

Type Theory Theory exercises

Alberto Lazari

II Semester A.Y. 2022-2023

Exercise 1

3.1 Singleton type and exercises

- 3. Show that the rule E-S) is derivable in the type theory T_1 replacing the rule E-S) elimination with the E-N_{1prog}) rule and adding the substitution and weakening rules and the sanitary checks rules set out in the previous sections.
- Rule E-S)

$$\text{E-S)} \ \frac{t \in \mathsf{N}_1 \ [\Gamma] \qquad M(z) \ type \ [\Gamma, z \in \mathsf{N}_1] \qquad c \in M(\star) \ [\Gamma]}{\mathsf{El}_{\mathsf{N}_1}(t,c) \in M(t) \ [\Gamma]}$$

• Rule E- N_{1prog})

$$\text{E-N}_{1prog}) \ \frac{D(w) \ type \ [\Sigma, w \in \mathsf{N}_1] \qquad d \in D(\star) \ [\Sigma]}{\mathsf{El}_{\mathsf{N}_1}(w, d) \in D(w) \ [\Sigma, w \in \mathsf{N}_1]}$$

Solution

Assuming:

$$a_1$$
) $t \in \mathbb{N}_1 [\Gamma]$

$$a_2)\ M(z)\ type\ [\Gamma,z\in \mathsf{N}_1]$$

$$a_3$$
) $c \in M(\star) [\Gamma]$

The rule E-S) is derivable:

$$\begin{array}{c} a_2 & a_3 \\ \text{E-N}_{1prog}) & \underline{M(z) \ type \ [\Gamma,z \in \mathbb{N}_1] \quad c \in M(\star) \ [\Gamma]} \\ \text{sub-ter}) & \underline{\operatorname{El}_{\mathbb{N}_1}(z,c) \in M(z) \ [\Gamma,z \in \mathbb{N}_1]} \\ \end{array} \qquad \begin{array}{c} a_1 \\ t \in \mathbb{N}_1 \ [\Gamma] \end{array}$$

Exercise 2

3.2 Natural Numbers Type

3. Define the addition operation using the rules of the natural number type

$$x + y \in \mathsf{Nat} [x \in \mathsf{Nat}, y \in \mathsf{Nat}]$$

such that $x + 0 = x \in \mathsf{Nat}\ [x \in \mathsf{Nat}].$

Solution

The addition x + y can be defined as:

$$\mathsf{El}_{\mathsf{Nat}}(y,x,(w,z).\,\mathsf{succ}(z))$$

Let $\Gamma = x \in \mathsf{Nat}, y \in \mathsf{Nat};$

 $x + y \in \mathsf{Nat} \ [x \in \mathsf{Nat}, y \in \mathsf{Nat}]$ is derivable:

$$\text{var}) \frac{\Gamma \ cont}{y \in \text{Nat} \ [\Gamma]} \quad \text{F-Nat}) \frac{\Gamma \ cont}{\text{Nat} \ type \ [\Gamma]} \quad \text{var}) \frac{\Gamma \ cont}{x \in \text{Nat} \ [\Gamma]} \quad \text{var}) \frac{\Gamma \ cont}{x \in \text{Nat} \ [\Gamma]} \quad \text{Var}) \frac{\Gamma \ cont}{z \in \text{Nat} \ [\Gamma, w \in \text{Nat}, z \in \text{Nat}]} \frac{\Gamma \ cont}{z \in \text{Nat} \ [\Gamma, w \in \text{Nat}, z \in \text{Nat}]} \frac{\Gamma \ cont}{z \in \text{Nat} \ [\Gamma, w \in \text{Nat}, z \in \text{Nat}]} \frac{\Gamma \ cont}{z \in \text{Nat} \ [\Gamma, w \in \text{Nat}, z \in \text{Nat}]} \frac{\Gamma \ cont}{z \in \text{Nat} \ [\Gamma, w \in \text{Nat}, z \in \text{Nat}]} \frac{\Gamma \ cont}{z \in \text{Nat} \ [\Gamma, w \in \text{Nat}, z \in \text{Nat}]} \frac{\Gamma \ cont}{z \in \text{Nat} \ [\Gamma, w \in \text{Nat}, z \in \text{Nat}]} \frac{\Gamma \ cont}{z \in \text{Nat} \ [\Gamma, w \in \text{Nat}, z \in \text{Nat}]} \frac{\Gamma \ cont}{z \in \text{Nat} \ [\Gamma, w \in \text{Nat}, z \in \text{Nat}]} \frac{\Gamma \ cont}{z \in \text{Nat} \ [\Gamma, w \in \text{Nat}, z \in \text{Nat}]} \frac{\Gamma \ cont}{z \in \text{Nat} \ [\Gamma, w \in \text{Nat}, z \in \text{Nat}]} \frac{\Gamma \ cont}{z \in \text{Nat} \ [\Gamma, w \in \text{Nat}, z \in \text{Nat}]} \frac{\Gamma \ cont}{z \in \text{Nat} \ [\Gamma, w \in \text{Nat}, z \in \text{Nat}]} \frac{\Gamma \ cont}{z \in \text{Nat} \ [\Gamma, w \in \text{Nat}, z \in \text{Nat}]} \frac{\Gamma \ cont}{z \in \text{Nat} \ [\Gamma, w \in \text{Nat}, z \in \text{Nat}]} \frac{\Gamma \ cont}{z \in \text{Nat} \ [\Gamma, w \in \text{Nat}, z \in \text{Nat}]} \frac{\Gamma \ cont}{z \in \text{Nat} \ [\Gamma, w \in \text{Nat}, z \in \text{Nat}]} \frac{\Gamma \ cont}{z \in \text{Nat} \ [\Gamma, w \in \text{Nat}, z \in \text{Nat}]} \frac{\Gamma \ cont}{z \in \text{Nat}} \frac{\Gamma \ cont}$$

Where Γ cont derivable, because:

- $\Gamma = x \in \mathsf{Nat}, y \in \mathsf{Nat}$
- $x \in \mathsf{Nat}, y \in \mathsf{Nat}\ cont\ derivable$:

$$\begin{aligned} & \text{F-Nat)} & \frac{ \quad [\] \ cont }{ \quad \text{Nat} \ type \ [\] } \\ & \text{F-Nat)} & \frac{ x \in \text{Nat} \ cont }{ \quad \text{Nat} \ type \ [x \in \text{Nat}] } \\ & \text{F-c)} & \frac{ \quad \text{Nat} \ type \ [x \in \text{Nat}] }{ \quad x \in \text{Nat}, y \in \text{Nat} \ cont } \end{aligned}$$

Correctness

The definition is correct, in fact:

Base case

$$y = 0 \Rightarrow x + y = x + 0 = x$$

This is true, because:

