The comonadic smell of spreadsheets

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D2	▼ fx =B2	*C2			
	A	В	С	D	E
1	Fruit	Weight (kg)	Price / kg	Total price	% price
2	Apples	1.5	1.75	2.625	0.4738267148
3	Bananas	0.7	1.29	0.903	0.1629963899
4	Tangerines	0.6	2.5	1.5	0.2707581227
5					
6					
7					
8		Total	5.54		
9					
10					

Figure 1: Example of a 2D spreadsheet

H8	▼ fx					
	А	В	С	D	E	F
1	Fruit	Weight (kg)	Price / kg	Total price	% price	
2	Apples	1.5	1.75	=B2*C2	=D2/\$C\$8	
3	Bananas	0.7	1.29	=B3*C3	=D3/\$C\$8	
4	Tangerines	0.6	2.5	=B4*C4	=D4/\$C\$8	
5						
6						
7						
8		Total	=SUM(\$C\$2:\$C			
9						
10						

Figure 2: Formulas

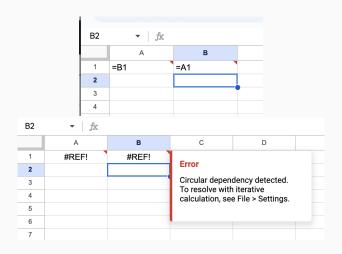


Figure 3: Circular dependencies

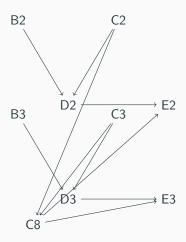


Figure 4: DAG

Associated to a spreadsheet is a DAG where

- nodes are cells
- there is a directed edge from a to b if and only if b has a formula that depends on a (i. e., b depends on a)

Spreadsheet $\it evaluation$ is analogous to $\it dependency \it resolution$

A modern spreadsheet application provides many additional capabilities, not in scope in this talk, such as

- Parsing of cell formulas
- Built-in functions
- Graphics (e.g., histograms, piecharts, etc)
- Formatting, exporting, drag-and-drop, copy-pasting, etc
- Programmable functions (e.g., in VBA)
- Data connectors
- Collaborative online spreadsheets
- And many more

					_	
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Figure 5: Recalculation and support graph



POST MVP: incremental recalculation

Traditional technique:

• Topological sorting of the support graph¹ (e.g., Kahn's algorithm)

 $^{^{1} \}verb|https://en.wikipedia.org/wiki/Topological_sorting|$

Production applications need to take into account more complicated challenges:²

- Performance tradeoffs (e.g., calculating the topological order also has an associated cost)
- Resource utilisation (e.g., compact representations)
- Volatile functions (e.g., now(), rand())
- DAG caching
- Parallelisation algorithms, etc

²https://learn.microsoft.com/en-us/office/client-developer/excel/ excel-recalculation

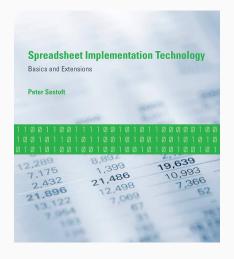


Figure 6: P. Sestoft, *Spreadsheet Implementation Technology*, MIT press (2014).⁴

⁴https://direct.mit.edu/books/book/3071/ Spreadsheet-Implementation-TechnologyBasics-and

Observation

Spreadsheet-like evaluation can be expressed as a *fixed-point of higher dimensional comonads*

Agenda

- 1. A taste of comonads
- 2. Comonadic vibes meet spreadsheets
- 3. Down the rabbit hole

A taste of comonads

monads are burritos?

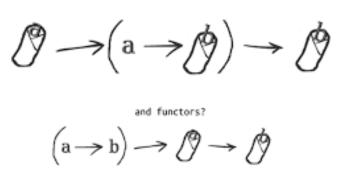


Figure 7: Monads are burritos

Definition

Comonads are co-burritos

Functor hierarchy:

```
class Functor f where
  fmap :: (a -> b) -> f a -> f b
together with some laws.
```

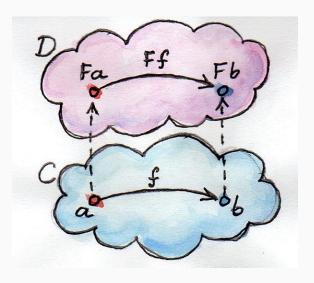


Figure 8: Functors: *Lifting* of 1-ary pure functions to the *f*-world

```
f1 :: [Int] -> [String]
-- vs

f2 :: Int -> String

f3 :: [Int] -> [String]
f3 = fmap f2
```

```
class (Functor f) => Applicative f where
  pure :: Applicative f => a -> f a
  ap :: Applicative f => f (a ->b) -> f a -> f b

together with laws governing properties of fmap, ap, and pure.
Applicatives provide lifting of N-ary pure functions in the f-context
```

```
class (Applicative f) => Monad f where
  return :: Monad f => a -> f a
  join :: Monad f => f (f a) -> f a

-- flatMap
bind :: Monad f => m a -> (a -> m b) -> m b
```

In short:

- ullet Functor \sim map
- Applicative ~ mapN (or equivalent)
- Monad ~ flatMap (or equivalent)

(and the caveat that the implementation must satisfies certain "reasonable" laws that rule out "unexpected" behaviours)

This seems very abstract but from a software engineer perspective lifting provides

