

Scientific Computation

Spring, 2019

Lecture 8

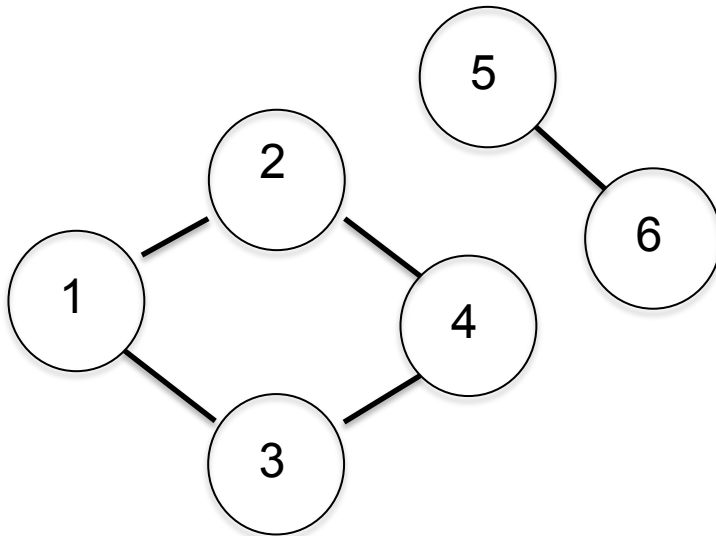
Today

- **Depth-first search**
- **Dijkstra's algorithm for weighted networks**

Graph search

- **Basic idea:** Given a graph, G , and a source node, s , find all other nodes that can be reached from s
- **Basic approach:**
 - Label s as “explored” and all other nodes as “unexplored”

While there is at least one edge between an explored and unexplored node:
Select one such edge and re-label the unexplored node as explored

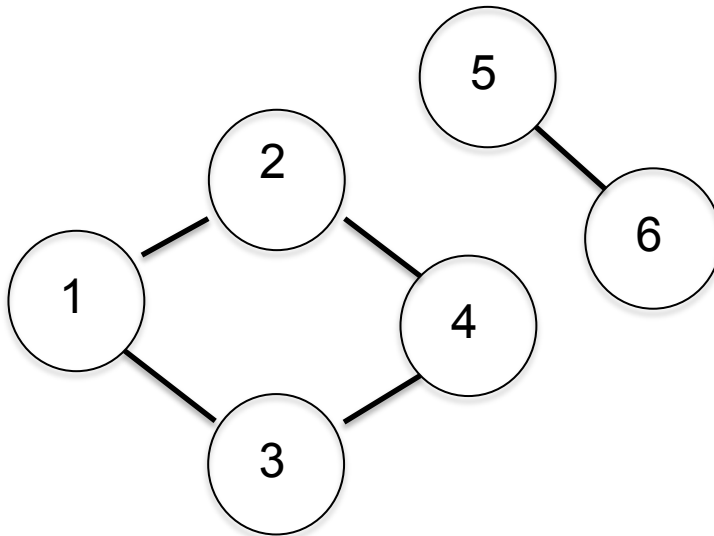


Graph search

- **Basic approach:**
 - Label s as “explored” and all other nodes as “unexplored”
- While there is at least one edge between an explored and unexplored node:
Select one such edge and re-label the unexplored node as explored*

Implementation: Depends on how edges are selected each iteration

- **Depth-first search** – aggressively move into the graph (1-2, 2-4)
- **Breadth-first search** – consider one “layer” at a time (1-2, 1-3)
- **Today: DFS**



Breadth-first search

- Python implementation
 - Specify: Graph and source node
 - Maintain: 1) list of nodes, 2) list of labels for nodes and 3) *queue* of nodes to-be explored
 - Initialize the queue with the source node and mark it as explored
 - Remove nodes from the queue in the order they were added (first in, first out)
 - Search through edges of removed node and add unexplored nodes to queue
 - Label added nodes as explored
 - Terminate search when queue is empty

