

ICPC TEMPLATE



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This template is a supplementary version.

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1 Standard Solution Template

1.1 support bits/stdc++.h

```
#include <bits/stdc++.h>
#define lc (o<<1)</pre>
#define rc ((o<<1)|1)
#define PB push_back
#define MK make_pair
using namespace std;
#define DebugP(x) cout << "Line" << __LINE__</pre>
    << " " << #x << "=" << x << endl
const int maxn = 3000 + 5;
const int modu = 998244353; // 1e9 + 7
const int inf = 0x3f3f3f3f;
const double eps = 1e-5;
const int dx[4] = \{1, 0, -1, 0\};
const int dy[4] = \{0, 1, 0, -1\};
typedef long long LL;
void read(LL &x) {
    x=0;int f=0;char ch=getchar();
    while (ch<'0'||ch>'9') \{f|=(ch=='-'); ch=
        getchar();}
    while (ch \ge 0' \& ch \le 9') \{x = (x < 1) + (x < 3) + (x < 1) \}
        ch^48);ch=getchar();}
    x=f?-x:x;
    return;
void read(int &x) { LL y; read(y); x = (int)y
void read(char &ch) { char s[3]; scanf("%s",
    s); ch = s[0]; }
void read(char *s) { scanf("%s", s); }
template < class T, class ... U> void read (T &x,
     U& ... u) { read(x); read(u...); }
int main() {
    // freopen ("my. txt ", "w", stdout );
    ios::sync_with_stdio(0);
    cin.tie(0);
    return 0;
```

1.2 unsupport bits/stdc++.h

```
#include <cstdio>
#include <cstring>
#include <iostream>
#include <algorithm>
#include <cmath>
#include <set>
#include <map>
#include <queue>
#include <cstdlib>
#include <string>

using namespace std;

int main() {
    // freopen ("my.txt", "w", stdout);
    return 0;
}
```

1.3 Python Template

```
for T in range(0, int(input())): #T组数据
   N= int( input())
   n,m= map( int, input().split())
   s= input()
   s=[ int(x) for x in
       input().split()] #一行输入的数组
   a[1:]=[int(x) for x in
       input().split()] #从下标1开始读入一行
   for i in range(0, len(s)):
       a,b= map( int, input().split())
while True: #未知多组数据
   try:
       \#n,m=map(int,input(),split())
       \#print(n+m,end="\n")
   except EOFError: #捕获到异常
       break
```

2 GRAPH

2 Graph

2.1 Network Flow

2.1.1 Dinic

```
struct Edge {
   int from, to, cap, flow;
   Edge(int from=0, int to=0, int cap=0, int
        flow=0):
       from(from), to(to), cap(cap), flow(
           flow) {}
};
vector<Edge> edges;
vector<int> g[maxn];
int d[maxn], cur[maxn];
void addedge(int u, int v, int cap) {
    edges.push_back(Edge(u, v, cap));
   g[u].push_back(edges.size()-1);
    edges.push_back(Edge(v, u, 0));
   g[v].push_back(edges.size()-1);
 * 时间复杂度: O(n^2*m)
 * 对于特殊的图,所有容量为1的: O(min(n^0.67,
    m^0.5)*m)
               二分图最大匹配: O(n^0.5*m)
bool BFS(int s, int t) {
   memset(d, -1, sizeof(d));
   queue<int> Q;
   Q.push(s);
   d[s] = 0;
   while (!Q.empty()) {
       int x = Q.front(); Q.pop();
       for (int i = 0; i < g[x].size(); ++i)</pre>
           Edge &e = edges[g[x][i]];
           if (d[e.to] == -1 \&\& e.cap > e.
               flow) {
               d[e.to] = d[x] + 1;
               Q.push(e.to);
           }
```

```
}
   return d[t] > -1;
int DFS(int x, int a, int t) {
   if (x == t || a == 0) return a;
   int flow = 0, f;
   for (int &i = cur[x]; i < g[x].size(); ++</pre>
        i) {
       Edge &e = edges[g[x][i]];
        if (d[x] + 1 == d[e.to] && (f = DFS(e
            .to, min(a, e.cap-e.flow), t)) >
           0) {
           flow += f;
           e.flow += f;
           edges[g[x][i]^1].flow -= f;
           a -= f;
           if (a == 0) break;
       }
   return flow;
int MaxFlow(int s, int t) {
   int res = 0;
   while (BFS(s, t)) {
       for (int i = 0; i <= n; ++i) cur[i] =</pre>
       res += DFS(s, inf, t);
   return res;
```

2.1.2 Edmonds-Karp

```
queue<int> Q;
int imp[maxn], p[maxn];

// 不断寻找增广路增广

int MaxFlow(int s, int t) {
    int res = 0;
    for (;;) {
        memset(imp, 0, sizeof(imp));
        while (!Q.empty()) Q.pop();
        Q.push(s);
```

