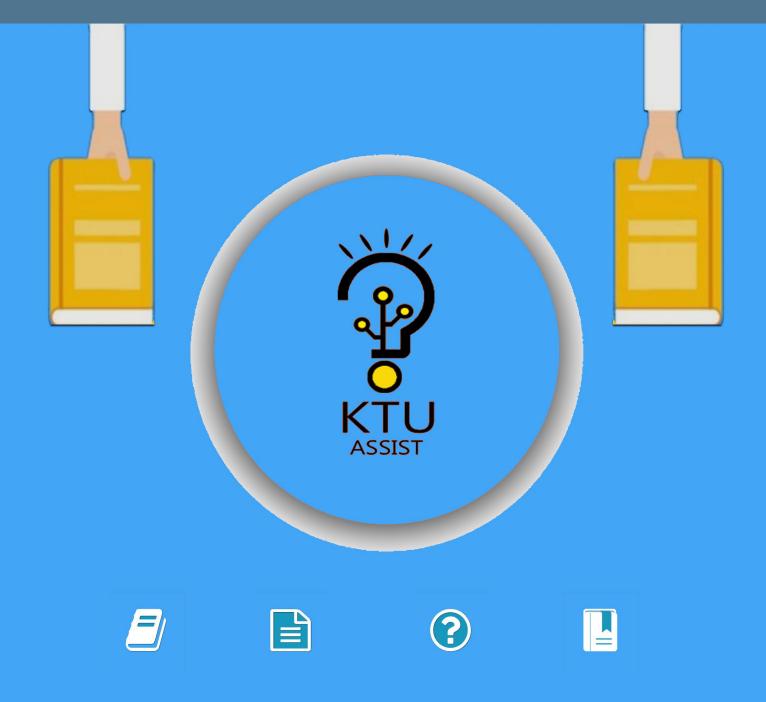
APJ ABDUL KALAM TECHNOLOGICAL UNIVERSITY

STUDY MATERIALS





a complete app for ktu students

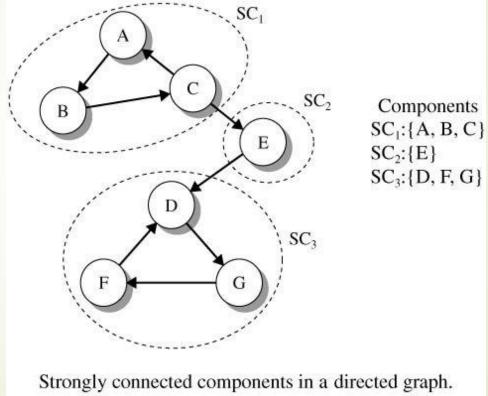
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Definition: Strongly connected component (SCC) of a directed graph G=(V,E) is a maximal set of vertices $C \subseteq V$ such that

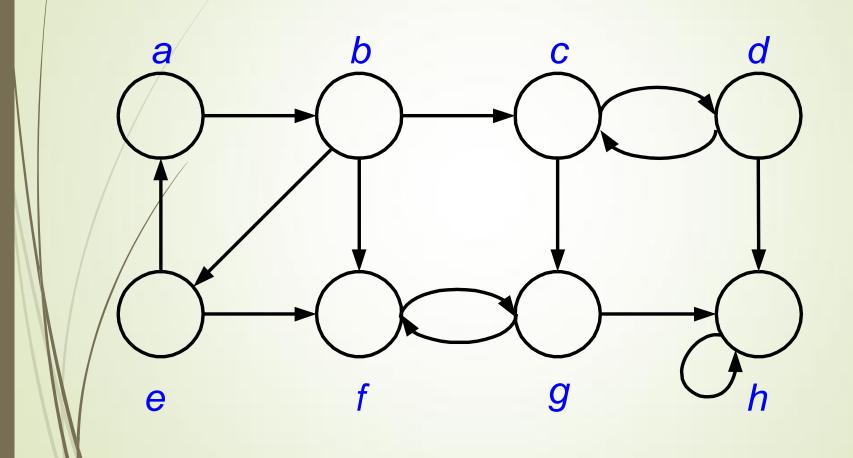
- For every pair of vertices u and v in C (u ~-> v and v~->u)
- \rightarrow i.e., u and v are mutually reachable from each other (u = v)
- Let $G^T = (V, E^T)$ be the transpose of G = (V, E) where $E^T = \{(U, V): (U, V) \in E\}$
 - i.e., E^T consists of edges of G with their directions reversed

Any graph can be partitioned into a unique set of strong components.

