry = 0
$1 + 1 = 10$	Sum = 0, carry = 1

When an input carry = 1 due to a previous result, the rules are

$1 + 0 + 0 = 01$	Sum = 1, carry = 0
$1 + 0 + 1 = 10$	Sum = 0, carry = 1
$1 + 1 + 0 = 10$	Sum = 0, carry = 1
$1 + 1 + 1 = 11$	Sum = 1, carry = 1

Ch.2 Summary

Binary Subtraction

The rules for binary subtraction are

$$0 - 0 = 0$$

$$1 - 1 = 0$$

$$1 - 0 = 1$$

$$10 - 1 = 1 \text{ with a borrow of } 1$$

Subtract the binary number 00111 from 10101 and show the equivalent decimal subtraction.

$$\begin{array}{r} 111 \\ 111 \\ \underline{00111} \\ 01110 \end{array} \quad \begin{array}{r} 21 \\ - 7 \\ \hline 14 \end{array}$$

Ch.2 Summary

Binary Multiplication

The rules for binary multiplication are:

$$0 \times 0 = 0$$

$$0 \times 1 = 0$$

$$1 \times 0 = 0$$

$$1 \times 1 = 1$$

Perform the following binary multiplications:

(a) 11×11

(b) 101×111

SOLUTION

(a)

$$\begin{array}{r} 11 \\ \times 11 \\ \hline 11 \\ + 11 \\ \hline 1001 \end{array}$$

Partial products {

$$\begin{array}{r} 3 \\ \times 3 \\ \hline 9 \end{array}$$

(b)

$$\begin{array}{r} 111 \\ \times 101 \\ \hline 111 \\ 000 \\ + 111 \\ \hline 100011 \end{array}$$

Partial products {

$$\begin{array}{r} 7 \\ \times 5 \\ \hline 35 \end{array}$$

Ch.2 Summary

Binary Division

The rules for binary division are:

(always the same with decimal operations)

Perform the following binary divisions:

(a) $110 \div 11$

(b) $110 \div 10$

SOLUTION

$$\begin{array}{r} \mathbf{10} \quad 2 \\ \text{(a) } 11 \overline{)110} \quad 3 \overline{)6} \\ \underline{11} \quad \underline{6} \\ 000 \quad 0 \end{array}$$

$$\begin{array}{r} \mathbf{11} \quad 3 \\ \text{(b) } 10 \overline{)110} \quad 2 \overline{)6} \\ \underline{10} \quad \underline{6} \\ 10 \quad 0 \\ \underline{10} \\ 00 \end{array}$$

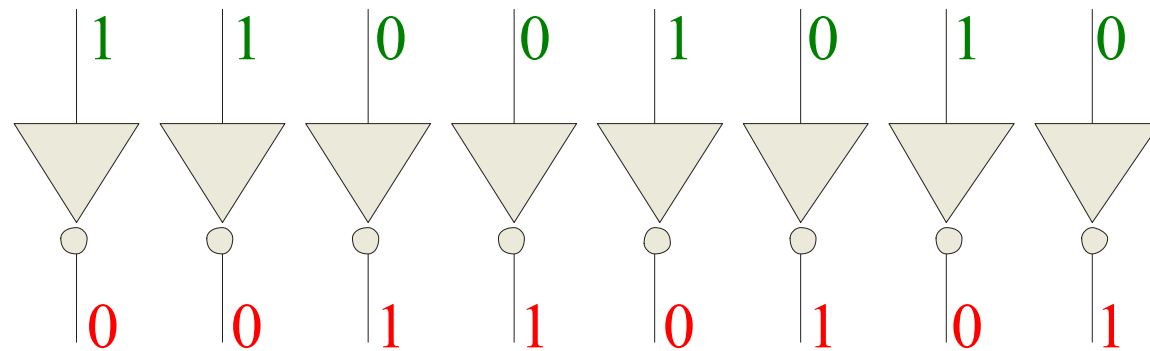
Ch.2 Summary

1's Complement

1's complement : reversing (inverting) all the bits.