- Separation of concerns
- Reusability
- Composability

flatMap ensures no "callback hell" of nested effects
f (f (f (... (f a)))) = f a

Duality:

```
class Functor f => Comonad f where
  coreturn :: f a -> a
  cojoin :: f a -> f ( f a )

(and, as usual, there are important laws)
```

```
class Functor w => Comonad w where
  extract :: f a -> a
  duplicate :: w a -> w ( w a )

-- coKleisli composition / extract
coFlatMap :: Comonad w => (w a -> b) -> w a -> w b
coFlatMap f = fmap f . duplicate

(and, as usual, there are important laws)
```

Monadic values are typically produced in effectful computations
 a -> m b

• Comonadic values are typically *consumed* in context-sensitive computations (e. g., "queries")

w a -> b

Maybe and [] are not comonads (why?)
class Copointed f where
 extract :: f a -> a

instance Comonad NonEmpty where
 extract = head
 duplicate = tails1

```
{-# LANGUAGE OverloadedLists #-}
-- return all suffices including itself
tails [1,2,3] 'shouldBe'
        [[1,2,3], [2, 3], [3]]
```

Simple moving average of a time series $[\phi_1,\ldots,\phi_N]$ with window size k at the time t is

$$SMA_k(t) = \frac{1}{k} \sum_{j=t-k+1}^t \phi_j \tag{1}$$

```
{-# LANGUAGE OverloadedLists #-}
sma :: Window -> NonEmpty Int -> [Double]
sma 3 [1,2,3,4,5,6,7,8,9,10] 'shouldBe'
[2.0,3.0,4.0,5.0,6.0,7.0,8.0,9.0]
```

because

- (1+2+3)/3 = 6/3 = 2
- (2+3+4)/3 = 9/3 = 3
- (3+4+5)/3 = 12/3 = 4

```
smaLocal :: Window -> NonEmpty Double
   -> Maybe Double
smaLocal k s = (/ (fromIntegral k)) <$>
        sumK k s
sumFirstK :: Num a =>
   Window -> NonEmpty a -> Maybe a
sumFirstK (Window k) _ | k <= 0 = Nothing</pre>
sumFirstK (Window 1) (MyNonEmpty a _) = Just a
sumFirstK k (MyNonEmpty a as) =
   case nonEmpty as of
     Just as 2 \rightarrow (+a) <  sumFirstK (k-1) as
     Nothing -> Nothing
```

Folkloristic understanding of some comonads usages

- the function passed to coFlatMap represents a local computation (can explore the "neighbourhood" of the focus point, but produces a single value)
- cojoin produces a (lazy) view of the structure from all perspectives
- coFlatMap applies the local computation to all perspectives

Similar examples:

- Numerical derivation
- Reconciliation in a list of events
- etc

Interesting usages:

- Cellular automaton (e.g., Conway's game of life)
- Store, Moore machine and comonadic UIs

Comonadic vibes meet

spreadsheets

```
data Stream a = Stream a (Stream a)
 deriving (Eq, Show, Functor)
iterate :: (a -> a) -> a
 -> Stream a
iterate f seed =
  Stream seed (iterate f (f seed))
instance Comonad Stream where
   extract (Stream a _) = a
   duplicate s'@(Stream _ as) =
     Stream s' (duplicate as)
```

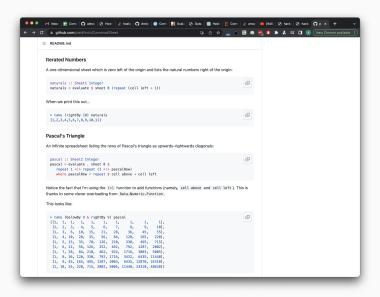
```
iterate :: (a -> a) -> a
   -> Stream a
iterate f seed =
   Stream seed (iterate f (f seed))
```

```
data Sheet1 = Sheet1 (Stream a) a (Stream a)
  deriving (Eq,Show, Functor)
```

```
focus :: Sheet1 a -> a
focus (Sheet1 _ a _ ) = a
moveL :: Sheet1 a -> Sheet1 a
moveL (Sheet1 (Stream a as) focus s)
   = Sheet1 as a (Stream focus s)
moveR :: Sheet1 a -> Sheet1 a
moveR (Sheet1 s focus (Stream a as))
   = Sheet1 (Stream focus s) a as
```

```
instance Comonad Sheet1 where
  extract = focus
duplicate s =
    Sheet1 (allLeft s) s (allRight s) where
    allLeft = iterate moveL
    allRight = iterate moveR
```

Down the rabbit hole



About comonads:

- https://bartoszmilewski.com/2017/01/02/comonads/
- https://blog.higher-order.com/blog/2015/10/04/ scala-comonad-tutorial-part-2/
- https:
 //reasonablypolymorphic.com/blog/cofree-comonads/
- https://reasonablypolymorphic.com/blog/ comonadic-physics/index.html

About comonadic UIs:

- https:
 - //functorial.com/the-future-is-comonadic/main.pdf
- https://arthurxavierx.github.io/ComonadsForUIs.pdf

From Löb's Theorem in modal logic⁵

$$\Box (\Box \phi \to \phi) \to \Box \phi \tag{2}$$

to spreadsheet evaluation:

- Original post from Piponi (2006): http://blog.sigfpe.com/ 2006/11/from-l-theorem-to-spreadsheet.html
- Functional pearl paper from Kenneth Foner (2015):
 https://dl.acm.org/doi/10.1145/2887747.2804310
- ComonadSheet hackage package (unmaintained):
 https://hackage.haskell.org/package/ComonadSheet

⁵In this context, $\Box \phi$ is a formal provability predicate which reads " ϕ is provable"