- $x + 0 = \mathsf{El}_{\mathsf{Nat}}(0, x, (w, z). \operatorname{succ}(z))$
- $\mathsf{El}_{\mathsf{Nat}}(0,x,(w,z).\operatorname{succ}(z)) = x \in \mathsf{Nat}\ [x \in \mathsf{Nat}]$ derivable:

$$\begin{array}{c} \text{F-Nat}) & \frac{\left[\ \right] \ cont}{\text{Nat type } \left[\ \right]} \\ \text{F-Nat}) & \frac{\left[\ \right] \ cont}{x \in \text{Nat cont}} \\ \text{F-Nat}) & \frac{\left[\ \right] \ cont}{x \in \text{Nat cont}} \\ \text{F-Nat}) & \frac{\left[\ \right] \ cont}{x \in \text{Nat type } \left[\ \right]} \\ \text{F-Nat}) & \frac{\left[\ \right] \ cont}{x \in \text{Nat type } \left[\ \right]} \\ \text{F-Nat}) & \frac{\left[\ \right] \ cont}{x \in \text{Nat type } \left[\ \right]} \\ \text{F-C}) & \frac{\left[\ \right] \ cont}{x \in \text{Nat type } \left[\ \right]} \\ \text{F-C}) & \frac{\left[\ \right] \ cont}{x \in \text{Nat type } \left[\ \right]} \\ \text{F-C}) & \frac{\left[\ \right] \ cont}{x \in \text{Nat cont}} \\ \text{F-C}) & \frac{\left[\ \right] \ cont}{x \in \text{Nat type } \left[\ \right]} \\ \text{Var}) & \frac{\left[\ \right] \ cont}{x \in \text{Nat type } \left[\ \right]} \\ \text{Var}) & \frac{\left[\ \right] \ cont}{x \in \text{Nat type } \left[\ \right]} \\ \text{C}_{1}\text{-Nat}) & \frac{\left[\ \right] \ cont}{x \in \text{Nat type } \left[\ \right]} \\ \text{El}_{\text{Nat}}(0,x,(w,z). \text{succ}(z)) = x \in \text{Nat } \left[x \in \text{Nat} \right] \\ \text{Succ}(z) \in \text{Nat} \left[x \in \text{Nat}, w \in \text{Nat}, x \in \text{Nat}, x \in \text{Nat} \right]} \\ \text{El}_{\text{Nat}}(0,x,(w,z). \text{succ}(z)) = x \in \text{Nat} \left[x \in \text{Nat} \right] \\ \text{Nat} & \frac{\left[\ \right] \ cont}{x \in \text{Nat}} \\ \text{Nat} & \frac{\left[\ \right] \ cont}{x \in \text{Nat}} \\ \text{Nat} & \frac{\left[\ \right] \ cont}{x \in \text{Nat}} \\ \text{Nat} & \frac{\left[\ \right] \ cont}{x \in \text{Nat}} \\ \text{Nat} & \frac{\left[\ \right] \ cont}{x \in \text{Nat}} \\ \text{Nat} & \frac{\left[\ \right] \ cont}{x \in \text{Nat}} \\ \text{Nat} & \frac{\left[\ \right] \ cont}{x \in \text{Nat}} \\ \text{Nat} & \frac{\left[\ \right] \ cont}{x \in \text{Nat}} \\ \text{Nat} & \frac{\left[\ \right] \ cont}{x \in \text{Nat}} \\ \text{Nat} & \frac{\left[\ \right] \ cont}{x \in \text{Nat}} \\ \text{Nat} & \frac{\left[\ \right] \ cont}{x \in \text{Nat}} \\ \text{Nat} & \frac{\left[\ \right] \ cont}{x \in \text{Nat}} \\ \text{Nat} & \frac{\left[\ \right] \ cont}{x \in \text{Nat}} \\ \text{Nat} & \frac{\left[\ \right] \ cont}{x \in \text{Nat}} \\ \text{Nat} & \frac{\left[\ \right] \ cont}{x \in \text{Nat}} \\ \text{Nat} & \frac{\left[\ \right] \ cont}{x \in \text{Nat}} \\ \text{Nat} & \frac{\left[\ \right] \ cont}{x \in \text{Nat}} \\ \text{Nat} & \frac{\left[\ \right] \ cont}{x \in \text{Nat}} \\ \text{Nat} & \frac{\left[\ \right] \ cont}{x \in \text{Nat}} \\ \text{Nat} & \frac{\left[\ \right] \ cont}{x \in \text{Nat}} \\ \text{Nat} & \frac{\left[\ \right] \ cont}{x \in \text{Nat}} \\ \text{Nat} & \frac{\left[\ \right] \ cont}{x \in \text{Nat}} \\ \text{Nat} & \frac{\left[\ \right] \ cont}{x \in \text{Nat}} \\ \text{Nat} & \frac{\left[\ \right] \ cont}{x \in \text{Nat}} \\ \text{Nat} & \frac{\left[\ \right] \ cont}{x \in \text{Nat}} \\ \text{Nat} & \frac{\left[\ \right] \ cont}{x \in \text{Nat}} \\ \text{Nat} & \frac{\left[\ \right] \ cont}{x \in \text{Nat}$$

Inductive case

$$y = \operatorname{succ}(v) \ [v \in \operatorname{Nat}] \Rightarrow x + y = x + \operatorname{succ}(v) = \operatorname{succ}(x + v)$$

This is true, because:

- $x + \operatorname{succ}(v) = \operatorname{El}_{\operatorname{Nat}}(\operatorname{succ}(v), x, (w, z). \operatorname{succ}(z))$
- $\operatorname{succ}(x+v) = \operatorname{succ}(\operatorname{El}_{\operatorname{Nat}}(v,x,(w,z).\operatorname{succ}(z)))$
- Let $\Gamma = x \in \mathsf{Nat}, v \in \mathsf{Nat};$

$$\mathsf{El}_{\mathsf{Nat}}(\mathsf{succ}(v), x, (w, z).\,\mathsf{succ}(z)) = \mathsf{succ}(\mathsf{El}_{\mathsf{Nat}}(v, x, (w, z).\,\mathsf{succ}(z))) \in \mathsf{Nat}\ [\Gamma]\ \mathsf{derivable} :$$

$$\text{var}) \ \frac{\Gamma \ cont}{V \in \mathsf{Nat} \ [\Gamma]} \\ \ F-\mathsf{Nat}) \ \frac{\Gamma \ cont}{\mathsf{Nat} \ type \ [\Gamma]} \\ \ \mathsf{var}) \ \frac{\Gamma \ cont}{x \in \mathsf{Nat} \ [\Gamma]} \\ \ \mathsf{var}) \ \frac{\Gamma \ cont}{x \in \mathsf{Nat} \ [\Gamma]} \\ \ \mathsf{var}) \ \frac{\Gamma \ cont}{x \in \mathsf{Nat} \ [\Gamma]} \\ \ \mathsf{var}) \ \frac{\Gamma \ cont}{x \in \mathsf{Nat} \ [\Gamma]} \\ \ \mathsf{var}) \ \frac{\Gamma \ cont}{x \in \mathsf{Nat} \ [\Gamma]} \\ \ \mathsf{var}) \ \frac{\Gamma \ cont}{x \in \mathsf{Nat} \ [\Gamma, w \in \mathsf{Nat}, z \in \mathsf{Nat}]} \\ \ \mathsf{var}) \ \frac{\Gamma \ cont}{x \in \mathsf{Nat} \ [\Gamma, w \in \mathsf{Nat}, z \in \mathsf{Nat}]} \\ \ \mathsf{El}_{\mathsf{Nat}}(\mathsf{succ}(v), x, (w, z). \mathsf{succ}(z)) = \mathsf{succ}(\mathsf{El}_{\mathsf{Nat}}(v, x, (w, z). \mathsf{succ}(z))) \in \mathsf{Nat} \ [\Gamma]$$

Where Γ cont derivable, because:

- $\Gamma = x \in \mathsf{Nat}, v \in \mathsf{Nat}$
- $x \in \mathsf{Nat}, v \in \mathsf{Nat}\ cont\ derivable$:

$$\begin{array}{c} \text{F-Nat)} & \underline{ \begin{array}{c} \left[\ \right] \ cont \\ \text{Nat} \ type \left[\ \right] \end{array} } \\ \text{F-Nat)} & \underline{ \begin{array}{c} x \in \text{Nat} \ cont \\ \text{Nat} \ type \left[x \in \text{Nat} \right] \end{array} } \\ \text{F-c)} & \underline{ \begin{array}{c} x \in \text{Nat} \ cont \\ x \in \text{Nat}, v \in \text{Nat} \ cont \end{array} } \end{array} }$$