Breadth-first search

- Specify: Graph and source node

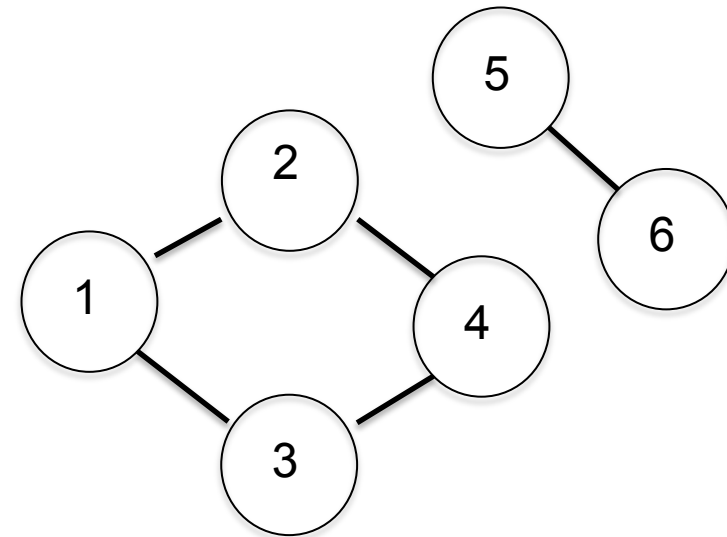
```
G = nx.Graph()
edges = [[1,2],[1,3],[1,2],[2,4],[3,4],[5,6]]
G.add_edges_from(edges)
s = 1
Q = [s] #Nodes to be explored
```

- Create list of nodes and labels

```
a_list = G.adjacency_list()
nodes = G.nodes()
z = [0 for i in nodes] #labels
L = [nodes,z]
L[1][s-1]=1 #mark source node as explored
```

- Iterate through nodes in queue (updating Q as appropriate)

```
while len(Q)>0:
    n = Q.pop(0)
    for v in a_list[n-1]: #iterate through neighbors of n
        if L[1][v-1]==0:
            L[1][v-1]=1
            Q.append(v)
```



Breadth-first search

- Adding `print("n=%d,Q=" % (n),Q)` to the while loop and running the code:

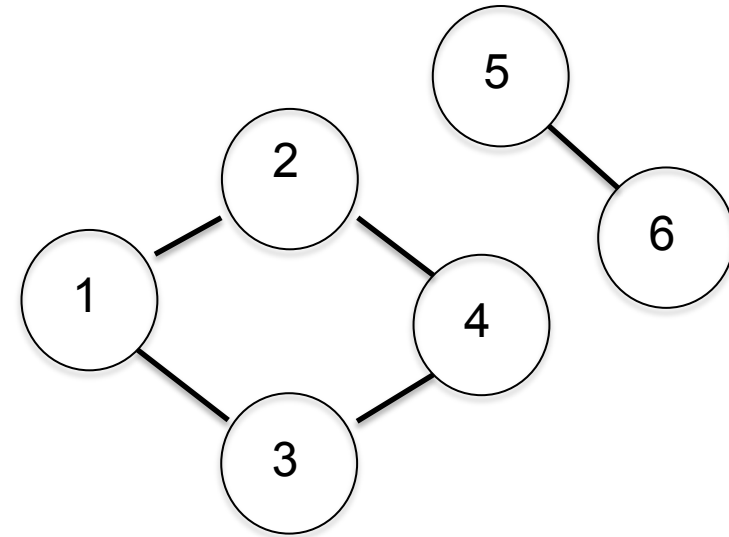
n=1,Q= [2, 3]

n=2,Q= [3, 4]

n=3,Q= [4]

