

# **IF232**

# **ALGORITHMS**

# **&**

# **DATA STRUCTURES**

06  
STACKS

DENNIS GUNAWAN

# REVIEW

## Linked Lists:

Double Linked Lists

Circular Linked Lists

# OUTLINE

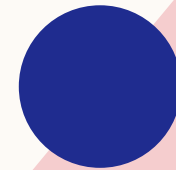
Array Representation of Stacks

Operations on a Stack

Linked Representation of Stacks

Operations on a Linked Stack

Applications of Stacks

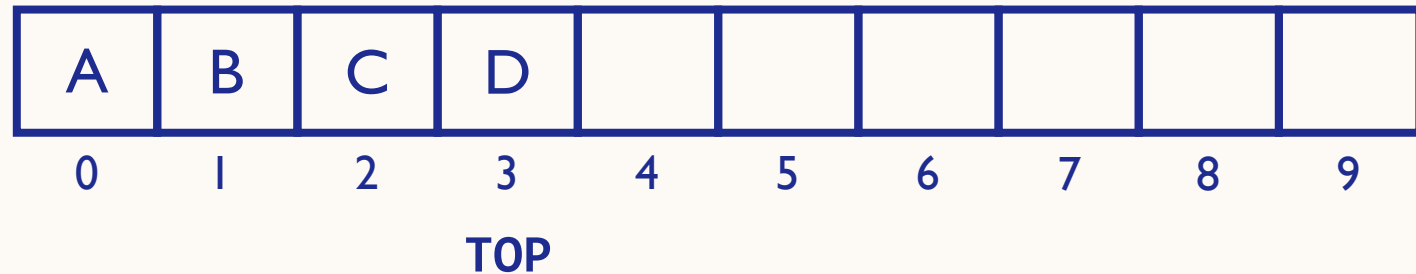


# STACKS

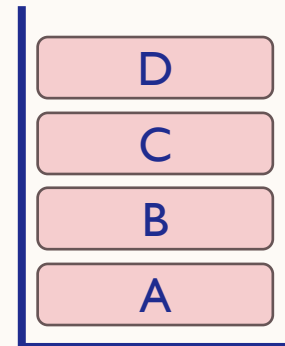
- LIFO: Last In First Out
- **TOP** is used to store the address of the topmost element of the stack
  - It is this position where the element will be added to or deleted from
- **SIZE** or **MAX** is used to store the maximum number of elements that the stack can hold

# ARRAY REPRESENTATION OF STACKS

MAX = 10

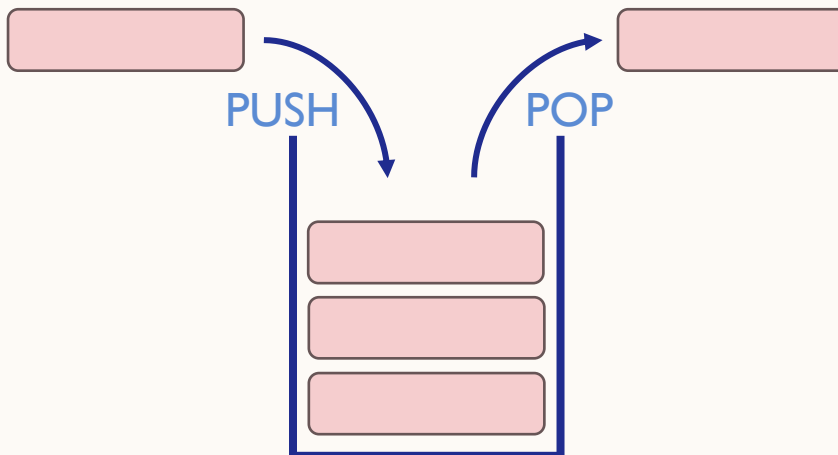


- The stack is empty:       $TOP = -1$
- The stack is full:       $TOP = MAX - 1$



