

# ECON 4673 Assignment #1

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## Problem 1

#	Statement	True/False
1	Every finite normal-form game has a pure strategy Nash equilibrium.	false
2	Every dominant strategy equilibrium is also a Nash equilibrium.	true
3	It is possible that there might exist some pure strategy Nash equilibria which did not survive the iterated elimination of strictly dominated strategies.	true
4	A strictly dominated strategy can also be a Nash equilibrium.	false
5	Every Nash equilibrium is also a dominant strategy equilibrium.	false
5	Every rationalizable strategy profile is also a Nash equilibrium.	false

## Problem 2

		Player 2	
		<i>C</i>	<i>D</i>
Player 1	<i>A</i>	1, 1	2, 0
	<i>B</i>	0, 2	$x, y$

### 2.1)

These values will lead to a dominant strategy equilibrium of  $(A, C)$ :

$$x < 2$$

$$y < 2$$

What follows is justification for the solution ...

Player 1's strategy  $A$  will be dominant if  $A$  has a greater payoff than strategy  $B$  for all strategies of Player 2.

$$u_1(A, s_2) > u_1(B, s_2) \quad \forall \quad s_2 \in S_2 = \{C, D, \sigma_2\}$$

where  $\sigma_2$  is a mixed strategy with probability  $0 < p < 1$  that  $s_2 = D$ .

When  $s_2 = C$ , it is given that  $s_1 = A$  has a higher payoff than  $s_2 = B$ .

$$u_1(A, C) = 1 > u_1(B, C) = 0$$

To satisfy the condition of dominance same must be true for  $s_2 = D$ :

$$u_1(A, D) = 2 > u_1(B, D) = x$$

$$x < 2$$

For the mixed strategy case:

$$u_1(A, \sigma_2) = u_1(A, C)(1 - p) + u_1(A, D)p = 1(1 - p) + 2p = 1 + p$$

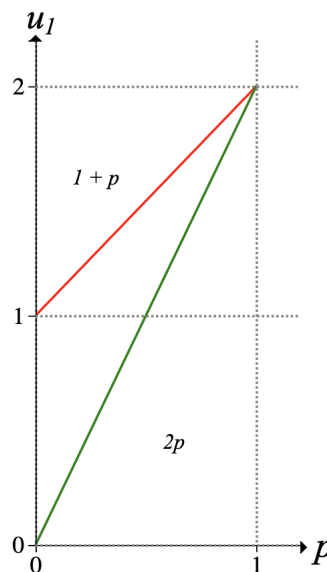
$$u_1(B, \sigma_2) = u_1(B, C)(1 - p) + u_1(B, D)p = 0(1 - p) + xp = xp$$

To satisfy the dominance condition:

$$u_1(A, \sigma_2) > u_1(B, \sigma_2)$$

$$1 + p > xp$$

The relationship is satisfied for  $0 < p < 1$  if  $x < 2$ , which can be seen graphically:



The same analysis for player 2:

$$u_2(s_1, C) > u_2(s_1, D) \quad \forall \quad s_1 \in S_1 = \{A, B, \sigma_1\}$$

( $\sigma_1$  is a mixed strategy with probability  $0 < p < 1$  that  $s_1 = B$ )

$$u_2(A, C) = 1 > u_2(A, D) = 0$$

$$u_2(B, C) = 2 > u_2(B, D) = y$$

$$y < 2$$

$$u_2(\sigma_1, C) = 1 + p$$

$$u_2(\sigma_1, D) = yp$$

$$u_2(\sigma_1, C) > u_2(\sigma_1, D)$$

$$1 + p > yp \quad , \quad 0 < p < 1 \implies y < 2$$

**2.2)**

These values will make this game into a *Prisoner's Dilemma*:

$$1 < x < 2$$

$$1 < y < 2$$

In a *Prisoner's Dilemma* there are three scenarios that can occur:

*Scenario 1:* Both criminals defect

The authorities offer each of the criminals the opportunity to testify on the other in exchange for a reduced sentence. In this scenario, both criminals accept the offer and testify against the other (strategy  $D$ ). They receive a payoff value  $x_1$  representing the reduced sentence.

$$u_1(D, D) = u_2(D, D) = x_1$$

*Scenario 2:* One criminal cooperates & the other defects

Criminal  $i$  decides to testify against the other ( $s_i = D$ ) and is freed, which means this criminal's payoff  $x_{2i}$  is higher than in scenario 1. The other criminal  $-i$  does not testify ( $s_{-i} = C$ ), which means they get the full sentence for the crime so their payoff  $x_{2ii}$  is lower than the payoff from scenario 1.

$$u_i(C, D) = x_{2i} > x_1$$

$$u_{-i}(C, D) = x_{2ii} < x_1$$

*Scenario 3: Criminals Cooperation*

Each criminal decides not to testify on the other so the authorities are only able to convict them of a minor offense. The sentence is less harsh than the sentence for the main crime, which means the payoff value  $x_3$  is higher than in scenario 1, but not as high as the case where the criminal goes free.

$$x_1 < u_1(C, C) = u_2(C, C) = x_3 < x_{2i}$$

The prisoner's dilemma in normal form:

		Criminal 1	
		$D$	$C$
Criminal 1	$D$	$x_1, x_1$	$x_{2i}, x_{2ii}$
	$C$	$x_{2i}, x_{2ii}$	$x_3, x_3$

$$x_{2ii} < x_1 < x_3 < x_{2i}$$

Substituting the payoff from the game in the problem:

$$x_1 = 1, \quad x_{2i} = 2, \quad x_{2ii} = 0$$

		Criminal 1	
		$D$	$C$
Criminal 1	$D$	1, 1	2, 0
	$C$	0, 2	$x_3, x_3$

$$0 < 1 < x_3 < 2$$

### Problem 3

		Player 2	
		X	Y
Player 1	A	3, 3	0, 1
	B	5, 5	2, 2

#### 3.1)

The only pure strategy Nash equilibrium is  $(B, X)$ .