Exercise 3

3.2 Natural Numbers Type

4. Define the addition operation using the rules of the natural number type

$$x + y \in \mathsf{Nat} \ [x \in \mathsf{Nat}, y \in \mathsf{Nat}]$$

such that
$$0 + x = x \in \mathsf{Nat} \ [x \in \mathsf{Nat}].$$

Solution

The addition x + y can be defined as:

$$\mathsf{El}_{\mathsf{Nat}}(x,y,(w,z).\operatorname{\mathsf{succ}}(z))$$

Let $\Gamma = x \in \mathsf{Nat}, y \in \mathsf{Nat};$ $x + y \in \mathsf{Nat} \ [x \in \mathsf{Nat}, y \in \mathsf{Nat}] \ \text{is derivable}:$

$$\begin{array}{c} \operatorname{F-C}) \frac{\Gamma \ cont}{\Gamma, w \in \operatorname{Nat} \ cont} \\ \operatorname{Var}) \frac{\Gamma \ cont}{x \in \operatorname{Nat} \ [\Gamma]} & \operatorname{F-Nat}) \frac{\Gamma \ cont}{\operatorname{Nat} \ type \ [\Gamma]} & \operatorname{var}) \frac{\Gamma \ cont}{y \in \operatorname{Nat} \ [\Gamma]} & \operatorname{Var}) \frac{\Gamma \ cont}{y \in \operatorname{Nat} \ [\Gamma]} & \operatorname{Var}) \frac{\Gamma \ cont}{z \in \operatorname{Nat} \ [\Gamma, w \in \operatorname{Nat}, z \in \operatorname{Nat}]} \\ \operatorname{E-Nat}) & \operatorname{El}_{\operatorname{Nat}}(x, y, (w, z). \operatorname{succ}(z)) \in \operatorname{Nat} \ [\Gamma] \end{array}$$

Where Γ *cont* derivable, because:

- $\Gamma = x \in \mathsf{Nat}, y \in \mathsf{Nat}$
- $x \in \mathsf{Nat}, y \in \mathsf{Nat}\ cont\ derivable$:

$$\begin{array}{c} \text{F-Nat)} & \cfrac{ \left[\ \right] \ cont }{ \text{Nat} \ type} \left[\ \right] }{ x \in \text{Nat} \ cont } \\ \text{F-Nat)} & \cfrac{ x \in \text{Nat} \ cont }{ \text{Nat} \ type} \left[x \in \text{Nat} \right] }{ x \in \text{Nat}, y \in \text{Nat} \ cont } \end{array}$$

Correctness

The definition is correct, in fact:

Base case

$$x = 0 \Rightarrow x + y = 0 + y = y$$

Note that the exercise requires that $0+x=x\in \mathsf{Nat}\ [x\in \mathsf{Nat}]$, but that is equivalent to proving that $0+y=y\in \mathsf{Nat}\ [y\in \mathsf{Nat}]$, by renaming y to x in the latter, and this is true, because:

- $0 + y = \mathsf{El}_{\mathsf{Nat}}(0, y, (w, z). \operatorname{succ}(z))$
- $\mathsf{El}_{\mathsf{Nat}}(0,y,(w,z).\,\mathsf{succ}(z)) = y \in \mathsf{Nat}\ [y \in \mathsf{Nat}]$ derivable:

$$\begin{array}{c} \text{F-Nat}) & \frac{ \left[\ \right] \ cont}{\text{Nat } \ type \ \left[\ \right]} \\ \text{F-Nat}) & \frac{ \left[\ \right] \ cont}{y \in \text{Nat } \ cont} \\ \text{F-Nat}) & \frac{ \text{F-Nat}) \frac{ \left[\ \right] \ cont}{y \in \text{Nat } \ cont} \\ \text{F-Nat}) & \frac{ \text{F-Nat}) \frac{ \left[\ \right] \ cont}{y \in \text{Nat } \ type \ \left[\ \right]} \\ \text{F-Nat}) & \frac{ \left[\ \right] \ cont}{y \in \text{Nat } \ type \ \left[\ \right]} \\ \text{F-C}) & \frac{ \text{Nat } \ type \ \left[\ \right]}{y \in \text{Nat } \ cont} \\ \text{F-C}) & \frac{ \text{Nat } \ type \ \left[\ \right]}{y \in \text{Nat } \ cont} \\ \text{Var}) & \frac{ \text{F-C}) \frac{ \left[\ \right] \ cont}{y \in \text{Nat } \ cont} \\ \text{Var}) & \frac{ \text{F-C}) \frac{ \left[\ \right] \ cont}{y \in \text{Nat } \ type \ \left[\ \right]} \\ \text{Var}) & \frac{ \left[\ \right] \ cont}{y \in \text{Nat } \ type \ \left[\ \right]} \\ \text{Var}) & \frac{ \left[\ \right] \ cont}{y \in \text{Nat } \ type \ \left[\ \right]} \\ \text{Var}) & \frac{ \left[\ \right] \ cont}{y \in \text{Nat } \ type \ \left[\ \right]} \\ \text{El}_{\text{Nat}}(0,y,(w,z). \text{succ}(z)) = y \in \text{Nat } \ \left[y \in \text{Nat} \right] \\ \text{Succ}(z) \in \text{Nat} \ \left[y \in \text{Nat}, w \in \text{Nat}, z \in \text{Nat} \right]} \\ \text{El}_{\text{Nat}}(0,y,(w,z). \text{succ}(z)) = y \in \text{Nat } \ \left[y \in \text{Nat} \right] \\ \text{Nat} & \frac{ \left[\ \right] \ cont}{y \in \text{Nat}, w \in \text{Nat}, z \in \text{Nat} } \\ \text{Succ}(z) \in \text{Nat} \ \left[y \in \text{Nat}, w \in \text{Nat}, z \in \text{Nat} \right]} \\ \text{Succ}(z) \in \text{Nat} \ \left[y \in \text{Nat}, w \in \text{Nat}, z \in \text{Nat} \right]} \\ \text{Succ}(z) \in \text{Nat} \ \left[y \in \text{Nat}, w \in \text{Nat}, z \in \text{Nat} \right]} \\ \text{Succ}(z) \in \text{Nat} \ \left[y \in \text{Nat}, w \in \text{Nat}, z \in \text{Nat} \right]} \\ \text{Succ}(z) \in \text{Nat} \ \left[y \in \text{Nat}, w \in \text{Nat}, z \in \text{Nat} \right]} \\ \text{Succ}(z) \in \text{Nat} \ \left[y \in \text{Nat}, w \in \text{Nat}, z \in \text{Nat} \right]} \\ \text{Succ}(z) \in \text{Nat} \ \left[y \in \text{Nat}, w \in \text{Nat}, z \in \text{Nat} \right]} \\ \text{Succ}(z) \in \text{Nat} \ \left[y \in \text{Nat}, w \in \text{Nat}, z \in \text{Nat} \right]} \\ \text{Succ}(z) \in \text{Nat} \ \left[y \in \text{Nat}, w \in \text{Nat}, z \in \text{Nat} \right]} \\ \text{Succ}(z) \in \text{Nat} \ \left[y \in \text{Nat}, w \in \text{Nat}, z \in \text{Nat} \right]} \\ \text{Succ}(z) \in \text{Nat} \ \left[y \in \text{Nat}, w \in \text{Nat}, z \in \text{Nat} \right]} \\ \text{Succ}(z) \in \text{Nat} \ \left[y \in \text{Nat}, w \in \text{Nat}, z \in \text{Nat} \right]}$$

Inductive case

$$x = \mathsf{succ}(v) \; [v \in \mathsf{Nat}] \Rightarrow x + y = \mathsf{succ}(v) + y = \mathsf{succ}(v + y)$$

This is true, because:

- $\bullet \ \operatorname{succ}(v) + y = \operatorname{El}_{\operatorname{Nat}}(\operatorname{succ}(v), y, (w, z).\operatorname{succ}(z))$
- $\bullet \ \operatorname{succ}(v+y) = \operatorname{succ}(\operatorname{El}_{\mathsf{Nat}}(v,y,(w,z).\operatorname{succ}(z)))$
- Let $\Gamma = v \in \mathsf{Nat}, y \in \mathsf{Nat};$