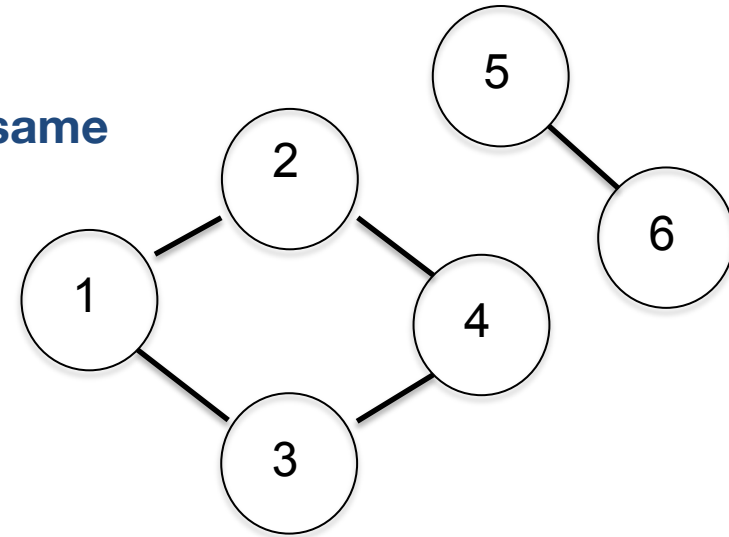
n=4,Q= []

- n is the node removed from the queue and Q is the queue after edges from n have been added (if n is unexplored)



Depth-first search

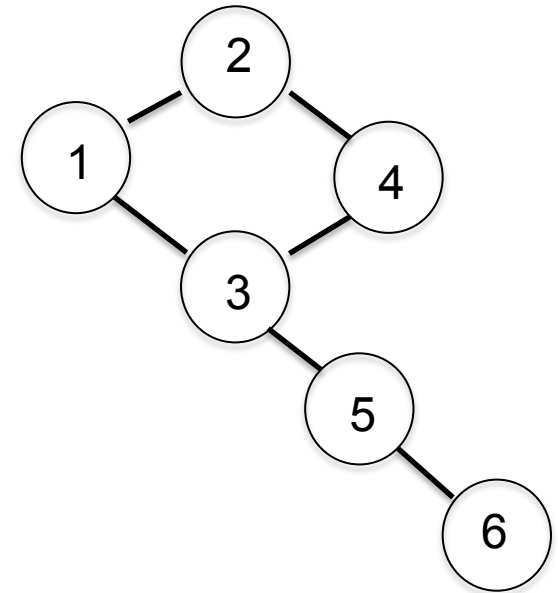
- What changes for DFS?
- How much does the code change?
- The problem setup and initialization will be the same
- Can focus on the search through the graph
- **BFS:** Remove node from front of queue, append unexplored neighbors to back of queue
- **DFS:** Remove node from end of *stack*, append unexplored neighbors to end of stack



```
#Depth-first search
while len(Q)>0:
    n = Q.pop() #Only change from BFS
    for v in a_list[n-1]: #iterate through neighbors of n
        if L[1][v-1]==0:
            L[1][v-1]=1
            Q.append(v)
```


Depth-first search

- Let's compare DFS and BFS on this modified graph
- Output from codes is provided below
- Compare the order in which the nodes are visited



In [16]: run bfs1

n=1,Q= [2, 3]

n=2,Q= [3, 4]

n=3,Q= [4, 5]

n=4,Q= [5]

n=5,Q= [6]

n=6,Q= []

In [17]: run dfs1

n=1,Q= [2, 3]

n=3,Q= [2, 4, 5]

n=5,Q= [2, 4, 6]

n=6,Q= [2, 4]

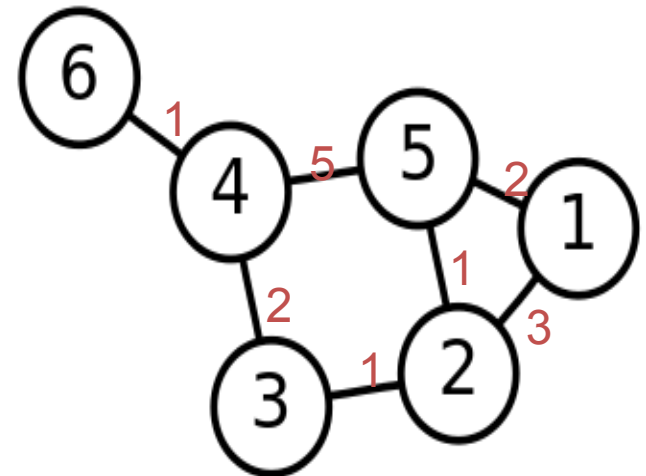
n=4,Q= [2]

n=2,Q= []

- We have left the length calculation in the DFS code the same as in BFS – is this correct?