For example, the 1's complement of **11001010**
is: **00110101**

We can use inverter to represent this operation



Ch.2 Summary

2's Complement

The 2's complement : 1's complement + 1

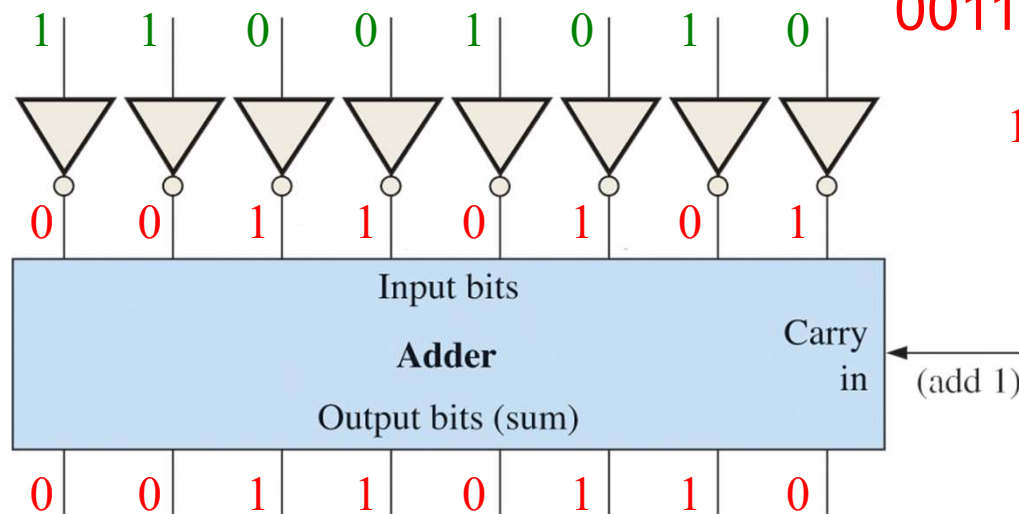
Recall that the 1's complement of **11001010** is

00110101 (1's complement)

For the 2's complement, add 1:

+1

00110110 (2's complement)



Ch.2 Summary

Signed Binary Numbers

MSB (most significant bit) represents the sign.

any symbol other than 0, 1 is not allowed!

Computers use a *modified 2's complement* for signed numbers.

Positive numbers: sign bit = 0

Negative numbers: sign bit = 1

For example, the positive number 58 is written using 8-bits
as 00111010 (true form).

Sign bit

Magnitude bits

Ch.2 Summary

Signed Binary Numbers

Negative numbers = 2's complement of the positive number.

The number -58 is written as:

-58 = 11000110 (complement form)

Sign bit Magnitude bits

MSB is a bit with weight -128.

We have 11000110 = -58, because

Column weights:	-128	64	32	16	8	4	2	1.
	1	1	0	0	0	1	1	0
	-128	+64				+4	+2	= -58

Ch.2 Summary

Floating Point Numbers

Floating point numbers: use scientific notation.

A 32-bit single precision number:



Express the speed of light, c , in single precision floating point notation.
($c = 0.2998 \times 10^9$)

In binary, $c = 0001\ 0001\ 1101\ 1110\ 1001\ 0101\ 1100\ 0000_2$.

In scientific notation, $c = 1.0001\ 1101\ 1110\ 1001\ 0101\ 1100\ 0000 \times 2^{28}$.

$S = 0$ because the number is positive. $E = 28 + 127 = 155_{10} = 1001\ 1011_2$. F is the next 23 bits after the first 1 is dropped.

In floating point notation, $c =$

0	10011011	0001 1101 1110 1001 0101 110
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Ch.2 Summary

Arithmetic Operations With Signed Numbers

Use a *modified 2's complement* for addition and subtraction.

Rules for **addition**: Add the two signed numbers. Discard any final carries. The result is in signed form.

Examples:

$$\begin{array}{r} 00011110 = +30 \\ 00001111 = +15 \\ \hline 00101101 = +45 \end{array}$$

$$\begin{array}{r} 00001110 = +14 \\ 11101111 = -17 \\ \hline 11111101 = -3 \end{array}$$

$$\begin{array}{r} 11111111 = -1 \\ 11111000 = -8 \\ \hline 11110111 = -9 \end{array}$$

Discard carry



Ch.2 Summary

Arithmetic Operations With Signed Numbers

Note that if the number of bits required for the answer is exceeded, overflow will occur. This occurs only if both numbers have the same sign. The overflow is indicated by an incorrect sign bit.

Two examples are:

$$\begin{array}{r} 01000000 = +64 \\ 01000001 = +65 \\ \hline 10000001 = -127 \end{array}$$

$$\begin{array}{r} 10000001 = -127 \\ 10000001 = -127 \\ \hline 100000010 = +2 \end{array}$$

Discard carry →

Wrong! The answer is incorrect
and the sign bit has changed.

Ch.2 Summary



Arithmetic Operations With Signed Numbers

Rules for **subtraction**: use 2's complement and add the numbers.
Discard any final carries. The result is in signed form.

