

LEARN CODING

ale66

STRUCTURED DATA

Structured data is the *gist* of *institutions*

Data organised in *tables* is structured

Silent conventions make such tables easy to recognise and work with

SPREADSHEETS

- broadly support working with struc. data on a personal level
- fail to separate data and code: impossible to maintain!

RELATIONAL DB MANAGEMENT SYSTEMS

- software that creates *isolated*, permanent databases: data is safe
- One RDBMS can host hundreds of databases, which remain independent from each other
- we create, access and modify data only via the SQL language

STRUCTURED DATA

- Table organisation can be cumbersome: see the *publishing company* example
- A single table shall only represent either
 - a single class of homogenous objects, e.g., writers, catalog entries...
 - a record of the relationship between two objects that are described in the db

THE RELATIONAL DATA MODEL

- columns (attribute) names shall be unique in each table
- rows may in principle be repeated.
- special columns called keys use ids to make rows unique
- A *primary key* is an ‘ID’ column that is usually specified for search

SQL TABLES

Careful planning of the table:

```
1 CREATE TABLE mywriters (
2     writer_id INTEGER(3) PRIMARY KEY,
3     writer_fname varchar(20) NOT NULL,
4     writer_lname varchar(20) NOT NULL,
5     writer_gender 'M' or 'F',
6     writer_birthyear INTEGER(4)
7 );
```

A table may have a lifespan of decades!

QUERYING TABLES

SQL naturally expressess queries on structured data:

```
1 SELECT gender, SUM(no_of_books), AVG(no_of_books)  
2  
3 FROM writers  
4  
5 GROUP BY gender;
```

QUERYING TABLES, 1

```
1 SELECT City, PostalCode, CustomerName  
2 FROM Customers  
3 WHERE Country = 'Germany' AND City = 'Berlin';
```

SELECT defines the format of the output, columns and their order and names

QUERYING TABLES, 2

```
1 SELECT City, PostalCode, CustomerName  
2 FROM Customers  
3 WHERE Country = 'Germany' AND City = 'Berlin';
```

SELECT defines the format of the output, columns and their order and names

FROM defines the table(s) to be queried

QUERYING TABLES, 3

```
1 SELECT City, PostalCode, CustomerName  
2 FROM Customers  
3 WHERE Country = 'Germany' AND City = 'Berlin';
```

SELECT defines the format of the output, columns and their order and names

FROM defines the table(s) to be queried

WHERE defines the conditions to be met *by the rows*

Column names act as variables

THE CONCEPT OF CURSOR

an invisible finger that iteratively will point at each row of the table

The column names-variables take up the values of the row pointed at by the cursor

QUERYING TABLES, 4

```
1 SELECT gender, SUM(no_of_books), AVG(no_of_books)
2
3 FROM writers
4
5 GROUP BY gender;
```

GROUP BY divides up the table in several minitable with
the same value of the grouping variable

HAVING selects the minitable according to a condition

QUERYING TABLES, 4

```
1 SELECT gender, SUM(no_of_books), AVG(no_of_books)
2
3 FROM writers
4
5 GROUP BY gender
6
7 HAVING SUM(no_of_books) > 10;
```

if male/female writers are not prolific enough we exclude them from the output

CURSOR

SQL executes row-by-row (mostly)

Cursor: an invisible finger pointing at the current row

Python can create cursors on tables and make them execute operations

KEY

Forces rows to be unique

Exists even in paper archive

```
1 CREATE TABLE students(  
2     student_id: PRIMARY KEY int(6),  
3     ...  
4 );
```

FOREIGN KEY

Key values used in another table

Supports the concept of *normalization*:

- each cell has an atomic value
-
- each table is strictly about either
 - a stand-alone class of objects: students, modules, lectures
 - a relationship, possibly with multiple instances, between said entities

FK EXAMPLE, 1

We include a FK column in order to access extra data from an external table

```
1 CREATE TABLE students(  
2     student_id: PRIMARY KEY int(6),  
3     ...  
4     // navigate to extra personal info  
5     nino: varchar(16) REFERENCES residents(nino)  
6 );
```

only inserts and updates involving a *valid* Nat'l insurance numbers are allowed

POSTGRESQL

```
1 CREATE TABLE students(  
2     student_id: PRIMARY KEY int(6),  
3     ...  
4     cod_fis: varchar(16),  
5     CONSTRAINT fiscal  
6         FOREIGN KEY(cod_fis)  
7             REFERENCES residents(cf)  
8 );
```

