1 Definitions

1.1 Alphabet

Let Σ be a finite set of symbols (alternatively called characters), called the alphabet. No assumption is made about the nature of the symbols.

1.2 Strings

A string (or word) over Σ is any finite sequence of symbols from Σ . For example if $\Sigma = \{0, 1\}$ then 01011 is a string over Σ .

The **length of a string** S **is the number of symbols** in s (the length of the sequence) and can be any non-negative integer. It is often denoted as |S|.

The **empty string** is the unique string over Σ of length 0, and is denoted as ε .

The set of all strings over Σ of length n is denoted Σ^n . For example, if $\Sigma = 0, 1$, then $\Sigma^2 = \{00, 01, 10, 11\}$. Note that $\Sigma^0 = \{\varepsilon\}$ for any alphabet Σ .

The set of all strings of any length over Σ is the Kleene closure of Σ and is denoted as Σ^* . Also:

$$\Sigma^* = \bigcup_{n \in N \cup \{0\}} \Sigma^n$$

The set of all non-empty strings over Σ is denoted by Σ^+

1.3 Subsequences

A subsequence of S is a string that can be obtained by deleting any number (from 0 to n) of non-consecutive characters from S, including S and the empty string too.

1.3.1 Let us count: subsequences

A subsequence of S can be described by a binary string B of n elements, telling us for each position i if the corresponding character is kept in the subsequence (1) or deleted (0).

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$$S = apple$$
, $B = 11010 \rightarrow subsequence = apl$

Thus, the number of subsequences corresponds to the number of binary strings of length n, hence

$$subsequences = 2^n$$

Notice that if two subsequences produced by different choices are **identical**, **they will be counted twice**.

1.4 Substrings

A string s is said to be a **substring** (or factor) of t if there exist (possibly empty) strings u and v such that t = usv.

Given a string t, suffixes and prefixes are special substrings of t.

1.4.1 Let us count: substrings

Given a string of length n:

- We have 1 substring of length n (starting at 1)
- We have 2 substrings of length n-1 (starting at 1 and 2)
- We have 3 substrings of length n-2 (starting at 1,2 and 3)
- ...
- We have n-i+1 substrings of length i (starting at all positions from 1 to n-i+1)
- We have n substrings of length 1 (starting at every position)

Plus, 1 substring of length 0.

This brings the **total number of substrings** of a string of length n to

$$substrings = \sum_{i=1}^{n} i = \frac{n(n+1)}{2} + 1$$

1.5 Prefixes and suffixes

A string s is said to be a **prefix** of t if there exists a string u such that t = su. If u is nonempty, s is said to be a *proper* prefix of t.

Simmetrically, a string s is said to be a **suffix** of t if there exists a string u such that t = us. If u is nonempty, s is said to be a *proper* suffix of t.

1.5.1 Let us count: prefixes and suffixes

The number of all posssible prefixes/suffixes of a string of length n is

$$prefixes/suffixes = n + 1$$

n are non-empty, +1 because we include also the empty string as prefix/suffix.

1.6 Reverse and palindrome

The **reverse** of a string is a string with the same symbols but in reverse order. For example, if s = abc (where a, b, and c are symbols of the alphabet), then the reverse of s is cba.

A string that is the reverse of itself is called a **palindrome**, which also includes the empty string and all strings of length 1.

1.7 Rotations and permutations

A string s = uv is said to be a **rotation** of t if t = vu. For example, if $\Sigma = \{0, 1\}$ the string 0011001 is a rotation of 0100110, where u = 00110 and v = 01.

1.7.1 Let us count: rotations

Given a string S, with |S| = n, we denote **S=AB**, where **A** is a prefix substring, and **B** a suffix substring.

For every possible pair A, B the string **R=BA** is a rotation of **S**, including S itself, that is, the case where $A = \varepsilon$ and B = S.

Since we have n different non empty suffixes B, the number of rotations is

$$rotations = n$$

It is not n+1 because the two suffixes B=S and $B=\varepsilon$ produce exactly the same rotation, given by S itself.

1.7.2 Let us count: permutations

Given a string S, with |S| = n, the number of permutations of the characters in S is

$$permutations = n!$$

Notice that the formula does not consider the case of indentical permutations: that is, a string of n identical characters, will have n identical permutations.