 $\mathsf{El}_{\mathsf{Nat}}(\mathsf{succ}(v),y,(w,z).\,\mathsf{succ}(z)) = \mathsf{succ}(\mathsf{El}_{\mathsf{Nat}}(v,y,(w,z).\,\mathsf{succ}(z))) \in \mathsf{Nat}\ [\Gamma]\ \mathsf{derivable} : \mathsf{Proposition}(v,y,(w,z)) = \mathsf{Proposition}$

$$\begin{aligned} & \text{Var)} & \frac{\Gamma \ cont}{V_2 - \text{Nat}} & \frac{\Gamma \ cont}{V \in \text{Nat} \ [\Gamma]} & \frac{\Gamma \ cont}{\text{Nat} \ type \ [\Gamma]} & \text{Var)} & \frac{\Gamma \ cont}{V_2 + \text{Nat} \ [\Gamma]} & \frac{\Gamma \ cont}{V_2 + \text{$$

Where Γ cont derivable, because:

- $\Gamma = v \in \mathsf{Nat}, y \in \mathsf{Nat}$
- $v \in \mathsf{Nat}, y \in \mathsf{Nat}\ cont\ derivable$:

$$\begin{aligned} & \text{F-Nat)} \frac{}{-\text{Nat } type \ [\]} \\ & \text{F-c)} \frac{}{v \in \text{Nat } cont} \\ & \text{F-Nat)} \frac{}{-\text{Nat } type \ [v \in \text{Nat}]} \\ & \text{F-c)} \frac{}{v \in \text{Nat}, y \in \text{Nat } cont} \end{aligned}$$

Exercise 4

3.2 Natural Numbers Type

6. Define the predecessor operator

$$\mathbf{p}(x) \in \mathsf{Nat} \ [x \in \mathsf{Nat}]$$

such that

$$\mathbf{p}(0) = 0$$
$$\mathbf{p}(\mathsf{succ}(\mathbf{n})) = \mathbf{n}$$

Solution

The predecessor $\mathbf{p}(x)$ can be defined as:

$$\mathsf{EI}_{\mathsf{Nat}}(x,0,(w,z).\,w)$$

 $\mathbf{p}(x) \in \mathsf{Nat} \ [x \in \mathsf{Nat}] \ \text{is derivable:}$

Correctness

The definition is correct, in fact:

Base case

$$x = 0 \Rightarrow \mathbf{p}(x) = \mathbf{p}(0) = 0$$

This is true, because:

- $\mathbf{p}(0) = \mathsf{El}_{\mathsf{Nat}}(0, 0, (w, z). w)$
- $\mathsf{El}_{\mathsf{Nat}}(0,0,(w,z).\,w) = 0 \in \mathsf{Nat}\,[\,\,]\,$ derivable:

$$\begin{aligned} \text{F-Nat}) & \frac{\left[\; \right] \; cont}{\text{Nat} \; type} \; \left[\; \right]}{\text{F-Nat}} \\ \text{F-Nat}) & \frac{\text{F-Nat}}{w \in \text{Nat} \; cont}}{\text{Nat} \; type} \; \left[\; \right]} \\ \text{F-Nat}) & \frac{\left[\; \right] \; cont}{\text{Nat} \; type} \; \left[\; \right]}{\text{Nat} \; type} \; \left[\; \right]} & \text{I}_{1}\text{-Nat}) & \frac{\left[\; \right] \; cont}{0 \in \text{Nat} \; \left[\; \right]}} & \text{var}) & \frac{\text{Nat} \; type}{w \in \text{Nat}, z \in \text{Nat}} \\ \hline & w \in \text{Nat}, z \in \text{Nat} \; cont} \\ \hline & w \in \text{Nat} \; \left[\; w \in \text{Nat}, z \in \text{Nat} \right]} \\ & \text{El}_{\text{Nat}}(0,0,(w,z),w) = 0 \in \text{Nat} \; \left[\; \right] \end{aligned}$$

Inductive case

$$x = \mathsf{succ}(y) \ [y \in \mathsf{Nat}] \Rightarrow \mathbf{p}(x) = \mathbf{p}(\mathsf{succ}(y)) = y$$

This is true, because:

- $\mathbf{p}(\mathsf{succ}(y)) = \mathsf{El}_{\mathsf{Nat}}(\mathsf{succ}(y), 0, (w, z), w)$
- $\mathsf{El}_{\mathsf{Nat}}(\mathsf{succ}(y), 0, (w, z). \, w) = y \in \mathsf{Nat} \ [y \in \mathsf{Nat}] \ \mathsf{derivable} :$

 $\mathsf{EI}_{\mathsf{Nat}}(\mathsf{succ}(y), 0, (w, z), w) = y \in \mathsf{Nat} \ [y \in \mathsf{Nat}]$

Exercise 5

3.6 Martin-Löf's Intensional Propositional Equality

7. Prove that there exists a proof-term **pf** such that.

$$\mathbf{pf} \in \mathsf{Id}(\mathsf{N}_1, \star, w) \ [w \in \mathsf{N}_1]$$

is derivable.

Solution

There exists a proof-term $\mathbf{pf} = \mathsf{El}_{\mathsf{N}_1}(w,\mathsf{id}(\star)),(x).\,\mathsf{id}(x),$ such that

$$\mathbf{pf} \in \mathsf{Id}(\mathsf{N}_1, \star, w) \ [w \in \mathsf{N}_1]$$

is derivable, in fact $\mathsf{El}_{\mathsf{N}_1}(w,\mathsf{id}(\star)) \in \mathsf{Id}(\mathsf{N}_1,\star,w)$ $[\Gamma]$ is derivable:

$$\begin{array}{c} \text{var)} & \Gamma \ \text{cont} \\ \text{E-S)} & W \in \mathbb{N}_1 \ [\Gamma] \end{array} \quad \begin{array}{c} \Gamma \cdot \text{S} \cdot \frac{\Gamma \ \text{cont}}{\mathbb{N}_1 \ \text{type} \ [\Gamma]} & \text{I-S} \cdot \frac{\Gamma \ \text{cont}}{\star \in \mathbb{N}_1 \ [\Gamma]} & \text{var)} \cdot \frac{\Gamma \ \text{cont}}{w \in \mathbb{N}_1 \ [\Gamma]} \\ & \text{I-Id} \cdot \frac{\Gamma \cdot \text{S} \cdot \frac{\Gamma \ \text{cont}}{\star \in \mathbb{N}_1 \ [\Gamma]}}{\text{id}(\star) \in \mathbb{Id}(\mathbb{N}_1, \star, w) \ [\Gamma]} \end{array}$$

Where Γ cont derivable, because:

- $\Gamma = w \in \mathsf{N}_1$
- $w \in \mathsf{N}_1 \ cont \ derivable$:

F-S)
$$\frac{[\]\ cont}{\mathsf{N}_1\ type\ [\]}$$
$$\overline{w\in \mathsf{N}_1\ cont}$$

Exercise 6

3.6 Martin-Löf's Intensional Propositional Equality

8. Prove that there exists a proof-term **pf** such that.

$$\mathbf{pf} \in \mathsf{Id}(\mathsf{N}_1, x, w) \ [x \in \mathsf{N}_1, w \in \mathsf{N}_1]$$

is derivable.

