

Outline

1 What is a Binary Search Tree (BST)?

2 Implementation of the Ordered Symbol Table API Using a BST

3 Binary Tree Traversal

4 Performance Characteristics



What is a Binary Search Tree (BST)?	
A binary tree is either empty or a node with a key (and associated value) and links (left and r subtrees	ight) to two disjoint binary

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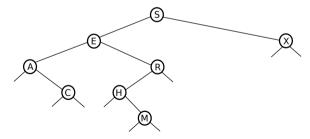
A binary tree is in symmetric order if each node's key is larger than all keys in its left subtree and smaller than all keys in its right subtree

What is a Binary Search Tree (BST)?

A binary tree is either empty or a node with a key (and associated value) and links (left and right) to two disjoint binary subtrees

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A binary search tree (BST) is a binary tree in symmetric order





A BST representation in Java is a reference to a root node, which is composed of five fields: a key, a value, a reference to the left subtree, a reference to the right subtree, and the number of nodes in the subtree

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A BST representation in Java is a reference to a root node, which is composed of five fields: a key, a value, a reference to the left subtree, a reference to the right subtree, and the number of nodes in the subtree

```
private class Node {
   private Key key;
   private Value val;
   private int size;
   private Node left, right;

public Node(Key key, Value value) {
      this.key = key;
      this.val = value;
      this.size = i;
   }
}
```



```
☑ BinarySearchTreeST.java

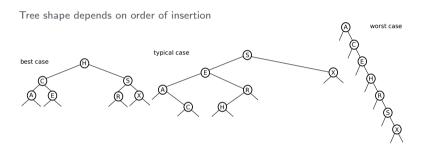
package dsa:
import java.util.NoSuchElementException;
import stdlib.StdIn:
import stdlib.StdOut;
public class BinarySearchTreeST<Key extends Comparable<Key>, Value>
        implements OrderedST < Key, Value > {
    private Node root:
    public BinarySearchTreeST() {
        root = null;
    public boolean isEmpty() {
        return size() == 0;
    public int size() {
        return size(root):
    public void put (Kev kev. Value value) {
        if (kev == null) {
            throw new IllegalArgumentException("kev is null"):
        if (value == null) {
            throw new IllegalArgumentException("value is null"):
        root = put(root, kev, value):
    public Value get(Kev kev) {
        if (kev == null) {
            throw new IllegalArgumentException("kev is null"):
        }
```

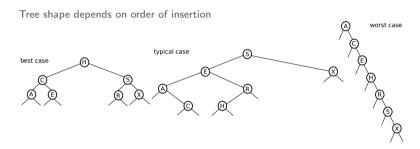


```
☑ BinarySearchTreeST.java

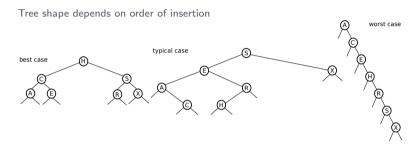
    public boolean contains (Kev kev) {
        if (key == null) {
            throw new IllegalArgumentException("key is null");
        return get(key) != null;
    private Node put(Node x. Kev kev. Value value) {
        if (x == null) {
            return new Node(key, value);
        int cmp = kev.compareTo(x.kev);
        if (cmp < 0) {
            x.left = put(x.left, key, value);
        } else if (cmp > 0) {
            x.right = put(x.right, key, value);
        } else {
            x.val = value:
        x.size = size(x.left) + size(x.right) + 1;
        return x:
    private Value get (Node x. Kev kev) {
        if (x == null) {
            return null:
        int cmp = kev.compareTo(x.kev);
        if (cmp < 0) {
            return get(x.left, kev):
        } else if (cmp > 0) {
            return get(x,right, kev);
        } else {
            return x.val:
```







Number of comparisons for search/insert is \sim depth of node

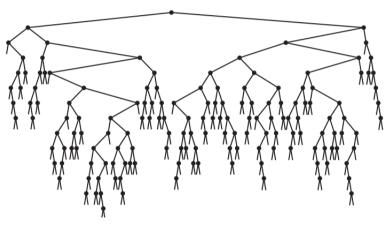


Number of comparisons for search/insert is \sim depth of node

If n distinct keys are inserted into a BST in random order, the expected number of comparisons for a search/insert is $\sim \lg n$

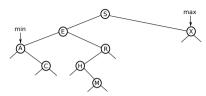


Typical BST, built from 256 random keys

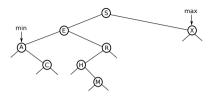




Minimum and maximum



Minimum and maximum

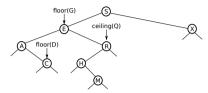


```
☑ BinarySearchTreeST.java

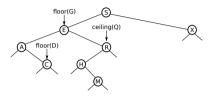
    public Key min() {
        if (isEmpty()) { return null; }
        return min(root).kev:
    private Node min(Node x) {
        if (x.left == null) { return x; }
        else
                  { return min(x.left); }
    public Kev max() {
        if (isEmpty()) { return null; }
        return max(root).key;
    private Node max(Node x) {
        if (x.right == null) { return x: }
        else
                            { return max(x,right); }
```



