

Stat4DS / Homework 02

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Due Thursday, December 24 (on Moodle)

General Instructions

I expect you to upload your solutions on Moodle as a **single running R Markdown** file (.rmd) + its html output, **named with your surnames**. Alternatively, a zip-file with all the material inside will be fine too.

R Markdown Test

To be sure that everything is working fine, start **RStudio** and create an empty project called **HW1**. Now open a new **R Markdown** file (File > New File > R Markdown...); set the output to **HTML mode**, press OK and then click on **Knit HTML**. This should produce a web page with the knitting procedure executing the default code blocks. You can now start editing this file.

1. Background: Cooperative Games (more info)

1.1 Generalities

A *game* in the sense of *game theory* is an abstract mathematical model of a scenario in which some sort of agents *interact*. It is abstract in the sense that irrelevant detail is omitted: the game aims to capture only those features of the scenario that are relevant to the decisions that must be made by players within the game.

Agents can be anything depending on the application: typically real humans (e.g. investors, political parties, etc.), but also signals/towers in a communication network, genes in a biological setup or even variables in a predictive/learning problem.

The form of games we are interested in is the most basic and widely-studied model of **cooperative games**.

More specifically, our games are populated by a (non-empty) set $\mathcal{P} = \{1, \dots, p\}$ of agents: the **players** of the game. A **coalition** C is simply *any* subset of the player-set \mathcal{P} . The **grand coalition** is the set \mathcal{P} of *all* players.

All this said, let's define what we mean by a *characteristic function game*. In the following, with $2^{\mathcal{P}}$ we will denote the **power-set**, that is, the set of all subsets, of \mathcal{P} .

Definition 1. A **characteristic function game** G is given by a pair (\mathcal{P}, ν) , where $\mathcal{P} = \{1, \dots, p\}$ is a finite, non-empty set of agents, and $\nu : 2^{\mathcal{P}} \mapsto \mathbb{R}$ is a **characteristic function**, which maps each coalition $C \subseteq \mathcal{P}$ to a real number $\nu(C)$.

The number $\nu(C)$ is usually referred to as **the value of the coalition** C .

There are many possible interpretations for the *characteristic function*, but note right away that characteristic function games assign values to a coalition as a whole, and **not** to its individual members. That is, the model behind a characteristic function game does **not** tell you how the coalition value $\nu(C)$ should be **divided** among the members of C .

In fact, the question of **how to divide** the coalition value is a fundamental research topic in cooperative game theory¹.

Notice also that, from a computational point of view, the naïve representation of a characteristic function game that consists in explicitly listing every coalition C together with the associated value $\nu(C)$, being of the order 2^p in size, is **not** practical unless the number of players is very small. On the other hand, **most real-life interactions** admit an encoding of size polynomial in p ; such an encoding provides an **implicit** description of the characteristic function.

As an example, consider modeling the decision-making process in voting bodies.

Example 1. (*weighted voting games*)

- A country has a 101 seats in its parliament, and each representative belongs to one of four parties: Liberal (L \rightsquigarrow 40 seats), Moderate (M \rightsquigarrow 22 seats), Conservative (C \rightsquigarrow 30 seats), or Green (G \rightsquigarrow 9).
- The parliament needs to decide how to allocate “1 billion euros” of discretionary spending.

¹An implicit assumption here: the coalition value $\nu(C)$ can be **divided** among the members of C in *any* way that the members of C choose. Formally, games with this property are said to be **transferable utility** games (TU games)

- The decision is made by a **simple majority vote**, and we assume that all representatives vote along the party lines.
- Parties can form coalitions; a coalition has **value** “1 billion euros” **IF** it can win the vote no matter what the other parties do, and value 0 otherwise.

Hence, after some thinking, we see that the associated 4-players game has $\mathcal{P} = \{L, M, C, G\}$ and characteristic function

$$\nu(S) = \begin{cases} 0 & \text{if } (\#S \leq 1) \text{ or } (G \in S) \\ 10^9 & \text{otherwise} \end{cases} \quad \text{where} \quad \#S = \{\text{cardinality of the coalition } S\}.$$

1.2 Solution Concepts: the Shapley Value

A key problem in game theory is to try to understand what the **outcomes** of a game will be. In our cooperative framework, an *outcome* of a game G consists of two parts:

1. a **coalition structure**, that is, a **partition** $CS = \{C_1, \dots, C_m\}$ of the player-set $\mathcal{P} = \{1, \dots, p\}$ into coalitions;
2. a **payoff vector** $\mathbf{x} = (x_1, \dots, x_p) \in \mathbb{R}^p$ for a coalition structure $CS = \{C_1, \dots, C_m\}$, which distributes the value $\nu(C_j)$ of each coalition among its members. Any legit payoff vector \mathbf{x} must satisfy the following natural conditions:
 - $x_j \geq 0$ for all $j \in \mathcal{P}$;
 - $\sum_{r \in C_j} x_r \leq \nu(C_j)$ for any $j \in \{1, \dots, m\}$

Now, any partition of agents into coalitions and any payoff vector that respects this partition corresponds to a *plausible* outcome of a characteristic function game. However, not all outcomes are equally desirable.

For instance, if all agents contribute equally to the value of a coalition, a payoff vector that allocates the entire payoff to just one of the agents is less appealing than the one that shares the profits equally among all agents.

Similarly, an outcome that push all agents to work together is (typically) preferable to an outcome that some of the agents want to deviate from.

More broadly, one can evaluate outcomes according to two criteria: **fairness** (i.e., how well each agent's payoff reflects his contribution), and **stability** (i.e., what are the incentives for the agents to stay in the coalition structure). These criteria give rise to two families of solution concepts having as most notable members the **Shapley Value** and the **Core** respectively.

