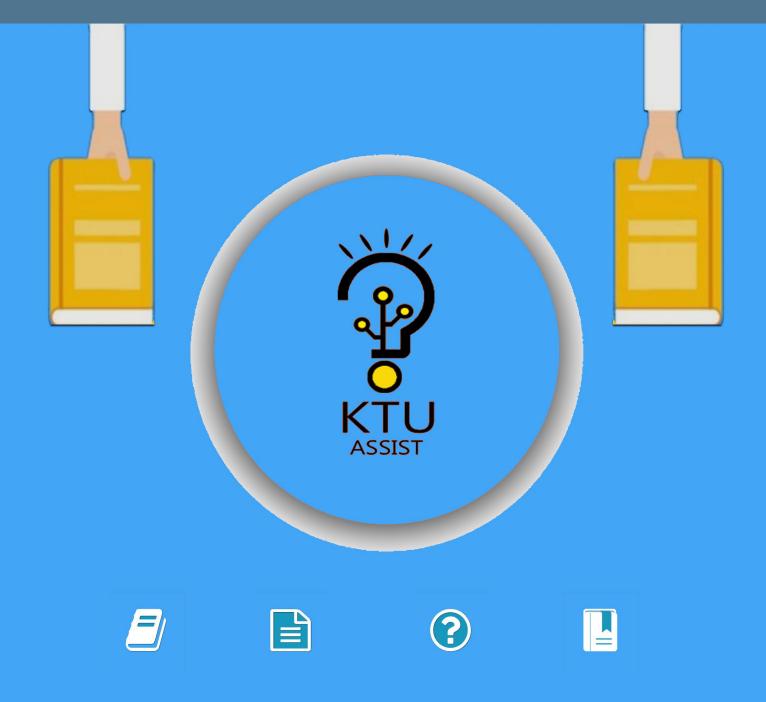
#### **APJ ABDUL KALAM TECHNOLOGICAL UNIVERSITY**

#### STUDY MATERIALS





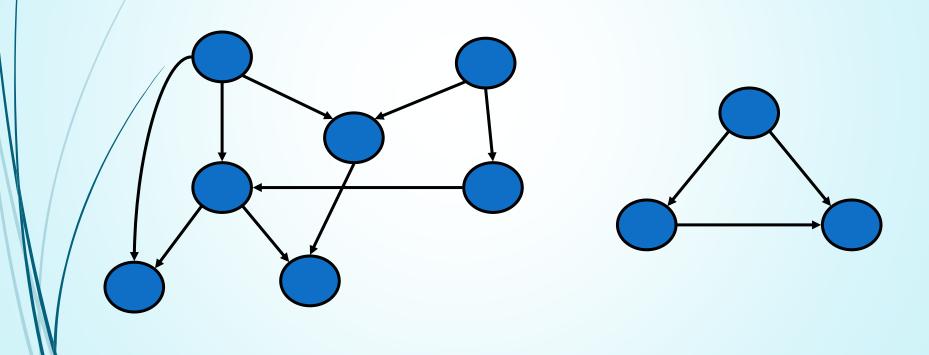
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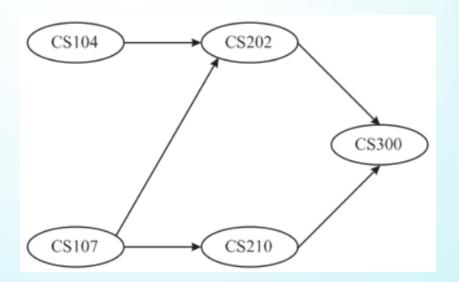
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### Directed Acyclic Graphs

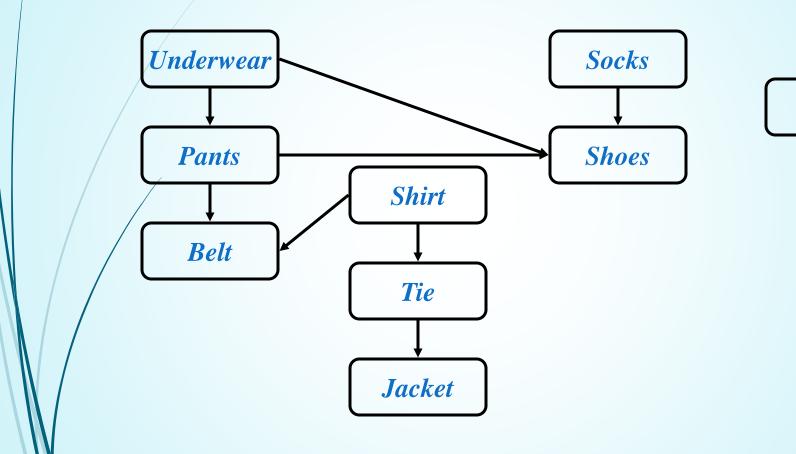
A directed acyclic graph or DAG is a directed graph with no directed cycles:



- Topological sort of a DAG:
  - Linear ordering of all vertices in graph G such that vertex u comes before vertex v if edge (u, v) ∈ G
  - Topological Sorting for a graph is not possible if the graph is not a DAG

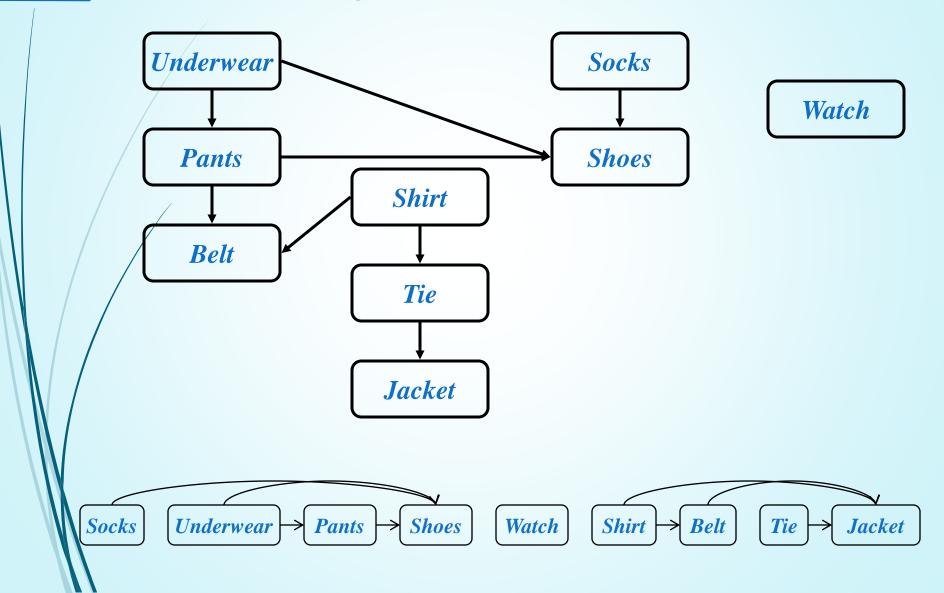


### Getting Dressed



Watch

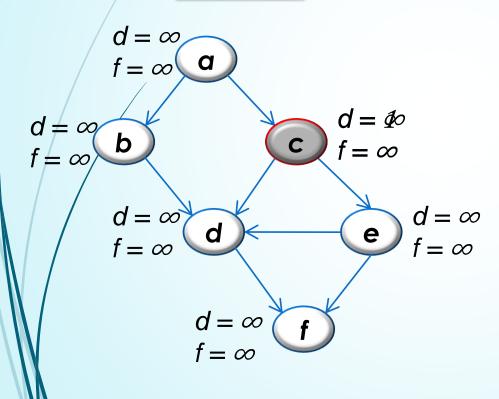
### Getting Dressed



### Topological Sort Algorithm

```
Topological-Sort()
  Run DFS
  When a vertex is finished, output it
  Output
         vertices in reverse topological
    order
Time: O(V+E)
```

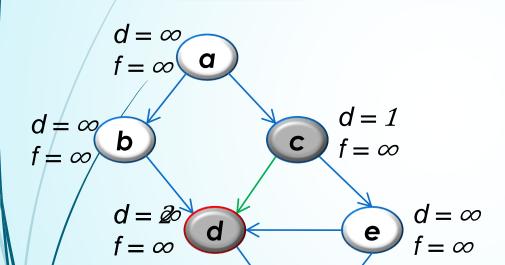
#### Time = 2



1) Call DFS(**G**) to compute the finishing times **f**[**v**]

Let's say we start the DFS from the vertex **c** 

Next we discover the vertex **d** 



 $d = \infty$ 

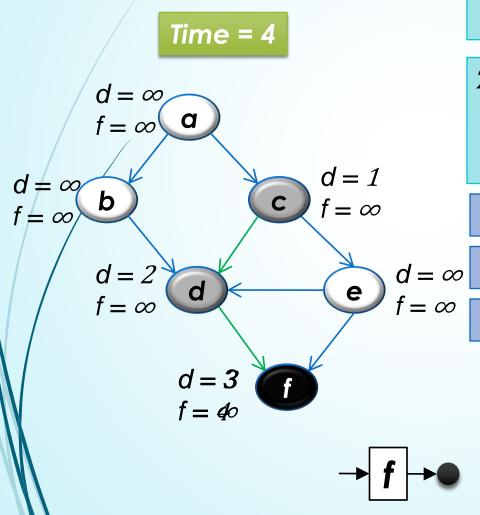
 $f = \infty$ 

Time = 3

1) Call DFS(**G**) to compute the finishing times **f**[**v**]

