

Ranking and unranking Dyck paths

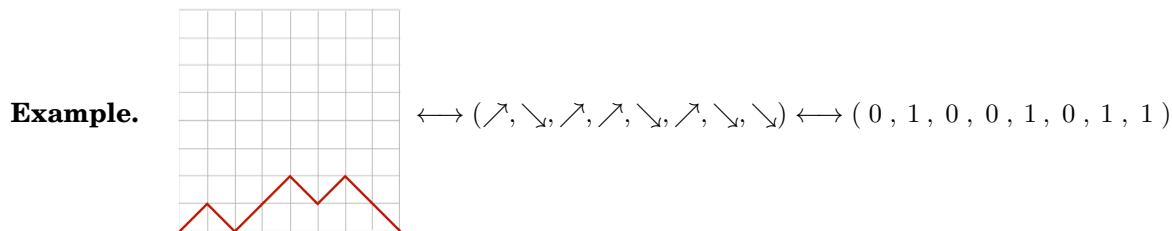
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1 Ranking and unranking Dyck paths

1.1 Dyck paths to words

Recall from Lecture 5 that a Dyck path of length $2n$ is a lattice walk from $(0, 0)$ to $(2n, 0)$ using **up-steps** $\nearrow = (1, 1)$ and **down-steps** $\searrow = (1, -1)$. We can represent a Dyck path as a binary string of length $2n$ by writing \nearrow as 0 and \searrow as 1.



Which binary strings do we get?

Definition. A binary string of length $2n$ is **totally balanced** if

- The string contains n zeros and n ones.
- For $T1 \leq i \leq 2n$, the first i elements of the string include at least as many zeros and ones.

Dyck paths correspond to totally balanced binary strings. The first condition ensures that a Dyck path finishes on the x -axis. The second condition ensures that it never goes strictly below the x -axis.

Viewing Dyck paths as binary strings gives them a natural lexicographic order. In this lecture, we find how to lexicographically rank and unrank Dyck paths (i.e. totally balanced binary words).

Example. When $n = 3$, the five Dyck paths of length 6 correspond to balanced binary words lexicographically ordered as follows. In particular, we have $\text{Rank}(\nearrow, \nearrow, \searrow, \searrow, \nearrow, \searrow) = 3$.

$$(0, 0, 0, 1, 1, 1), (0, 0, 1, 0, 1, 1), (0, 0, 1, 0, 1, 1), (0, 0, 1, 1, 0, 1), (0, 1, 0, 1, 0, 1)$$

We can rank Dyck paths by using the general lexicographic ranking formula from Lecture 8. This formula requires that, for any totally balanced word $W = (w_1, w_2, \dots, w_{2n})$, any index $1 \leq i \leq 2n$ and any entry a with $0 \leq a \leq w_i - 1$, we know to compute the number $P(W; i, a)$ of totally balanced words that lexicographically precede W and begin with the sequence $(w_1, w_2, \dots, w_{i-1}, a)$. The condition $0 \leq a < w_i \leq 1$ implies $a = 0$ and $w_i = 1$, so we can simplify the inner sum of the general ranking formula.

$$\text{Rank}(W) = \sum_{i=1}^{2n} \sum_{a=0}^{w_i-1} P(W; i, a) = \sum_{i=1}^{2n} w_i P(W; i, 0). \quad (1)$$

We turn to the problem of evaluating $P(W; i, 0)$. We must count the Dyck paths p that correspond to a totally balanced word of the form $(w_1, w_2, \dots, w_{i-1}, 0, *, *, \dots, *)$. The first i steps of p will bring us to a lattice point (x, y) where $x = i$. We need to be able to count the paths from (x, y) to $(2n, 0)$ that do not go strictly below the x -axis.

1.2 Counting suffixes of Dyck paths

Definition. Let $\mathcal{D}_{2n}(x, y)$ be the set of paths from (x, y) to $(2n, 0)$ using the steps $(1, 1)$ and $(1, -1)$ and which never go strictly below the x -axis.

Let $d_{2n}(x, y) = |\mathcal{D}_{2n}(x, y)|$.

Such paths are ends (suffixes if you think of them as words) of Dyck paths.