# OPERATIONS ON A STACK

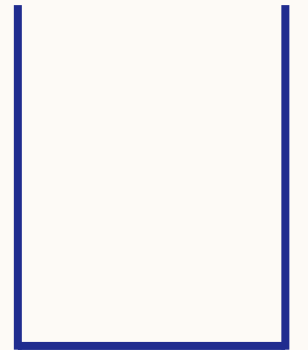
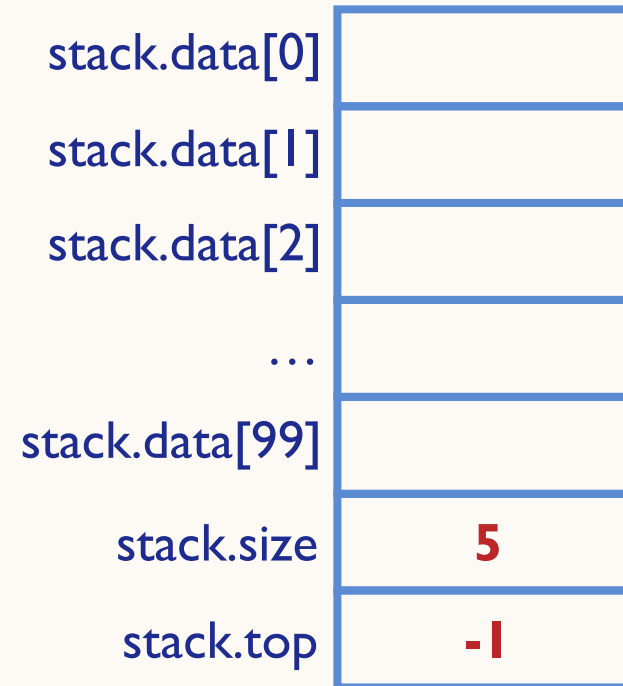
- **Push:** adds an element to the top of the stack
- **Pop:** removes the element from the top of the stack
- **Peek:** returns the value of the topmost element of the stack



# DECLARATION

```
struct tstack{  
    int data[100];  
    int size;  
    int top;  
};
```

```
int main()  
{  
    struct tstack stack;  
    int number;  
    ...  
    stack.size = 5;  
    stack.top = -1;  
    ...  
}
```

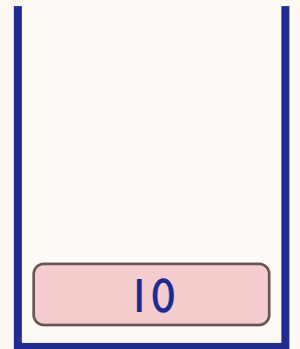


# PUSH OPERATION

```
scanf("%d", &number); //10

if(stack.top == stack.size - 1)
    printf("OVERFLOW");
else{
    stack.top++;
    stack.data[stack.top] = number;
}
```

stack.data[0]	10
stack.data[1]	
stack.data[2]	
...	
stack.data[99]	
stack.size	5
stack.top	0



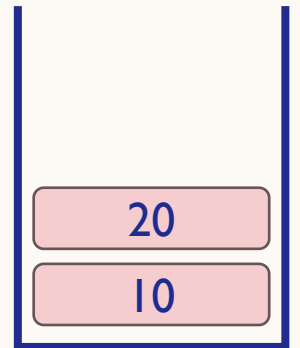


# PUSH OPERATION

```
scanf("%d", &number); //20

if(stack.top == stack.size - 1)
    printf("OVERFLOW");
else{
    stack.top++;
    stack.data[stack.top] = number;
}
```

stack.data[0]	10
stack.data[1]	20
stack.data[2]	
...	
stack.data[99]	
stack.size	5
stack.top	1