Solution

There exists a proof-term $\mathbf{pf} = \mathsf{El}_{\mathsf{N}_1} \big(x, \mathsf{El}_{\mathsf{N}_1} (w, \mathsf{id}(\star)) \big), (y). \, \mathsf{id}(y),$ such that

$$\mathbf{pf} \in \mathsf{Id}(\mathsf{N}_1, x, w) \ [x \in \mathsf{N}_1, w \in \mathsf{N}_1]$$

is derivable, in fact $\mathsf{El}_{\mathsf{N}_1} ig(x, \mathsf{El}_{\mathsf{N}_1} (w, \mathsf{id}(\star)) ig) \in \mathsf{Id}(\mathsf{N}_1, x, w) \ [\Gamma]$ is derivable:

$$\begin{array}{c} \pi_1 \\ \text{var)} \\ \frac{\Gamma \ cont}{x \in \mathbb{N}_1 \ [\Gamma]} \end{array} \begin{array}{c} \pi_1 \\ \text{F-Id)} \\ \hline \\ \text{E-S)} \end{array} \begin{array}{c} \pi_1 \\ \text{var)} \\ \hline \\ \text{EI}_{\mathbb{N}_1}(x, \text{El}_{\mathbb{N}_1}(w, \text{id}(\star))) \in \text{Id}(\mathbb{N}_1, x, w) \ [\Gamma] \end{array} \end{array} \begin{array}{c} \pi_1 \\ \hline \\ \\ \text{var)} \\ \hline \\ \hline \\ \\ \text{EI}_{\mathbb{N}_1}(w, \text{id}(\star)) \in \text{Id}(\mathbb{N}_1, x, w) \ [\Gamma] \end{array} \end{array}$$

Where:

 π_1) Γ cont derivable, because:

- $\Gamma = x \in \mathbb{N}_1, w \in \mathbb{N}_1$
- $x \in \mathsf{N}_1, w \in \mathsf{N}_1 \ cont \ derivable$:

$$\begin{aligned} & \text{F-S)} & \frac{\text{[]} cont}{\mathsf{N}_1 \ type \ []} \\ & \text{F-c)} & \frac{x \in \mathsf{N}_1 \ cont}{x \in \mathsf{N}_1 \ type \ [x \in \mathsf{N}_1]} \\ & \text{F-c)} & \frac{\mathsf{N}_1 \ type \ [x \in \mathsf{N}_1]}{x \in \mathsf{N}_1, w \in \mathsf{N}_1 \ cont} \end{aligned}$$

 π_2) $\mathsf{El}_{\mathsf{N}_1}(w,\mathsf{id}(\star)) \in \mathsf{Id}(\mathsf{N}_1,\star,w)$ $[\Gamma]$ derivable:

$$\underbrace{ \begin{array}{c} \pi_1 \\ \text{var)} \\ \text{E-S)} \end{array}}_{\text{E-S)}} \underbrace{ \begin{array}{c} \pi_1 \\ \text{F-S)} \\ \hline \text{F-Id)} \end{array}}_{\text{F-Id)} \underbrace{ \begin{array}{c} \pi_1 \\ \Gamma \ cont \\ \hline \text{N}_1 \ type \ [\Gamma] \end{array}}_{\text{F-Id)} \underbrace{ \begin{array}{c} \Gamma \ cont \\ \hline \text{N}_1 \ type \ [\Gamma] \end{array}}_{\text{I-S)} \underbrace{ \begin{array}{c} \pi_1 \\ \Gamma \ cont \\ \hline \star \in \text{N}_1 \ [\Gamma] \end{array}}_{\text{Var)} \underbrace{ \begin{array}{c} \Gamma \ cont \\ \hline w \in \text{N}_1 \ [\Gamma] \end{array}}_{\text{I-Id)} \underbrace{ \begin{array}{c} \Pi_1 \\ \Gamma \ cont \\ \hline \star \in \text{N}_1 \ [\Gamma] \end{array}}_{\text{I-Id)} \underbrace{ \begin{array}{c} \Gamma \ cont \\ \hline \star \in \text{N}_1 \ [\Gamma] \end{array}}_{\text{I-Id}) \underbrace{ \begin{array}{c} \Gamma \ cont \\ \hline \text{I-Id} \end{array}}_{\text{I-Id}} \underbrace{ \begin{array}{c} \Pi_1 \\ \Gamma \ cont \\ \hline \text{I-Id} \end{array}}_{\text{I-Id}} \underbrace{ \begin{array}{c} \Pi_1 \\ \Gamma \ cont \\ \hline \text{I-Id} \end{array}}_{\text{I-Id}} \underbrace{ \begin{array}{c} \Pi_1 \\ \Gamma \ cont \\ \hline \text{I-Id} \end{array}}_{\text{I-Id}} \underbrace{ \begin{array}{c} \Pi_1 \\ \Gamma \ cont \\ \hline \text{I-Id} \end{array}}_{\text{I-Id}} \underbrace{ \begin{array}{c} \Pi_1 \\ \Gamma \ cont \\ \hline \text{I-Id} \end{array}}_{\text{I-Id}} \underbrace{ \begin{array}{c} \Pi_1 \\ \Gamma \ cont \\ \hline \text{I-Id} \end{array}}_{\text{I-Id}} \underbrace{ \begin{array}{c} \Pi_1 \\ \Gamma \ cont \\ \hline \text{I-Id} \end{array}}_{\text{I-Id}} \underbrace{ \begin{array}{c} \Pi_1 \\ \Gamma \ cont \\ \hline \text{I-Id} \end{array}}_{\text{I-Id}} \underbrace{ \begin{array}{c} \Pi_1 \\ \Gamma \ cont \\ \hline \text{I-Id} \end{array}}_{\text{I-Id}} \underbrace{ \begin{array}{c} \Pi_1 \\ \Gamma \ cont \\ \hline \text{I-Id} \end{array}}_{\text{I-Id}} \underbrace{ \begin{array}{c} \Pi_1 \\ \Gamma \ cont \\ \hline \text{I-Id} \end{array}}_{\text{I-Id}} \underbrace{ \begin{array}{c} \Pi_1 \\ \Gamma \ cont \\ \hline \text{I-Id} \end{array}}_{\text{I-Id}} \underbrace{ \begin{array}{c} \Pi_1 \\ \Gamma \ cont \\ \hline \text{I-Id} \end{array}}_{\text{I-Id}} \underbrace{ \begin{array}{c} \Pi_1 \\ \Gamma \ cont \\ \hline \text{I-Id} \end{array}}_{\text{I-Id}} \underbrace{ \begin{array}{c} \Pi_1 \\ \Gamma \ cont \\ \hline \text{I-Id} \end{array}}_{\text{I-Id}} \underbrace{ \begin{array}{c} \Pi_1 \\ \Gamma \ cont \\ \hline \text{I-Id} \end{array}}_{\text{I-Id}} \underbrace{ \begin{array}{c} \Pi_1 \\ \Gamma \ cont \\ \hline \text{I-Id} \end{array}}_{\text{I-Id}} \underbrace{ \begin{array}{c} \Pi_1 \\ \Gamma \ cont \\ \hline \text{I-Id} \end{array}}_{\text{I-Id}} \underbrace{ \begin{array}{c} \Pi_1 \\ \Gamma \ cont \\ \hline \text{I-Id} \end{array}}_{\text{I-Id}} \underbrace{ \begin{array}{c} \Pi_1 \\ \Gamma \ cont \\ \hline \text{I-Id} \end{array}}_{\text{I-Id}} \underbrace{ \begin{array}{c} \Pi_1 \\ \Gamma \ cont \\ \hline \text{I-Id} \end{array}}_{\text{I-Id}} \underbrace{ \begin{array}{c} \Pi_1 \\ \Gamma \ cont \\ \hline \text{I-Id} \end{array}}_{\text{I-Id}} \underbrace{ \begin{array}{c} \Pi_1 \\ \Gamma \ cont \\ \hline \text{I-Id} \end{array}}_{\text{I-Id}} \underbrace{ \begin{array}{c} \Pi_1 \\ \Gamma \ cont \\ \hline \text{I-Id} \end{array}}_{\text{I-Id}} \underbrace{ \begin{array}{c} \Pi_1 \\ \Gamma \ cont \\ \hline \text{I-Id} \end{array}}_{\text{I-Id}} \underbrace{ \begin{array}{c} \Pi_1 \\ \Gamma \ cont \\ \hline \text{I-Id} \end{array}}_{\text{I-Id}} \underbrace{ \begin{array}{c} \Pi_1 \\ \Gamma \ cont \\ \hline \text{I-Id} \end{array}}_{\text{I-Id}} \underbrace{ \begin{array}{c} \Pi_1 \\ \Gamma \ cont \\ \hline \text{I-Id} \end{array}}_{\text{I-Id}} \underbrace{ \begin{array}{c} \Pi_1 \\ \Gamma \ cont \\ \hline \text{I-Id} \end{array}}_{\text{I-Id}} \underbrace{ \begin{array}{c} \Pi_1 \\ \Gamma \ cont \\ \hline \text{I-Id} \end{array}}_{\text{I-Id}} \underbrace{ \begin{array}{c} \Pi_1 \\ \Gamma \ cont \\ \hline \text{I-Id} \end{array}}_{\text{I-Id}} \underbrace{ \begin{array}{c} \Pi_1 \\ \Gamma \ cont \\ \hline \text{I-Id} \end{array}}_{\text{I-Id}} \underbrace{ \begin{array}{c} \Pi_1 \\ \Gamma \ cont \\ \hline \text{I-Id} \end{array}}_{\text{I-I$$

