

Alen

Last one::  
Assignment 6

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i)  $\hat{H} = \theta_1(Z \otimes X) + \theta_2(X \otimes Z)$

ii) Show that  $(Z \otimes X)(X \otimes Z) = (X \otimes Z)(Z \otimes X)$

$$\begin{aligned}(Z \otimes X)(X \otimes Z) &= ZX \otimes XZ \\&= -XZ \otimes -ZX \\&= (-1)(-1)XZ \otimes ZX \\&= XZ \otimes ZX \\&= (X \otimes Z)(Z \otimes X)\end{aligned}$$

iii)  $e^{A+B} = e^A e^B$  when A & B commute

$$\begin{aligned}U(t) &= e^{-i\hat{H}t} = e^{-i(\theta_1(Z \otimes X) + \theta_2(X \otimes Z))t} \\&= e^{-i\theta_1 t(Z \otimes X) - i\theta_2 t(X \otimes Z)} \\&= e^{-i\theta_1 t(Z \otimes X)} e^{-i\theta_2 t(X \otimes Z)}\end{aligned}$$

$i\omega$  are just constants  
so  $Z \otimes X$  &  $X \otimes Z$  commute

Are  $(Z \otimes X)$  &  $(X \otimes Z)$  normal? ( $AA^\dagger = A^\dagger A$ )

$$\begin{array}{ll}(Z \otimes X)(Z \otimes X)^\dagger & (Z \otimes X)^\dagger (Z \otimes X) \\(Z \otimes X)(Z^\dagger \otimes X^\dagger) & (Z^\dagger \otimes X^\dagger)(Z \otimes X) \\ZZ^\dagger \otimes XX^\dagger & Z^\dagger Z \otimes X^\dagger X \\I \otimes I = I_4 & I \otimes I = I_4\end{array}\checkmark \quad (Z \otimes X) \text{ is normal}$$

$$\begin{array}{ll}(X \otimes Z)(X \otimes Z)^\dagger & (X \otimes Z)^\dagger (X \otimes Z) \\XX^\dagger \otimes ZZ^\dagger & X^\dagger X \otimes Z^\dagger Z \\I \otimes I & I \otimes I\end{array}\checkmark \quad (X \otimes Z) \text{ is normal}$$

Meaning we can use matrix exponentials on  $e^{-i\theta_1 t(Z \otimes X)} e^{-i\theta_2 t(X \otimes Z)}$

$$\Rightarrow P_{ZX} \begin{pmatrix} e^{-i\theta_1 \lambda_{1ZX}} & & \\ & \ddots & \\ & & e^{-i\theta_1 \lambda_{nZX}} \end{pmatrix} P_{ZX}^\dagger P_{XZ} \begin{pmatrix} e^{-i\theta_2 \lambda_{1XZ}} & & \\ & \ddots & \\ & & e^{-i\theta_2 \lambda_{nXZ}} \end{pmatrix} P_{XZ}^\dagger \quad \text{where}$$

$P_{\#}$  has columns that are eigenvectors of #  
 $\lambda$  are eigenvalues of their individual matrices ( $ZX$  or  $XZ$ )

ii) Putting all this together, finding  $P_{XZ}$  &  $P_{ZX}$

$$X \otimes Z = \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix} \otimes \begin{pmatrix} 1 & 0 \\ 0 & -1 \end{pmatrix} = \begin{pmatrix} 0 & 0 & 0 & -1 \\ 0 & 0 & 0 & 0 \\ 1 & 0 & 0 & 0 \\ 0 & -1 & 0 & 0 \end{pmatrix}$$

