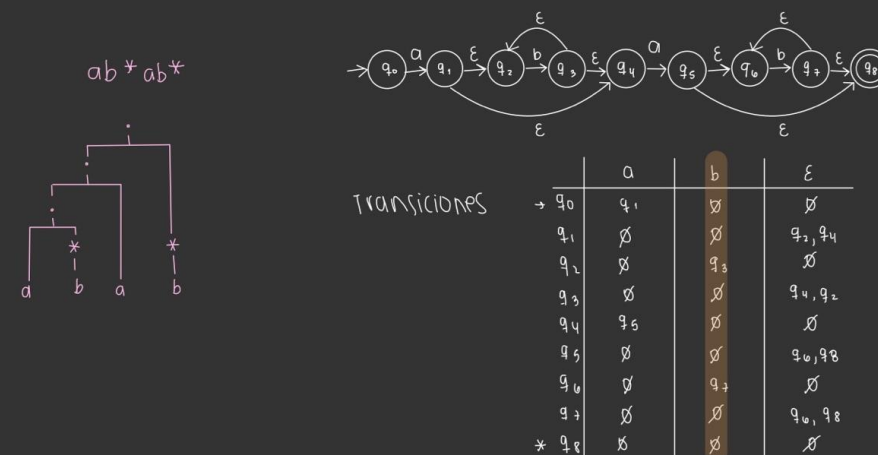


Pre Laboratorio No.1

- ab^*ab^*

Construya un AFN con el algoritmo de Thompson:



construya un AFD utilizando el algoritmo de subconjuntos:

calcular cerraduras ε:

ε closure (q₀) = {q₀}

(q₁) = {q₁, q₂, q₄}

(q₂) = {q₂}

(q₃) = {q₃, q₅, q₆}

(q₄) = {q₄}

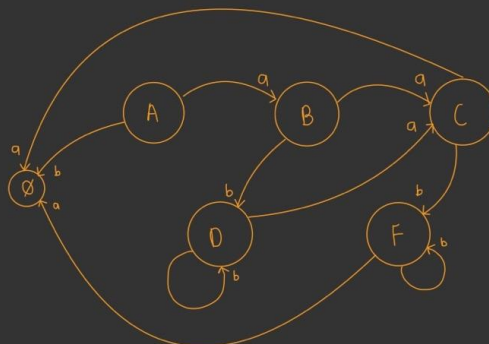
(q₅) = {q₅, q₆, q₈}

(q₆) = {q₆}

(q₇) = {q₆, q₇, q₈}

(q₈) = {q₈}

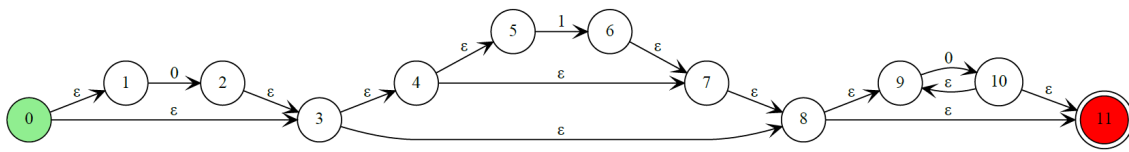
	a	b
A → {q ₁ }	{q ₁ , q ₂ , q ₄ , q ₅ , q ₆ , q ₈ } A	∅
B → {q ₂ , q ₃ , q ₄ , q ₅ , q ₆ , q ₈ }	{q ₅ , q ₆ , q ₈ } C	{q ₂ , q ₃ , q ₄ , q ₆ , q ₇ , q ₈ } D
C → {q ₅ , q ₆ , q ₈ }	∅	{q ₆ , q ₇ , q ₈ } F
D → {q ₂ , q ₃ , q ₄ , q ₆ , q ₇ , q ₈ }	{q ₅ , q ₆ , q ₈ } C	{q ₂ , q ₃ , q ₄ , q ₆ , q ₇ , q ₈ } D
F → {q ₆ , q ₇ , q ₈ }	∅	{q ₆ , q ₇ , q ₈ } F



- $0?(1?)^?0^*$

- AFN

- Diagrama



- Tabla de transiciones

	0	1	ϵ
0	\emptyset	\emptyset	1,3
1	2	\emptyset	\emptyset
2	\emptyset	\emptyset	3
3	\emptyset	\emptyset	4,8
4	\emptyset	\emptyset	5,7
5	\emptyset	6	\emptyset
6	\emptyset	\emptyset	7
7	\emptyset	\emptyset	8
8	\emptyset	\emptyset	9,11
9	10	\emptyset	\emptyset
10	\emptyset	\emptyset	11,9
11	\emptyset	\emptyset	\emptyset