Exercise 7

- 5 How to translate predicative logic with equality into type theory
- 15. Show that by using the **Propositional Equality with Path Induction**, for any type A and $a \in A$, there exists a proof-term \mathbf{q}

$$\mathbf{q} \in \Sigma_{z \in \Sigma_{x \in A} \; \mathsf{Id}_{\mathsf{p}}(A, a, x)} \; \forall_{w \in \Sigma_{x \in A} \; \mathsf{Id}_{\mathsf{p}}(A, a, x)} \; \mathsf{Id}_{\mathsf{p}} \big(\Sigma_{x \in A} \; \mathsf{Id}_{\mathsf{p}}(A, a, x), z, w \big)$$

First, I transform the universal quantifier into a dependent product, in order to be able to derive it in type theory. The original judgment so becomes

$$\Sigma_{z \in \Sigma_{x \in A} \ \operatorname{Id}_{\mathsf{p}}(A, a, x)} \ \Pi_{w \in \Sigma_{x \in A} \ \operatorname{Id}_{\mathsf{p}}(A, a, x)} \ \operatorname{Id}_{\mathsf{p}} \left(\Sigma_{x \in A} \ \operatorname{Id}_{\mathsf{p}}(A, a, x), z, w\right)$$

I will assume the following rules, that are essential to use the **Propositional Equality with Path Induction**:

$$\begin{split} \operatorname{F-ld_p}) & \xrightarrow{A \ type \ [\Gamma]} \ a \in A \ [\Gamma] \ b \in A \ [\Gamma] \\ & \operatorname{I-ld_p}) \xrightarrow{a \in A \ [\Gamma]} \ \operatorname{I-ld_p}) \xrightarrow{a \in A \ [\Gamma]} \\ \operatorname{E-ld_p}) & \xrightarrow{C(y,z) \ type \ [\Gamma,y \in A,z \in \operatorname{Id_p}(A,a,y)]} \ a \in A \ [\Gamma] \ b \in A \ [\Gamma] \ p \in \operatorname{Id_p}(A,a,b) \ [\Gamma] \ c \in C\big(a,\operatorname{id_p}(a)\big) \ [\Gamma] \\ & \operatorname{E-ld_p}) & \xrightarrow{E(y,z) \ type \ [\Gamma,y \in A,z \in \operatorname{Id_p}(A,a,y)]} \ a \in A \ [\Gamma] \ b \in A \ [\Gamma] \ p \in \operatorname{Id_p}(A,a,b) \ [\Gamma] \ c \in C\big(a,\operatorname{id_p}(a)\big) \ [\Gamma] \end{split}$$

These are different than the ones written in page 34 of the course notes, in fact a subscript p was added to F-Id) and I-Id), in order to distinguish the ones related to Propositional Equality with Path Induction from the ones defined for Martin-Löf's Propositional Equality. The subscript was also added to the term $\mathrm{id}(x)$, for the same reason $(\mathrm{id}(x) \in \mathrm{Id}(A,x,x))$, but $\mathrm{id}(x) \notin \mathrm{Id}_p(A,x,x)$. Another small change in $\mathrm{E}\mathrm{-Id}_p$) was that $z \in \mathrm{Id}_p(A,a,y)$, not $z \in \mathrm{Id}(A,a,y)$.

Propositional Equality with Path Induction

$$\begin{aligned} \operatorname{F-Id}) & \frac{A \; type \; [\Gamma] \quad a \in A \; [\Gamma] \quad b \in A \; [\Gamma]}{\operatorname{Id}_{\mathsf{p}}(A,a,b) \; type \; [\Gamma]} & \operatorname{I-Id}) \; \frac{a \in A \; [\Gamma]}{\operatorname{id}(a) \in \operatorname{Id}_{\mathsf{p}}(A,a,a) \; [\Gamma]} \\ \operatorname{E-Id}_{\mathsf{p}}) & \frac{C(y,z) \; type \; [\Gamma,y \in A,z \in \operatorname{Id}(A,a,y)]}{\operatorname{E-Id}_{\mathsf{p}}(p,c) \in C(b,p) \; [\Gamma]} & a \in A \; [\Gamma] \quad b \in A \; [\Gamma] \quad p \in \operatorname{Id}_{\mathsf{p}}(A,a,b) \; [\Gamma] \quad c \in C(a,\operatorname{id}(a)) \\ & \operatorname{E-Id}_{\mathsf{p}}(p,c) \in C(b,p) \; [\Gamma] \end{aligned}$$

Solution

Assuming:

$$a_1$$
) A type []

$$a_2$$
) $a \in A$ []

• Let
$$\alpha = \langle a, id_p(a) \rangle$$

• Let
$$\mathbf{q} = \langle \alpha, \lambda w, \mathsf{El}_{\Sigma} \Big(w, \mathsf{El}_{\mathsf{Id}_{\mathsf{p}}} \big(x_2, \mathsf{id}_{\mathsf{p}}(\alpha) \big) \Big) \rangle$$

• Let
$$\phi = \Sigma_{x \in A} \operatorname{Id}_{\mathbf{p}}(A, a, x)$$

• Let
$$\psi(z,w) = \operatorname{Id}_{\mathbf{p}}(\phi,z,w)$$

$$\langle \alpha, \lambda w. \, \mathsf{El}_{\Sigma} \Big(w, \mathsf{El}_{\mathsf{Id}_{\, \mathsf{p}}} \big(x_2, \mathsf{id}_{\, \mathsf{p}} (\alpha) \big) \Big) \rangle \in \Sigma_{z \in \phi} \ \Pi_{w \in \phi} \ \psi(z, w) \ [\] \ \mathsf{derivable} :$$

$$\text{I-}\Sigma) \xrightarrow{\begin{array}{c} \pi_1 \\ \text{I-}\Pi \end{array}} \xrightarrow{\text{I-}\Pi} \xrightarrow{\begin{array}{c} \text{El}_{\Sigma} \Big(w, \text{El}_{\mathsf{Id}_{\mathsf{p}}} \big(x_2, \mathsf{id}_{\mathsf{p}} (\alpha) \big) \Big) \in \psi(\alpha, w) \; [w \in \phi] \\ \hline \lambda w. \; \text{El}_{\Sigma} \Big(w, \text{El}_{\mathsf{Id}_{\mathsf{p}}} \big(x_2, \mathsf{id}_{\mathsf{p}} (\alpha) \big) \Big) \in \Pi_{w \in \phi} \; \psi(\alpha, w) \; [\;] \end{array}} \xrightarrow{\text{F-}\Pi} \xrightarrow{\begin{array}{c} \pi_3 \\ \psi(z, w) \; type \; [z \in \phi, w \in \phi] \\ \hline \Pi_{w \in \phi} \; \psi(z, w) \; type \; [z \in \phi] \\ \hline \langle \alpha, \lambda w. \; \text{El}_{\Sigma} \Big(w, \text{El}_{\mathsf{Id}_{\mathsf{p}}} \big(x_2, \mathsf{id}_{\mathsf{p}} (\alpha) \big) \Big) \Big\rangle \in \Sigma_{z \in \phi} \; \Pi_{w \in \phi} \; \psi(z, w) \; [\;] \end{array}$$

