

Writing Assignment 1

Alec Merdler
CS 444: Operating Systems II
Spring 2017

1 INTRODUCTION

1.1 Processes

The explicit purpose of computers is to execute sets of instructions, called programs. As computers have increased in power and speed, the complexity and number of these programs has also increased. A simple program could count to the number ten, printing out each number along the way, then stop. The resources needed for a computer to execute this program would be minimal (barring any fancy recursion), and has only one explicit process: to count numbers. However, a more complex example would be a program that receives user input, performs extensive calculations, updates a database, and returns a result to the user. This type of program would be more aptly classified as a computer application, in that it contains many different components and processes, versus just counting numbers. Many of these processes can happen at the same time, and there can be duplicate processes in order to divide the work needed to be done. We can therefore think of a computer process as a program that is being executed, from which applications are built.

1.2 Threads

Each computer process requires resources in order to run, including the CPU, RAM, and hard disk storage. The operating system allocates these resources using threads.

1.3 CPU Scheduling

Any operating system that claims to be a multiprocessing operating system must have a way of managing all of the running processes and allocating system resources to them. One of the most critical resources is the CPU's time. Modern CPU's have multiple physical cores, which can execute their own process. However, in order to run even more processes at the same time, the operating system can quickly switch between different processes, giving them time to use the CPU, and creating the appearance of running many programs simultaneously.

2 LINUX

2.1 Processes

Each process (also known as a task) in Linux is represented by a `task_struct` data structure, which is naturally large and complex. Outlined below are the main basic functional areas and fields:

- State - Can be one of following: running, waiting, stopped, zombie
- Scheduling Information - Data regarding process precedence in the system
- Identifiers - Simply a unique number for the process
- Inter-Process Communication - Standard *nix information for processes to share data
- Links - Pointers to parent, sibling, and child processes
- Timers - Data about when the process was created and how much CPU time it has consumed
- File System - Pointers to working directories and other files
- Virtual Memory - How much virtual memory assigned to the process
- Processor Specific Context - Information that is needed to restart the process in the same state it was in

The tasks vector contains pointers to every `task_struct`. The default size of the task vector is 512 entries, and the currently running process can be found using the current pointer.

3 FREEBSD

3.1 Processes

Processes in FreeBSD are similar to those in Linux in many ways. Each process has a unique process ID (PID). Each process has one owner and group whose permissions determine which files and devices the process can open. Processes also have

a tree structure, with parent, sibling, and child processes.

New processes are created using the fork family of system calls. There are a few variations of the fork system call: fork creates a complete copy of the parent process; rfork creates a copy that shares a selected set of resources from its parent; vfork handles virtual memory differently than fork.

4 WINDOWS

4.1 Processes

Processes in Windows are similar to those in Linux in many ways. Each process contains a virtual address space, executable code, open handles to system objects, a security context, a unique process ID, environment variables, a priority class, minimum and maximum working set sizes, and at least one thread of execution. Processes also have a tree structure, with parent, sibling, and child processes.

Below is a code snippet showing how to create a child process in Windows using C++:

```
#include <windows.h>
#include <stdio.h>
#include <tchar.h>

void _tmain( int argc, TCHAR *argv[] )
{
    STARTUPINFO si;
    PROCESS_INFORMATION pi;

    ZeroMemory( &si, sizeof(si) );
    si.cb = sizeof(si);
    ZeroMemory( &pi, sizeof(pi) );

    if( argc != 2 )
    {
        printf("Usage: %s [cmdline]\n", argv[0]);
        return;
    }

    // Start the child process.
    if( !CreateProcess( NULL,          // No module name (use command line)
        argv[1],                    // Command line
        NULL,                        // Process handle not inheritable
        NULL,                        // Thread handle not inheritable
        FALSE,                       // Set handle inheritance to FALSE
        0,                           // No creation flags
        NULL,                        // Use parent's environment block
        NULL,                        // Use parent's starting directory
        &si,                          // Pointer to STARTUPINFO structure
        &pi )                        // Pointer to PROCESS_INFORMATION structure
    )
    {
        printf( "CreateProcess failed (%d).\n", GetLastError() );
        return;
    }

    // Wait until child process exits.
    WaitForSingleObject( pi.hProcess, INFINITE );

    // Close process and thread handles.
    CloseHandle( pi.hProcess );
}
```

```
        CloseHandle( pi.hThread );  
    }
```

4.2 *Threads*

Each Windows process is started with a single thread, but can create more threads which all share its virtual address space and system resources. Threads can be created by using the `CreateThread` function, which returns a handle to the thread.