Proposition. Let x, y , and n be integers with $x + y$ even and $x + y \leq 2n$. Then

$$d_{2n}(x, y) = \binom{2n - x}{n - \frac{x+y}{2}} - \binom{2n - x}{n - 1 - \frac{x+y}{2}}$$

Proof. Let $\mathcal{P}_{2n}(x, y)$ be the set of all paths from (x, y) to $(2n, 0)$ that use steps $(1, 1)$ and $(1, -1)$, including those that go below the x -axis. Let \mathcal{B}_{2n} be the set of paths in $\mathcal{P}_{2n}(x, y)$ that go strictly below the x -axis. Then

$$\mathcal{D}_{2n}(x, y) = \mathcal{P}_{2n}(x, y) - \mathcal{B}_{2n}(x, y).$$

We aim to count the paths in $\mathcal{P}_{2n}(x, y)$ and the paths in $\mathcal{B}_{2n}(x, y)$.

We start with $\mathcal{P}_{2n}(x, y)$. Consider a path $p \in \mathcal{P}_{2n}(x, y)$ go from (x, y) to $(2n, 0)$. Let u be the number of up-steps in p , and let d be the number of down-steps in p . Since p goes from (x, y) to $(2n, 0)$ we must have

$$d + u = 2n - x \quad \text{and} \quad d - u = y.$$

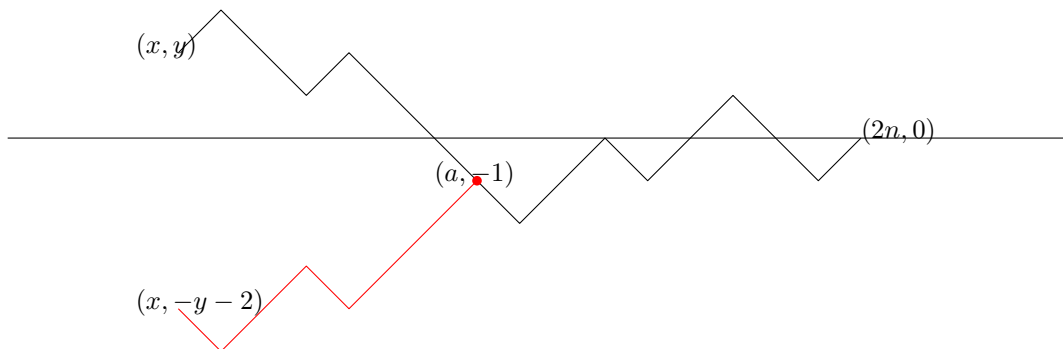
Solving these two equations, we find that

$$u = n - \frac{x + y}{2} \quad \text{and} \quad d = n - \frac{x - y}{2}.$$

Every path in $\mathcal{P}_{2n}(x, y)$ corresponds bijectively to a sequence of length $2n - x$ containing exactly u up-steps (the remaining steps are down-steps). There are $\binom{2n-x}{u}$ ways to choose the locations of the up-steps. Therefore

$$\mathcal{P}_{2n}(x, y) = \binom{2n - x}{u} = \binom{2n - x}{n - \frac{x+y}{2}}. \quad (2)$$

We now count $\mathcal{B}_{2n}(x, y)$. To determine $d_{2n}(x, y)$ we just need to subtract from (2) the number paths which *do* go strictly below the x axis. Let w be such a path, and let $(a, -1)$ be the first point along w which touches the line $y = -1$. We reflect the portion of w before $(a, -1)$ across the line $y = -1$, as illustrated, to obtain a new path w' .



Note that w' goes from $(x, -y - 2)$ to $(2n, 0)$. Note also that the construction $w \mapsto w'$ is reversible. That is, for every path w' from $(x, -y - 2)$ to $(2n, 0)$ there is a unique path w which produces w' in this way. Specifically, there is a unique point where it may have been flipped according to this rule: namely the first place where the path reaches $(a, -1)$ for some a .

Thus the number of paths from (x, y) to $(2n, 0)$ which go strictly below the x -axis is the same as the number of paths from $(x, -y - 2)$ to $(2n, 0)$ with no restrictions. This number is given by (2) with y replaced by $-y - 2$. we leave as a little exercise to show this equals

$$\binom{2n-x}{n-1-\frac{x+y}{2}}$$

Therefore the number of paths from (x, y) to $(2n, 0)$ which never go strictly below the x -axis is

$$\binom{2n-x}{n-\frac{x+y}{2}} - \binom{2n-x}{n-1-\frac{x+y}{2}}$$

□

Corollary. The number of Dyck paths of length $2n$ is the n -th Catalan number $\frac{1}{n+1} \binom{2n}{n}$.