Where:

 π_1) $\alpha \in \phi$ [] derivable, because:

•
$$\alpha = \langle a, \mathrm{id}_{p}(a) \rangle$$

$$\bullet \ \phi = \Sigma_{x \in A} \ \mathrm{Id}_{\mathbf{p}}(A,a,x)$$

•
$$\langle a, \mathsf{id}_{\mathsf{p}}(a) \rangle \in \Sigma_{x \in A} \ \mathsf{Id}_{\mathsf{p}}(A, a, x)$$
 [] derivable:

$$\text{I-}\Sigma) \xrightarrow{\begin{array}{c} a_2 \\ a \in A \ [\] \end{array}} \text{F-Id}_{\mathsf{p}}) \xrightarrow{\begin{array}{c} a_2 \\ \text{id}_{\mathsf{p}}(a) \in \text{Id}_{\mathsf{p}}(A,a,a) \ [\] \end{array}} \begin{array}{c} a_1 \\ \text{ind-ty}) \xrightarrow{\begin{array}{c} A \ type \ [\] \\ x \in A \ cont \end{array}} \begin{array}{c} a_1 \\ x \in A \ cont \end{array} \begin{array}{c} a_2 \\ \text{ind-ter}) \xrightarrow{\begin{array}{c} A \ type \ [\] \\ x \in A \ cont \end{array}} \begin{array}{c} A \ type \ [\] \\ x \in A \ cont \end{array} \begin{array}{c} x \in A \ cont \end{array} \begin{array}{c} a_2 \\ \text{ind-ter}) \xrightarrow{\begin{array}{c} A \ type \ [\] \\ x \in A \ cont \end{array}} \begin{array}{c} x \in A \ cont \end{array} \\ \text{Var)} \xrightarrow{\begin{array}{c} A \ type \ [\] \\ x \in A \ cont \end{array}} \begin{array}{c} x \in A \ cont \end{array} \\ \text{Id}_{\mathsf{p}}(A,a,x) \ type \ [x \in A] \end{array}$$

$$\begin{split} \pi_2) & \ \mathsf{El}_{\Sigma}\Big(w, \mathsf{El}_{\mathsf{Id}_{\mathsf{p}}}\big(x_2, \mathsf{id}_{\mathsf{p}}(\alpha)\big)\Big) \in \psi(\alpha, w) \ [w \in \phi] \ \mathsf{derivable}, \mathsf{because}: \\ & \bullet \ \psi(\alpha, w) = \mathsf{Id}_{\mathsf{p}}(\phi, \alpha, w) \\ & \bullet \ \mathsf{El}_{\Sigma}\Big(w, \mathsf{El}_{\mathsf{Id}_{\mathsf{p}}}\big(x_2, \mathsf{id}_{\mathsf{p}}(\alpha)\big)\Big) \in \mathsf{Id}_{\mathsf{p}}(\phi, \alpha, w) \ [w \in \phi] \ \mathsf{derivable}: \end{split}$$

•
$$\psi(\alpha, w) = \mathrm{Id}_{\mathbf{p}}(\phi, \alpha, w)$$

•
$$\mathsf{El}_{\Sigma} (w, \mathsf{El}_{\mathsf{Id}_{\mathsf{p}}}(x_2, \mathsf{id}_{\mathsf{p}}(\alpha))) \in \mathsf{Id}_{\mathsf{p}}(\phi, \alpha, w) \ [w \in \phi] \ derivable:$$

$$\text{E-}\Sigma) \begin{array}{c} \pi_{3} \\ \text{E-}\Sigma) \end{array} \underbrace{\begin{array}{c} \pi_{3} \\ \text{Id}_{\mathsf{p}}(\phi,z,w) \ type \ [w \in \phi,z \in \phi] \end{array}}_{\begin{array}{c} \pi_{3} \\ \text{var)} \end{array} \underbrace{\begin{array}{c} \varphi \ type \ [\] \\ w \in \phi \ cont \\ \hline w \in \phi \ cont \\ \hline w \in \phi \ [w \in \phi] \end{array}}_{\begin{array}{c} \text{El}_{\mathsf{Id}_{\mathsf{p}}}(x_{2},\mathsf{id}_{\mathsf{p}}(\alpha)) \in \mathsf{Id}_{\mathsf{p}}(\phi,\alpha,\langle x_{1},x_{2}\rangle) \ [w \in \phi,x_{1} \in A,x_{2} \in \mathsf{Id}_{\mathsf{p}}(A,a,x_{1})] \\ \hline \\ \text{El}_{\Sigma}\Big(w,\mathsf{El}_{\mathsf{Id}_{\mathsf{p}}}(x_{2},\mathsf{id}_{\mathsf{p}}(\alpha))\Big) \in \mathsf{Id}_{\mathsf{p}}(\phi,\alpha,w) \ [w \in \phi] \end{array}$$

Where:

$$\pi_{2.1}) \ \ \mathsf{El}_{\mathsf{Id}_{\mathsf{p}}}\big(x_2,\mathsf{id}_{\mathsf{p}}(\alpha)\big) \in \mathsf{Id}_{\mathsf{p}}(\phi,\alpha,\langle x_1,x_2\rangle) \ \big[w \in \phi, x_1 \in A, x_2 \in \mathsf{Id}_{\mathsf{p}}(A,a,x_1)\big] \ \text{derivable, because:}$$

• Let
$$\Gamma = w \in \phi, x_1 \in A, x_2 \in \mathsf{Id}_{\mathsf{n}}(A, a, x_1)$$

$$\begin{array}{l} \bullet \ \ {\rm Let} \ \Gamma = w \in \phi, x_1 \in A, x_2 \in {\rm Id_p}(A,a,x_1) \\ \bullet \ \ {\rm El_{Id_p}} \big(x_2, {\rm id_p}(\alpha)\big) \in {\rm Id_p}(\phi,\alpha,\langle x_1,x_2\rangle) \ [\Gamma] \ {\rm derivable:} \end{array}$$

$$\text{E-Id}_{\mathsf{p}}) \xrightarrow{\pi_{2.2}} \text{var}) \xrightarrow{\pi_{2.3}} \text{var}) \xrightarrow{\pi_{2.3}} \text{var}) \xrightarrow{\pi_{2.3}} \text{var}) \xrightarrow{\pi_{2.3}} \text{var}) \xrightarrow{\Gamma \ cont} \text{var}) \xrightarrow{\Gamma \ cont} \text{var}) \xrightarrow{\Gamma \ cont} \xrightarrow{x_2 \in \mathsf{Id}_{\mathsf{p}}(\phi, a, x_1) \ [\Gamma]} \text{F-Id}_{\mathsf{p}}) \xrightarrow{\mathsf{ind-te})} \xrightarrow{\alpha \in \phi \ [\Gamma]} \xrightarrow{\mathsf{rad}_{2.3}} \text{var} \xrightarrow{\pi_{2.3}} \text{var}$$

$$\xrightarrow{\Gamma \ cont} \xrightarrow{x_2 \in \mathsf{Id}_{\mathsf{p}}(\phi, a, x_1) \ [\Gamma]} \text{F-Id}_{\mathsf{p}}) \xrightarrow{\mathsf{ind-te})} \xrightarrow{\alpha \in \phi \ [\Gamma]} \xrightarrow{\mathsf{rad}_{2.3}} \text{var}$$

$$\xrightarrow{\mathsf{E-Id}_{\mathsf{p}}} \xrightarrow{\mathsf{rad}_{2.3}} \text{var} \xrightarrow{\mathsf{r$$