Using symbolab for my own sanity

$$\lambda_1 = 1, \lambda_2 = -1, \text{ eigenvectors } V_1 = \begin{pmatrix} 1 \\ 0 \\ 1 \\ 0 \end{pmatrix}, V_2 = \begin{pmatrix} -1 \\ 0 \\ 0 \\ 1 \end{pmatrix}, V_3 = \begin{pmatrix} 0 \\ 1 \\ 0 \\ 1 \end{pmatrix}, V_4 = \begin{pmatrix} 0 \\ 1 \\ 0 \\ 0 \end{pmatrix}$$

has  $\lambda_1 = 1 \quad \lambda_2 = -1 \quad \lambda_1 = 1 \quad \lambda_2 = -1$

$$Z \otimes X = \begin{pmatrix} 1 & 0 \\ 0 & -1 \end{pmatrix} \otimes \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix} = \begin{pmatrix} 0 & 1 & 0 & 0 \\ 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & -1 \\ 0 & 0 & -1 & 0 \end{pmatrix}$$

$$\lambda_1 = 1, \lambda_2 = -1, \text{ eigenvectors } V_1 = \begin{pmatrix} 1 \\ 1 \\ 0 \\ 0 \end{pmatrix}, V_2 = \begin{pmatrix} -1 \\ 1 \\ 0 \\ 0 \end{pmatrix}, V_3 = \begin{pmatrix} 0 \\ 0 \\ 1 \\ 1 \end{pmatrix}, V_4 = \begin{pmatrix} 0 \\ 0 \\ 1 \\ 1 \end{pmatrix}$$

has  $\lambda_1 = 1, \lambda_2 = -1, \lambda_1 = 1, \lambda_2 = -1$

Putting this all together,  $P_{ZX} \left( \begin{array}{c|cc} e^{-i\omega_1 t} \lambda_{1ZX} & \dots & P_{ZX} + P_{XZ} \left( \begin{array}{c|cc} e^{-i\omega_2 t} \lambda_{1XZ} & \dots & e^{-i\omega_2 t} \lambda_{2XZ} \\ e^{-i\omega_2 t} \lambda_{2XZ} & \dots & e^{-i\omega_2 t} \lambda_{2XZ} \end{array} \right) \end{array} \right)$

$$= \begin{pmatrix} 1 & -1 & 0 & 0 \\ 1 & 1 & 0 & 0 \\ 0 & 0 & -1 & 1 \\ 0 & 0 & 1 & 1 \end{pmatrix} \begin{pmatrix} e^{-i\omega_1 t} & 0 & 0 & 0 \\ 0 & e^{-i\omega_1 t} & 0 & 0 \\ 0 & 0 & e^{-i\omega_2 t} & 0 \\ 0 & 0 & 0 & e^{i\omega_1 t} \end{pmatrix} \begin{pmatrix} 1 & 1 & 0 & 0 \\ 1 & 1 & 0 & 0 \\ 0 & 0 & -1 & 1 \\ 0 & 0 & 1 & 0 \end{pmatrix} \begin{pmatrix} 1 & -1 & 0 & 0 \\ 1 & 1 & 0 & 0 \\ 0 & 0 & e^{-i\omega_2 t} & 0 \\ 0 & 0 & 0 & e^{i\omega_2 t} \end{pmatrix} \begin{pmatrix} 1 & 0 & 1 & 0 \\ 1 & 0 & 1 & 0 \\ 0 & 1 & 0 & 1 \\ 0 & 1 & 0 & 1 \end{pmatrix}$$

$V_1 V_2 V_3 V_4$

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is our two qubit circuit for  $U(t)$

iii) After approximating to Clifford+T gate set,  $C \approx 1.5$  with error  $\epsilon = 10^{-17}$

$$\Rightarrow \left| \log_2 \left( \frac{1}{10^{-17}} \right) \right|^{1.5} \approx 1 \text{ (very close to 1)}$$

iv)  $(Z \otimes X)|V_i\rangle = \lambda_{aj,i}|V_i\rangle$   
 $(X \otimes Z)|V_i\rangle = \lambda_{bj,i}|V_i\rangle, \langle V_i | V_j \rangle = 0 \forall i \neq j$

$$(Z \otimes X)|V_i\rangle = \lambda_{aj,i}|V_i\rangle$$

$$(X \otimes Z)(Z \otimes X)|V_i\rangle = (X \otimes Z)\lambda_{aj,i}|V_i\rangle$$

$$(X \otimes Z)(Z \otimes X)|V_i\rangle = \lambda_{aj,i}(X \otimes Z)|V_i\rangle$$