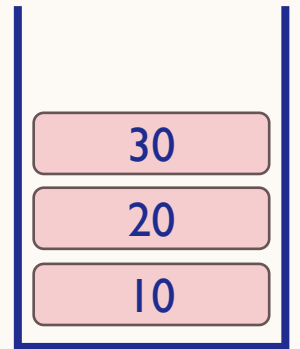


# PUSH OPERATION

```
scanf("%d", &number); //30

if(stack.top == stack.size - 1)
    printf("OVERFLOW");
else{
    stack.top++;
    stack.data[stack.top] = number;
}
```

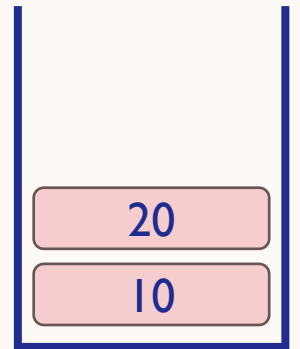
stack.data[0]	10
stack.data[1]	20
stack.data[2]	30
...	
stack.data[99]	
stack.size	5
stack.top	2



# POP OPERATION

```
if(stack.top == -1)
    printf("UNDERFLOW");
else{
    number = stack.data[stack.top]; //30
    stack.top--;
    //process number
}
```

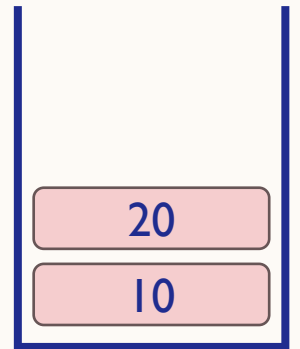
stack.data[0]	10
stack.data[1]	20
stack.data[2]	
...	
stack.data[99]	
stack.size	5
stack.top	1



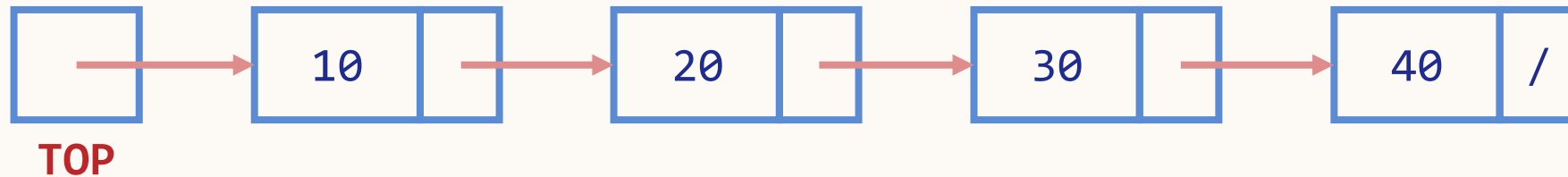
# PEEK OPERATION

```
if(stack.top == -1)
    printf("EMPTY");
else{
    //process stack.data[stack.top]; 20
}
```

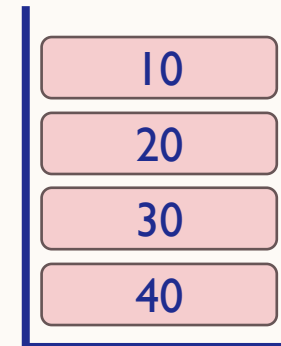
stack.data[0]	10
stack.data[1]	20
stack.data[2]	
...	
stack.data[99]	
stack.size	5
stack.top	1



# LINKED REPRESENTATION OF STACKS



- All insertions and deletions are done at the node pointed by **TOP**
- The stack is empty: **TOP** = **NULL**