Weighted networks

- Have considered unweighted networks so far
- Often have a weight associated with each edge
 - E.g. distance between nodes, time to travel between nodes
 - More generally, nodes can represent states in a state space, and weights can represent the cost of moving between states
- What is the best way to compute shortest paths in a weighted network?
- Can we use BFS or DFS? If weights are small, can replace a weight- n link with n weight-1 links
 - But weights can be arbitrarily large
- Dijkstra's algorithm is a graph-search method for networks with positive weights



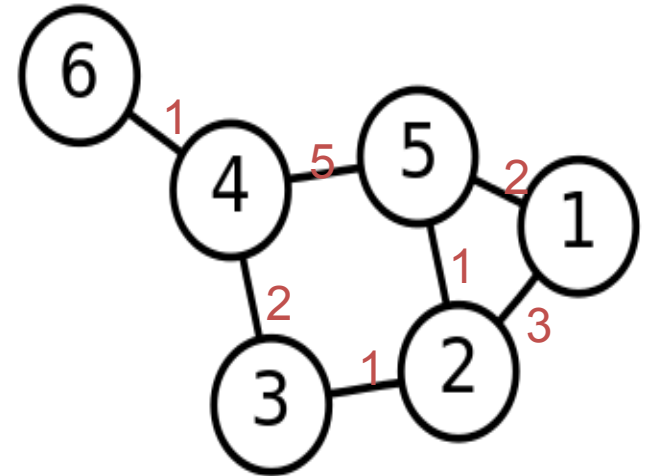
Shortest paths:

1→2: 3

1→4: 6

Dijkstra's algorithm

- Dijkstra's algorithm is a graph-search method for networks with positive weights
- Input: Graph (nodes, edges, weights) and source node
- Output: Shortest paths from source node to all nodes reachable from source



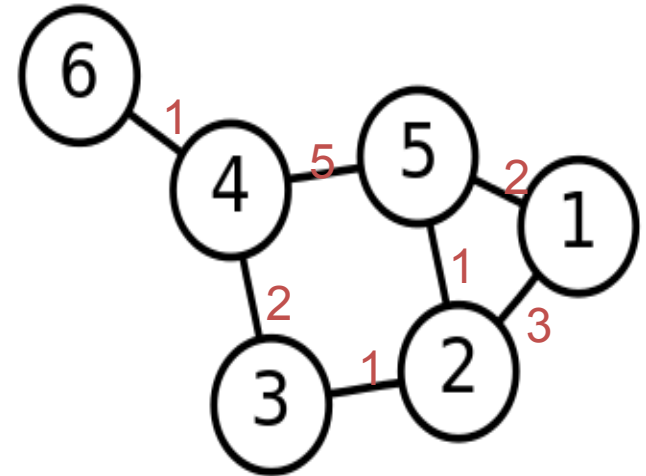
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Dijkstra's algorithm

- Dijkstra's algorithm is a graph-search method for networks with positive weights
- Input: Graph (nodes, edges, weights) and source node
- Output: Shortest paths from source node to all nodes reachable from source
- Basic idea:
 - Maintain lists of explored and unexplored nodes (E and U)
 - For each explored node, assume that we have calculated the shortest distance to the source
 - Each iteration, move one node from U to E
 - Which node? The neighbor of an explored node that is closest to the source.