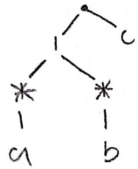
- AFD (lab B)

- $(a^*|b^*)c$

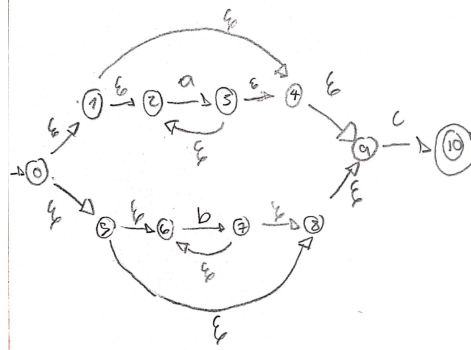
Construcción por Thomson

1. Postfix: $a * b * | c$.

2. Árbol:



3. Dibujo:



4. Tabla de transiciones

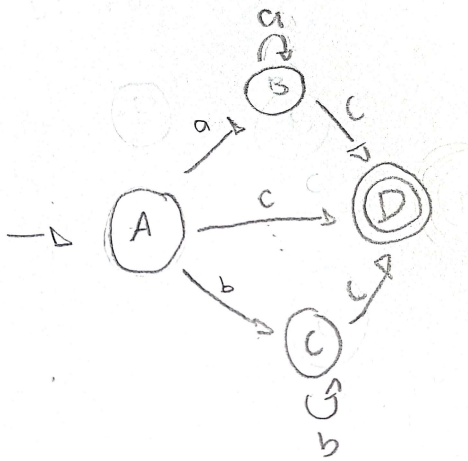
	A	B	C	ϵ
0	\emptyset	\emptyset	\emptyset	{1, 5}
1	\emptyset	\emptyset	\emptyset	{2, 4}
2	{3}	\emptyset	\emptyset	\emptyset
3	\emptyset	\emptyset	\emptyset	{2, 4}
4	\emptyset	\emptyset	\emptyset	{9}
5	\emptyset	\emptyset	\emptyset	{6, 8}
6	\emptyset	{7}	\emptyset	\emptyset
7	\emptyset	\emptyset	\emptyset	{6, 8}
8	\emptyset	\emptyset	\emptyset	{9}
9	\emptyset	\emptyset	{10}	\emptyset
10	\emptyset	\emptyset	\emptyset	\emptyset

Construcción de subconjuntos

1. Tabla de transiciones

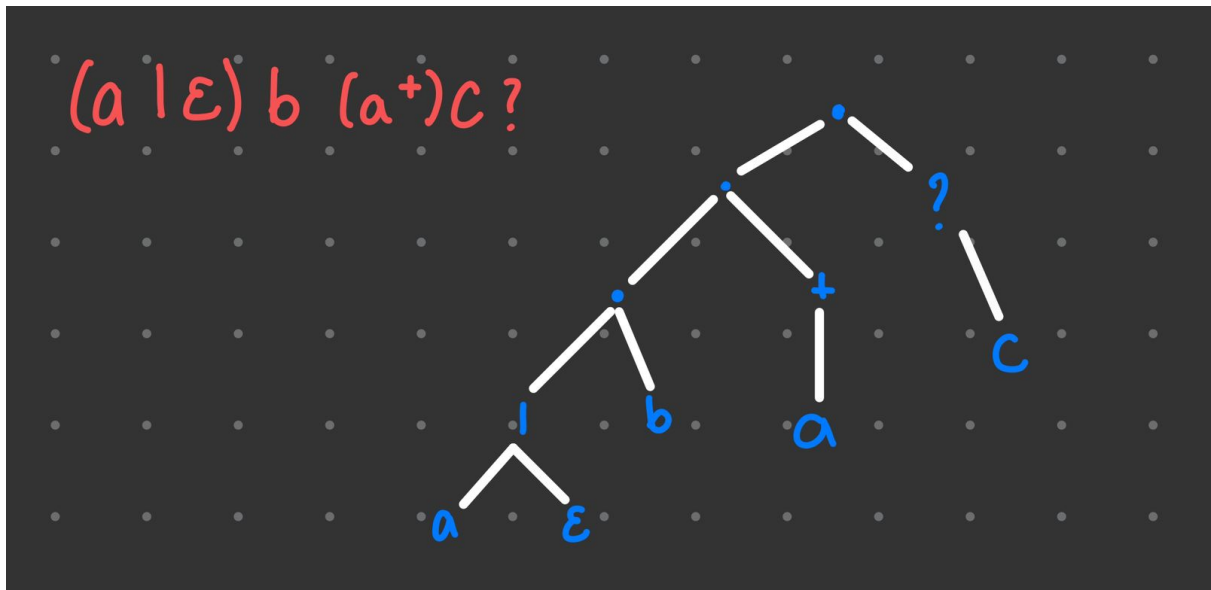
	Transiciones		
Subconjunto (estados)	a	b	c
A = {0, 1, 2, 4, 5, 6, 8, 9}	B	C	D
B = {2, 3, 4, 9}	B	\emptyset	D
C = {6, 7, 8, 9}	\emptyset	C	D
D = {10}	\emptyset	\emptyset	\emptyset