# DECLARATION

```
struct tstack{  
    int data;  
    struct tstack *next;  
};
```

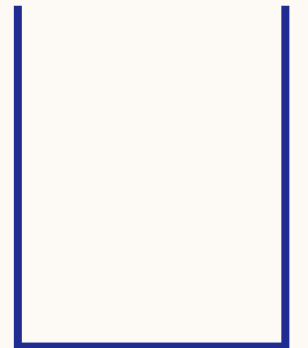
```
int main()  
{  
    struct tstack *top, *node;  
    int number;  
    ...  
    top = NULL;  
    ...  
}
```



top



node

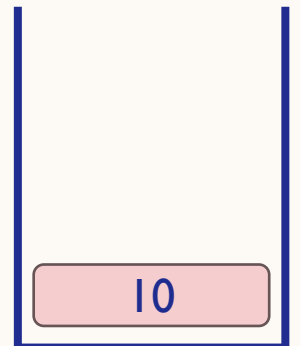
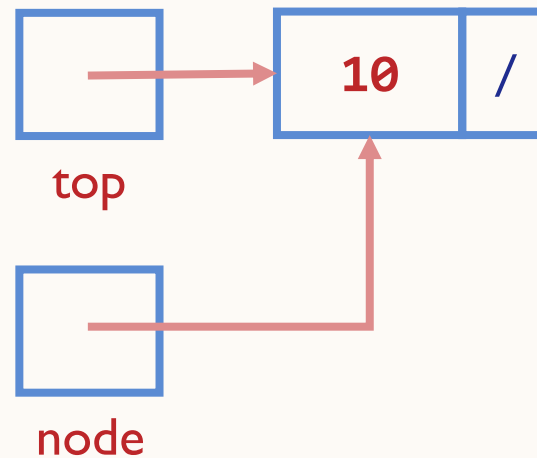


# PUSH OPERATION

```
scanf("%d", &number); //10

node = (struct tstack *) malloc
        (sizeof(struct tstack));
node->data = number;

if(top == NULL)
    node->next = NULL;
else
    node->next = top;
top = node;
```

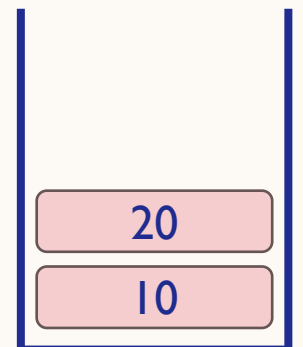
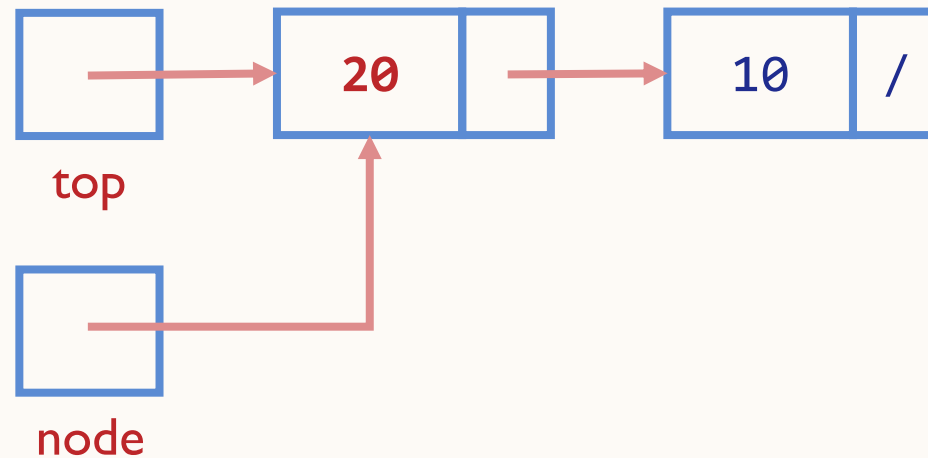


# PUSH OPERATION

```
scanf("%d", &number); //20

node = (struct tstack *) malloc
      (sizeof(struct tstack));
node->data = number;

if(top == NULL)
    node->next = NULL;
else
    node->next = top;
top = node;
```