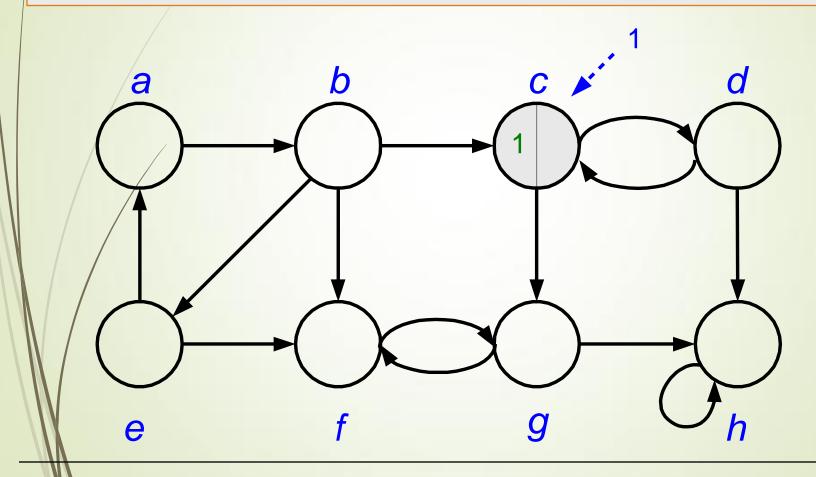


Algorithm

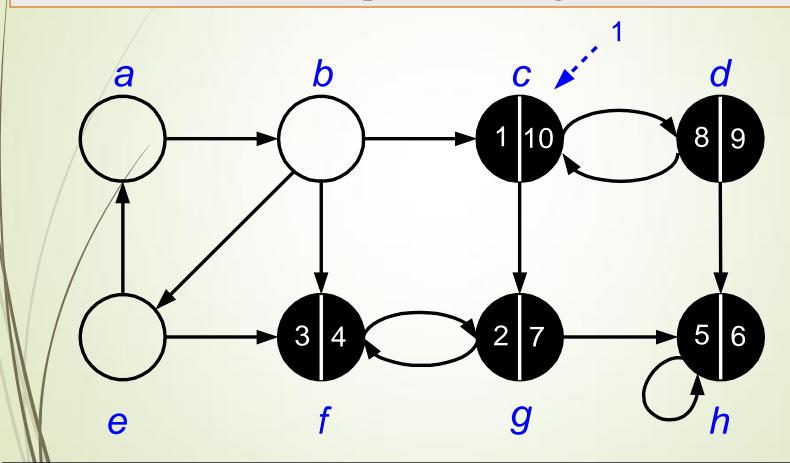
- (1) Run DFS(G) to compute finishing times for all $u \in V$
- (2) Compute G^T
- 3) Call $DFS(G^T)$ processing vertices in main loop in decreasing f[u] computed in Step (1)
- (4) Output vertices of each DFT in DFF of Step (3) as a separate SCC



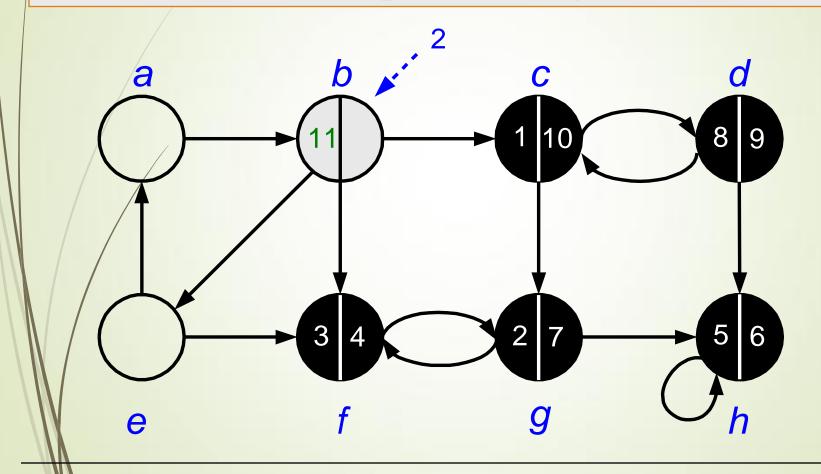
(1) Run **DFS**(G) to compute finishing times for all $u \in V$