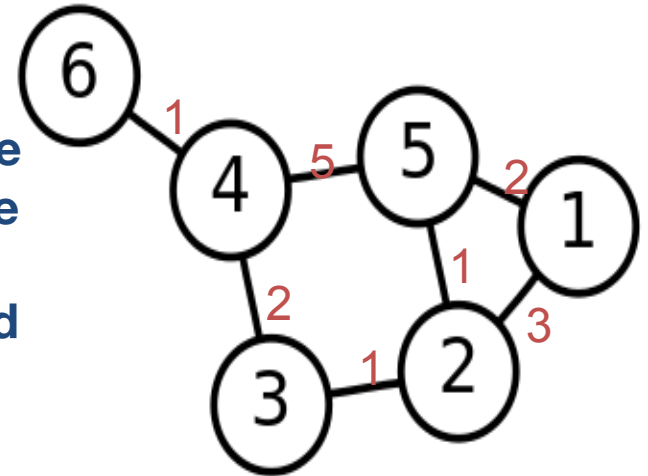
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Dijkstra's algorithm

- **Basic idea:**
 - Maintain lists of explored and unexplored nodes (E and U)
 - For each explored node, assume that we have calculated the shortest distance to the source
 - Each iteration, move one node from U to E
 - Which node? The neighbor of an explored node that is closest to the source, s
- **Why is this the right node to move?**
 - Let w_{ij} be the weight of the edge between nodes i and j
 - Let d_i be the length of the shortest path between node i and the source, s
 - Then $d_i + w_{ij}$ is the length of the path between the source and node j which includes edge i,j
 - If we choose the “unexplored neighbor” with the minimum value of $(d_i + w_{ij})$, then all other paths from s to j will be longer (why?)



Shortest paths:

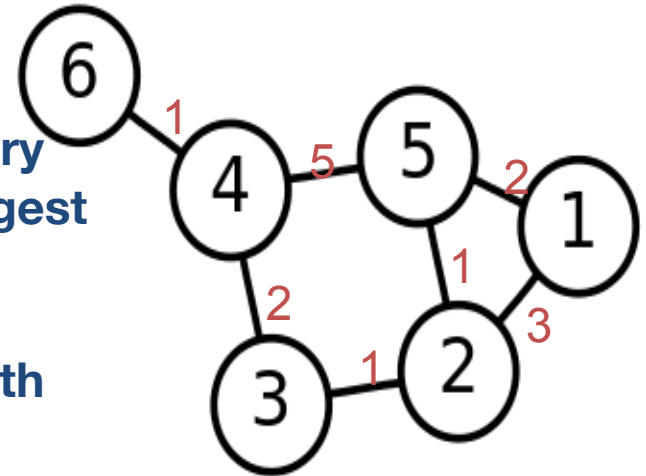
1→2: 3

1→4: 6

Dijkstra's algorithm

Algorithm (sketch)

- 1) Set the distance for the source to be zero
- 2) Set the distances for all other nodes to an arbitrary large number (much larger than the expected largest path lengths)
- 3) Label source as explored
- 4) Set provisional distances for all neighbors of s with weights of edges between s and neighbors
- 5) Each iteration, find the unexplored node (node i) that is closest to the source (smallest provisional distance)
- 6) Label this node as explored
- 7) Update the provisional distance of all neighbors of i with $\min(d_j, d_i + w_{ij})$
- 8) Continue “exploring” until no reachable nodes remain unexplored.



Shortest paths:

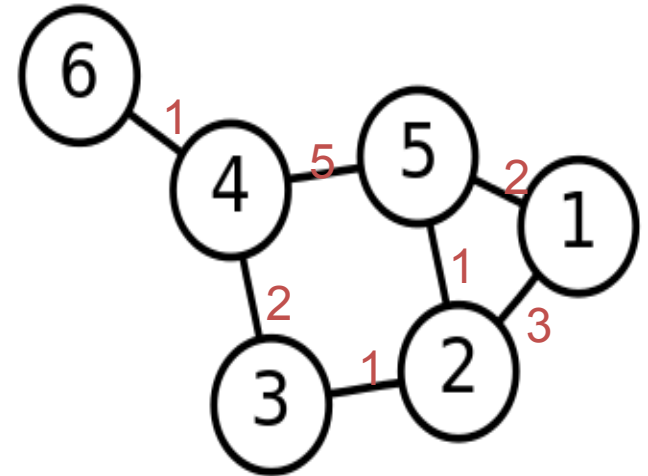
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1→4: 6

Dijkstra's algorithm

Python implementation

- What data type(s) should we use?
 - list, array, dictionary, something else?
- Basic operations needed:
 - find min
 - insert/delete
- Let's try dictionaries