Let's say we start the DFS from the vertex **c** 

Next we discover the vertex **d** 

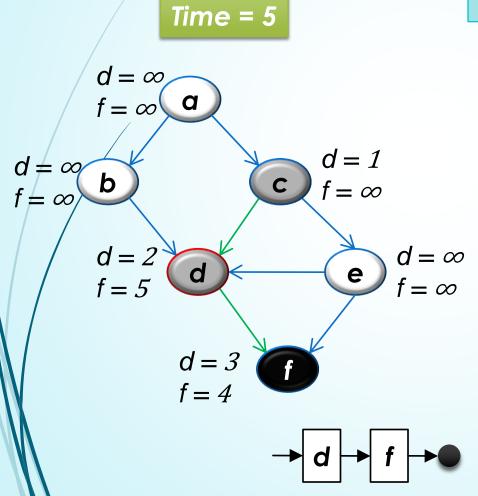


- 1) Call DFS(**G**) to compute the finishing times **f**[**v**]
- as each vertex is finished, insert it onto the **front** of a linked list

Next we discover the vertex **d** 

 $d = \infty$  Next we discover the vertex **f** 

**f** is done, move back to **d** 



1) Call DFS(**G**) to compute the finishing times **f**[**v**]

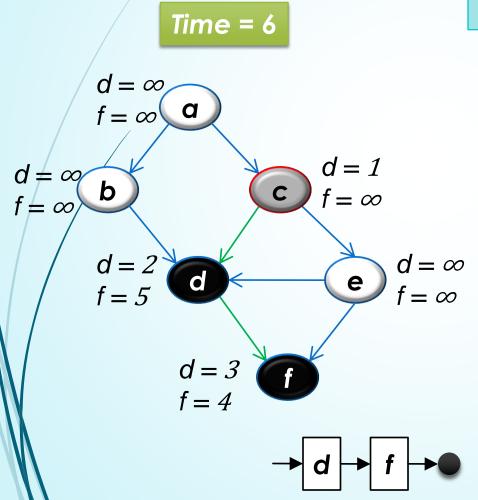
Let's say we start the DFS from the vertex **c** 

Next we discover the vertex **d** 

Next we discover the vertex f

**f** is done, move back to **d** 

**d** is done, move back to **c** 



1) Call DFS(**G**) to compute the finishing times **f**[**v**]

Let's say we start the DFS from the vertex **c** 

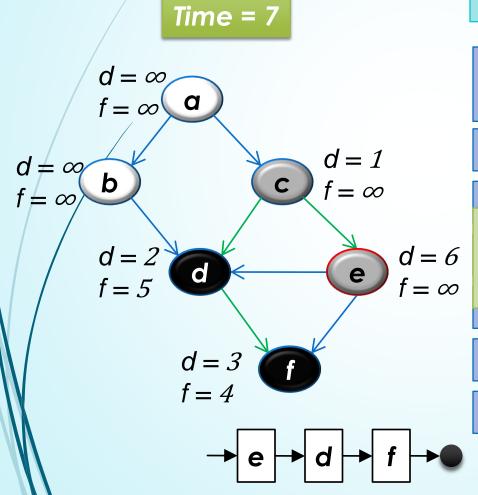
Next we discover the vertex **d** 

Next we discover the vertex f

**f** is done, move back to **d** 

**d** is done, move back to **c** 

Next we discover the vertex **e** 



1) Call DFS(**G**) to compute the finishing times **f**[**v**]

Let's say we start the DFS from the vertex **c** 

Next we discover the vertex d

Both edges from e

are cross edges

Next we discover the vertex **e** 

e is done, move back to c

**Time = 8**  $d = \infty$  $d \neq \infty$ b d = 6d = 2d = 3f = 4

1) Call DFS(**G**) to compute the finishing times **f**[**v**]

Let's say we start the DFS from the vertex **c** 

Just a note: If there was (**c**,**f**) edge in the graph, it would be classified as a **forward edge** (in this particular DFS run)