2.2 DSU on Tree

```
imp[s] = inf;
    while (!Q.empty()) {
        int x = Q.front(); Q.pop();
       for (int i = 0; i < g[x].size();</pre>
           ++i) {
           Edge &e = edges[g[x][i]];
           if (!imp[e.to] && e.cap > e.
               flow) {
               p[e.to] = g[x][i];
               imp[e.to] = min(imp[x], e.
                   cap-e.flow);
               Q.push(e.to);
           }
       if (imp[t]) break;
    if (!imp[t]) break;
    for (int u = t; u != s; u = edges[p[u
       ]].from) {
        edges[p[u]].flow += imp[t];
        edges[p[u]^1].flow -= imp[t];
   }
   res += imp[t];
return res;
```

2.2 DSU on Tree

4 3 MATHEMATICS

3 Mathematics

3.1 Number Theory

3.1.1 Linear Inverse Modulo

```
const int maxn = 2e5 + 5;
const int modu = 1e9 + 7;
long long inv[maxn]; // k在模modu的意义下的逆
元是inv[k]
inv[1] = 1;
for (int i = 2; i < maxn; ++i)
    inv[i] = (modu - (modu/i))*inv[modu%i]%
    modu;
```

3.1.2 Quick Power

```
LL powmod(LL a, LL b, LL modu) {
    LL res = 1;
    for (a %= modu; b; b >>= 1, a = a*a%modu)
        if (b&1) res = res * a % modu;
    return res;
}
```

3.1.3 Baby-Step-Giant-Step Algorithm

$$a^x \equiv b \pmod{n}$$

```
// no solution = -1
int log_mod(int a, int b, int n) {
   LL m, v, e = 1, i;
   m = S; // S = int(sqrt(n) + 0.5)
   v = powmod(powmod(a, m), n-2);
   unordered_map<int, int> x;
   x[1] = 0;
   for (i = 1; i < m; ++i) {
        e = e*a %n;
        if (!x.count(e)) x[e] = i;
   }
   for (i = 0; i < m; ++i) {
        if (x.count(b)) return (LL)i*m + x[b]
        ];
        b = (LL)b*v%n;
   }
   return -1;
}</pre>
```

3.1.4 Möbius Inversion

3.1.5 Dujiao Sieve

常见积性函数:

$$\epsilon(n) = [n = 1]$$

$$I(n) = 1$$

$$id(n) = n$$

$$d(x) = \sum_{i|n} 1$$

$$\sigma(x) = \sum_{i|n} i$$

$$\phi(i) = \sum_{i|n} [gcd(x, i) = 1]$$

$$\mu(x) = \begin{cases} 1 & x = 1 \\ (-1)^k & \prod_{i=1}^k q_i = 1 \\ 0 & \max\{q_i\} > 1 \end{cases}$$

狄利克雷卷积: 设 f,g 是两个数论函数,则有 $(f*g)(n) = \sum_{d|n} f(d)g(\frac{n}{d})$

常见性质:

1.
$$\mu * I = \epsilon$$

2. $\phi * I = id$
3. $\mu * id = \phi$

最后要有形如:

$$g(1)S(1) = \sum_{i=1}^{n} (f * g)(i) - \sum_{i=2}^{n} g(i)S(\lfloor \frac{n}{i} \rfloor)$$

```
#include <algorithm>
#include <cstdio>
#include <cstring>
#include <map>
using namespace std;
const int maxn = 2000010;
typedef long long 11;
11 T, n, pri[maxn], cur, mu[maxn], sum mu[
   maxn];
bool vis[maxn];
map<11, 11> mp_mu;
11 S_mu(11 x) {
 if (x < maxn) return sum_mu[x];</pre>
 if (mp_mu[x]) return mp_mu[x];
 11 ret = 111;
 for (11 i = 2, j; i <= x; i = j + 1) {
    j = x / (x / i);
```

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```
ret -= S_mu(x / i) * (j - i + 1);
 }
 return mp_mu[x] = ret;
11 S_phi(11 x) {
 ll ret = 011;
 for (11 i = 1, j; i <= x; i = j + 1) {</pre>
   j = x / (x / i);
   ret += (S_mu(j) - S_mu(i - 1)) * (x / i)
        * (x / i);
 return ((ret - 1) >> 1) + 1;
int main() {
 scanf("%11d", &T);
 mu[1] = 1;
 for (int i = 2; i < maxn; i++) {</pre>
   if (!vis[i]) {
     pri[++cur] = i;
     mu[i] = -1;
   for (int j = 1; j <= cur && i * pri[j] <</pre>
       maxn; j++) {
      vis[i * pri[j]] = true;
      if (i % pri[j])
        mu[i * pri[j]] = -mu[i];
      else {
       mu[i * pri[j]] = 0;
       break;
     }
 for (int i = 1; i < maxn; i++) sum_mu[i] =</pre>
      sum_mu[i - 1] + mu[i];
 while (T--) {
   scanf("%lld", &n);
   printf("%lld %lld\n", S_phi(n), S_mu(n));
 }
 return 0;
```

3.1.6 Chinese Remainder Theory

```
// x mod a[i] == b[i] mod a[i]

// gcd(a[i], a[j]) == 1 (i != j)