Shortest paths:

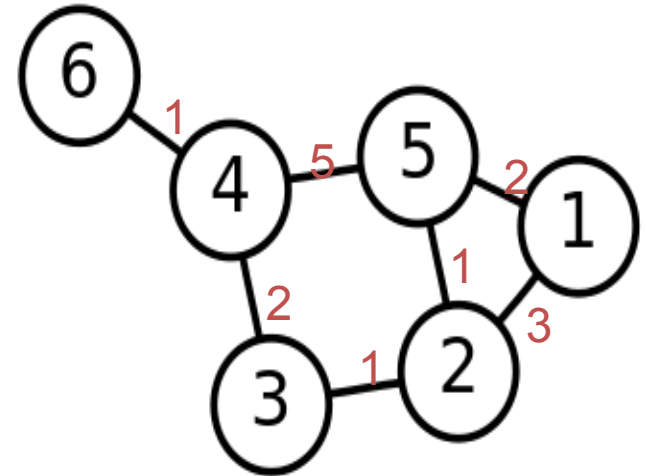
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Dijkstra's algorithm

Python implementation

- What data type(s) should we use?
 - list, array, dictionary, something else?
- Basic operations needed:
 - find min
 - insert/delete
- Let's try dictionaries
- We'll work with two dictionaries – one for explored nodes (where the shortest path length has been determined) and one for unexplored neighbors of explored nodes (where provisional shortest paths have been specified)
- And we'll also have the adjacency list of the graph containing the edges and their weights



Shortest paths:

1→2: 3

1→4: 6

Dijkstra's algorithm

Python implementation

1. Create the graph, initialize the dicts

```
e=[[1,2,3],[1,5,2],[2,5,1],[2,3,1],[3,4,2],[4,5,5],[4,6,1]]
```

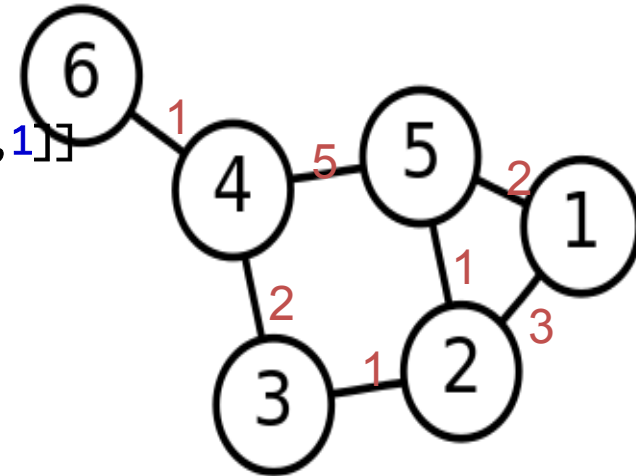
```
G = nx.Graph()
```

```
G.add_weighted_edges_from(e)
```

```
Udict = {} #Unexplored neighbor nodes
```

```
Udict[5]=0
```

```
Edict={} #Explored nodes
```



Shortest paths:

1→2: 3

1→4: 6

Dijkstra's algorithm

Python implementation

1. Create the graph, initialize the dicts

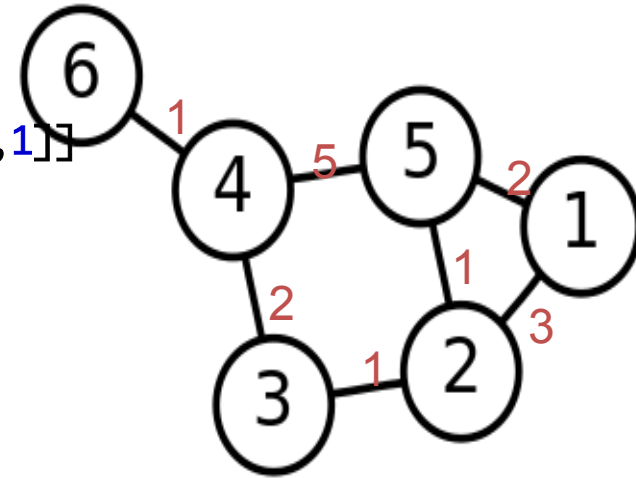
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e=[[1,2,3],[1,5,2],[2,5,1],[2,3,1],[3,4,2],[4,5,5],[4,6,1]]
G = nx.Graph()
G.add_weighted_edges_from(e)
```