Repeat the examples done previously, but subtract:

$$\begin{array}{r} 00011110 \quad (+30) \\ - 00001111 \quad -(+15) \\ \hline \end{array} \quad \begin{array}{r} 00001110 \quad (+14) \\ - 11101111 \quad -(-17) \\ \hline \end{array} \quad \begin{array}{r} 11111111 \quad (-1) \\ - 11111000 \quad -(-8) \\ \hline \end{array}$$

2's complement, then add:

$\begin{array}{r} 00011110 = +30 \\ 11110001 = -15 \\ \hline 100001111 = +15 \end{array}$	$\begin{array}{r} 00001110 = +14 \\ 00010001 = +17 \\ \hline 00011111 = +31 \end{array}$	$\begin{array}{r} 11111111 = -1 \\ 00001000 = +8 \\ \hline 100000111 = +7 \end{array}$
 Discard carry		 Discard carry

Ch.2 Summary

Hexadecimal Numbers

Hexadecimal: 16 symbols
0 through 9 and A through F.

Large binary number can easily be converted to hexadecimal by dividing it into 4-bit groups and converting each into its equivalent hexadecimal character.

Express $1001\ 0110\ 0000\ 1110_2$ in hexadecimal:

Group the binary number by 4-bits starting from the right. Thus, $1001011000001110 = 960E$

Decimal	Hexadecimal	Binary
0	0	0000
1	1	0001
2	2	0010
3	3	0011
4	4	0100
5	5	0101
6	6	0110
7	7	0111
8	8	1000
9	9	1001
10	A	1010
11	B	1011
12	C	1100
13	D	1101
14	E	1110
15	F	1111

Ch.2 Summary

Hexadecimal Numbers

Weights: powers of 16

Column weights $\begin{cases} 16^3 & 16^2 & 16^1 & 16^0 \\ 4096 & 256 & 16 & 1 \end{cases}$

Ex) express $1A2F_{16}$ in decimal.

4096 256 16 1
1 A 2 F_{16}

$$1(4096) + 10(256) + 2(16) + 15(1) = 6703_{10}$$

Decimal	Hexadecimal	Binary
0	0	0000
1	1	0001
2	2	0010
3	3	0011
4	4	0100
5	5	0101
6	6	0110
7	7	0111
8	8	1000
9	9	1001
10	A	1010
11	B	1011
12	C	1100
13	D	1101
14	E	1110
15	F	1111

Ch.2 Summary

Octal Numbers

Octal: 8 symbols
0 through 7

Binary number can easily be converted to octal by grouping bits 3 at a time and writing the equivalent octal character for each group.

Express 1 001 011 000 001 110₂ in octal:

Group the binary number by 3-bits starting from the right. Thus, 1001011000001110 = **113016₈**

Decimal	Octal	Binary
0	0	0000
1	1	0001
2	2	0010
3	3	0011
4	4	0100
5	5	0101
6	6	0110
7	7	0111
8	10	1000
9	11	1001
10	12	1010
11	13	1011
12	14	1100
13	15	1101
14	16	1110
15	17	1111

Ch.2 Summary

Octal Numbers

Octal is also a weighted number system.
The column weights are powers of 8,
which increase from right to left.

Column weights $\left\{ \begin{array}{cccc} 8^3 & 8^2 & 8^1 & 8^0 \\ 512 & 64 & 8 & 1 \end{array} \right.$

Ex) express 3702_8 in decimal.

512	64	8	1
3	7	0	2

$$3(512) + 7(64) + 0(8) + 2(1) = 1986_{10}$$

Decimal	Octal	Binary
0	0	0000
1	1	0001
2	2	0010
3	3	0011
4	4	0100
5	5	0101
6	6	0110
7	7	0111
8	10	1000
9	11	1001
10	12	1010
11	13	1011
12	14	1100
13	15	1101
14	16	1110
15	17	1111

Ch.2 Summary

Binary Coded Decimal(BCD)

Convert each decimal digits to binary number

BCD represents each decimal digit with a 4-bit code.

Codes 1010 through 1111 are not used in BCD.

Decimal	Binary	BCD
0	0000	0000
1	0001	0001
2	0010	0010
3	0011	0011
4	0100	0100
5	0101	0101
6	0110	0110
7	0111	0111
8	1000	1000
9	1001	1001
10	1010	0001 0000
11	1011	0001 0001
12	1100	0001 0010
13	1101	0001 0011
14	1110	0001 0100
15	1111	0001 0101

Ch.2 Summary

BCD

Weights of BCD: 80 40 20 10 8 4 2 1.

Example: What are the weights for the BCD number
1000 0011 0101 1001?

