# Structured concurrency with Kotlin coroutines

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## **Agenda**

- 1. Warm up
- 2. Before structured concurrency
- 3. Welcome structured concurrency
- 4. Exceptions

Warm up

```
interface Contract { // mvp contract
   interface View {
    interface Presenter {
```

```
interface Contract { // higher level contract
    interface View {
        fun render(state : ViewState)
    }
    interface Presenter {
        fun perform(action : ViewAction)
    }
    sealed class ViewAction
    sealed class ViewState
}
```

```
interface Contract { // more granular contract
    interface View {
        fun showLoading()
        fun hideLoading()
        fun showError(error: CharSequence)
        fun showResults(items: List<Item>)
    }
    interface Presenter {
        fun onRefresh()
```

```
// usually
class SomeActivity : Activity(), Contract.View
class SomeFragment: Fragment(), Contract.View
class SomeCustomView : ViewGroup(), Contract.View
```

```
class Presenter : Contract.Presenter {
    override fun onRefresh() = TODO()
}
```

- Two-way bindings
- Passive view

- View instance holds a reference to its presenter
- Presenter instance holds a reference to the associated view instance<sup>1</sup>

 $<sup>{}^{1}\</sup>mathsf{Circular}\ \mathsf{dependencies};\ \mathsf{see}\ \mathsf{https://www.martinfowler.com/eaaDev/uiArchs.html}$ 

```
class Presenter : Contract.Presenter {
    override fun onRefresh() {
       view.showLoading() // <-- view?
    }
}</pre>
```

```
interface Contract {
    interface View /* ... */
    interface Presenter {
        fun bind(view : View) // <---</pre>
        fun unbind()
        fun onRefresh()
```

```
class Presenter : Contract.Presenter {
    private var view : Contract.View? = null
    override fun bind(view: Contract.View) {
        this.view = view
    }
    override fun unbind() {
        this.view = null
    override fun onRefresh() {
        view?.showLoading()
   }
```

java.lang.IllegalStateException(s)

override fun onSaveInstanceState(outState: Bundle)

```
class SomeFragment : Fragment(), Contract.View {
  @Inject
  lateinit var presenter: Contract.Presenter
  override fun onViewCreated(view: View,
          savedInstanceState: Bundle?) {
     // ...
     presenter.bind(this)
     presenter.onRefresh()
  override fun onDestroyView() {
     presenter.unbind()
     super.onDestroyView()
```

```
// why not this?
class SomeFragment : Fragment(), Contract.View {
    // ...

    override fun onSaveInstanceState(outState: Bundle)
        presenter.unbind()
        super.onSaveInstanceState(outState)
    }
}
```

```
// why not this?
class Presenter(view : Contract.View) :
        Contract Presenter {
  private var isAttached : Boolean = false
  override fun bind() {isAttached = true}
  override fun unbind() {isAttached = false}
  override fun onRefresh() {
    if(isAttached) {
      view.showLoading()
```

```
// what about memory leaks ?
class Presenter(view : Contract.View) :
        Contract.Presenter {
 private val executor : Executor = /*...*/
  override fun onRefresh() {
    view.showLoading()
    executor.execute {
        // long-running operation
        updateUI()
  private fun updateUI() {
  view.hideLoading()
```

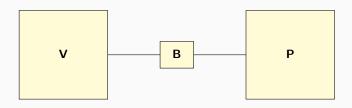
```
class Presenter(view : Contract.View) :
        Contract.Presenter {
 private val executor : Executor = /*...*/
  override fun onRefresh() {
    view.showLoading()
    executor.execute {
        // long-running operation
        updateUI()
  private fun updateUI() {
   // do nothing!
```

Can we handle this wiring somewhere else?

```
class SomeFragment : Fragment(), Contract.View {
    // ...

    override fun showLoading() {
        if (isAdded() ) {
            // update the android view
        }
    }
}
```

```
class SomeFragment : Fragment(), Contract.View {
 // ...
  override fun showA() = if(isAdded()){/* ... */}
  override fun hideA() = if(isAdded()){/* ... */}
  override fun showB() = if(isAdded()){/* ... */}
  override fun hideB() = if(isAdded()){/* ... */}
  override fun showC() = if(isAdded()){/* ... */}
  override fun hideC() = if(isAdded()){/* ... */}
```



### Options are:

- Presenter
- View
- "Binder" between view and presenter

### Pros:

Testability

Cons:2

- Verbosity
- Noise in the contracts

 $<sup>^2 \</sup>mbox{Base class/type class can help though}$ 

# Moving things out of the UI thread

Who is responsible?