LL crt(LL *a, LL *b, int n) {

   LL res = 0;
```

```
LL tota = 1;
    for (int i = 0; i < n; ++i) tota *= a[i];</pre>
    for (int i = 0; i < n; ++i) {</pre>
        LL m = tota / a[i];
        LL r, tmp;
        exgcd(m, a[i], r, tmp);
        r = (r\%a[i]+a[i])\%a[i];
        res = (res + b[i]*m%tota*r%tota) %
            tota;
    return res;
// extended version
// x \mod a[i] == b[i] \mod a[i]
// \gcd(a[i], a[j]) >= 1
LL excrt(LL *a, LL *b, int n) {
    for (int i = 1; i < n; ++i) {</pre>
        LL x, y;
        LL g = exgcd(a[0], a[i], x, y);
        x \%= a[i];
        a[0] /= g;
        x = ((_int128)(b[i] - b[0])%a[i]*x%a
            [i] + a[i])%a[i];
        LL lcm = a[0]*a[i];
        b[0] = ((_int128)x*a[0] + b[0])%lcm;
        a[0] = lcm;
    return b[0];
```

3.1.7 Floor Sum

```
// sigma(floor (a*i+b/c)) 0 <= i <= n
LL floor_sum(LL a, LL b, LL c, LL n) {
   if (a == 0) return (n+1)*(b/c);
   if (n == 0) return b/c;
   if (n < 0) return 0;
   LL res = (a/c)*n*(n+1)/2 + (n+1)*(b/c);
   a %= c;
   b %= c;
   LL m = (a*n+b)/c;
   return res + n*m - floor_sum(c, c-b-1, a, m-1);
}
// atcoder</pre>
```

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```
using ll = long long;
ll floor_sum(ll n, ll m, ll a, ll b) {
   11 \text{ ans} = 0;
   if (a >= m) {
        ans += (n - 1) * n * (a / m) / 2;
        a \%= m;
   if (b >= m) {
        ans += n * (b / m);
        b \% = m;
   }
   11 y_max = (a * n + b) / m, x_max = (
        y_max * m - b);
   if (y_max == 0) return ans;
    ans += (n - (x_max + a - 1) / a) * y_max;
    ans += floor_sum(y_max, a, m, (a - x_max
        % a) % a);
   return ans;
```

3.1.8 Min 25 Sieve

3.1.9 $\pi(x)$ $\pi(x) \sim \frac{x}{\ln x}$

```
LL n;
vector<LL> p, val;
bool vis[maxn+_];
LL f[maxn+_];
int id1[maxn+_], id2[maxn+_];
inline int getid(LL k) {
    if (k <= n/k) return id1[k];</pre>
    return id2[n/k];
int main() {
    scanf("%lld", &n);
    memset(vis, 0, sizeof(vis));
    int m = sqrt(n);
    for (LL i = 2; i <= m; ++i) {</pre>
        if (!vis[i]) p.PB(i);
        for (int j = 0; j < SZ(p) && i*p[j]</pre>
            <= m; ++j) {
            vis[i*p[j]] = 1;
            if (i % p[j] == 0) break;
```

```
}
}
for (LL L = 1, R; L <= n; L = R+1) {
    R = n / (n/L);
    val.PB(R);
    if (R <= n/R) id1[R] = SZ(val)-1;
    else id2[n/R] = SZ(val)-1;
}

for (int i = 0; i < SZ(val); ++i) f[i] =
    val[i] - val[i]/2 - (val[i] == 1);

for (int i = 1; i < SZ(p); ++i)
    for (int j = SZ(val)-1; j >= 0 && val
        [j] >= p[i]*p[i]; --j)
        f[j] += - f[getid(val[j]/p[i])] +
        i;
    printf("%11d\n", f[SZ(val)-1]);
    return 0;
}
```

3.1.10 Lucas' Theorem

3.2 Karatsuba Multiply

3.3 Fast Fourier Transform 7

```
z0 = karatsuba_polymul(m - 1, a, b);
 z2 = karatsuba_polymul(n - m, a + m, b + m)
 // 计算 z1
 // 临时更改, 计算完毕后恢复
 for (int i = 0; i + m <= n; ++i) a[i] += a[</pre>
 for (int i = 0; i + m <= n; ++i) b[i] += b[
     i + m];
 z1 = karatsuba_polymul(m - 1, a, b);
 for (int i = 0; i + m <= n; ++i) a[i] -= a[</pre>
     i + m];
 for (int i = 0; i + m <= n; ++i) b[i] -= b[</pre>
     i + m];
 for (int i = 0; i <= (m - 1) * 2; ++i) z1[i
     ] = z0[i];
 for (int i = 0; i \le (n - m) * 2; ++i) z1[i
     ] = z2[i];
 // 由 z0、z1、z2 组合获得结果
 for (int i = 0; i <= (m - 1) * 2; ++i) r[i]</pre>
      += z0[i];
 for (int i = 0; i <= (m - 1) * 2; ++i) r[i</pre>
     + m] += z1[i];
 for (int i = 0; i <= (n - m) * 2; ++i) r[i</pre>
     + m * 2] += z2[i];
 delete[] z0;
 delete[] z1;
 delete[] z2;
 return r;
// 计算a*b=c, 时间复杂度是O(n^1.585)
void karatsuba_mul(int a[], int b[], int c[])
 int *r = karatsuba_polymul(LEN - 1, a, b);
 memcpy(c, r, sizeof(int) * LEN);
 for (int i = 0; i < LEN - 1; ++i)</pre>
   if (c[i] >= 10) {
     c[i + 1] += c[i] / 10;
     c[i] %= 10;
   }
 delete[] r;
```

