

Data Structures and Algorithms – Project

Statement of the problem & Why is this ADT suitable?

An Art Gallery has to store a list of artists and the number of pieces they created, in an alphabetical order. Clients must be able to browse through the Portfolio in whatever direction they like (they can go forward or back). A **Map** is a suitable option for solving this problem mainly because we can store pairs of ArtistName and nrOfPieces, the name being the key, and the number of pieces the associated value. In addition, making the map a **sorted** one solves the order problem. Lastly, the **doubly linked list** allows us to browse through the list with the help of the next/previous operations and the iterator, just like flipping the pages of a real Portfolio.

ADT Sorted Map

I. Specification

A Map is a container with <key, value> pairs. If it is a Sorted Map, we also have a relation(order) on the set of possible keys.

Domain:

$M = \{m \mid m \text{ is a map with elements } e = (k, v), \text{ where } k \in TComp \text{ and } v \in TValue\}$

II. Interface

init(sm, r)
"""

creates a new empty sorted map
pre: $r \in R$ is a relation
post: $sm \in M$, sm is an empty sorted map
"""

destroy(sm)
"""

destroys a sorted map
pre: $sm \in M$
post: sm was destroyed
"""

add(sm, k, v)
"""

adds a new key-value pair to the sorted map
pre: $sm \in M, k \in TComp, v \in TValue$
post: (k, v) is added to the sorted map sm
"""

delete(sm, k, v)
"""

deletes a pair with a given key from the sorted map
pre: sm ∈ M, k ∈ TComp
post: v ∈ TValue

where v is v' , if k was found, or 0 otherwise
"""

search(sm, k, v)
"""

searches for the value associated with the given key in the sorted map
pre: sm ∈ M, k ∈ TComp
post: v ∈ TValue

where v is v' , if k was found, or 0 otherwise
"""

iterator(sm, it)
"""

returns an iterator for a sorted map
pre: sm ∈ M
post: it ∈ I is an iterator over sm
"""

size(sm)
"""

returns the number of pairs from the sorted map
pre: sm ∈ M
post: size <- the number of pairs from the sorted map
"""

keys(sm, s)
"""

returns the sorted set of keys from the sorted map
pre: sm ∈ M
post: s ∈ S is the set of keys from sm
"""

values(sm, b)
"""

returns a bag with all the values from the sorted map
pre: sm ∈ M
post: b ∈ B is the bag of all values from sm
"""

pairs(sm, s)
"""

returns the sorted set of all the pairs from the sorted map

pre: $sm \in M$

post: $s \in S$ is the set of all the pairs from sm

"""

III. Representation

We implement the Sorted Map over a doubly linked list with dynamic allocation.

Pair

key : String

value: Integer

Node

info: Pair

next: \uparrow Node

prev: \uparrow Node

SM

head: \uparrow Node

tail: \uparrow Node

r: Relation

IV. Implementation

- **subalgorithm init**(sm, r) Complexity: $O(1)$
 sm.head <- NIL
 sm.tail <- NIL
 sm.r <- r
end subalgorithm
- **subalgorithm destroy**(sm) Complexity: $O(n)$
 while sm.head != NIL do
 aux <- sm.head
 sm.head <- [sm.head].next
 free(aux)
 end-while
end subalgorithm
- **function search**(sm, key) Complexity: BC: $O(1)$,
WC: $O(n)$, AC: $O(n)$
 currentNode <- sm.head
 found <- false
 while currentNode != NIL and found == false
 if [currentNode].info.key == key
 found <- true
 else
 currentNode <- [currentNode].next
 end-if
 end-while
 if found == true
 search <- currentNode
 else
 search <- NIL
 end-if
end function

- **subalgorithm add**(sm, key, value)