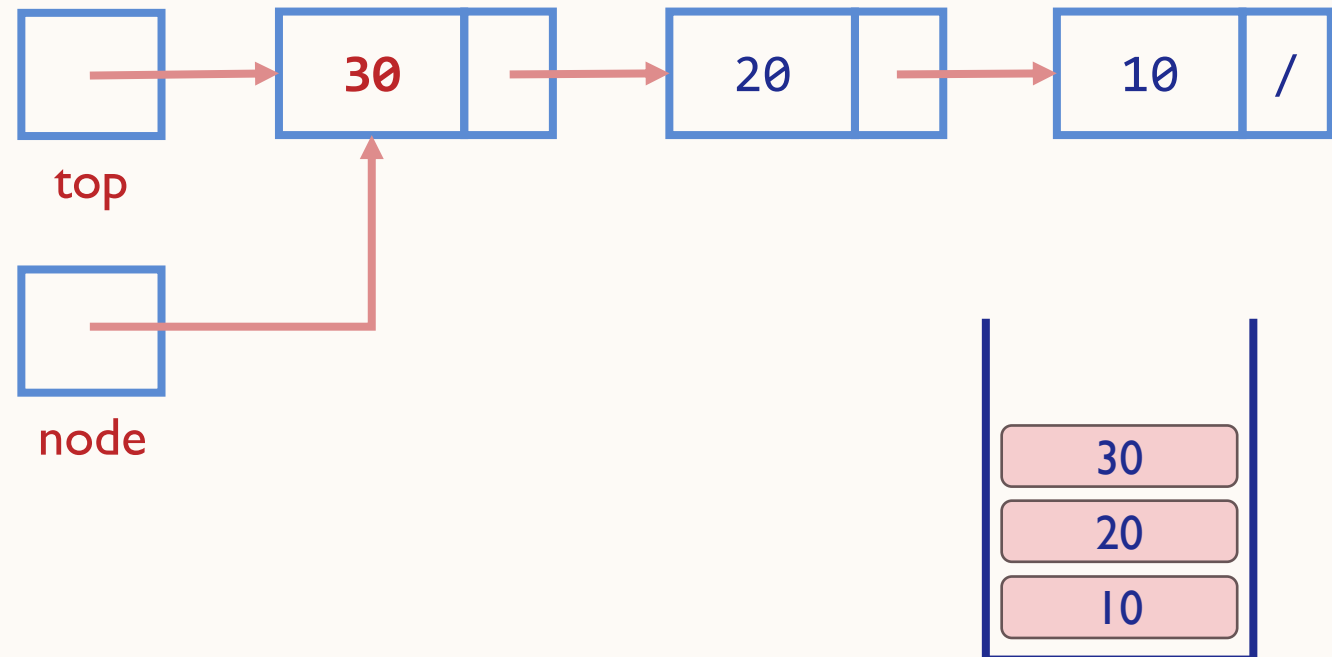


# PUSH OPERATION

```
scanf("%d", &number); //30

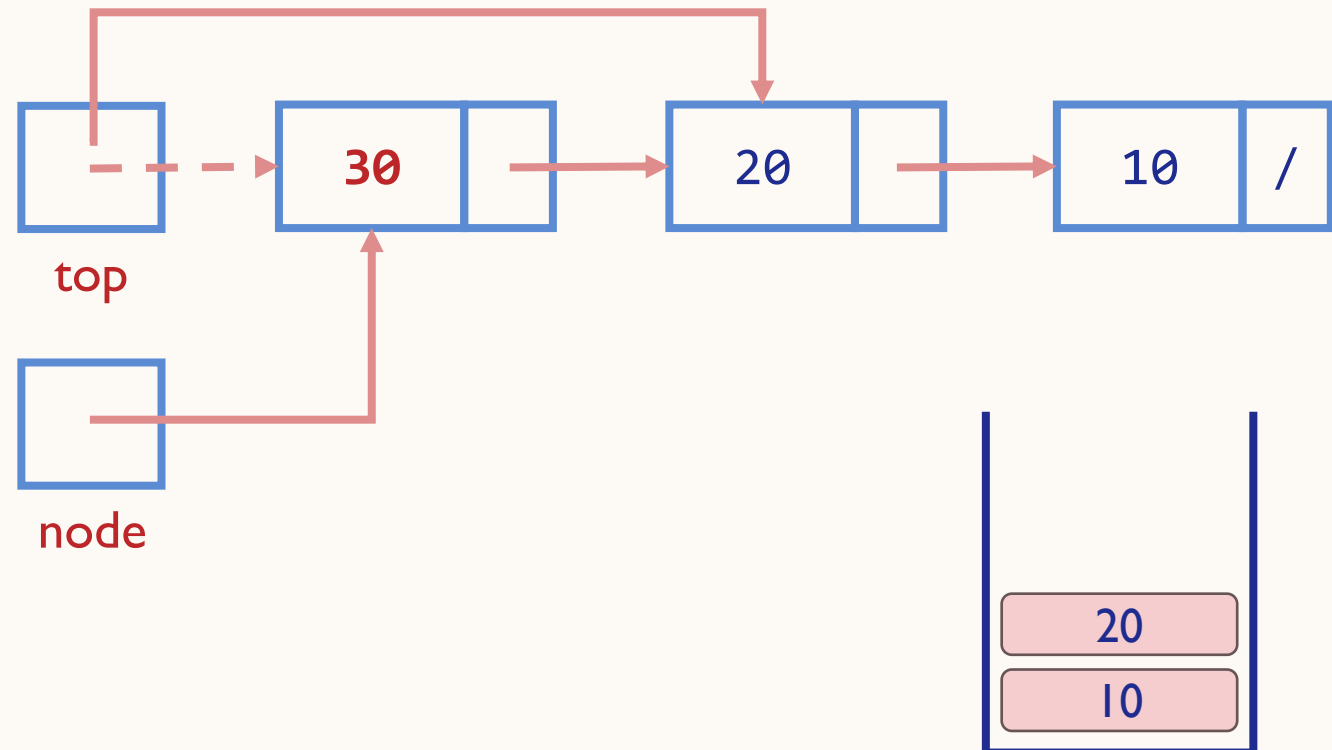
node = (struct tstack *) malloc
        (sizeof(struct tstack));
node->data = number;

if(top == NULL)
    node->next = NULL;
else
    node->next = top;
top = node;
```



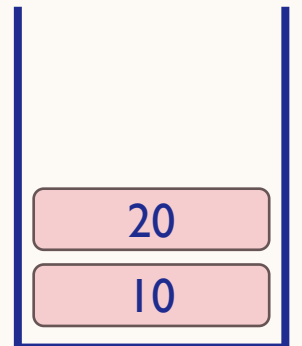
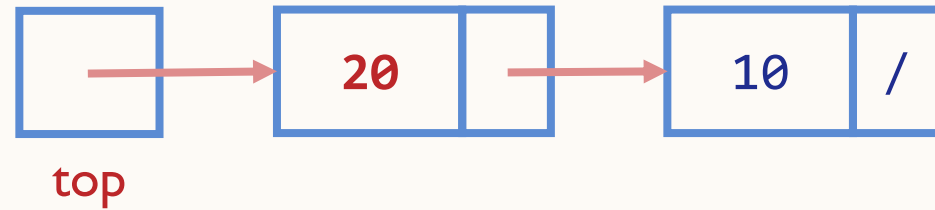
# POP OPERATION

```
if(top == NULL)
    printf("UNDERFLOW");
else{
    node = top;
    top = top->next;
    //process node->data 30
    free(node);
}
```



# PEEK OPERATION

```
if(top == NULL)
    printf("EMPTY");
else{
    //process top->data 20
}
```



# APPLICATIONS OF STACKS

- Reversing a list
- Parentheses checker
- Palindrome checker
- Conversion of an infix expression into a postfix expression
- Evaluation of a postfix expression
- Conversion of an infix expression into a prefix expression
- Evaluation of a prefix expression

# CONVERSION OF AN INFIX EXPRESSION INTO A POSTFIX EXPRESSION

- Convert the following infix expressions into postfix expressions.