3.3 Fast Fourier Transform

```
// f是系数数组,处理完后,f表示:
// rev=1,是点表示法
// rev=-1,除N后是系数
// N=2^n
typedef complex <double > Comp; // 先导入头文件
    complex
void DFT(Comp *f, int N, int rev) {
    if (N == 1) return;
    for (int i = 0; i < N; ++i) tmp[i] = f[i</pre>
    for (int i = 0; i < N; ++i)</pre>
       if (i%2) f[i/2+N/2] = tmp[i];
       else f[i/2] = tmp[i];
    Comp *g = f, *h = f + \mathbb{N}/2;
    DFT(g, N/2, rev); DFT(h, N/2, rev);
    // c[N] = cos(2*pi/N), s[N] = sin(2*pi/N)
    Comp w(c[N], s[N]*rev), cur(1, 0);
    for (int k = 0; k < N/2; ++k) {
        tmp[k] = g[k] + cur*h[k];
       tmp[k+N/2] = g[k] - cur*h[k];
       cur *= w;
    for (int i = 0; i < N; ++i) f[i] = tmp[i</pre>
       ];
// Important !!!!
//for (N = 1; N < n+m; N < < = 1) 
      s[N] = sin(2*pi/N);
     c[N] = cos(2*pi/N);
//s[N] = \sin(2*pi/N);
//c[N] = cos(2*pi/N);
// no recursion
inline void change(Comp y[], int len) {
    int i, j, k;
    for (int i = 1, j = len / 2; i < len - 1;</pre>
        i++) {
       if (i < j) swap(y[i], y[j]);</pre>
       // 交换互为小标反转的元素, i < j 保证交换
            一次
       //i 做正常的 +1, j 做反转类型的 +1,
            始终保持 i 和 j 是反转的
       k = len / 2;
       while (j \ge k) {
           j = j - k;
```

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```
k = k / 2;
       }
       if (j < k) j += k;
   }
* 做 FFT
 *len 必须是 27k 形式
*on == 1 时是 DFT, on == -1 时是 IDFT
inline void fft(Comp y[], int len, int on) {
   change(y, len);
   for (int h = 2; h <= len; h <<= 1) {</pre>
       Comp wn(c[h], on*s[h]);
       for (int j = 0; j < len; j += h) {</pre>
           Comp w(1, 0);
           for (int k = j; k < j + h / 2; k
               ++) {
               Comp u = y[k];
               Comp t = w * y[k + h / 2];
               y[k] = u + t;
               y[k + h / 2] = u - t;
               w = w * wn;
           }
       }
   if (on == -1) {
       for (int i = 0; i < len; ++i)</pre>
           y[i] = Comp(floor(y[i].real()/len
               +0.5), 0);
   }
```

3.4 Number Theory Transform

```
// convolution for F(x) * G(x)

/* N = 2 \hat{\ }k

* init : f[i] = ... 0 <= i < N

* g[i] = ... 0 <= i < N

* call :

* NTT(f, N, 1);

* NTT(g, N, 1);

* for (int \ i = 0; \ i < N; ++i) \ f[i] = (LL) f[i] *

g[i];

* NTT(f, N, -1);

*/
```

```
inline void change(int y[], int len) {
    for (int i = 1, j = len / 2, k; i < len -</pre>
         1; i++) {
        if (i < j) swap(y[i], y[j]);</pre>
        k = len / 2;
        while (j \ge k) {
            j = j - k;
            k = k / 2;
        }
        if (j < k) j += k;
    }
// g is the primitive root of modu
inline void NTT(int y[], int len, int on) {
    int g = 3; // modu = 998244353 or
        1004535809
    change(y, len);
    for (int h = 2; h <= len; h <<= 1) {</pre>
        LL wn = powmod(g, (modu-1)/h);
        if (on == -1) wn = powmod(wn, modu-2)
        for (int j = 0; j < len; j += h) {</pre>
            LL w = 1;
            for (int k = j; k < j + h / 2; k
                ++) {
                LL u = y[k];
                LL t = (w * 111 * y[k + h /
                    2])%modu;
                y[k] = (u + t) \text{mod}u;
                y[k + h / 2] = (u - t + modu)
                    %modu;
                w = w * wn \% modu;
        }
    }
    if (on == -1) {
        LL INV = powmod(len, modu-2);
        for (int i = 0; i < len; ++i)</pre>
            y[i] = 111 * y[i] * INV % modu;
    }
```

3.5 Lagrange Insertion Value Method

$$f(k) = \sum_{i=1}^{n} y_i \prod_{i \neq j} \frac{k - x_j}{x_i - x_j}$$

