

# ICPC TEMPLATE



### $\mathbf{BY}$

## albertxwz

School of Computer Science & Engineering, South China University of Technology

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This template is a supplementary version.

# Contents

1	Star	ndard Solution Template	1
	1.1	$support\ bits/stdc++.h\ .\ .\ .\ .\ .\ .\ .$	1
	1.2	unsupport bits/stdc++.h $\dots$	1
	1.3	Python Template	1
2	Gra	ph	2
	2.1	Network Flow	2
		2.1.1 Dinic	2
		2.1.2 Edmonds-Karp	2
	2.2	DSU on Tree	3
3	Mathematics		4
	3.1	Number Theory	4
		3.1.1 Linear Inverse Modulo	4
		3.1.2 M $\ddot{o}$ bius Inversion	4
		3.1.3 Dujiao Sieve	4
		3.1.4 Min_25 Sieve	4
		3.1.5 Lucas' Theorem	4
	3.2	Karatsuba Multiply	4
	3.3	Fast Fourier Transform	4
4	Data Structure		6
	4.1	Treap	6
	4.2	Splay Tree	6
	4.3	Two-dimensional Segment Tree $\ \ldots \ \ldots$	6
5	Stri	ng	7
	5.1	KMP	7
A	The	orem	7
	A.1	Lucas' Theorem	7
	A.2	Betty's Theorem	7
В	<b>C</b> +-	+ STL: set	7
	B.1	Basic Method	7
	B.2	Advanced Method	8

# 1 Standard Solution Template

### 1.1 support bits/stdc++.h

### 1.2 unsupport bits/stdc++.h

```
#include <cstdio>
#include <cstring>
#include <iostream>
#include <algorithm>
#include <cmath>
#include <iterator>
#include <map>
#include <queue>
#include <cstdlib>
#include <string>

using namespace std;

int main() {
    // freopen ("my.txt", "w", stdout);
    return 0;
}
```

### 1.3 Python Template

```
for T in range(0, int(input())): #T组数据
   N= int( input())
   n,m= map( int, input().split())
   s= input()
   s=[ int(x) for x in
       input().split()] #一行输入的数组
   a[1:]=[int(x) for x in
       input().split()] #从下标1开始读入一行
   for i in range(0, len(s)):
       a,b= map( int, input().split())
while True: #未知多组数据
   try:
       \#n,m=map(int,input(). split())
       \#print(n+m,end="\n")
   except EOFError: #捕获到异常
       break
```

2 GRAPH

## 2 Graph

### 2.1 Network Flow

#### 2.1.1 Dinic

```
struct Edge {
   int from, to, cap, flow;
   Edge(int from, int to=0, int cap=0, int
       flow=0):
       from(from), to(to), cap(cap), flow(
           flow) {}
};
vector<Edge> edges;
vector<int> g[maxn];
int d[maxn], cur[maxn];
void addedge(int u, int v, int cap) {
    edges.push_back(Edge(v, cap));
   g[u].push_back(edges.size()-1);
    edges.push_back(Edge(u, 0));
   g[v].push_back(edges.size()-1);
 * 时间复杂度: O(n^2*m)
 * 对于特殊的图,所有容量为1的: O(min(n^0.67,
    m^0.5)*m)
               二分图最大匹配: O(n^0.5*m)
bool BFS(int s, int t) {
   memset(d, -1, sizeof(d));
   queue<int> Q;
   Q.push(s);
   d[s] = 0;
   while (!Q.empty()) {
       int x = Q.front(); Q.pop();
       for (int i = 0; i < g[x].size(); ++i)</pre>
           Edge &e = edges[g[x][i]];
           if (d[e.to] == -1 \&\& e.cap > e.
               flow) {
               d[e.to] = d[x] + 1;
               Q.push(e.to);
           }
```

```
}
    return d[t] > -1;
int DFS(int x, int a, int t) {
    if (x == t || a == 0) return a;
    int flow = 0, f;
    for (int &i = cur[x]; i < g[x].size(); ++</pre>
        i) {
        Edge &e = edges[g[x][i]];
        if (d[x] + 1 == d[e.to] && (f = DFS(e
            .to, min(a, e.cap-e.flow), t)) >
            0) {
            flow += f;
            e.flow += f;
            edges[g[x][i]^1].flow -= f;
            a -= f;
            if (a == 0) break;
        }
    return flow;
}
int MaxFlow(int s, int t) {
    int res = 0;
    while (BFS(s, t)) {
        for (int i = 0; i <= n; ++i) cur[i] =</pre>
        res += DFS(s, inf, t);
    return res;
```