- View?
- Binder (if any)?
- Presenter?
- Interactors/usecases?
- Repository/service layer?

```
sealed class Model
interface UseCase {
    suspend fun fetch(): Model
}
```

```
class Presenter(private val useCase: UseCase) :
      Contract.Presenter {
// ...
  override fun onRefresh() { // deprecated way
    launch(Dispatchers.IO) {
      val model = useCase.fetch()
      launch(Dispatchers.Main) {
        view?.showResults(model)
```

```
class Presenter(private val useCase: UseCase) :
      Contract.Presenter {
// ...
  override fun onRefresh() { // more on this later!
    GlobalScope.launch(Dispatchers.IO) {
      val model = useCase.fetch()
      GlobalScope.launch(Dispatchers.Main) {
        view?.showResults(model)
```

```
class Presenter(
  private val useCase: UseCase
  private val ui: CoroutineDispatcher,
  private val io: CoroutineDispatcher
) : Contract.Presenter {
 // ...
  override fun onRefresh() {
    launch(io) {
      val model = useCase.fetch()
      launch(ui) {
        view?.showResults(model)
```

### What's the difference between

```
launch(ui) {
      /*...*/
      launch(io) {
       /*...*/
and
    launch(io) {
      /*...*/
      launch(ui) {
        /*...*/
```

```
class Presenter(
  private val useCase: UseCase
  private val ui: CoroutineDispatcher,
  private val io: CoroutineDispatcher
) : Contract.Presenter {
 // ...
  override fun onRefresh() {
    launch(ui) {
      val model = withContext(io){
         useCase.fetch()
      view?.showResults(model)
```

```
class Presenter(
  private val useCase: UseCase
  private val ui: CoroutineDispatcher,
  private val io: CoroutineDispatcher
) : Contract.Presenter {
 // ...
  override fun onRefresh() {
    launch(ui) {
      try {
        val model = /*...*/
        view?.showResults(model)
      } catch(e : Exception) { // more on this later!
        view?.showError(/*...*/)
```

#### What's the difference between

```
fun onRefresh() {
  launch(dispatcher) {
      /*...*/
and
suspend fun onRefresh() {
  launch(dispatcher) {
      /*...*/
```

```
suspend fun f(a : A) : B

suspend fun g(b : B) : C

suspend fun h(a : A) : C = g(f(a)) // no concurrency
```

```
suspend fun f(a : A, callback : Callback < A >)
suspend fun g(b : B, callback : Callback < C >)
interface Callback < T > {
  fun onSuccess(t : T)
  fun onError(throwable : Throwable)
}
```

```
fun h(a:A, callback: Callback<C>) {
  f(a, object : Callback < B > {
       override fun onSuccess(t: B) {
         g(t,object : Callback < C > {
            override fun onSuccess(t: C) {
               callback.onSuccess(t)
            }
            override fun onError(throwable: Throwable
               callback.onError(throwable)
         })
       override fun onError(throwable: Throwable) {
         callback.onError(throwable)
 })
```

```
class Presenter(dispatcher : CoroutineDispatcher) {
  fun doA() {
     launch(dispatcher) { /*...*/ }
  }
  fun doB() {
     launch(dispatcher) { /*...*/ }
  fun doC() {
     launch(dispatcher) { /*...*/ }
```

### **Concurrency**

Two operations are concurrent if they are not ordered by *happens* before relation.<sup>3</sup>.

https://www.microsoft.com/en-us/research/uploads/prod/2016/12/ Time-Clocks-and-the-Ordering-of-Events-in-a-Distributed-System.pdf

<sup>&</sup>lt;sup>3</sup>Leslie Lamport's paper

"Concurrency is the composition of independently executing processes, typically functions, but they don't have to be."