Complexity: BC: $O(1)$,
WC: $O(n)$, AC: $O(n)$

```
ok <- 0
find <- search(sm, key)
currentNode <- sm.head

if currentNode == NIL
  //the list is empty
  [toAdd].info <- Pair(key, value)
  [toAdd].next <- NIL
  [toAdd].prev <- NIL
  sm.head <- toAdd
  sm.tail <- toAdd
end-if

if sm.r([currentNode].info.key, key) == 0
  //insert it first
  [toAdd].info <- Pair(key, value)
  [toAdd].next <- currentNode
  [currentNode].prev <- toAdd
  [toAdd].prev <- NIL
  sm.head <- toAdd
  ok <- 1
end-if

while currentNode != NIL and ok == 0
  //the list is not empty

  if [currentNode].next != NIL
    //it has more than one element
    if sm.r([currentNode].next.info.key, key) == 0
      and sm.r([currentNode].info.key, key) == 1
      [toAdd].info <- Pair(key, value)
      [toAdd].next <- [currentNode].next
      [currentNode].next <- toAdd
      [[toAdd].next].prev <- toAdd
      [toAdd].prev <- currentNode
      ok <- 1
      currentNode <- [currentNode].next
    end-if
  else
    //it has only one element
    if sm.r([currentNode].info.key, key) == 1
      //insert it after the currentNode
      [toAdd].info <- Pair(key, value)
      [toAdd].next <- NIL
      [currentNode].next <- toAdd
      [toAdd].prev <- currentNode
```

```

                                sm.tail <- toAdd
                                ok <- 1
                                currentNode <- [currentNode].next
                            else
                                //insert it before the currentNode
                                [toAdd].info <- Pair(key, value)
                                [toAdd].next <- currentNode
                                [currentNode].prev <- toAdd
                                [toAdd].prev <- NIL
                                sm.head <- toAdd
                                ok <- 1
                                currentNode <- [currentNode].next
                            end-if
                        end-if
                    currentNode <- [currentNode].next
                end-while
            end subalgorithm
```

- **subalgorithm delete**(sm, key) Complexity: BC: $O(1)$
WC: $O(n)$, AC: $O(n)$

```

    ok <- 0
    find <- search(sm, key)
    currentNode <- sm.head

    if currentNode == find
        //the first element is the one we want
        if [currentNode].next == NIL
            //it's the only element in the list
            sm.head <- NIL
            sm.tail <- NIL
            delete find
        else
            //there are more elements
            [[find].next].prev <- NIL
            sm.head <- [find].next
            delete find
        end-if
    end-if

    while currentNode != NIL
        if currentNode == find
            if [currentNode].next == NIL
                //the searched element is the last one
                [[find].prev].next <- NIL
                sm.tail <- [find].prev
                delete find
                currentNode = NIL
            else
                //there are more elements after it
```

```

                                aux <- [currentNode].next
                                [aux].prev <- [find].prev
                                [[find].prev].next <- aux
                                delete find
                                currentNode <- aux
                            end-if
                        else
                            currentNode <- [currentNode].next
                        end-if
                    end-while
end subalgorithm

```

- **function size**(sm) Complexity: $O(n)$

```

    nr <- 0
    currentNode <- sm.head
    while currentNode != NIL
        nr <- nr + 1
        currentNode <- [currentNode].next
    end-while
    size <- nr
end function

```

- **function keys**(sm, s) Complexity: $O(n)$

```

    currentNode <- sm.head
    while currentNode != NIL
        add(s, [currentNode].info.key)
        currentNode <- [currentNode].next
    end-while
    keys <- s
end function

```

- **function values**(sm, b) Complexity: $O(n)$

```

    currentNode <- sm.head
    while currentNode != NIL
        add(b, [currentNode].info.key)
        currentNode <- [currentNode].next
    end-while
    values <- b
end function

```

- **function pairs**(sm, s) Complexity: $O(n)$

```

    currentNode <- sm.head
    while currentNode != NIL
        add(s, [currentNode].info.key)
        currentNode <- [currentNode].next
    end-while
    pairs <- s
end function