#### 2.1.2 Edmonds-Karp

```
queue<int> Q;
int imp[maxn], p[maxn];

// 不断寻找增广路增广

int MaxFlow(int s, int t) {
   int res = 0;
   for (;;) {
      memset(imp, 0, sizeof(imp));
      while (!Q.empty()) Q.pop();
      Q.push(s);
```

2.2 DSU on Tree

```
imp[s] = inf;
    while (!Q.empty()) {
        int x = Q.front(); Q.pop();
       for (int i = 0; i < g[x].size();</pre>
           ++i) {
           Edge &e = edges[g[x][i]];
           if (!imp[e.to] && e.cap > e.
               flow) {
               p[e.to] = g[x][i];
               imp[e.to] = min(imp[x], e.
                   cap-e.flow);
               Q.push(e.to);
           }
       if (imp[t]) break;
    if (!imp[t]) break;
    for (int u = t; u != s; u = edges[p[u
       ]].from) {
        edges[p[u]].flow += imp[t];
        edges[p[u]^1].flow -= imp[t];
   }
   res += imp[t];
return res;
```

## 2.2 DSU on Tree

4 3 MATHEMATICS

### 3 Mathematics

## 3.1 Number Theory

#### 3.1.1 Linear Inverse Modulo

- 3.1.2 Möbius Inversion
- 3.1.3 Dujiao Sieve
- 3.1.4 Min 25 Sieve
- 3.1.5 Lucas' Theorem

### 3.2 Karatsuba Multiply

```
z2 = karatsuba_polymul(n - m, a + m, b + m)
 // 计算 z1
 // 临时更改, 计算完毕后恢复
 for (int i = 0; i + m <= n; ++i) a[i] += a[</pre>
     i + m];
 for (int i = 0; i + m <= n; ++i) b[i] += b[</pre>
     i + m];
 z1 = karatsuba_polymul(m - 1, a, b);
 for (int i = 0; i + m <= n; ++i) a[i] -= a[</pre>
     i + m];
 for (int i = 0; i + m <= n; ++i) b[i] -= b[</pre>
 for (int i = 0; i <= (m - 1) * 2; ++i) z1[i
     ] = z0[i];
 for (int i = 0; i \le (n - m) * 2; ++i) z1[i
     ] = z2[i];
 // 由 z0、z1、z2 组合获得结果
 for (int i = 0; i <= (m - 1) * 2; ++i) r[i]</pre>
      += z0[i];
 for (int i = 0; i \le (m - 1) * 2; ++i) r[i
     + m] += z1[i];
 for (int i = 0; i \le (n - m) * 2; ++i) r[i
     + m * 2] += z2[i];
 delete[] z0;
 delete[] z1;
 delete[] z2;
 return r;
// 计算a*b=c, 时间复杂度是O(n^1.585)
void karatsuba_mul(int a[], int b[], int c[])
 int *r = karatsuba_polymul(LEN - 1, a, b);
 memcpy(c, r, sizeof(int) * LEN);
 for (int i = 0; i < LEN - 1; ++i)</pre>
   if (c[i] >= 10) {
     c[i + 1] += c[i] / 10;
     c[i] %= 10;
 delete[] r;
```

### 3.3 Fast Fourier Transform

3.3 Fast Fourier Transform 5

```
// f是系数数组,处理完后,f表示:
// rev=1,是点表示法
// rev=-1,除N后是系数
// N=2^n
typedef complex<double> Comp; // 先导入头文件
   complex
void DFT(Comp *f, int N, int rev) {
   if (N == 1) return;
   for (int i = 0; i < N; ++i) tmp[i] = f[i</pre>
       ];
   for (int i = 0; i < N; ++i)</pre>
       if (i\%2) f[i/2+N/2] = tmp[i];
       else f[i/2] = tmp[i];
   Comp *g = f, *h = f + N/2;
   DFT(g, N/2, rev); DFT(h, N/2, rev);
   // c[N] = cos(2*pi/N), s[N] = sin(2*pi/N)
   Comp w(c[N], s[N]*rev), cur(1, 0);
   for (int k = 0; k < N/2; ++k) {
       tmp[k] = g[k] + cur*h[k];
       tmp[k+N/2] = g[k] - cur*h[k];
       cur *= w;
   for (int i = 0; i < N; ++i) f[i] = tmp[i</pre>
       ];
```