2. Dibujo

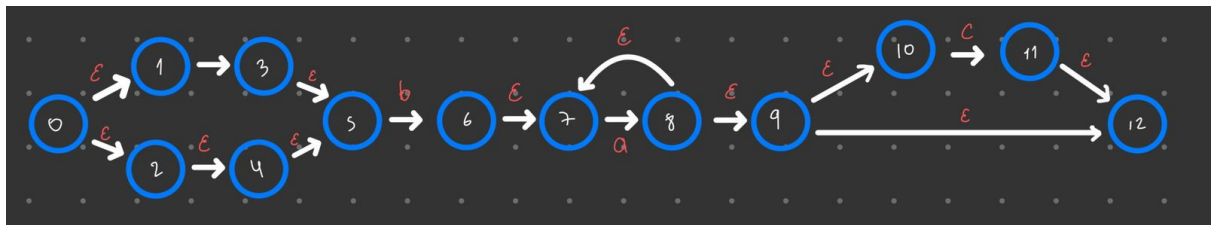


- $(b|b)^*abb(a|b)^*$

1. Arbol de expresión



2. AFN (Dibujo)



3. Tabla de transición

	A	B	C	ϵ
0	\emptyset	\emptyset	\emptyset	{1, 2}
1	{3}	\emptyset	\emptyset	\emptyset
2	\emptyset	\emptyset	\emptyset	{4}
3	\emptyset	\emptyset	\emptyset	{5}
4	\emptyset	\emptyset	\emptyset	{5}
5	\emptyset	{6}	\emptyset	\emptyset
6	\emptyset	\emptyset	\emptyset	{7}

7	{8}	\emptyset	\emptyset	\emptyset
8	\emptyset	\emptyset	\emptyset	{ 7,9 }
9	\emptyset	\emptyset	\emptyset	{ 10,12 }
10	\emptyset	\emptyset	{11}	\emptyset
11	\emptyset	\emptyset	\emptyset	{12}
12	\emptyset	\emptyset	\emptyset	\emptyset

4.

- $(a|b)^*a(a|b)(a|b)$

$$-(a|b)^* \cdot a \cdot (a|b) \cdot (a|b)$$

$$\rightarrow ab^*a \cdot ab| \cdot ab|.$$

árbol sintáctico:



autómata AFN

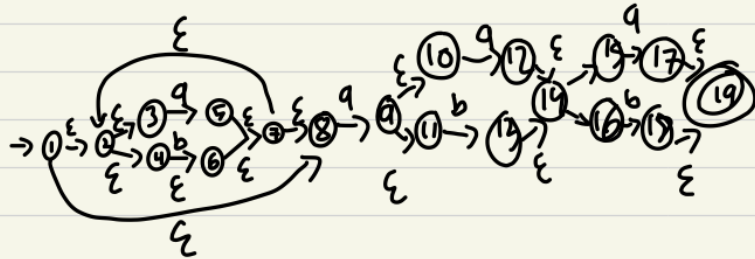


tabla de transiciones

	A	B	ε
1	∅	∅	2,8
2	∅	∅	3,4
3	5	∅	∅
4	∅	6	∅
5	∅	∅	7
6	∅	∅	7
7	∅	∅	8
8	9	∅	∅
9	∅	∅	10,11
10	12	∅	∅
11	∅	13	∅
12	∅	∅	14
13	∅	∅	14
14	∅	∅	15,16
15	17	∅	∅
16	∅	18	∅
17	∅	∅	19
18	∅	∅	19
19	∅	∅	∅