$$(X \otimes Z)(Z \otimes X)|V_i\rangle = \lambda_{aj,i}\lambda_{bj,i}|V_i\rangle$$

$$(X \otimes Z)(Z \otimes X) = \begin{vmatrix} 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & -1 \\ 1 & 0 & 0 & 0 \\ 0 & -1 & 0 & 0 \end{vmatrix} \begin{vmatrix} 0 & 1 & 0 & 0 \\ 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & -1 \\ 0 & 0 & -1 & 0 \end{vmatrix} = \begin{vmatrix} 0 & 0 & 0 & -1 \\ 0 & 0 & 1 & 0 \\ 0 & 1 & 0 & 0 \\ -1 & 0 & 0 & 0 \end{vmatrix}$$

eigenvectors of this matrix:  $V_1 = \begin{pmatrix} 0 \\ 1 \\ 1 \\ 0 \end{pmatrix}, V_2 = \begin{pmatrix} 0 \\ -1 \\ 1 \\ 0 \end{pmatrix}, V_3 = \begin{pmatrix} -1 \\ 0 \\ 0 \\ 1 \end{pmatrix}, V_4 = \begin{pmatrix} 0 \\ 0 \\ 0 \\ 1 \end{pmatrix}$

$$\lambda_1 = 1 \quad \lambda_2 = -1 \quad \lambda_3 = 1 \quad \lambda_4 = -1$$

Make them unit: multiply by  $\frac{1}{\sqrt{2}}$ .  $V_i \cdot V_j = 0 \forall i \neq j \checkmark$

Therefore, orthonormal basis for eigenspace:

$$\left\{ \frac{1}{\sqrt{2}} \begin{pmatrix} 0 \\ 1 \\ 1 \\ 0 \end{pmatrix}, \frac{1}{\sqrt{2}} \begin{pmatrix} 0 \\ -1 \\ 1 \\ 0 \end{pmatrix}, \frac{1}{\sqrt{2}} \begin{pmatrix} -1 \\ 0 \\ 0 \\ 1 \end{pmatrix}, \frac{1}{\sqrt{2}} \begin{pmatrix} 0 \\ 0 \\ 0 \\ 1 \end{pmatrix} \right\}$$

V)  $\Lambda_X$  is matrix of eigenvectors of  $X$

$$Z \otimes X = U \Lambda_{ZX} U^+, \quad X \otimes Z = U \Lambda_{XZ} U^+$$

"as  $Z \otimes I$  &  $I \otimes Z$ "

$$Z \otimes I = \begin{pmatrix} 1 & 0 \\ 0 & -1 \end{pmatrix} \otimes \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix} = \begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & -1 & 0 \\ 0 & 0 & 0 & -1 \end{pmatrix}$$

$$I \otimes Z = \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix} \otimes \begin{pmatrix} 1 & 0 \\ 0 & -1 \end{pmatrix} = \begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & -1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & -1 \end{pmatrix}$$

$$Z \otimes X =$$

2) Given 32-bit hash  $H$  of a message  $M$  with two-one protocol

$H = h(M)$  where  $h: \{0,1\}^{32} \rightarrow \{0,1\}^{32}$  &  $h$  has the property that for every  $H = h(M)$   $\exists$  exactly one collision  $M' \neq M$  s.t.  $h(M') = H$

ii) input:  $h: \{0,1\}^{32} \rightarrow \{0,1\}^{32}$  s.t.  $\forall H = h(M) \exists! M' \neq M$  s.t.  $h(M') = H$

goal: find  $M' \in \{0,1\}^{32}$  s.t.  $h(M') = M_j$ ,  $M' \neq M$   
here, our search function  $f = h$

Implementation: Prepare  $\frac{1}{\sqrt{2^{32}}} \sum_{X \in \{0,1\}^{32}} |X\rangle$ . All states have amplitude  $\frac{1}{2^{16}}$

Apply phase oracle  $U_f |X\rangle \rightarrow (-1)^{f(X)} |X\rangle$  and  
Apply  $U_{\text{diff}}$  for amplitude manipulation