The condition for Nash equilibrium

$$s_i \in BR_i(s_{-i}) \quad \forall i$$

Strategy Y is dominated by X, which means  $s_2 = X$  is always best response

$$BR_2(s_1) = \{X\} \quad \forall s_1$$

$s_1 = B$  is the best response to X

$$BR_1(X) = B$$

#### 3.2)

Let  $\sigma_1$  be a mixed strategy with probability  $0 < p < 1$  that  $s_1 = A$

Let  $\sigma_2$  be a mixed strategy with probability  $0 < q < 1$  that  $s_2 = X$

A Nash equilibrium will exist  $p = 0$  and  $q = 1$  which results in the strategy profile  $(B, X)$

Player 1's expected payoff

$$\begin{aligned}
 v_1(\sigma_1, \sigma_2) &= \sum_{(s_1, s_2)} u_1(s_1, s_2) P(s = (s_1, s_2)), \quad (s_1, s_2) \in S_1 \times S_2 \\
 &= u_1(A, X)pq + u_1(A, Y)p(1-q) + u_1(B, X)(1-p)q + u_1(B, Y)(1-p)(1-q) \\
 &= 3pq + 0p(1-q) + 5(1-p)q + 2(1-p)(1-q) \\
 &= 3q - 2p + 2
 \end{aligned}$$

It is strictly decreasing with  $p$

$$\frac{\partial}{\partial p} v_1(\sigma_1, \sigma_2) = -2$$

$v_1(\sigma_1, \sigma_2)$  is maximized when  $p$  is minimized, so for any given  $q$ ,  $p = 0$  is best.

$$BR_1(\sigma_2) = B, \quad 0 < q < 1$$

The same analysis to determine player 2's expected payoff

$$\begin{aligned} v_2(\sigma_1, \sigma_2) &= \sum_{(s_1, s_2)} u_2(s_1, s_2) P(s = (s_1, s_2)), \quad (s_1, s_2) \in S_1 \times S_2 \\ &= u_2(A, X)pq + u_2(A, Y)p(1 - q) + u_2(B, X)(1 - p)q + u_2(B, Y)(1 - p)(1 - q) \\ &= 3pq + 1p(1 - q) + 5(1 - p)q + 2(1 - p)(1 - q) \\ &= -pq - p + 3q + 2 \end{aligned}$$

It is strictly increasing with  $q$  over the interval  $0 < p < 1$

$$\frac{\partial}{\partial p} v_2(\sigma_1, \sigma_2) = -p + 3 > 0, \quad 0 < p < 1$$

Player 2 will maximize the payoff with the largest possible value  $q = 1$

$$BR_2(\sigma_1) = X, \quad 0 < p < 1$$

## Problem 4

		Player 2		
		$L$	$M$	$R$
Player 1	$T$	1, -1	5, 0	0, 0
	$C$	0, 5	4, 4	0, 0
	$B$	0, 0	0, 0	3, 3

### 4.1)

The relationalizable strategies are  $\{M, R\} \times \{T, B\}$ . Rationalizable strategies are those which can survive the process of iterated dominance. Strategies  $C$  and  $L$  can be dominated.

Strategy  $C$  can be dominated by a mixed strategy  $\sigma_1$  with of playing  $T$ ,  $C$  and  $B$  of

$$4/5 < p_T < 1$$

$$p_C = 0$$

$$p_B = 1 - p_T$$

For example  $p_T = 0.85 \implies p_B = 0.15$

$$u_1(\sigma_1, L) = 0.85(1) + 0 + 0 = 0.85 > u_1(C, L) = 0$$

$$u_1(\sigma_1, M) = 0.85(5) + 0 + 0 = 4.25 > u_1(C, M) = 4$$

$$u_1(\sigma_1, R) = 0 + 0 + .15(3) = 0.45 > u_1(C, R) = 0$$

The game with  $C$  removed:

		Player 2		
		$L$	$M$	$R$
Player 1	$T$	1, -1	5, 0	0, 0
	$B$	0, 0	0, 0	3, 3

In the new game, strategy  $L$  is dominated by a mixed strategy  $\sigma_2$  of playing  $L$ ,  $M$  and  $R$  of

$$p_L = 0$$

$$0 < p_M < p_R < 1$$

For example  $p_M = 0.5 \implies p_B = 0.5$

$$u_2(T, \sigma_2) = 0 + 0.5(5) + 0 = 0.25 > u_2(T, L) = -1$$

$$u_2(B, \sigma_2) = 0 + 0 + 3(0.5) = 1.5 > u_2(B, L) = 0$$

The resulting game with  $L$  removed:

		Player 2	
		$M$	$R$
Player 1	$T$	5, 0	0, 0
	$B$	0, 0	3, 3

There are no more dominated strategies.

**4.2)**

There are two Nash equilibria  $(T, M)$  and  $(B, R)$ .

The condition for Nash equilibrium

$$u_i(s_i, s_{-i}) > u_i(s'_i, s_{-i}) \quad \forall i, s'_i \in S_i$$

$(T, M)$  is a Nash equilibrium but  $(T, R)$  is not:

$$u_1(T, M) = 5 > u_1(T, R) = 0$$

$$u_2(T, M) = 0 \geq u_2(T, R) = 0$$

$(B, R)$  is a Nash equilibrium but  $(B, M)$  is not:

$$u_1(B, R) = 3 > u_1(B, M) = 0$$

$$u_2(B, R) = 3 > u_2(B, M) = 0$$

**4.3)**

There is not a dominant strategy equilibrium for the game.