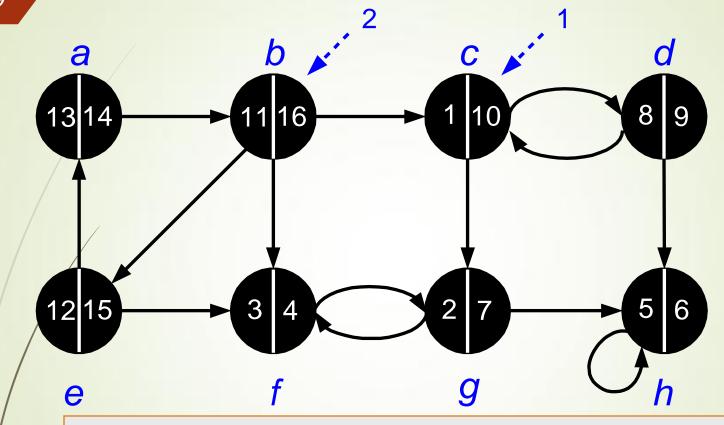


(1) Run **DFS**(G) to compute finishing times for all $u \in V$



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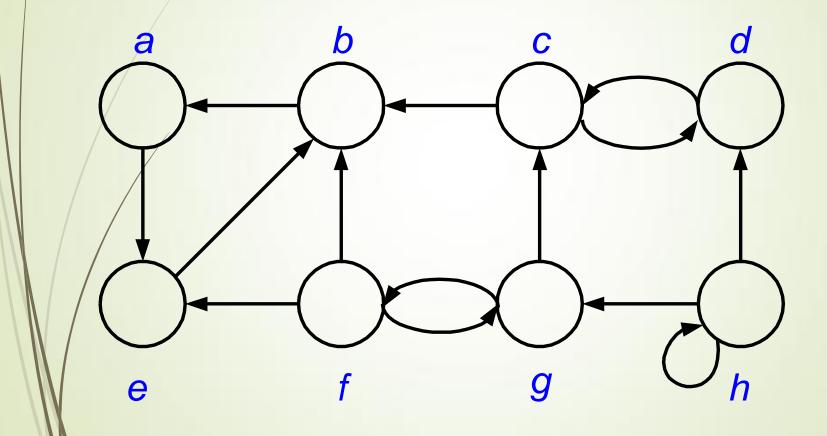