Measure when solution state has high prob

iii) Each  $M$  has a pair that has same image under  $h$

Meaning for any  $M \exists 2$  hashes, one of which is  $M_j$  the other is  $M'$

Our search function def "excludes"  $M' \neq M$  meaning we have one sol

$|S\rangle$  will have angle  $\theta$  with  $|\Psi_{\text{bad}}\rangle$ ,  $\sin \theta = \frac{1}{2^{16}}$  (very small meaning,  $\sin \theta \approx \theta$ )

$\rightarrow$  Opt. iterations are  $\frac{\pi}{4\sqrt{2^{32}}} - \frac{1}{2} \approx 51471$

iii) After  $M$  reps, angle with  $|\Psi_{bad}\rangle$  is  $(2M+1)\theta$ ,  $\theta \approx \frac{1}{2^{16}}$  small  
 $\Rightarrow (51471 \cdot 2 + 1) \left( \frac{1}{2^{16}} \right) = 1.570785522\dots$  is angle with  $|\Psi_{bad}\rangle$ . We have  $\frac{\pi}{2}$  angle  
 $\frac{\pi}{2} \approx 1.570796327$

Very high similarity with  $\frac{\pi}{2}$ , approximate similarity:  $\frac{1.570785522}{1.570796327} = 0.9999931213\dots$

Round down to about 99% success

$$3) X \otimes X = \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix} \otimes \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix} = \begin{pmatrix} 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \\ 0 & 1 & 0 & 0 \\ 1 & 0 & 0 & 0 \end{pmatrix}$$

eigenvectors of this matrix  $\begin{pmatrix} 1 \\ 0 \\ 0 \\ 1 \end{pmatrix}, \begin{pmatrix} 0 \\ 1 \\ 1 \\ 0 \end{pmatrix}$  by inspection

$$\begin{pmatrix} 1 \\ 0 \\ 0 \\ 1 \end{pmatrix} \cdot \begin{pmatrix} 0 \\ 1 \\ 1 \\ 0 \end{pmatrix} = 0 \rightarrow \text{orthogonal}$$

$$\begin{pmatrix} 1 \\ 0 \\ 0 \\ 1 \end{pmatrix} = |00\rangle + |11\rangle = |\Psi\rangle \quad \text{with } \alpha = \beta = 1$$

$$\begin{pmatrix} 0 \\ 1 \\ 1 \\ 0 \end{pmatrix} = |01\rangle + |10\rangle = |\Psi\rangle \quad \text{with } \alpha = \beta = 1$$

If bit flips are independent we could end up with a non-eigenvector like  $\begin{pmatrix} 0 \\ 0 \\ 1 \\ 1 \end{pmatrix}$   
meaning  $(X \otimes I)|\Psi\rangle \neq |\Psi\rangle \rightarrow \text{doesn't work}$

4) Quantum channel:  $p \rightarrow p'$

Kraus operators:  $\{K_i\}$  s.t.  $\sum_i K_i^\dagger K_i = I$

The channel then maps  $p \rightarrow \sum_i K_i p K_i^\dagger$ .  $\{\frac{1}{2}I, \frac{1}{2}X, \frac{1}{2}Y, \frac{1}{2}Z\}$

Assume  $p$  is  $2 \times 2$

Depolarizing channel sends  $p \rightarrow \sum_i K_i p K_i^\dagger$

$$\begin{aligned}\sum_i K_i p K_i^\dagger &= \frac{1}{2}I p \frac{1}{2}I + \frac{1}{2}X p \frac{1}{2}X + \frac{1}{2}Y p \frac{1}{2}Y + \frac{1}{2}Z p \frac{1}{2}Z \\ &= \frac{1}{4}I p I + \frac{1}{4}X p X + \frac{1}{4}Y p Y + \frac{1}{4}Z p Z\end{aligned}$$

Let's do this manually. let  $p = \begin{pmatrix} a & b \\ c & d \end{pmatrix}$ .  $\text{Tr}(p) = 1 \rightarrow a+d=1$   
density matrix