```
Udict = {} #Unexplored neighbor nodes
Udict[5]=0
Edict={} #Explored nodes
```

2. Find node in Udict with smallest distance

```
while len(Udict)>0:
    dmin=2000000
    for k,v in Udict.items():
        if v<dmin:
            dmin=v #min distance
            n=k #corresponding node
```

3. Remove 'smallest-distance' node and update provisional distances of its neighbors



Shortest paths:

1→2: 3

1→4: 6

Dijkstra's algorithm

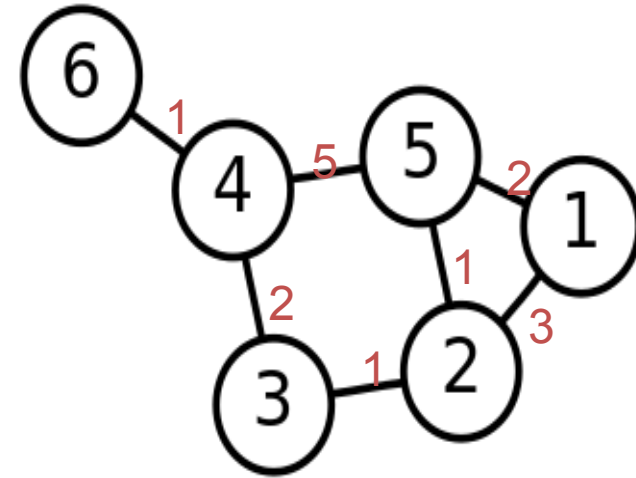
Python implementation

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    for k,v in Udict.items():
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            dmin=v #min distance
            n=k #corresponding node
```

3. Remove 'smallest-distance' node and update provisional distances of its unexplored neighbors

```
for m,en,wn in G.edges(n,data='weight'):
    ddist = dmin+wn
    if en in Edict:
        pass
    elif en in Udict:
        Udict[en] =
min(Udict[en],ddist)
    else:
        Udict[en] = ddist
    Edict[n] = Udict.pop(n)
```

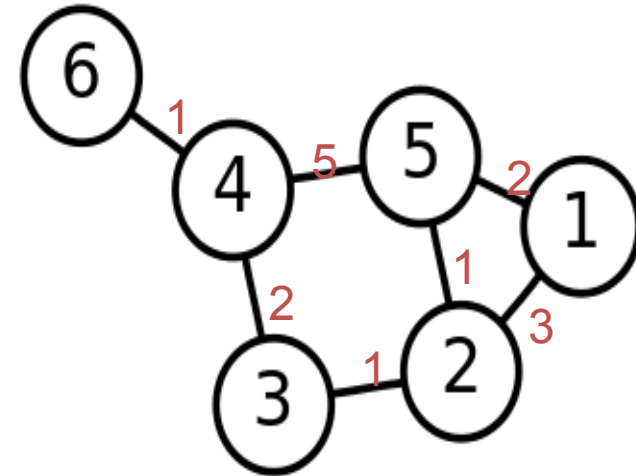


Shortest paths:

1→2: 3

1→4: 6

Dijkstra's algorithm



Setting the source to be node 5 and running the code:

```
In [56]: Edict
```

```
Out[56]: {1: 2, 2: 1, 3: 2, 4: 4, 5: 0, 6: 5}
```

The key-value pairs are node:distance

Dijkstra's algorithm

What is the cost for a graph with N nodes and M edges?

- Estimate cost of setting up graph to be $O(N+M)$
- Each edge is only considered maximum of two times $\rightarrow O(M)$
- dmin calculation $\rightarrow O(N^2)$ operations
- So it is the minimum calculation which is the bottleneck
- Can we do better?