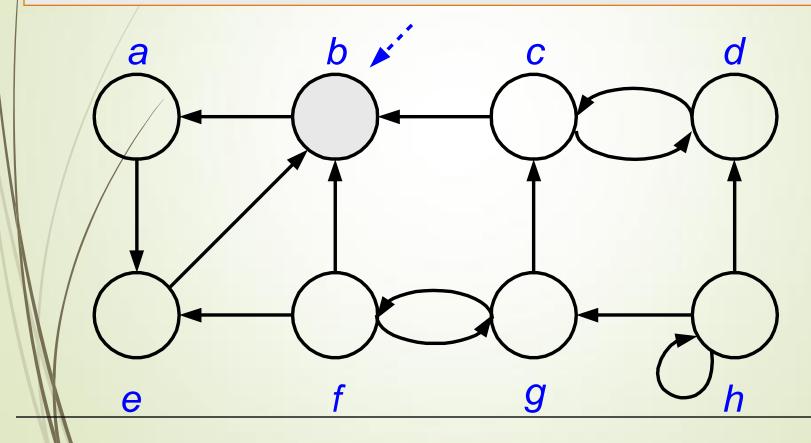


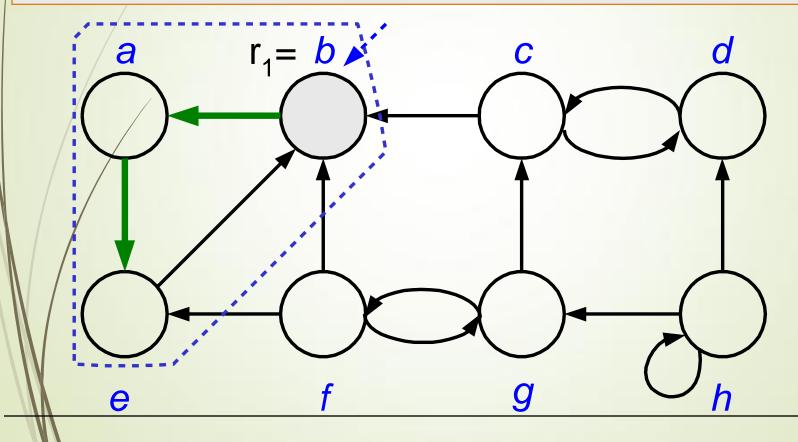
Vertices sorted according to the finishing times:

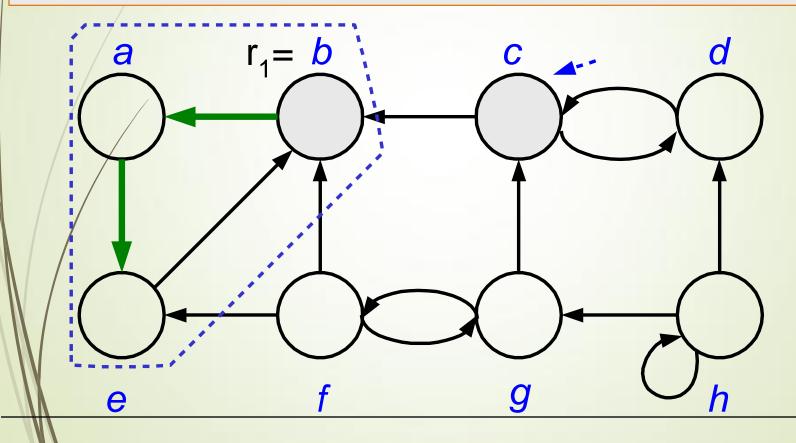
$$\langle b, e, a, c, d, g, h, f \rangle$$