$I$  is identity so expression simplifies to

$$\frac{1}{4}p + \frac{1}{4}|0\rangle\langle 0| + \frac{1}{4}|a\rangle\langle a| + \frac{1}{4}|b\rangle\langle b| + \frac{1}{4}|c\rangle\langle c| + \frac{1}{4}|d\rangle\langle d|$$

$$\frac{1}{4}p + \frac{1}{4}|0\rangle\langle 0| + \frac{1}{4}|b\rangle\langle a| + \frac{1}{4}|0\rangle\langle -i| + \frac{1}{4}|i\rangle\langle b| + \frac{1}{4}|i\rangle\langle -a| + \frac{1}{4}|1\rangle\langle 1| + \frac{1}{4}|a\rangle\langle -b| + \frac{1}{4}|b\rangle\langle a| + \frac{1}{4}|c\rangle\langle d| + \frac{1}{4}|c\rangle\langle -i| + \frac{1}{4}|i\rangle\langle c| + \frac{1}{4}|d\rangle\langle -c| + \frac{1}{4}|0\rangle\langle -1| + \frac{1}{4}|0\rangle\langle 1| + \frac{1}{4}|c\rangle\langle -d| + \frac{1}{4}|d\rangle\langle c|$$

$$\frac{1}{4}p + \frac{1}{4}|d\rangle\langle c| + \frac{1}{4}|d\rangle\langle -c| + \frac{1}{4}|a\rangle\langle -b| + \frac{1}{4}|b\rangle\langle a| + \frac{1}{4}|d\rangle\langle a| + \frac{1}{4}|a\rangle\langle d|$$

$$\frac{1}{4}|a\rangle\langle b| + \frac{1}{4}|d\rangle\langle c| + \frac{1}{4}|d\rangle\langle -c| + \frac{1}{4}|a\rangle\langle -b| + \frac{1}{4}|b\rangle\langle a| + \frac{1}{4}|b\rangle\langle -a| + \frac{1}{4}|c\rangle\langle d| + \frac{1}{4}|d\rangle\langle -d|$$

$$\frac{1}{4}(a+d+a-d+b+c-c-b)$$

$$\frac{1}{4}(c-b-b-c+d+a+d)$$

$$\frac{1}{4}\begin{pmatrix} 1 & 1 \\ 0 & 1 \end{pmatrix} + \frac{1}{4}\begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix} = \frac{1}{2}\begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix} = 2I$$

5) Steane code encodes one logical qubit through 7 physical qubits

5i) for  $X \otimes I \otimes I \otimes I \otimes I \otimes I \otimes I$  the matrix will be an off diagonal of 1's, 0 elsewhere meaning it will flip qubits no matter its position

Comparing  $|0\rangle_L \otimes |1\rangle_L$ , their qubits are just flipped meaning  $X|0\rangle_L = |1\rangle_L \Rightarrow X_L$  acts like  $X \otimes X \otimes X \otimes X \otimes X \otimes X \otimes X$  as it flips each qubit

for  $Z \otimes Z \otimes Z \otimes Z \otimes Z \otimes Z \otimes Z$  the matrix will be a diagonal of  $1, -1, 1, -1, \dots$  meaning it will flip every  $|1\rangle$  qubit, leave  $|0\rangle$  qubit alone

Comparing  $|0\rangle_L \otimes |1\rangle_L$ , each state in  $|0\rangle_L$  has even # of 1's each state in  $|1\rangle_L$  has odd # of 1's

meaning  $Z \otimes Z \otimes Z \otimes Z \otimes Z \otimes Z \otimes Z$  will make  $|0\rangle_L \rightarrow |0\rangle_L$  (even amount of -1's)  $|1\rangle_L \rightarrow -|1\rangle_L$  (odd amount of -1's multiplied)

$$\rightarrow Z_L |0\rangle_L \rightarrow |0\rangle_L, Z_L |1\rangle \rightarrow -|1\rangle_L$$