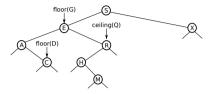
Floor and ceiling



Floor and ceiling



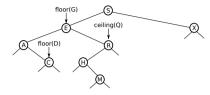
Floor and ceiling



Computing the floor of key k

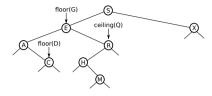
• Case 1 (k equals the key in the node): the floor of k is k

Floor and ceiling



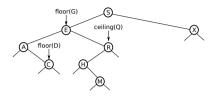
- Case 1 (k equals the key in the node): the floor of k is k
- Case 2 (k is less than the key in the node): the floor of k is in the left subtree

Floor and ceiling



- Case 1 (k equals the key in the node): the floor of k is k
- Case 2 (k is less than the key in the node): the floor of k is in the left subtree
- Case 3 (k is greater than the key in the node): the floor of k is in the right subtree if there is any key $\leq k$ in there; otherwise, it is the key in the node

Floor and ceiling



- Case 1 (k equals the key in the node): the floor of k is k
- Case 2 (k is less than the key in the node): the floor of k is in the left subtree
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Computing the ceiling of key k

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Computing the ceiling of key k

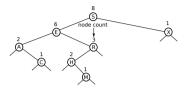
- Case 1 (k equals the key in the node): the ceiling of k is k
- Case 2 (k is greater than the key in the node): the ceiling of k is in the right subtree
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Computing the ceiling of key k

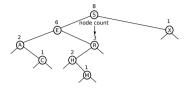
- Case 1 (k equals the key in the node): the ceiling of k is k
- Case 2 (k is greater than the key in the node): the ceiling of k is in the right subtree
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Rank and selection



Rank and selection



```
☑ BinarySearchTreeST.java

    public int rank(Kev kev) { return rank(kev, root); }
    private int rank(Key key, Node x) {
        if (x == null) { return 0; }
        int cmp = kev.compareTo(x.kev);
                (cmp < 0) { return rank(key, x.left); }
        else if (cmp > 0) { return 1 + size(x.left) + rank(key, x.right); }
        0100
                           { return size(x.left): }
    public Kev select(int k) {
        if (k < 0 \mid | k >= size()) { return null: }
        Node x = select(root, k):
        return x.key;
    private Node select(Node x, int k) {
        if (x == null) { return null: }
        int t = size(x.left):
                (t > k) { return select(x.left. k): }
        else if (t < k) { return select(x.right, k - t - 1); }</pre>
                         { return x: }
        0100
```



Range count and range search

```
☑ BinarySearchTreeST.java
    public int size(Kev lo. Kev hi) {
        if (lo.compareTo(hi) > 0) { return 0: }
        if (contains(hi)) { return rank(hi) - rank(lo) + 1; }
                           { return rank(hi) - rank(lo); }
        else
    public Iterable < Key > keys() {
        return keys(min(), max());
    public Iterable < Kev > kevs (Kev lo, Kev hi) {
        Queue < Kev > queue = new Queue < Kev > ():
        keys(root, queue, lo, hi);
        return queue:
    private void kevs(Node x. Queue<Kev> queue. Kev lo. Kev hi) {
        if (x == null) { return: }
        int cmplo = lo.compareTo(x.kev);
        int cmphi = hi.compareTo(x.key);
        if (cmplo < 0) { kevs(x.left, queue, lo, hi); }
        if (cmplo <= 0 && cmphi >= 0) { queue.enqueue(x.kev); }
        if (cmphi > 0) { kevs(x.right, queue, lo, hi); }
```



Deletion: to delete the minimum (maximum) key

• Go left (right) until you find a node with null left (right) link

- Go left (right) until you find a node with null left (right) link
- Replace that node by its right (left) link

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- Update subtree counts

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```
☑ BinarySearchTreeST.java
    public void deleteMin() {
        if (isEmpty()) { throw new NoSuchElementException(); }
        root = deleteMin(root);
    private Node deleteMin(Node x) {
        if (x.left == null) { return x.right; }
        x.left = deleteMin(x.left);
        x.N = size(x.left) + size(x.right) + 1;
        return x;
    public void deleteMax() {
        if (isEmptv()) { throw new NoSuchElementException(): }
        root = deleteMax(root);
    private Node deleteMax(Node x) {
        if (x.right == null) { return x.left; }
        x.right = deleteMax(x.right);
        x.N = size(x.left) + size(x.right) + 1:
        return x:
```



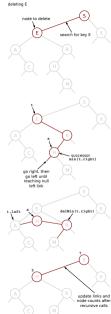
Deletion: to delete a node with key k (Hibbard deletion), search for the node t containing key ksuccessor min(t,right) go right, then go left until reaching null left link delMin(t.right) update links and node counts after

deleting E

recursive calls

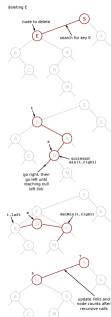
Deletion: to delete a node with key k (Hibbard deletion), search for the node t containing key k

 Case 1 (0 children): delete t by setting parent link to null



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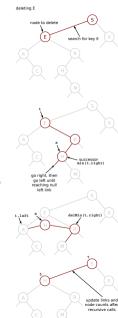
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Deletion: to delete a node with key k (Hibbard deletion), search for the node t containing key k

- Case 1 (0 children): delete t by setting parent link to null
- Case 2 (1 child): delete t by replacing parent link
- Case 3 (2 children): find successor x of t (x has no left child); delete the minimum in t's right subtree; and put x in t's spot

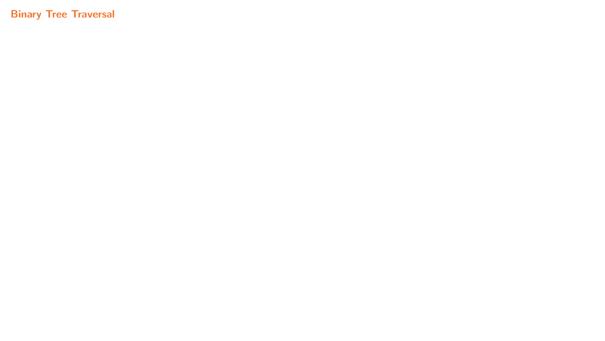
and update subtree counts