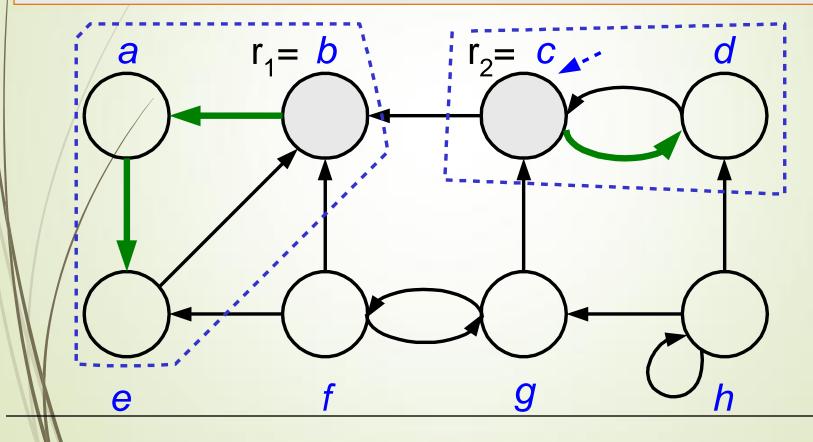
(2) Compute G^T

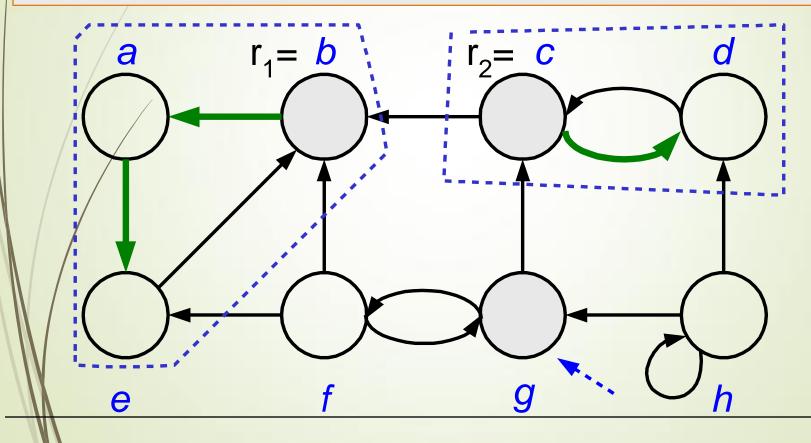


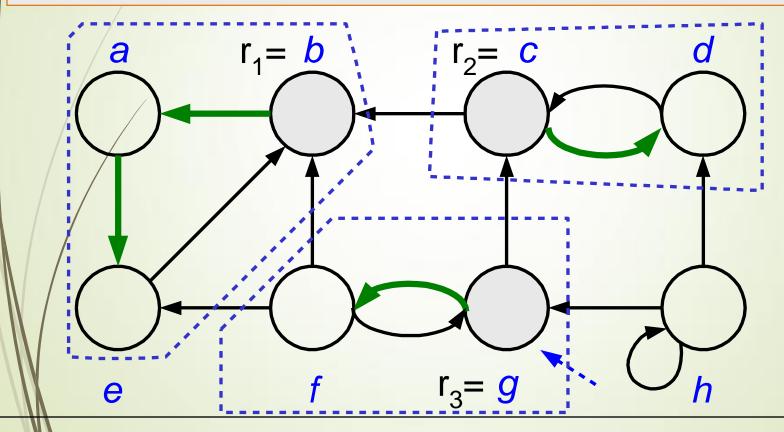


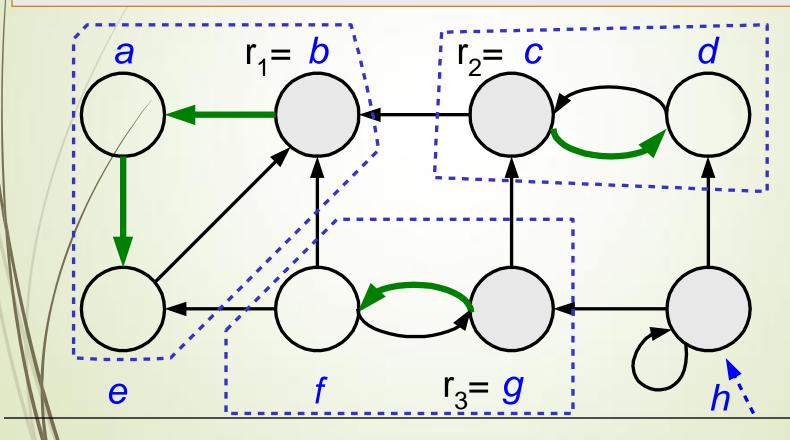


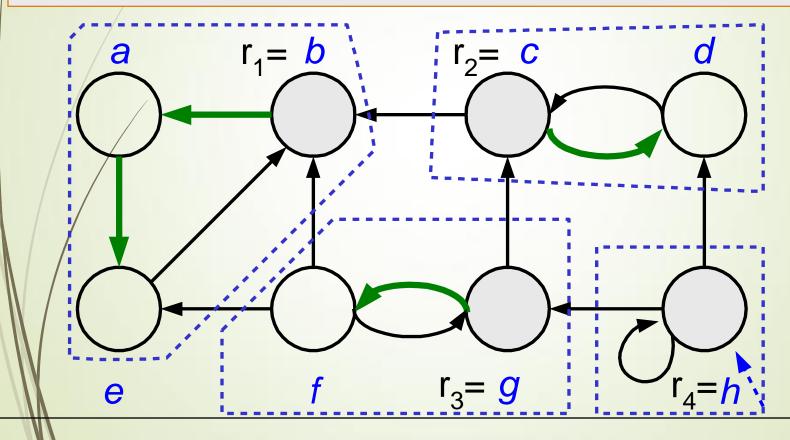