FK EXAMPLE, 2

A table recording relationship instances: the exam

```
1 CREATE TABLE exams(  
2     exam_date: date,  
3     mark: int(2),  
4     student_name: int(6) REFERENCES students(student_id),  
5     module_name: int(6) REFERENCES modules(module_id),  
6     programme: varchar(6) REFERENCES programmes(prog_id)  
7 );
```

Only inserts and updates involving three *valid* keys are allowed

POSTGRESQL

An example with the Italian fiscal code (16 characters)

```
1 CREATE TABLE students(
2     student_id: PRIMARY KEY int(6),
3     ...
4     cod_fis: varchar(16) REFERENCES residents(cf),
5     CONSTRAINT fiscal
6         FOREIGN KEY(cod_fis)
7         REFERENCES residents(cf)
8 );
```

SQL AT WORK

IEEE Spectrum: Top Programming languages 2023:
SQL is the most requested language in job advertisements
[worldwide]

MANY TABLES, MANY RELATIONSHIPS

- DBs are designed for the long-term
- Entity-Relationship diagrams: domain knowledge meets IT
- for maintenance, the more tables the better
- for querying, SELECTs become complex

JOINS

- Follows up the Foreign keys to create a giant summary table
- Row Selection follows
- column selection to improve readability

```
1 WHERE ATable.fk = AnotherTable.key
```

```
1 SELECT ATable.col, AnotherTable.anotherCol
2 FROM   ATable, AnotherTable
3 WHERE  ATable.fk = AnotherTable.key
4 AND    ATable.col > 10
5 AND    AnotherTable.aThirdCol = 0
6 ;
```

WITH ALIASES

```
1 SELECT T.col, AT.anotherCol
2 FROM ATable as T, AnotherTable as AT
3 WHERE T.fk = AT.key
4 AND T.col > 10
5 AND AT.aThirdCol = 0
6 ;
```

WITH COL. RENAMING

```
1 SELECT T.col as Name, AT.anotherCol as Income  
2 FROM ATable as T, AnotherTable as AT  
3 WHERE T.fk = AT.key  
4 AND AT.anotherCol > 10  
5 AND AT.aThirdCol = 0  
6 ;
```

EXAMPLE QUERIES

See again the E-R project for the Film studios

```
1 CREATE TABLE Cities (
2     city_id VARCHAR(30) PRIMARY KEY,
3     city_name VARCHAR(30) NOT NULL,
4     city_country VARCHAR(30) NOT NULL,
5     city_population INT(8)
6 );
```

```
1 CREATE TABLE Actors (
2     actor_id VARCHAR(30),
3     actor_name VARCHAR(30) NOT NULL,
4     actor_age INT(3),
5     actor_gender VARCHAR(1) NOT NULL,
6     born_in VARCHAR(30),
7     PRIMARY KEY (actor_id),
8     FOREIGN KEY (born_in) REFERENCES Cities(city_id)
9 );
```

```
1 CREATE TABLE Worked_in (
2     actor VARCHAR(30),
3     film VARCHAR(30),
4     salary INT(8),
5     PRIMARY KEY (actor_id, film_id),
6     FOREIGN KEY (actor) REFERENCES Actors(actor_id),
7     FOREIGN KEY (film) REFERENCES Films(film_id)
8 );
```

```
1 CREATE TABLE Films(  
2     film_id : PRIMARY KEY int(6),  
3     premiere : date,  
4     budget : int(10)  
5     ...  
6 );
```

Which actor is coming from a big city?

```
1 SELECT Actors.actor_name, Cities.city_name  
2 FROM Actors, Cities  
3 WHERE Actors.birth_city = Cities.city_id  
4 AND Cities.city_population > 1000000;
```

Are there pairs of actors who are from the same city?

We need two cursors running over the same table

```
1 SELECT A.actor_name, B.actor_name, A.birth_city  
2 FROM Actors AS A, Actors AS B  
3 WHERE A.birth_city = B.birth_city;
```

Select all pairs of actors who acted in the same film

We need **Worked_in**

```
1 SELECT A.actor_name, B.actor_name, W.film  
2 FROM Actors AS A, Actors AS B, Worked_in AS C, Worked_in AS D  
3 WHERE C.actor = A.actor_id  
4 AND D.actor = B.actor_id  
5 AND C.film = D.film  
6 ;
```

Has anyone ever acted in a film taking place in their own birthplace?

[we need to record the places where each film took place]

```
1 SELECT ...
2 FROM ...
3 WHERE ...
4 ;
```