```
#include <cstdio>
const int maxn = 2010;
using ll = long long;
11 \mod = 998244353;
11 n, k, x[maxn], y[maxn], ans, s1, s2;
ll powmod(ll x, ll n) {
 11 ret = 111;
 while (n) {
   if (n & 1) ret = ret * x % mod;
   x = x * x \% mod;
   n >>= 1;
 return ret;
11 inv(11 x) { return powmod(x, mod - 2); }
int main() {
 scanf("%lld%lld", &n, &k);
 for (int i = 1; i <= n; i++) scanf("%lld%</pre>
     11d'', x + i, y + i);
 for (int i = 1; i <= n; i++) {</pre>
   s1 = y[i] \% mod;
   s2 = 111;
   for (int j = 1; j <= n; j++)</pre>
     if (i != j) s1 = s1 * (k - x[j]) % mod,
           s2 = s2 * (x[i] - x[j]) % mod;
   ans += s1 * inv(s2) % mod;
 printf("%lld\n", (ans % mod + mod) % mod);
 return 0;
}
// Lagrange interpolation O(n^2)
// n is the highest degree of polynomial
// return f(x) where x = n
// x[i] = i-1, x[i] sorted from low to high
LL lagrange(LL n, LL *x, LL *y, int m) {
   LL res = 0;
   for (int i = 0; i < m; ++i) {</pre>
       LL tmp = 1;
```

10 4 DATA STRUCTURE

4 Data Structure

4.1 Treap

```
struct Treap
    struct Node
        Node *ch[2];
        int r, v, cnt;
        Node()
        {
            *this = NULL;
        Node(int v) : v(v)
            r = rand()/10*rand();
            ch[0] = ch[1] = NULL;
            cnt = 1;
        bool operator < (const Node &rhs)</pre>
            const
        {
            return r < rhs.r;</pre>
        }
        int cmp(int x) const
            if (x == v) return -1;
            return x > v;
        }
    }*root;
    void rotate(Node *&o, int d)
    {
        Node *k = o->ch[d^1];
        o->ch[d^1] = k->ch[d];
        k \rightarrow ch[d] = o;
        o = k;
    void insert(Node *&o, int x)
    {
```

```
if (o == NULL)
         o = new Node(x);
         return;
    int k = o \rightarrow cmp(x);
    if (k == -1) o->cnt++;
    else
    {
         insert(o->ch[k], x);
         if (o->ch[k]->r > o->r) rotate(o,
              k^1);
    }
}
void remove(Node *&o, int x)
    int d = o \rightarrow cmp(x);
    if (d == -1)
         if (o->cnt == 0 || (--o->cnt) ==
             0)
         {
             if (o->ch[0] == NULL)
                  o = o \rightarrow ch[1];
             else if (o->ch[1] == NULL)
                 o = o \rightarrow ch[0];
             else
             {
                  int d1 = o \rightarrow ch[0] \rightarrow r > o \rightarrow
                      ch[1]->r;
                 rotate(o, d1);
                 remove(o->ch[d1], x);
         }
         return;
    remove(o->ch[d], x);
}
int finds(Node *o, int x)
    if (o == NULL) return INF;
    int d = o \rightarrow cmp(x);
    if (d == -1) return 0;
    if (o->ch[d] == NULL) return abs(x-o
         ->v);
```

4.2 Splay Tree

4.2 Splay Tree

4.3 Two-dimensional Segment Tree

4.4 Mo's Algorithm

```
struct Interval {
   int 1, r, t, id;
   int k;
   Interval(int 1=0, int r=0, int t=0, int
       id=0, int k=0):
       l(1), r(r), t(t), id(id), k(k) {}
};
struct UpdateOp {
   int p, x;
   UpdateOp(int p=0, int x=0): p(p), x(x) {}
};
int n, q, S, st, ed, qt; // S is the size of
    one block
vector<Interval> intervals;
vector<UpdateOp> up;
bool cmp(Interval A, Interval B) {
   return (A.1-1)/S < (B.1-1)/S ||
           ((A.1-1)/S == (B.1-1)/S && (A.r)
               -1)/S < (B.r-1)/S) | |
           ((A.1-1)/S == (B.1-1)/S && (A.r)
               -1)/S == (B.r-1)/S && A.t < B
               .t);
   // 没有更新操作时需要修改
```

```
void add(int loc) {
void del(int loc) {
int main() {
    // freopen ("input . txt ", "r ", stdin );
    ios::sync_with_stdio(false);
    cin.tie(0);
    read(n); read(q);
    for (int i = 1; i <= n; ++i) {</pre>
        read(a[i]);
        tmpa[i] = a[i];
    for (int i = 0; i < q; ++i) {</pre>
        int cmd;
        read(cmd);
        if (cmd == 1) {
            int 1, r, k;
            read(1); read(r); read(k);
            intervals.push_back(Interval(1, r
                 , up.size(), i, k));
        }
        else {
            ans[i] = -2;
            int p, x;
            read(p); read(x);
            up.push_back(UpdateOp(p, x));
    S = (int)pow(2.0*n*n, 1.0/3); // without
        update \ operation \ , \ S=(int)(n/sqrt(q))
    sort(intervals.begin(), intervals.end(),
        cmp);
    tot = 0;
    qt = 0;
    st = 1; ed = 1;
    add(1);
    for (auto qj: intervals) {
        if (ed <= qj.r) {</pre>
            while (ed < qj.r) add(++ed);</pre>
            while (ed > qj.r) del(ed--);
            while (st < qj.l) del(st++);</pre>
            while (st > qj.l) add(--st);
```

12 4 DATA STRUCTURE

