Having fun with Kotlin coroutines

A first tour of concurrency models in Kotlin

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Agenda

- 1. We live in a concurrent world
- 2. Blocking vs non-blocking
- 3. Demystifying coroutines
- 4. Coroutines-powered concurrency models

We live in a concurrent world

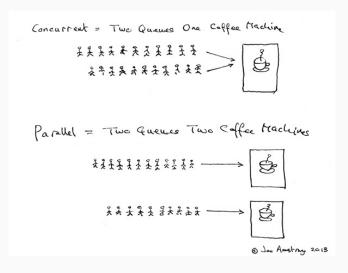


Figure 1: https://joearms.github.io/published/
2013-04-05-concurrent-and-parallel-programming.html

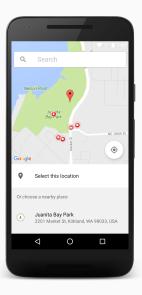
- IO (e.g., network, etc)
- sensors (e.g., gps, etc)
- UI events
- platform lifecycle

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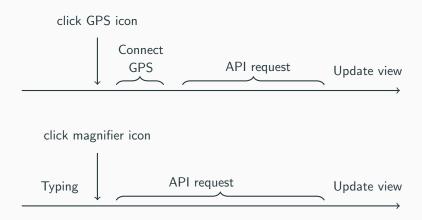
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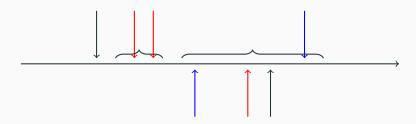


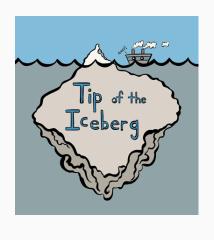
acceptance criteria:

- search by current location
- search by location name advanced
 - search suggestions when tying

Translate ACs into code: simple sequential state machine (simplified)







- Delays
- User inputs
- Failures
- Ordering of api response
- retry logic
- debouncing
- timeouts
- caching
- android lifecycle, etc

First approach: put constrains in place to restrict the range of possible options

- Conditionally forbid user events (disable buttons, loading spinners, etc)
- Boolean flags
- Be defensive (if/else)
- Bind/unbind from lifecycle, etc

(Or more technical constrains like single thread executors, queues, synchronization, etc)

"Concurrency is the composition of independently executing processes, typically functions, but they don't have to be." "Parallelism is the simultaneous execution of multiple things, possibly related, possibly not."

Rob Pike



Rob Pike - 'Concurrency Is Not Parallelism'

Figure 2: https://www.youtube.com/watch?v=cN_DpYBzKso&t=1061s

grams, three basic constructs have received widespread recognition and use: A repetitive construct (e.g. the while loop), an alternative construct (e.g. the conditional if then else), and normal sequential program composi-S. L. Graham, R. L. Rivest tion (often denoted by a semicolon). Less agreement has been reached about the design of other important program structures, and many suggestions have been made: Communicating Subroutines (Fortran), procedures (Algol 60 [15]), entries Sequential Processes (PL/I), coroutines (UNIX II7I), classes (SIMULA 67 ISI), processes and monitors (Concurrent Pascal [2]), clusters (CLU [13]), forms (ALPHARD [19]), actors (Hewitt [1]). C.A.R. Hoare The traditional stored program digital computer has The Queen's University been designed primarily for deterministic execution of a Belfast, Northern Ireland single sequential program. Where the desire for greater speed has led to the introduction of parallelism, every attempt has been made to disguise this fact from the This paper suggests that input and output are basic programmer, either by hardware itself (as in the multiple function units of the CDC 6600) or by the software (as primitives of programming and that parallel composition of communicating sequential processes is a in an I/O control package, or a multiprogrammed opfundamental program structuring method. When erating system). However, developments of processor combined with a development of Dijkstra's guarded technology suggest that a multiprocessor machine, concommand, these concepts are surprisingly versatile. structed from a number of similar self-contained proc-Their use is illustrated by sample solutions of a variety essors (each with its own store), may become more of familiar programming exercises. powerful, capacious, reliable, and economical than a Key Words and Phrases: programming machine which is disguised as a monoprocessor. programming languages, programming primitives. In order to use such a machine effectively on a single

1. Introduction

Among the primitive concept or company programs, and of the high religious in which program are expressed, the action of aniqueness is similar and expressed, the action of aniqueness is dismitted and expressed aniqueness of a measure of a law system good the internal aniqueness of a new value to come variable part of the aniqueness of a new value to come variable part of the aniqueness of a new value to come variable part of the aniqueness of a new value to come variable part of the aniqueness of a new value of the aniqueness of the aniqueness and the aniqueness of the aniqueness o

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This research was supported by a Senior Fellowship of the Science

program structures, parallel programming, concurrency,

coroutines, procedures, multiple entries, multiple exits,

input, output, guarded commands, nondeterminacy,

classes, data representations, recursion, conditional

critical regions, monitors, iterative arrays

CR Categories: 4.20, 4.22, 4.32

Research Council. (1) Simple forms of Author's present address: Programming Research Onough, 45, Ban-thery Road, Oxford, England. (1978 ACM 0001-4792/78,7000-6664 \$00.75 concurrent processes.

events (PL/I), conditional critical regions [10], monitors and queues (Concurrent Pascal [2]), and path expressions [3]. Most of these are demonstrably adequate for their purpose, but there is no widely recognized criterion for choosing between them. This paper makes an ambitious attempt to find a single simple solution to all these problems. The essential proposals are:

task, the component processors must be able to com-

municate and to synchronize with each other. Many

methods of achieving this have been proposed. A widely

adopted method of communication is by inspection and

updating of a common store (as in Algol 68 [18], PL/I,

and many machine codes). However, this can create severe problems in the construction of correct programs and it may lead to expense (e.g., crossbar switches) and

unreliability (e.g. glitches) in some technologies of hardware implementation. A greater variety of methods has

been recogned for synchronization; semanhores [6].

proposals are: (1) Dijkstra's guarded commands [8] are adopted (with a slight change of notation) as sequential control structures, and as the sole means of introducing and controlline nondeterminism.

(2) A parallel command, based on Dijistra's parhegin [6], specifies concurrent execution of its constituent sequernial commands (peccesses). All the peccesses start simultaneously, and the parallel command ends only when they are all finished. They may not communicate with each other by updating global variables. (3) Simple forms of input and output command are introduced. They are used for communication between

Communications of the ACM

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Figure 3: Tony Hoare's seminal paper

"The most obvious application of the new ideas is to the specification, design, and implementation of computer systems which continuously act and interact with their environment. The basic idea is that these systems can be readily decomposed into subsystems which operate concurrently and interact with each other as well as with their common environment. The parallel composition of subsystems is as simple as the sequential composition of lines or statements in a conventional programming language."

Tony Hoare (CSP book, 2015)

Java definition of concurrency and Leslie Lamport seminal paper

Traditional picture of concurrency vs parallelism: coffe machine

Better reppresentation: driving in the desert vs driving in London



The deep problem: communication

Communicating, orchestrating independent processing. Software development as a *dialog* between parts.

We will just sketch the surface of this deeper problem by focusing only on a smaller technical problem in this talk: *blocking* and *non wasting* unnecessary resources when waiting

Blocking vs non-blocking

Test

Demystifying coroutines

What rae

Coroutines-powered concurrency

models

- CSP (aka, channels)
- actors