$\Rightarrow Z_L$  acts like  $Z \otimes Z \otimes Z \otimes Z \otimes Z \otimes Z \otimes Z$  as it negates  $|1\rangle$  and leaves  $|0\rangle$

for  $S \otimes S \otimes S$ , the matrix will leave  $|0\rangle$  alone  $|1\rangle \rightarrow i|1\rangle$

Comparing  $|0\rangle_L \otimes |1\rangle_L$ , each state in  $|0\rangle_L$  has 4 1's (expect 1000000) each state in  $|1\rangle_L$  has 3 1's (expect 1111111)

meaning  $S \otimes S \otimes S$  will make  $|0\rangle_L \rightarrow |0\rangle_L$  as  $i^4 = 1$   
 $|1\rangle_L \rightarrow -i|1\rangle_L$  as  $i^3 = -i$

meaning for  $S^t = \begin{pmatrix} 1 & 0 \\ 0 & -i \end{pmatrix}$ ,  $S^t |0\rangle = |0\rangle, S^t |1\rangle = -i|1\rangle$

$\Rightarrow S_L$  acts like  $S \otimes S \otimes S \otimes S \otimes S \otimes S$  as it leaves  $|0\rangle$  & adds factor of  $-i$  to  $|1\rangle$

iii) for  $T_{0..0T}$ , the matrix will leave  $|0\rangle$  alone,  $|1\rangle \rightarrow e^{\frac{i\pi}{4}}|1\rangle$

$T_{0..0T}$  will make  $|0\rangle_L \rightarrow -|0\rangle$  as  $(e^{\frac{i\pi}{4}})^4 = -1$   
 $|1\rangle_L \rightarrow e^{\frac{3i\pi}{4}}|1\rangle_L$  as  $(e^{\frac{i\pi}{4}})^3 = e^{\frac{3i\pi}{4}}$

This is not equal to  $T_L$  as  $T_{0..0T}|0\rangle_L \rightarrow -|0\rangle_L$   
 $\neq T_L|0\rangle_L \rightarrow |0\rangle_L$

Different mapping means it is not traversal in Steane code  
as  $T_L \neq T_{0..0T}$

iii) let  $|4\rangle = \alpha|0\rangle + \beta|1\rangle$

Circuit 1

$$T|4\rangle = \alpha|0\rangle + \beta e^{\frac{i\pi}{4}}|1\rangle$$

Circuit 2

$$|4\rangle|A\rangle = (\alpha|0\rangle + \beta|1\rangle)\otimes \frac{1}{\sqrt{2}}(|0\rangle + e^{\frac{i\pi}{4}}|1\rangle)$$

$$= \frac{1}{\sqrt{2}}(\alpha|100\rangle + \beta|110\rangle + \alpha e^{\frac{i\pi}{4}}|101\rangle + \beta e^{\frac{i\pi}{4}}|111\rangle)$$

$$(CNOT|4\rangle|A\rangle) = \frac{1}{\sqrt{2}}(\alpha|100\rangle + \beta|111\rangle + \alpha e^{\frac{i\pi}{4}}|101\rangle + \beta e^{\frac{i\pi}{4}}|110\rangle)$$

measure second qubit  $a=0 \rightarrow \frac{1}{\sqrt{2}}(\alpha|10\rangle + \beta e^{\frac{i\pi}{4}}|11\rangle)$

$$\text{apply } S^0 = I \rightarrow \frac{1}{\sqrt{2}}(\alpha|10\rangle + \beta e^{\frac{i\pi}{4}}|11\rangle) = |4\rangle$$

$$a=1 \rightarrow \frac{1}{\sqrt{2}}(\beta|11\rangle + \alpha e^{\frac{i\pi}{4}}|10\rangle)$$

$$\text{apply } S^1 \rightarrow \frac{1}{\sqrt{2}}(\beta|11\rangle + \alpha e^{\frac{i\pi}{4}}|10\rangle)$$

$$\frac{1}{\sqrt{2}}(e^{\frac{i\pi}{4}}(\beta e^{\frac{i\pi}{4}}|11\rangle + \alpha|10\rangle)) = \frac{\theta}{\sqrt{2}}|4\rangle$$