$(A - B) * (C + D)$

$(A + B) / (C + D) - (D * E)$

$[AB-] * [CD+]$

$[AB+] / [CD+] - [DE*]$

$AB-CD+*$

$[AB+CD+ /] - [DE*]$

$AB+CD+ / DE*-$

# CONVERSION OF AN INFIX EXPRESSION INTO A POSTFIX EXPRESSION

- Algorithm to convert an infix notation to postfix notation

Step 1: Add “)” to the end of the infix expression

Step 2: Push “(” on the stack

Step 3: Repeat until each character in the infix notation is scanned

IF a “(” is encountered, push it on the stack

IF an operand (whether a digit or a character) is encountered, add it to the postfix expression

IF a “)” is encountered, then

a. Repeatedly pop from stack and add it to the postfix expression until a “(” is encountered

b. Discard the “(”. That is, remove the “(” from stack and do not add it to the postfix expression

IF an operator O is encountered, then

a. Repeatedly pop from stack and add each operator (popped from the stack) to the postfix expression which has the same precedence or a higher precedence than O

b. Push the operator O to the stack

[END OF IF]

Step 4: Repeatedly pop from the stack and add it to the postfix expression until the stack is empty

Step 5: EXIT

# CONVERSION OF AN INFIX EXPRESSION INTO A POSTFIX EXPRESSION

- Convert the following infix expression into postfix expression using the algorithm.

$(A - B) * (C + D)$

Infix Character Scanned	Stack	Postfix Expression
	(	
(	( (	
A	( (	A
-	( ( -	A
B	( ( -	A B
)	(	A B -

Infix Character Scanned	Stack	Postfix Expression
*	( *	A B -
(	( * (	A B -
C	( * (	A B - C
+	( * ( +	A B - C
D	( * ( +	A B - C D
)	( *	A B - C D +
)		A B - C D + *

# CONVERSION OF AN INFIX EXPRESSION INTO A POSTFIX EXPRESSION

- Convert the following infix expression into postfix expression using the algorithm.

**A - B \* C + D**

Infix Character Scanned	Stack	Postfix Expression
	(	
A	(	A
-	( -	A
B	( -	A B
*	( - *	A B

Infix Character Scanned	Stack	Postfix Expression
C	( - *	A B C
+	( +	A B C * -
D	( +	A B C * - D
)		A B C * - D +



# EVALUATION OF A POSTFIX EXPRESSION

- Algorithm to evaluate a postfix expression

Step 1: Scan every character of the postfix expression and repeat Step 2 until the end of the expression

Step 2: IF an operand is encountered, push it on the stack

IF an operator  $\theta$  is encountered, then

a. Pop the top two elements from the stack as A and B

b. Evaluate  $B \theta A$ , where A is the topmost element and B is the element below A

c. Push the result of evaluation on the stack

[END OF IF]

Step 3: SET RESULT equal to the topmost element of the stack

# EVALUATION OF A POSTFIX EXPRESSION

- Evaluate the following postfix expression using the algorithm.

9 3 4 \* 8 + 4 / -

Character Scanned	Stack
9	9
3	9 3
4	9 3 4
*	9 12
8	9 12 8
+	9 20
4	9 20 4
/	9 5
-	4

# CONVERSION OF AN INFIX EXPRESSION INTO A PREFIX EXPRESSION

- Convert the following infix expressions into prefix expressions.

$(A + B) * C$

$(A - B) * (C + D)$

$(A + B) / (C + D) - (D * E)$

$[+AB] * C$

$[-AB] * [+CD]$

$[+AB] / [+CD] - [*DE]$

$*+ABC$

$*-AB+CD$

$[/+AB+CD] - [*DE]$

$-/+AB+CD*DE$

# CONVERSION OF AN INFIX EXPRESSION INTO A PREFIX EXPRESSION

- Algorithm to convert an infix expression to prefix expression

Step 1: Reverse the infix string. Note that while reversing the string you must interchange left and right parentheses.