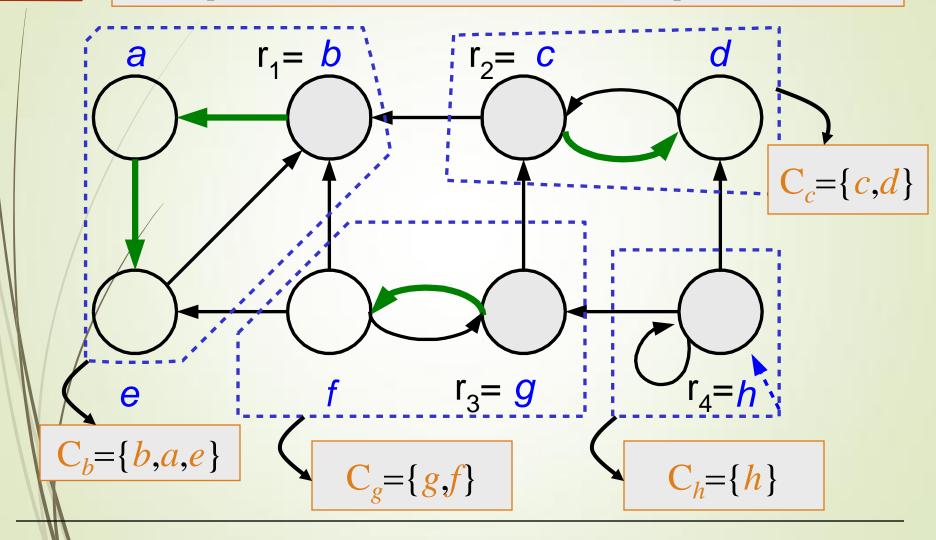


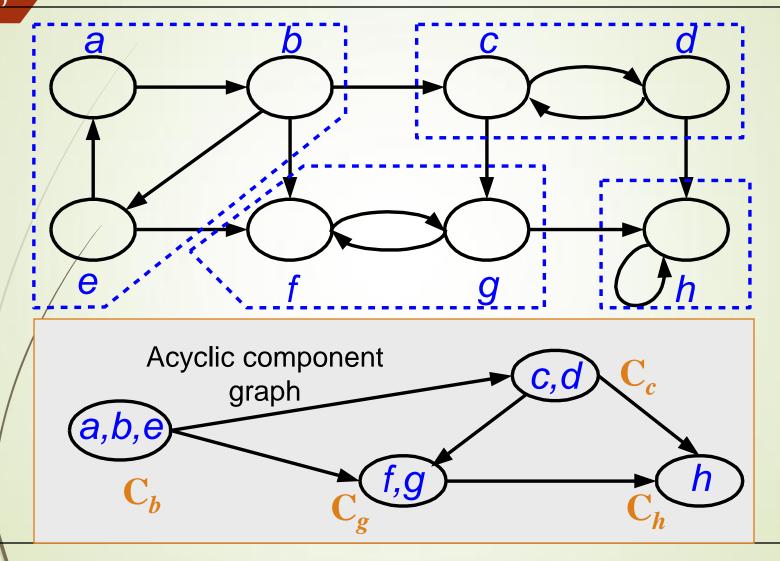


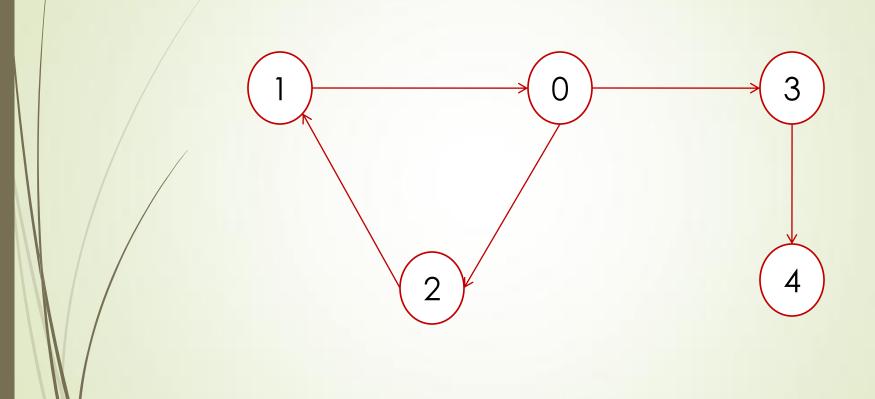




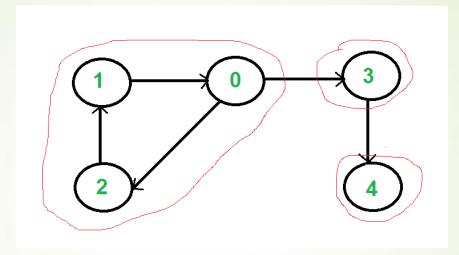
(4) Output vertices of each DFT in DFF as a separate SCC







22 EXAMPLE





END