```

- ***function iterator***(sm, it)
 iterator <- it(sm)
end function

Complexity: $O(1)$

Iterator

1) Interface

Domain: $I = \{it \mid it \text{ is an iterator over a sorted map with elements of type Pair } \}$

init(it, sm)
"""

creates a new iterator for a sorted map
pre: sm is a sorted map
post: $it \in I$ and it points to the first element in sm if sm is not empty or it is not valid
"""

getCurrent(it, e)
"""

returns the current element from the iterator
pre: $it \in I$, it is valid
post: e is a Pair, the current element from it
"""

next(it)
"""

moves the current element from the sorted map to the next element or makes the iterator invalid if no elements are left
pre: $it \in I$, it is valid
post: the current element from it points to the next element from the sorted map
"""

previous(it)
"""

moves the current element from the sorted map to the previous element or makes the iterator invalid if no elements are left
pre: $it \in I$, it is valid
post: the current element from it points to the previous element from the sorted map
"""

valid(it)
"""

verifies if the iterator is valid
pre: $it \in I$
post: valid <- True if it points to a valid element in the sorted map, False otherwise
"""

setFirst(it, sm)
"""

sets the iterator to the head of the sorted map

pre: $it \in I, sm \in M$

post: the iterator points to the head of the sm

"""

setLast(it, sm)
"""

sets the iterator to the tail of the sorted map

pre: $it \in I, sm \in M$

post: the iterator points to the tail of the sm

"""

2) Representation

Iterator

currentNode: \uparrow Node

sm: SM

3) Implementation

- **subalgorithm** *init*(it, sm) Complexity: $O(1)$
 it.sm <- sm
 it.currentNode <- sm.head
end subalgorithm
- **subalgorithm** *next*(it) Complexity: $O(1)$
 it.currentNode <- [it.currentNode].next
end subalgorithm
- **subalgorithm** *prev*(it) Complexity: $O(1)$
 it.currentNode <- [it.currentNode].prev
end subalgorithm
- **function** *getCurrent*(it) Complexity: $O(1)$
 getCurrent <- it.currentNode
end function
- **function** *valid*(it) Complexity: $O(1)$
 valid <- it.currentNode != NIL
end function

- ***subalgorithm setFirst***(it, sm) Complexity: O(1)
 it.currentNode <- sm.head
end subalgorithm
- ***subalgorithm setLast***(it, sm) Complexity: O(1)
 it.currentNode <- sm.tail
end subalgorithm

Solution of the problem

Menu:

1. Add an (artist, number of pieces) pair to the portfolio
2. See the whole portfolio
3. Delete an artist with a given name
4. Search for an artist with a given name
5. View the list from the start
6. View the list from the end
7. Next artist
8. Previous artist
0. Exit

"""

Method that prints a Node's information (Pair – the artist's name and the number of his pieces)

"""

- ***subalgorithm toString***(node) Complexity: O(1)
 print "Name: "
 print [Node].info.key
 print "Number of pieces: "
 print [Node].info.value
end subalgorithm

"""

Starts viewing the list from the beginning

Pre: $it \in I, sm \in M$

Post: the viewing starts and the first artist is being printed

"""

- ***subalgorithm first***(it, sm) Complexity: O(1)
 setFirst(it, sm)
 toString(getCurrent(it))
end subalgorithm

"""

Starts viewing the list from the end

Pre: $it \in I, sm \in M$

Post: the viewing starts and the last artist is being printed

"""

- **subalgorithm last**(it, sm) Complexity: $O(1)$
 setLast(it, sm)
 toString(getCurrent(it))
end subalgorithm

 """
 Moves the iterator to the next element in the list and prints it
 Pre: $it \in I, sm \in M$
 Post: the iterator points to the next artist(which is printed)
 """
- **subalgorithm next**(it, sm) Complexity: $O(1)$
 next(it)
 if valid(it)
 toString(getCurrent(it))
end subalgorithm

 """
 Moves the iterator to the previous element in the list and prints it
 Pre: $it \in I, sm \in M$
 Post: the iterator points to the previous artist (which is printed)
 """
- **subalgorithm prev**(it, sm) Complexity: $O(1)$
 previous(it)
 if valid(it)
 toString(getCurrent(it))
end subalgorithm

 """
 Function that prints the Portfolio using the iterator
 """
- **subalgorithm printList**(it) Complexity: $O(n)$
 while valid(it)
 aux <- getCurrent(it)
 toString(aux)
 next(it)
end subalgorithm