```
☑ BinarySearchTreeST.java

    public void delete(Kev kev) {
        root = delete(root, kev):
    private Node delete(Node x. Kev kev) {
        if (x == null) { return null; }
        int cmp = kev.compareTo(x.kev):
               (cmp < 0) { x.left = delete(x.left, kev); }
        else if (cmp > 0) { x.right = delete(x.right, key); }
        else {
            if (x.right == null) { return x.left: }
            if (x.left == null) { return x.right: }
            Node t = x:
            x = min(t.right);
            x.right = deleteMin(t.right);
            x.left = t.left:
        x.N = size(x.left) + size(x.right) + 1;
        return x:
```



Pre-order traversal

```
public void preorder() { preorder(root); }

private void preorder(Node x) {
    if (x == null) { return; }
    process(x);
    preorder(x.left);
    preorder(x.right);
}
```

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public void preorder() { preorder(root); }

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In-order traversal

```
public void inorder() { inorder(root); }

private void inorder(Node x) {
    if (x == null) { return; }
    inorder(x.left);
    process(x);
    inorder(x.right);
}
```

Pre-order traversal

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public void preorder() { preorder(root); }

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public void inorder() { inorder(root); }

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    if (x == null) { return; }
    inorder(x.left);
    process(x);
    inorder(x.right);
}
```

Post-order traversal

```
public void postorder() { postorder(root); }

private void postorder(Node x) {
   if (x == null) { return; }
   postorder(x.left);
   postorder(x.right);
   process(x);
}
```

Pre-order traversal

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public void preorder() { preorder(root); }

private void preorder(Node x) {
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```

In-order traversal

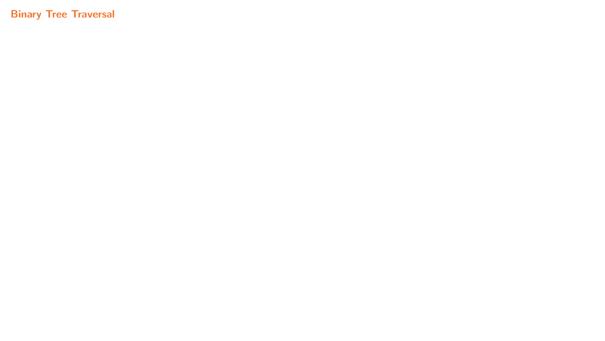
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    if (x == null) { return; }
    inorder(x.left);
    process(x);
    inorder(x.right);
}
```

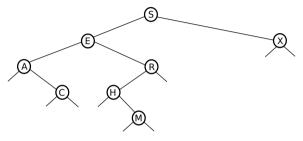
Post-order traversal

```
public void postorder() { postorder(root); }

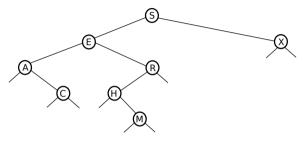
private void postorder(Node x) {
   if (x == null) { return; }
   postorder(x.left);
   postorder(x.right);
   process(x);
}
```



For example, let $_{\mathtt{root}}$ denote the following BST



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The calls preorder(), inorder(), and postorder() will process the tree as follows

```
s e a c R H M x (preorder)

a c e H M R s x (inorder)

c a M H R E x s (postorder)
```



Performance Characteristics

Symbol table operations summary

Operation	Unordered Linked List	Ordered Array	BST
search	n	lg n	h^{\dagger}
insert	n	n	h
delete	n	n	$\sqrt{n}^{\dagger\dagger}$
min/max	-	1	h
floor/ceiling	-	lg n	h
rank	-	lg n	h
select	-	1	h
ordered iteration	-	n	n

 $[\]dagger$ $\it h$ is the height of BST, proportional to $\lg \it n$ if keys inserted in random order

 $[\]dagger\dagger$ other operations also become \sqrt{n} if deletions are allowed