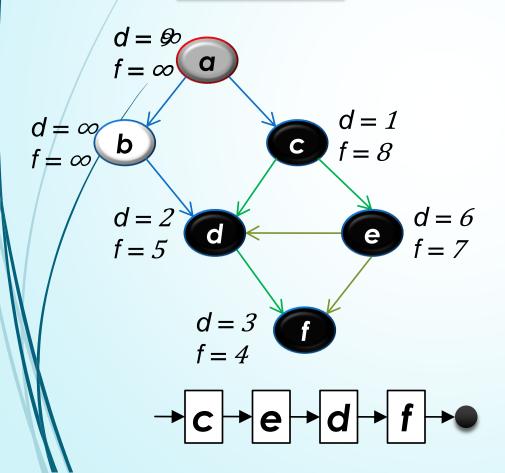
**d** is done, move back to **c** 

Next we discover the vertex **e** 

e is done, move back to c

**c** is done as well

Time = 10

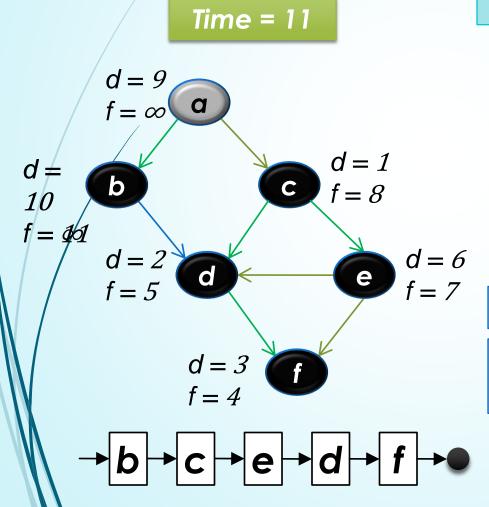


1) Call DFS(**G**) to compute the finishing times **f**[**v**]

Let's now call DFS visit from the vertex **a** 

Next we discover the vertex **c**, but **c** was already processed => (**a**,**c**) is a cross edge

Next we discover the vertex **b** 



1) Call DFS(**G**) to compute the finishing times **f**[**v**]

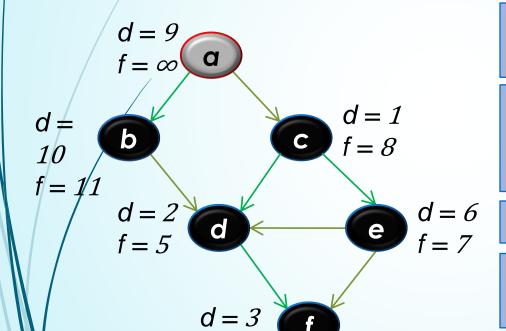
Let's now call DFS visit from the vertex **a** 

Next we discover the vertex **c**,

but **c** was already processed => (**a**,**c**) is a cross edge

Next we discover the vertex **b** 

**b** is done as (**b**,**d**) is a cross edge => now move back to **c** 



Time = 12

1) Call DFS(**G**) to compute the finishing times **f**[**v**]

Let's now call DFS visit from the vertex **a** 

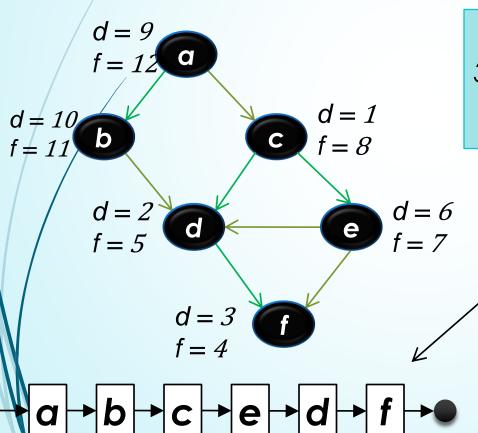
Next we discover the vertex **c**, but **c** was already processed => (**a**,**c**) is a cross edge

Next we discover the vertex **b** 

**b** is done as (**b**,**d**) is a cross edge => now move back to **c** 

a is done as well

Time = 13



1) Call DFS(**G**) to compute the finishing times **f**[**v**]

#### **WE HAVE THE RESULT!**

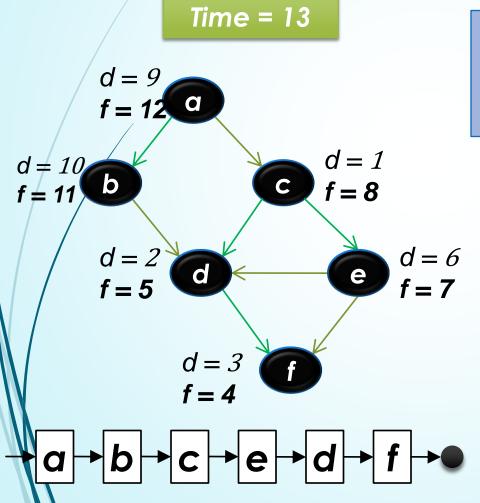
3) return the linked list of vertices

but **c** was already

Next we discover the

**b** is done as (**b**,**d**) is a cross edge => now

**a** is done as well



The linked list is sorted in **decreasing** order of finishing times **f**[]

## **END**