 """
 Starts the application and lets you choose an option from the menu
 """
- **subalgorithm start**(it, sm)
 printMenu() //prints the menu of the app
 initList() //initializes the Portfolio
 while option != 0
 print "Enter an option"
 read option
 if option == 1

```
//Add an artist and the number of his pieces (pair)
print "Name: "
read name

Print "Number of pieces: "
read nr
add(sm, name, nr)

else if option == 2
    //See the portfolio
    printList(it)

else if option == 3
    //Delete an artist
    print "Name: "
    read name

    delete(sm, name)

else if option == 4
    //Search for an artist
    print "Name: "
    read name

    node <- search(sm, name)
    if node != NIL
        toString(node)
    else print "No matches"
else if option == 5
    //From the start
    first(it, sm)
else if option == 6
    //From the end
    last(it, sm)
else if option == 7
    //Next artist
    next(it, sm)
else if option == 8
    //Previous artist
    prev(it, sm)

end subalgorithm
```

Tests

```
class Tests
{
public:
```

```
    /*
```

```
Default constructor for Tests
*/
Tests(){}

/*
Calls all the test functions
*/
void testAll()
{
    this->testSearch();
    this->testDelete();
    this->testAdd();
    this->testgetSize();
    this->testgetKeys();
    this->testgetValues();
    this->testgetPairs();
    this->testgetters();
}

/*
Initializes the list for the tests
Returns a DLL
*/
DLL initList()
{
    DLL l;
    l.add("Anne", 20);
    l.add("David", 1);
    l.add("Brianna", 10);
    return l;
}

/*
Tests the search function from the DLL
*/
void testSearch()
{
    DLL l = this->initList();
    Node *n = l.search("David");
    assert(n->getInfo().getKey() == "David");
    assert(n->getInfo().getValue() == 1);

    Node *m = l.search("Lily");
    assert(m == NULL);
}

/*
Tests the delete function from the DLL
*/
```

```
void testDelete()
{
    DLL l = this->initList();
    //delete the last element
    l.del("Brianna");
    assert(l.getSize() == 2);
    //delete the first element
    l.del("Anne");
    assert(l.getSize() == 1);

    DLL s{};
    //delete the only element
    s.add("Ana", 2);
    s.del("Ana");
    assert(s.getSize() == 0);

    DLL c = this->initList();
    //delete an element from the middle
    c.del("David");
    assert(c.getSize() == 2);

    //try to delete something that doesn't exist
    try
    {
        c.del("Serena");
    }
    catch (Exception& e)
    {
    }
}

/*
Tests the add function from the DLL
*/
void testAdd()
{
    DLL l{};
    //empty list
    l.add("Maria", 21);
    assert(l.getSize() == 1);

    DLL s{};
    s.add("Dave", 22);
    //add something before
    s.add("Andreea", 10);
    assert(s.getSize() == 2);
    assert(s.getHead()->getInfo().getKey() == "Andreea");
}
```

```
DLL a{};
a.add("Alina", 3);
//add something after
a.add("Matt", 8);
assert(a.getSize() == 2);
assert(a.getTail()->getInfo().getKey() == "Matt");

DLL b = this->initList();
//add something in the middle
b.add("Chris", 24);
assert(b.getSize() == 4);
assert(b.getKeys().at(2) == "Chris");
}

/*
Tests the getSize function from the DLL
*/
void testgetSize()
{
    DLL l = this->initList();
    assert(l.getSize() == 3);
}

/*
Tests the getKeys function from the DLL
*/
void testgetKeys()
{
    DLL l = this->initList();
    vector<string> s = l.getKeys();

    assert(s.at(0) == "Anne");
}

/*
Tests the getValues function from the DLL
*/
void testgetValues()
{
    DLL l = this->initList();
    vector<int> v = l.getValues();

    assert(v.at(0) == 20);
}

/*
Tests the getPairs function from the DLL
```

```
    */
    void testgetPairs()
    {
        DLL l = this->initList();
        vector<Pair> p = l.getPairs();

        assert(p.at(0).getKey() == "Anne");
        assert(p.at(0).getValue() == 20);

    }

    /*
    Tests the getHead and getTail functions from the DLL
    */
    void testgetters()
    {
        DLL l = this->initList();
        assert(l.getHead()->getInfo().getKey() == "Anne");
        assert(l.getTail()->getInfo().getKey() == "David");
    }

    /*
    Default destructor for Tests
    */
    ~Tests(){}

};
```