```
else {
        while (st < qj.l) del(st++);</pre>
        while (st > qj.l) add(--st);
        while (ed < qj.r) add(++ed);</pre>
        while (ed > qj.r) del(ed--);
    }
    if (qt > qj.t) { // recover the update
         operation
    }
    for (; qt < qj.t; ++qt) { // update</pre>
         operation
        if (st <= up[qt].p && up[qt].p <=</pre>
             ed)
            del(up[qt].p);
        a[up[qt].p] = up[qt].x;
        if (st <= up[qt].p && up[qt].p <=</pre>
             ed)
            add(up[qt].p);
    ans[qj.id] = solve();
for (int i = 0; i < q; ++i)</pre>
    if (ans[i] > -2) cout << ans[i] << "\</pre>
        n";
return 0;
```

5 String

5.1 KMP

```
int f[maxn];
void getfail(char *P, int *f, int m) {
   f[0] = 0;
   f[1] = 0;
   for (int i = 1; i < m; ++i) {</pre>
        int j = f[i];
       while (j && P[i] != P[j]) j = f[j];
       if (P[i] == P[j]) f[i+1] = j+1;
       else f[i+1] = 0;
   }
int find(char *T, char *P, int *f, int n, int
    m) {
   int res = 0;
   getfail(P, f);
   for (int i = 0, j = 0; i < n; ++i) {</pre>
       while (j && P[j] != T[i]) j = f[j];
       if (P[j] == T[i]) j++;
       if (j == m) { // 出现一次
           res++;
           j = f[m];
       }
   }
   return res;
```

5.2 Trie

```
void insert(const string &str, int v) {
        int i, u;
        for (i = 0, u = 0; i < str.length();</pre>
            int c = str[i]-'a';
            if (ch[c][u] == -1) {
                newNode();
                ch[c][u] = ch[c].size()-1;
            }
            u = ch[c][u];
        }
        val[u] = v;
    }
    int find(const string &str) {
        int i, u;
        for (i = 0, u = 0; i < str.length();</pre>
            ++i) {
            int c = str[i] - 'a';
            if (ch[c][u] == -1) return -1;
            u = ch[c][u];
        }
        return val[u];
    }
};
```

6 Computational Geometry

6.1 2D Template

```
using Point = pair<double, double>;
Point operator+(Point &A, Point &B) { return
   Point(A.X+B.X, A.Y+B.Y); }
Point operator-(Point &A, Point &B) { return
   Point(A.X-B.X, A.Y-B.Y); }
Point operator*(double a, Point A) { return
   Point(a*A.X, a*A.Y); }
Point operator*(Point A, double a) { return a
   *A; }
double Dot(Point A, Point B) { return A.X*B.X
     + A.Y*B.Y; }
double Cross(Point A, Point B) { return A.X*B
    .Y - A.Y * B.X; }
double Distance(Point A, Point B) { return
    sqrt(Dot(A-B, A-B)); }
bool Circle(Point A, Point B, Point C, Point
   &cir) {
   double a1 = 2*(B.X - A.X), b1 = 2*(B.Y -
       A.Y), c1 = Dot(B, B) - Dot(A, A);
   double a2 = 2*(C.X - A.X), b2 = 2*(C.Y -
       A.Y), c2 = Dot(C, C) - Dot(A, A);
   // a1*x + b1*x = c1
   // a2*x + b2*x = c2
   double D = a1 * b2 - a2 * b1;
   if (dcmp(D) == 0) return false;
   cir = Point(c1 * b2 - c2 * b1, a1 * c2 -
       a2 * c1) * (1 / D);
   return true;
struct Line {
   Point A, B;
   double rad;
   Line(Point A=MK(0, 0), Point B=MK(0, 0)):
        A(A), B(B) { rad = getAngle(); }
   double getAngle() { return atan2((B-A).Y,
         (B-A).X); }
   Point getdir() { return B-A; }
   bool operator<(const Line &op) const {</pre>
       return dcmp(rad-op.rad) < 0;</pre>
```

6.2 Smallest Covering Circle

```
for (int i = 0; i < n; ++i) swap(p[rand()%n],</pre>
     p[rand()%n]);
r = 0;
cir = p[0];
for (int i = 0; i < n; ++i) {</pre>
    if (dcmp(Distance(cir, p[i]) - r) <= 0)</pre>
        continue;
    cir = (p[0] + p[i]) * 0.5;
    r = Distance(p[0], p[i]) * 0.5;
    for (int j = 1; j < i; ++j) {</pre>
        if (dcmp(Distance(cir, p[j]) - r) <=</pre>
            0) continue;
        cir = (p[i] + p[j]) * 0.5;
        r = Distance(p[i], p[j]) * 0.5;
        for (int k = 0; k < j; ++k) {
            if (dcmp(Distance(cir, p[k]) - r)
                 <= 0) continue;
            Circle(p[i], p[j], p[k], cir);
            r = Distance(cir, p[i]);
        }
    }
```

6.3 Half Plane Intersection