Step 2: Obtain the postfix expression of the infix expression obtained in Step 1.

Step 3: Reverse the postfix expression to get the prefix expression.

# CONVERSION OF AN INFIX EXPRESSION INTO A PREFIX EXPRESSION

- Convert the following infix expression into prefix expression using the algorithm.

$(A - B / C) * (A / K - L)$

Step 1:  $(L - K / A) * (C / B - A)$

Step 2:  $(L - [KA/]) * ([CB/] - A)$

$[LKA/-] * [CB/A-]$

$L K A / - C B / A - *$

Step 3:  $* - A / B C - / A K L$

# EVALUATION OF A PREFIX EXPRESSION

- Algorithm for evaluation of a prefix expression

Step 1: Repeat until all the characters in the prefix expression have been scanned

- a. Scan the prefix expression from right, one character at a time
- b. IF the scanned character is an operand, push it on the operand stack
- c. IF the scanned character is an operator, then
  - (i) Pop two values from the operand stack
  - (ii) Apply the operator on the popped operands
  - (iii) Push the result on the operand stack

# EVALUATION OF A PREFIX EXPRESSION

- Apply the algorithm to evaluate the following prefix expression.

**+ - 7 2 \* 8 / 12 4**

Character Scanned	Operand Stack
4	4
12	4 12
/	3
8	3 8
*	24
2	24 2
7	24 2 7
-	24 5
+	29



**PRACTICE**



# EXERCISES

- I. Draw the stack structure in each case when the following operations are performed on an empty stack.
  - a. Add A, B, C, D, E, F
  - b. Delete two letters
  - c. Add G
  - d. Add H
  - e. Delete four letters
  - f. Add I

# EXERCISES

2. Convert the following infix expressions to their postfix equivalents.

a.  $A - B + C$

b.  $A * B + C / D$

c.  $(A - B) + C * D / E - C$

d.  $(A * B) + (C / D) - (D + E)$

e.  $((A - B) + D / ((E + F) * G))$

f.  $(A - 2 * (B + C) / D * E) + F$

g.  $A - (B / C + (D \% E * F) / G) * H$

# EXERCISES

3. Find the infix equivalents of the following postfix equivalents.

a.  $A\ B\ +\ C\ *\ D\ -$

b.  $A\ B\ C\ *\ +\ D\ -$

4. Give the infix expression of the following prefix expressions.

a.  $\*\ -\ +\ A\ B\ C\ D$

b.  $\+\ -\ A\ *\ B\ C\ D$

# EXERCISES

5. Convert the expression given below into its corresponding postfix expression and then evaluate it.

a.  $10 + ((7 - 5) + 10) / 2$

b.  $14 / 7 * 3 - 4 + 10 / 2$

# REFERENCES

- Deitel, P. and Harvey Deitel (2022), C How to Program (9th Edition), Pearson Education.
- Thareja, R. (2014), Data Structures Using C (2nd Edition), India: Oxford University Press.

# NEXT

## Queues:

Array Representation of Queues

Operations on Queues

Linked Representation of Queues

Operations on Linked Queues

Applications of Queues

# VISION

To become an **outstanding** undergraduate Computer Science program that produces **international-minded** graduates who are **competent** in software engineering and have **entrepreneurial spirit** and **noble character**.

# MISSION

1. To conduct studies with the best technology and curriculum, supported by professional lecturer
2. To conduct research in Informatics to promote science and technology
3. To deliver science-and-technology-based society services to implement science and technology

Without hard work,  
nothing grows but weeds.



**if** INFORMATIKA  
UMN

Have patience.

All things are difficult before they become easy.