"Parallelism is the simultaneous execution of multiple things, possibly related, possibly not."

Rob Pike



Rob Pike - 'Concurrency Is Not Parallelism'

Figure 1: https://www.youtube.com/watch?v=cN\_DpYBzKso&t=1061s

grams, three basic constructs have received widespread recognition and use: A repetitive construct (e.g. the while loop), an alternative construct (e.g. the conditional if then else), and normal sequential program composi-S. L. Graham, R. L. Rivest tion (often denoted by a semicolon). Less agreement has been reached about the design of other important program structures, and many suggestions have been made: Communicating Subroutines (Fortran), procedures (Algol 60 [15]), entries Sequential Processes (PL/I), coroutines (UNIX II7I), classes (SIMULA 67 ISI), processes and monitors (Concurrent Pascal [2]), clusters (CLU [13]), forms (ALPHARD [19]), actors (Hewitt [1]). C.A.R. Hoare The traditional stored program digital computer has The Queen's University been designed primarily for deterministic execution of a Belfast, Northern Ireland single sequential program. Where the desire for greater speed has led to the introduction of parallelism, every attempt has been made to disguise this fact from the This paper suggests that input and output are basic programmer, either by hardware itself (as in the multiple function units of the CDC 6600) or by the software (as primitives of programming and that parallel composition of communicating sequential processes is a in an I/O control package, or a multiprogrammed opfundamental program structuring method. When erating system). However, developments of processor technology suggest that a multiprocessor machine, constructed from a number of similar self-contained proc-

> In order to use such a machine effectively on a single task, the component processors must be able to communicate and to synchronize with each other. Many methods of achieving this have been proposed. A widely adopted method of communication is by inspection and updating of a common store (as in Algol 68 [18], PL/I, and many machine codes). However, this can create severe problems in the construction of correct programs and it may lead to expense (e.g. crossbar switches) and unreliability (e.g. glitches) in some technologies of hardware implementation. A greater variety of methods has been recogned for synchronization; semanhores [6]. events (PL/I), conditional critical regions [10], monitors and queues (Concurrent Pascal [2]), and path expressions [3]. Most of these are demonstrably adequate for their

essors (each with its own store), may become more

powerful, capacious, reliable, and economical than a

machine which is disguised as a monoprocessor.

purpose, but there is no widely recognized criterion for choosing between them. This paper makes an ambitious attempt to find a single simple solution to all these problems. The essential (1) Dijkstra's guarded commands [8] are adopted (with a slight change of notation) as sequential control struc-

tures, and as the sole means of introducing and controlline nondeterminism (2) A parallel command, based on Dijkstra's authoris (6), specifies concurrent execution of its constituent sequential commands (processes). All the processes start simultaneously, and the parallel command ends only when they are all finished. They may not communicate with each other by updating global variables

(3) Simple forms of input and output command are introduced. They are used for communication between concurrent processes.

Communications the ACM

combined with a development of Dijkstra's guarded command, these concepts are surprisingly versatile.

Their use is illustrated by sample solutions of a variety of familiar programming exercises. Key Words and Phrases: programming programming languages, programming primitives.

program structures, parallel programming, concurrency, input, output, guarded commands, nondeterminacy, coroutines, procedures, multiple entries, multiple exits, classes, data representations, recursion, conditional critical regions, monitors, iterative arrays CR Categories: 4.20, 4.22, 4.32

### 1. Introduction

Among the primitive concepts of computer programming, and of the high level languages in which programs are expressed, the action of assignment is familiar and well understood. In fact, any change of the internal state of a machine executing a program can be modeled as an assignment of a new value to some variable part of that machine. However, the operations of input and output, which affect the external environment of a machine, are not nearly so well understood. They are often added to a programming language only as an afterthought. Among the structuring methods for computer pro-General permission to make fair use in teaching or research of all or part of this material is granted to individual readers and to nonpools libraries acting for them provided that ACM's copyright notice is given to the fact that reprinting privileges were granted by permission of the Association for Computing Machinery. To otherwise reprint a figure, table, other substantial except, or the entire work requires specific

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Figure 2: Tony Hoare's seminal paper

"The most obvious application of the new ideas is to the specification, design, and implementation of computer systems which continuously act and interact with their environment. The basic idea is that these systems can be readily decomposed into subsystems which operate concurrently and interact with each other as well as with their common environment. The parallel composition of subsystems is as simple as the sequential composition of lines or statements in a conventional programming language."