 $\pi_{2.2}$) $\mathsf{Id}_{\mathsf{p}}(\phi, \alpha, \langle y, z \rangle) \ type \ [\Gamma, y \in A, z \in \mathsf{Id}_{\mathsf{p}}(A, a, y)]$ derivable, because:

- Let $\Sigma = \Gamma, y \in A, z \in \mathsf{Id}_{\mathsf{p}}(A, a, y)$
- $\mathsf{Id}_{\mathsf{p}}(\phi, \alpha, \langle y, z \rangle) \ type \ [\Sigma]$ derivable:

$$\begin{array}{c} \operatorname{ind-ty}) \frac{\pi_{3.1}}{\phi \ type \ [\]} \frac{\pi_{2.2.1}}{\Sigma \ cont} & \operatorname{ind-ter}) \frac{\pi_1}{\alpha \in \phi \ [\]} \frac{\pi_{2.2.1}}{\Sigma \ cont} & \pi_{2.2.2} \\ \operatorname{F-Id_p}) \frac{\phi \ type \ [\Sigma]}{\operatorname{Id_p}(\phi,\alpha,\langle y,z\rangle) \ type \ [\Sigma]} \end{array}$$

Where:

 $\pi_{2,2,1}$) Σ cont derivable, because:

- $\Sigma = \Gamma, y \in A, z \in \mathrm{Id}_{\mathbf{p}}(A, a, y)$
- $\Gamma, y \in A, z \in \mathsf{Id}_{\mathsf{p}}(A, a, y)$ cont derivable:

 $\pi_{2,2,2}$) $\langle y,z\rangle\in\phi$ [Σ] derivable, because:

- $\phi = \Sigma_{x \in A} \operatorname{Id}_{\mathbf{p}}(A, a, x)$
- $\langle y,z\rangle\in \Sigma_{x\in A}\ \operatorname{Id}_{\mathbf{p}}(A,a,x)\ [\Sigma]$ derivable:

$$\begin{array}{c} \text{var}) & \frac{\pi_{2.2.1}}{\sum cont} & \pi_{2.2.1} \\ \text{I-}\Sigma) & \frac{\sum cont}{y \in A \; [\Sigma]} & \text{var}) & \frac{\sum cont}{z \in \mathsf{Id}_{\mathsf{p}}(A,a,y) \; [\Sigma]} & \mathsf{Id}_{\mathsf{p}}(A,a,x) \; [\Sigma,x \in A] \\ \hline & \langle y,z \rangle \in \Sigma_{x \in A} \; \mathsf{Id}_{\mathsf{p}}(A,a,x) \; [\Sigma] \end{array}$$

 $\pi_{2,2,3}$) $\mathsf{Id}_{\mathsf{p}}(A,a,x)$ $[\Sigma,x\in A]$ derivable:

 $\pi_{2,3}$) Γ cont derivable, because:

- $\Gamma=w\in\phi, x_1\in A, x_2\in \mathrm{Id}_{\mathbf{p}}(A,a,x_1)$
- Let $\Delta = w \in \phi, x_1 \in A$
- $\Delta, x_2 \in \mathsf{Id}_{\mathsf{p}}(A, a, x_1)$ cont derivable:

$$\begin{array}{c} a_1 & \pi_{2.4} \\ \text{ind-ty)} & \frac{A \ type \ [\] \quad \Delta \ cont}{F \text{-Id}_{\mathsf{p}})} & \frac{a_2}{A \ type \ [\Delta]} & \text{ind-ter)} & \frac{a_2}{a \in A \ [\] \quad \Delta \ cont}}{a \in A \ [\Delta]} & \text{var)} & \frac{\Delta \ cont}{x_1 \in A \ [\Delta]} \\ \hline F\text{-c}) & \frac{\mathsf{Id}_{\mathsf{p}}(A,a,x_1) \ type \ [\Delta]}{\Delta, x_2 \in \mathsf{Id}_{\mathsf{p}}(A,a,x_1) \ cont} \\ \hline \\ A \ cont \ derivable \ because: \end{array}$$

 $\pi_{2,4}$) Δ cont derivable, because:

- $\Delta = w \in \phi, x_1 \in A \ cont$
- $w \in \phi, x_1 \in A \ cont \ derivable$:

$$\operatorname{ind-ty}) \ \frac{ \begin{matrix} a_1 \\ A \ type \ [\] \end{matrix} }{ \begin{matrix} F-c) \ \dfrac{\phi \ type \ [\] }{w \in \phi \ cont} \end{matrix} }$$

$$\overline{ \begin{matrix} F-c) \ \dfrac{A \ type \ [w \in \phi] }{w \in \phi, x_1 \in A \ cont} \end{matrix} }$$

 $\pi_3) \ \ \psi(z,w) \ type \ [z \in \phi, w \in \phi]$ derivable, because:

- $\psi(z,w) = \operatorname{Id}_{\mathbf{p}}(\phi,z,w)$
- $\mathsf{Id}_{\mathtt{p}}(\phi, z, w) \ type \ [z \in \phi, w \in \phi]$ derivable:

$$\begin{array}{c} \text{ind-ty)} \frac{\pi_{3.1}}{\phi \ type \ [\]} \quad x \in \phi, w \in \phi \ cont \\ \text{F-Id}_{\mathbf{p}}) \quad \frac{\phi \ type \ [z \in \phi, w \in \phi \]}{\text{Id}_{\mathbf{p}}(\phi, z, w) \ type \ [z \in \phi, w \in \phi \]} \end{array} \quad \text{var)} \quad \begin{array}{c} \pi_{3.2} \\ z \in \phi, w \in \phi \ cont \\ \hline z \in \phi, w \in \phi \ cont \\ \hline w \in \phi \ [z \in \phi, w \in \phi \] \end{array}$$

Where:

 $\pi_{3,1}$) ϕ type [] derivable, because:

- $\bullet \ \phi = \Sigma_{x \in A} \ \operatorname{Id}_{\mathbf{p}}(A,a,x)$
- $\Sigma_{x \in A} \operatorname{Id}_{p}(A, a, x) \ type []$ derivable:

$$\operatorname{ind-ty}) \begin{array}{c} a_{1} \\ A \ type \ [\] \\ \hline F-\operatorname{Id}_{\mathsf{p}}) \xrightarrow{A \ type \ [x \in A]} \end{array} \begin{array}{c} a_{1} \\ x \in A \ cont \\ \hline \\ F-\Sigma) \xrightarrow{\begin{bmatrix} \operatorname{Id}_{\mathsf{p}}(A,a,x) \ type \ [x \in A] \\ \Sigma_{x \in A} \ \operatorname{Id}_{\mathsf{p}}(A,a,x) \ type \ [x \in A] \\ \hline \\ \Sigma_{x \in A} \ \operatorname{Id}_{\mathsf{p}}(A,a,x) \ type \ [x \in A] \\ \hline \end{array} \begin{array}{c} a_{1} \\ F-c) \xrightarrow{A \ type \ [\]} \\ \hline x \in A \ cont \\ \hline x \in A \ cont \\ \hline x \in A \ [x \in A] \\ \hline \end{array}$$

 $\pi_{3,2}$) $z \in \phi, w \in \phi \ cont \ derivable$:

ind-ty)
$$\frac{\pi_{3.1}}{\phi \ type \ [\]} \text{F-c}) \frac{\phi \ type \ [\]}{z \in \phi \ cont}$$
$$F-c) \frac{\phi \ type \ [z \in \phi]}{z \in \phi, w \in \phi \ cont}$$