6 4 DATA STRUCTURE

# 4 Data Structure

- 4.1 Treap
- 4.2 Splay Tree
- 4.3 Two-dimensional Segment Tree

# 5 String

### 5.1 KMP

```
int f[maxn];
void getfail(char *P, int *f) {
   f[0] = 0;
   f[1] = 0;
   int m = strlen(P);
   for (int i = 1; i < m; ++i) {</pre>
       int j = f[i];
       while (j && P[i] != P[j]) j = f[j];
       if (P[i] == P[j]) f[i+1] = j+1;
       else f[i+1] = 0;
   }
}
int find(char *T, char *P, int *f) {
   int res = 0;
   int n = strlen(T), m = strlen(P);
   getfail(P, f);
   for (int i = 0, j = 0; i < n; ++i) {</pre>
       while (j && P[j] != T[i]) j = f[j];
       if (P[j] == T[i]) j++;
       if (j == m) { // 出现一次
           res++;
           j = f[m];
       }
   return res;
```

 $B \quad C++ STL: SET$ 

## A Theorem

### A.1 Lucas' Theorem

对于质数 p, 有

$$\binom{n}{m} \ mod \ p = \binom{\lfloor n/p \rfloor}{\lfloor m/p \rfloor} \cdot \binom{n \ mod \ p}{m \ mod \ p} \ mod \ p$$

### A.2 Betty's Theorem

如果两个无理数 a,b 满足:

$$\frac{1}{a} + \frac{1}{b} = 1$$

那么对于两个集合 A, B:

$$A = \{[na]\}, B = \{[nb]\}$$

有下面两个结论:

$$A \cap B = \emptyset, A \bigcup B = \mathbb{N}^+$$

### B C++ STL: set

set 与 undered\_set 的区别在于有无在内部存储时有无顺序。

### B.1 Basic Method

begin() 返回 set 容器的第一个元素

end() 返回 set 容器的最后一个元素 clear() 删除 set 容器中的所有的元素

empty() 判断 set 容器是否为空

max\_size() 返回 set 容器可能包含的元素最大个数

size() 返回当前 set 容器中的元素个数

rbegin()返回的值和 end()相同

rend() 返回的值和 rbegin() 相同

count() 用来查找 set 中某个某个键值出现的次数。(在 set 中只有 0 或 1 次)

equal\_range()返回一对定位器,分别表示第一个大于或等于给定关键值的元素和第

一个大于给定关键值的元素,这个返回值是一个 pair 类型,如果这一对定位器中哪个返回失败,就会等于 end()的值

erase(iterator) 删除定位器 iterator 指向的值

erase(first,second) 删除定位器 first 和 second 之间的值

erase(key\_value) 删除键值 key\_value 的值

insert(key\_value) 将 key\_value 插入到 set 中,返回值是 pair<set<int>::iterator,bool>, bool 标志着插入是否成功,而 iterator 代表 插入的位置,若 key\_value 已经在 set 中,则 iterator 表示的 key value 在 set 中的位置

lower\_bound(key\_value) 返回第一个大 于等于 key\_value 的定位器

upper\_bound(key\_value) 返回最后一个 大于等于 key\_value 的定位器

### **B.2** Advanced Method

注: 必须导入 algorithm 头文件 set intersection(first1,last1,first2,last2,d first,comp)

first1, last1 - 要检验的第一元素范围 first2, last2 - 要检验的第二元素范围

d first - 输出范围的起始

comp - 比较函数对象(即满足比较 (Compare) 概念的对象),若第一参数小于(即先序于)第二参数则返回 true

Example:

std::set\_intersection(v1.begin(),v1.
end(),v2.begin(),v2.end(),std::
back\_inserter(v\_intersection));

// v\_intersection 就是交集, back\_insecter ()用于 vector

set\_union(first1,last1,first2,last2,d\_first,comp) 同 intersection