```
/*
* input: Array of lines
```

6.3 Half Plane Intersection

15

```
* output: the points of intersection
 * Time: O(nlogn)
 */
deque<Line> Q;
deque<Point> Qp;
vector<Point> halfPlaneCross(vector<Line> &ls
    ) {
   while (!Q.empty()) Q.pop_front();
   while (!Qp.empty()) Qp.pop_front();
    sort(ALL(ls), cmp);
   Q.push_back(ls[0]);
   for (int i = 1; i < SZ(ls); ++i) {</pre>
        if (dcmp(Cross(ls[i].getdir(), ls[i
           -1].getdir())) != 0) {
           while (!Qp.empty() && !onLeft(ls[
               i], Qp.back())) Q.pop_back(),
                Qp.pop_back();
           while (!Qp.empty() && !onLeft(ls[
                i], Qp.front())) Q.pop_front
                (), Qp.pop_front();
           Point tmp;
            intersect(Q.back(), ls[i], tmp);
           Q.push_back(ls[i]);
           Qp.push_back(tmp);
       }
   }
   while (!Qp.empty() && !onLeft(Q.front(),
       Qp.back())) Q.pop_back(), Qp.pop_back
        ();
   vector<Point> ansp;
   if (Qp.size() <= 1) return ansp;</pre>
    ansp.PB(Point());
    intersect(Q.front(), Q.back(), ansp[0]);
   while (!Qp.empty()) ansp.PB(Qp.front()),
       Qp.pop_front();
   return ansp;
```

16 A ENV CONF

A Env Conf

```
" user configuration
"set cindent
"set autoindent
"set nu rnu
"set tabstop=4
"set backspace=2
"set vb t_vb=
"se shiftwidth=4
" set background=dark
" inoremap ( ()<LEFT>
" inoremap [ [] < LEFT >
" colorsheme evening
syntax on
color molokai
se ai nu rnu bs=2 ts=4 sw=4
se guifont=Consolas:b:h12
imap {<CR> {<CR>}<ESC>0
imap <F5> <ESC>:call Run()<CR>
imap <C-S> <ESC>:w<CR>
map <C-A>c ggvG$"+y
map <C-A>v ggvG$"+p
map <C-S> :w<CR>
map <F5> :call Run()<CR>
" if os is Linux, replace %<.exe with ./%<
func! Run()
   exec "w"
   exec "!g++ -Wall -g % -o %<.exe"</pre>
   exec "silent !%<.exe < my.in > my.out"
endfunc
map <F10> :call CaR()<CR>
func! CaR()
   exec "w"
   exec "!g++ -Wall -g % -o %<.exe"</pre>
   exec "!%<.exe"
endfunc
" Non
map <F4> :call PY()<CR>
func! PY()
    exec "w"
   exec "!python %"
endfunc
map <F6> :call Debug()<CR>
```

```
func! Debug()
    exec "w"
    exec "silent !g++ -Wall -g % -o %<.exe"
    exec "!gdb %<.exe"
endfunc
map <F7> :call Finderror()<CR>
func! Finderror()
    exec "w"
    exec "silent !g++ -Wall -g test.cpp -o
        test.exe"
    exec "silent !g++ -Wall -g % -o %<.exe"
    exec "silent !%<.exe < my.in > my.out"
    exec "silent !test.exe < my.in > ans.txt"
    exec "!fc ans.txt my.out"
endfunc
```

```
// launch
   // 使用 IntelliSense 了解相关属性。
   // 悬停以查看现有属性的描述。
   // 欲了解更多信息,请访问: https://go.
       microsoft.com/fwlink/?linkid = 830387
   "version": "0.2.0",
   "configurations": [
           "name": "g++.exe build and debug
              active file",
           "type": "cppdbg",
           "request": "launch",
           "program": "${fileDirname}\\${
              fileBasenameNoExtension}.exe"
           "args": [],
           "stopAtEntry": false,
           "cwd": "${workspaceFolder}",
           "environment": [],
           "externalConsole": false,
           "MIMode": "gdb",
           "miDebuggerPath": "C:\\mingw64\\
              bin\\gdb.exe",
           "setupCommands": [
              {
                  "description": "Enable
                      pretty-printing for
                      gdb",
                  "text": "-enable-pretty-
                      printing",
```

```
// tasks
// 有关 tasks.json 格式的文档,请参见
   // https://go.microsoft.com/fwlink/?LinkId
       =733558
   "version": "2.0.0",
   "tasks": [
           "type": "shell",
           "label": "g++.exe build active
               file",
           "command": "C:\\mingw64\\bin\\g
               ++.exe",
           "args": [
               "-Wall",
               "-g",
               "${file}",
               "-o",
               "${fileDirname}\\${
                   fileBasenameNoExtension}.
                   exe"
           ],
           "options": {
               "cwd": "C:\\mingw64\\bin"
           },
           "problemMatcher": [
               "$gcc"
           ],
           "group": {
               "kind": "build",
               "isDefault": true
       }
   ]
```