Tony Hoare (CSP book, 2015)

```
class Test {
  @Mock
  lateinit var useCase : UseCase
  lateinit var presenter : Contract.Presenter
  @Before
  fun setup() {
      presenter = Presenter(useCase,
              Dispatchers. Unconfined,
              Dispatchers. Unconfined
```

```
interface Contract {
    interface Presenter {
        suspend fun onRefresh()
        fun bind(view: View)
        fun unbind()
```

```
class SomeFragment : Fragment(), Contract.View {
 // ...
  override fun showA() = launch(ui){/* ... */}
  override fun hideA() = launch(ui){/* ... */}
  override fun showB() = launch(ui){/* ... */}
  override fun hideB() = launch(ui){/* ... */}
  override fun showC() = launch(ui){/* ... */}
  override fun hideC() = launch(ui){/* ... */}
```

Before structured concurrency

```
class Presenter( /*...*/ ) : Contract.Presenter {
  private var view: Contract.View? = null
  private val jobs = mutableListOf < Job > ()
  override fun onRefresh() {
    jobs += launch(ui) {/*...*/}
  override fun unbind() {
    view = null
    jobs.forEach { it.cancel() }
    jobs.clear()
```

```
class Presenter(
  private val useCase: UseCase,
 private val ui: Scheduler,
  private val io: Scheduler
) : Contract.Presenter {
 private val bag = CompositeDisposable()
  override fun onRefresh() {
    bag.add(
      useCase
        .observeOn( /*...*/ )
        .subscribeOn( /*...*/ )
        .subscribe { /* ... */ }
```

```
class Presenter(
  private val useCase: UseCase,
  private val ui: Scheduler,
 private val io: Scheduler
) : Contract.Presenter {
  private val bag = CompositeDisposable()
  override fun unbind() {
    view = null
    bag.clear() // not dispose!
```

# Question

Why disposing?

```
class Presenter( /*...*/ ) : Contract.Presenter {
  private var view: Contract.View? = null
  private val jobs = mutableListOf < Job > ()
  override fun onRefresh() {
    jobs += launch(ui) {/*...*/}
  override fun unbind() {
    view = null
    jobs.forEach { it.cancel() } // release resource
    jobs.clear() // avoid memory leaks
```

### Question

Why a list and not only one reference?

```
class Presenter( /*...*/ ) : Contract.Presenter {
 private var view: Contract.View? = null
  private val jobs = mutableListOf < Job > ()
  override fun onA(){ jobs += launch(ui){/*...*/}}
  override fun onB(){ jobs += launch(ui){/*...*/}}
  override fun onC(){ jobs += launch(ui){/*...*/}}
  override fun unbind() {
   view = null
    jobs.forEach { it.cancel() }
    jobs.clear()
```

## Question

Why is it error prone?

```
class Presenter( /*...*/ ) : Contract.Presenter {
 private var view: Contract.View? = null
  private val jobs = mutableListOf < Job > ()
  override fun onA(){ jobs += launch(ui){/*...*/}}
  override fun onB(){ launch(ui){/*...*/}}
  override fun onC(){ jobs += launch(ui){/*...*/}}
  override fun unbind() {
   view = null
    jobs.forEach { it.cancel() }
    jobs.clear()
```

Welcome structured concurrency

## **Exceptions**

```
interface UseCase {
    fun fetch(): List<Item>
```

```
class Presenter(
        private val view: Contract. View,
        private val useCase: UseCase
) : Contract.Presenter {
    override fun onRefresh() {
        val items = useCase.fetch()
        view.showItems(items)
    }
    override fun onSubmit() = TODO()
```