Proof. By its definition, $\mathcal{D}_{2n}(0, 0)$ is precisely the set of Dyck paths of length $2n$, so the number of Dyck paths of length $2n$ is

$$d_{2n}(0, 0) = \binom{2n}{n} - \binom{2n}{n-1} = \binom{2n}{n} - \frac{n}{n+1} \binom{2n}{n} = \left(1 - \frac{n}{n+1}\right) \binom{2n}{n}.$$

□

1.3 Ranking and unranking

The formula for counting suffixes tells us about the rank. Suppose $w = w_1 w_2 \cdots w_{2n}$ is a Dyck path represented as a totally balanced binary string. As discussed in Lecture 8, we step through $i = 1, 2, \dots, 2n$, each time accounting for those predecessors of begin with $w_1 w_2 \cdots w_{i-1} a$, where $a < w_i$. This only happens when $a = 0$ and $w_i = 1$. Suppose that after the first i steps the path is at the point (x_i, y_i) . Then $x_i = i$. If $w_i = 0$, then no accounting is needed. If $w_i = 1$, then the previous step was down, and there is exactly one totally balanced word which begins $w_1 w_2 \cdots w_{i-1} 0$ for every path from $(x_i, y_i + 2)$ to $(2n, 0)$ which does not go below the x -axis. There are exactly $d_{2n}(i, y_i + 2)$ such predecessors (the $+2$ comes because their i th step is an up step, instead of a down step). Thus we have

$$\text{Rank}(w) = \sum_{i=1}^{2n} w_i d_{2n}(i, y_i + 2) = \sum_{i=1}^{2n} w_i \left(\binom{2n-i}{n-\frac{i+y_i+2}{2}} - \binom{2n-i}{n-1-\frac{i+y_i+2}{2}} \right)$$

In the following we keep track of the pair $(x, y) = (i, y_i)$.

Algorithm: RankDyck

```

input: n, w.    w is a totally balanced word of length 2n
y = 0
r = 0
for x from 1 to 2n
  y = y+1
  if w(x) = 1
    r = r + binom( 2n-x, n-(x+y)/2 ) - binom( 2n-x , n-1-(x+y+2)/2 )
    y = y - 2
output: r

```

The following unranking algorithm is an adaptation of algorithm UnrankGeneral from Lecture 8.

Algorithm: UnrankDyck

```

input: n, r.
y = 0
R = 0
for x from 1 to 2n
  w(x) = 0

```

```

y = y+1
P = binom( 2n-x, n-(x+y)/2 ) - binom( 2n-x, n-(x+y)/2 )
if R + P <= r
    w(x) = 1
    y = y - 2
    R = R + P
output: w

```

Here's an example of each algorithm which is adapted from Kreher and Stinson, *Combinatorial Algorithms*, Section 3.4.

Suppose we want to compute the rank of $w = 0010110101$. We calculate as follows.

x	$w(x)$	y	$d_{10}(x, y+2)$	r
1	0	1		0
2	0	2		0
3	1	1	$\binom{7}{2} - \binom{7}{1} = 14$	14
4	0	2		14
5	1	1	$\binom{5}{1} - \binom{5}{0} = 4$	18
6	1	0	$\binom{4}{1} - \binom{4}{0} = 3$	21
7	0	1		21
8	1	0	$\binom{2}{1} - \binom{2}{0} = 4$	22
9	0	1		22
10	1	0	$\binom{0}{-1} - \binom{0}{-2} = 0$	22

So the rank is 22. Notice that steps $x = 1$ and $x = 2n$ can always be omitted.

Now we calculate $\text{Unrank}(22)$

$r = 22$					
x	y	$d_{10}(x, y)$	$= P$	R	$w(x)$
1	1	$\binom{9}{4} - \binom{9}{3} = 42$	0	0	0
2	2	$\binom{8}{3} - \binom{8}{2} = 28$	0	0	0
3	1	$\binom{7}{3} - \binom{7}{2} = 14$	14	1	1
4	2	$\binom{6}{2} - \binom{6}{1} = 9$	14	0	0
5	1	$\binom{5}{1} - \binom{5}{0} = 4$	18	1	1
6	0	$\binom{4}{1} - \binom{4}{0} = 3$	21	1	1
7	1	$\binom{3}{1} - \binom{3}{0} = 2$	21	0	0
8	0	$\binom{2}{0} - \binom{2}{-1} = 1$	22	1	1
9	1	$\binom{1}{0} - \binom{1}{-1} = 1$	22	0	0
10	1	$\binom{1}{-1} - \binom{0}{-2} = 0$	22	1	1

so $\text{Unrank}(22) = 0010110101$ as expected.