8000 4000 2000 1000 800 400 200 100 80 40 20 10 8 4 2 1

Get the decimal number by adding weights:

$$8000 + 200 + 100 + 40 + 10 + 8 + 1 = 8359_{10}$$

Ch.2 Summary

Gray Code

Gray code is an **unweighted** code

A single bit change between one code word and the next in a sequence.

To avoid problems in systems where an error can occur if more than one bit changes at a time.

Decimal	Binary	Gray code
0	0000	0000
1	0001	0001
2	0010	0011
3	0011	0010
4	0100	0110
5	0101	0111
6	0110	0101
7	0111	0100
8	1000	1100
9	1001	1101
10	1010	1111
11	1011	1110
12	1100	1010
13	1101	1011
14	1110	1001
15	1111	1000

Ch.2 Summary

ASCII

ASCII is a code for characters.

Originally, ASCII encoded 128 characters and symbols using 7-bits.

Ex) 35(0100011) -> #, 97(1100001) -> a

In 1981, IBM introduced extended ASCII, which is an 8-bit code and increased the character set to 256.

Other extended sets (such as Unicode) have been introduced to handle characters in languages other than English.

Ch.2 Summary

Parity Method for Error Detection

Method of error detection

parity bit : “extra” bit to force the number of 1’s to be either even (*even parity*) or odd (*odd parity*).

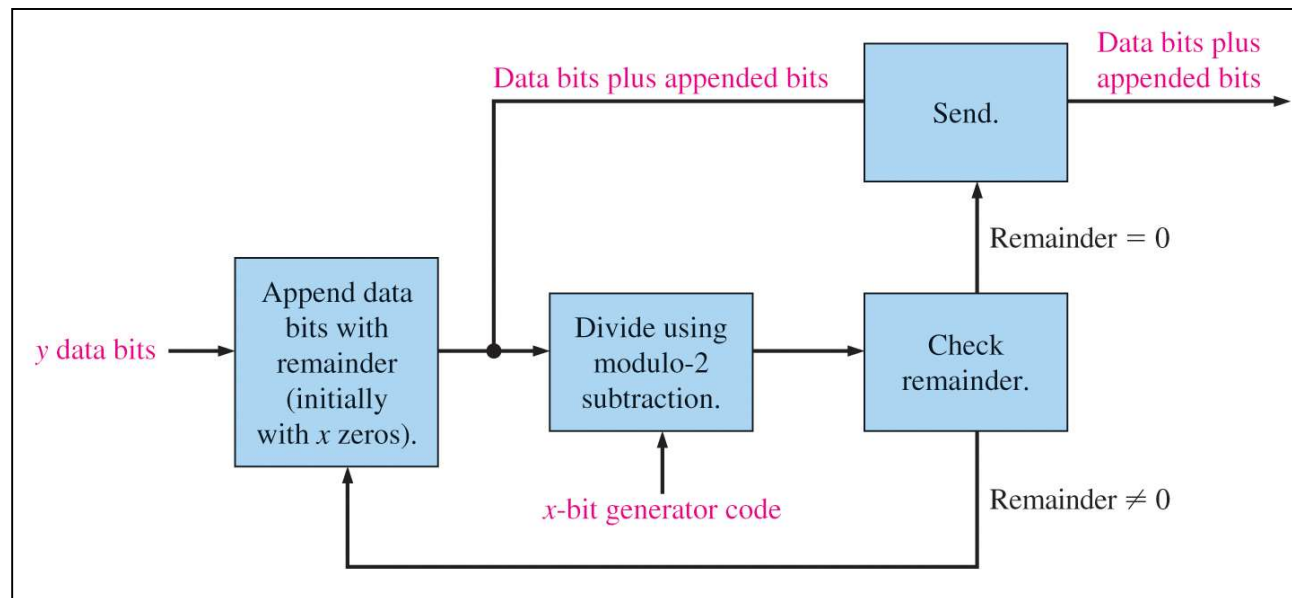
The ASCII character for “a” is 1100001 and for “A” is 1000001.
What is the correct bit to append to make both have odd parity?

The ASCII “a” has an odd number of bits that are equal to 1; therefore the parity bit is 0. The ASCII “A” has an even number of bits that are equal to 1; therefore the parity bit is 1.

Ch.2 Summary

Cyclic Redundancy Check

The cyclic redundancy check (CRC) is an error detection method that can detect multiple errors in larger blocks of data. At the sending end, a **checksum** is appended to a block of data. At the receiving end, the checksum is generated and compared to the sent checksum. If the checksums are the same, no error is detected.



Ch.2 Summary

Cyclic Redundancy Check

Example

11010011101100 with 1011