```
// properties {
```

```
"configurations": [
        "name": "Win32".
       "includePath": [
            "${workspaceFolder}/**"
       ],
        "defines": [
           " DEBUG".
            "UNICODE",
           " UNICODE"
       ],
        "compilerPath": "C:\\mingw64\\bin
            \\g++.exe",
        "cStandard": "gnu17",
        "cppStandard": "gnu++14",
        "intelliSenseMode": "gcc-x64"
   }
],
"version": 4
```

```
cmake_minimum_required(VERSION 3.16)
project(code)
set(CMAKE_CXX_STANDARD 14)

file (GLOB files *.cpp */*.cpp)
foreach ( file ${ files })
    string (REGEX REPLACE ".+/(.+)/(.+)\\..*"
        "\\1-\\2" exe ${file})
    add_executable (${exe} ${ file })
endforeach ()
```

B Theorem

B.1 Lucas' Theorem

对于质数 p, 有 $\binom{n}{m} \bmod p = \binom{\lfloor n/p \rfloor}{\lfloor m/p \rfloor} \cdot \binom{n \mod p}{m \mod p} \bmod p$

B.2 Betty's Theorem

如果两个无理数 a,b 满足:

$$\frac{1}{a} + \frac{1}{b} = 1$$

那么对于两个集合 A, B:

$$A = \{[na]\}, B = \{[nb]\}$$

C C++ STL: SET

有下面两个结论:

$$A \bigcap B = \emptyset, A \bigcup B = \mathbb{N}^+$$

B.3 Bernoulli Number

定义: $S(n,m) = \sum_{i=1}^{n} i^{m} \ (m > 0, n > 0)$ 则必有: $S(n,m) = \frac{1}{m+1} \sum_{i=0}^{m} {m+1 \choose i} B_{i} n^{m+1-i}$ 其中: $B_{n} = -\frac{1}{n+1} \sum_{i=0}^{n-1} {n+1 \choose i} B_{i}, \ B_{0} = 1, \ B_{1} = \pm \frac{1}{2}, \ B_{2} = \frac{1}{6}, \ B_{3} = 0, \ B_{4} = -\frac{1}{30}$ 一般来说,自然数幂求和要 B_{1} 取正数。另外,伯努利数和 $\zeta(s)$ 有关,

$$B_{2m} = \frac{2(-1)^{m-1}(2m)!}{(2\pi)^{2m}}\zeta(2m)$$

使用多项式求逆或者分治 ftt 可以快速计算伯努利数,因为伯努利数满足

$$\frac{x}{e^x - 1} = \sum_{n=0}^{\infty} \frac{B_n}{n!} x^n = \frac{1}{\frac{e^x - 1}{x}}$$

C C++ STL: set

set 与 unordered_set 的区别在于有无在内部存储时有无顺序。

C.1 Basic Method

begin() 返回 set 容器的第一个元素

end() 返回 set 容器的最后一个元素

clear() 删除 set 容器中的所有的元素

empty() 判断 set 容器是否为空

max_size() 返回 set 容器可能包含的元素最大个数

size() 返回当前 set 容器中的元素个数

rbegin() 返回的值和 end() 相同

rend() 返回的值和 rbegin() 相同

count() 用来查找 set 中某个某个键值出现的次数。(在 set 中只有 0 或 1 次)

equal_range()返回一对定位器,分别表示第一个大于或等于给定关键值的元素和第

一个大于给定关键值的元素,这个返回值是一个 pair 类型,如果这一对定位器中哪个返回 失败,就会等于 end()的值

erase(iterator) 删除定位器 iterator 指向的值

erase(first,second) 删除定位器 first 和 second 之间的值

erase(key_value) 删除键值 key_value 的值

insert(key_value) 将 key_value 插入到 set 中,返回值是 pair<set<int>::iterator,bool>, bool 标志着插入是否成功,而 iterator 代表 插入的位置,若 key_value 已经在 set 中,则 iterator 表示的 key_value 在 set 中的位置

lower_bound(key_value) 返回第一个大于等于 key_value 的定位器

upper_bound(key_value) 返回最后一个 大于等于 key_value 的定位器

C.2 Advanced Method

注: 必须导入 algorithm 头文件 set_intersection(first1,last1,first2,last2,d_first,comp)

first1, last1 - 要检验的第一元素范围 first2, last2 - 要检验的第二元素范围 d first - 输出范围的起始

comp - 比较函数对象(即满足比较 (Compare) 概念的对象),若第一参数小于(即先序于)第二参数则返回 true

Example:

std::set_intersection(v1.begin(),v1.
end(),v2.begin(),v2.end(),std::
back_inserter(v_intersection));
// v_intersection 就是交集,
back_insecter()用于 vector

set_union(first1,last1,first2,last2,d_first,comp)