Given its broad success for defining **variable importance** in supervised learning problems (see [here](#), [here](#), [here](#), [here](#), and also [here](#), just to mention a few) in the following we will focus on the former, the *Shapley Value*, **forged in the '50s** by the one and only, sir **Lloyd S. Shapley**.

The Shapley value is a solution concept that is usually formulated with respect to the grand coalition: it defines a way of distributing the value $\nu(\mathcal{P})$ that could be obtained by the grand coalition, and it is based on the intuition that the payment that each agent receives should be *proportional to his contribution*.

Idea v1.0: a naïve implementation of this idea would be to pay each agent according to how much he increases the value of the coalition of all other players when he joins it, i.e., set the payoff of player j to $\nu(\mathcal{P}) - \nu(\mathcal{P} \setminus \{j\})$.

Problem: Under this payoff scheme the total payoff assigned to the agents may differ from the value of the grand coalition.

Idea v2.0: we can **fix an ordering** of the agents and pay each agent according to how much he contributes to the coalition formed by his predecessors in this ordering: *Agent 1* receives $\nu(\{1\})$, *Agent 2* receives $\nu(\{1, 2\}) - \nu(\{1\})$, and so on.

It is easy to see that this payoff scheme distributes the value of the grand coalition among the agents.

Problem: agents that play symmetric roles in the game may receive different payoffs depending on their position in the order.

Shapley's insight: the dependence on the agents ordering can be eliminated by **averaging over all possible orderings** (or permutations) of the players.

Now, to formally define the Shapley value, we need some additional notation.

1. Fix a characteristic function game $G = (\mathcal{P}, \nu)$ and let $\Pi_{\mathcal{P}}$ be the set of all permutations of \mathcal{P} . Given a specific permutation $\pi \in \Pi_{\mathcal{P}}$, we denote by $S_{\pi}(j)$ the set of all predecessors of player j in π , i.e., we set $S_{\pi}(j) = \{r \in \mathcal{P} \text{ such that } \pi(r) < \pi(j)\}$. For example, if $\mathcal{P} = \{1, 2, 3\}$, then

$$\Pi_{\mathcal{P}} = \{(1, 2, 3), (1, 3, 2), (2, 1, 3), (3, 1, 2), (3, 2, 1)\},$$

moreover, if $\pi = (3, 1, 2)$, then $S_{\pi}(3) = \emptyset$, $S_{\pi}(1) = \{3\}$ and $S_{\pi}(2) = \{1, 3\}$.

2. The **marginal contribution** of an agent j with respect to a permutation π in a game $G = (\mathcal{P}, \nu)$ is denoted by $\Delta_{\pi}^G(j)$ and measures how much player j increases the value of the coalition consisting of its predecessor in π . More specifically:

$$\Delta_{\pi}^G(j) = \nu(S_{\pi}(j) \cup \{j\}) - \nu(S_{\pi}(j)). \quad (1)$$

We can now define the Shapley value $\psi^G(j)$ of player j : it is simply his **average marginal contribution**, where the average is taken over **all** permutations of the player-set (assumed to be **all equally likely**):

$$\psi^G(j) \stackrel{\text{def}}{=} \mathbb{E}_\pi(\Delta_\pi^G(j)) = \frac{1}{p!} \sum_{\pi \in \Pi_{\mathcal{P}}} \Delta_\pi^G(j) \stackrel{(\heartsuit)}{=} \sum_{C: j \notin C} \frac{(\#C)!(p-1-\#C)!}{p!} (\nu(C \cup \{j\}) - \nu(C)). \quad (2)$$

The second version (\heartsuit) can be derived by noting that the marginal contributions $\Delta_\pi^G(j)$ are of the form $\nu(C \cup \{j\}) - \nu(C)$ where C is a coalition not containing j . Now, for how many orderings does one have $S_\pi(j) = C$? There are $(\#C)!$ possible orderings of C and $(p-1-\#C)!$ orderings of $\mathcal{P} \setminus (C \cup \{j\}) \rightsquigarrow$ another nice probabilistic interpretation (see page 307 **here**).

It can be shown that the Shapley value is in fact **the only** payoff division scheme that has **all** these four properties:

- *Efficiency*: $\sum_j \psi^G(j) = \nu(\mathcal{P})$.
- *Dummy player*: if player j is *dummy* (i.e., $\nu(C \cup \{j\}) = \nu(C)$ for *any* coalition C), then $\psi^G(j) = 0$.
- *Symmetry*: if j_1 and j_2 are *symmetric* players (i.e., $\nu(C \cup \{j_1\}) = \nu(C \cup \{j_2\})$ for *any* coalition C), then $\psi^G(j_1) = \psi^G(j_2)$.
- *Additivity*: given any two characteristic function games $G_1 = (\mathcal{P}, \nu_1)$ and $G_2 = (\mathcal{P}, \nu_2)$, and their *sum* $G^+ = (\mathcal{P}, \nu_1 + \nu_2)$,

$$\psi^{G^+}(j) = \psi^{G_1}(j) + \psi^{G_2}(j) \quad \text{for all } j \in \mathcal{P}.$$

Example 2. Suppose that **Ciccio-Pharma** (C) is a small biotech company who discovered a new vaccine but, to manufacture and market it, needs to team up with a larger partner. Two candidates: **Aristo-Medical** (A) and **BruttiMaBuoni-Inc** (B). If A or B teams up with C , the big firm will split “1 billion” with C . Here is a *possible* characteristic function

$$\nu(A) = \nu(B) = \nu(C) = \nu(AB) = 0 \quad \text{and} \quad \nu(AC) = \nu(BC) = \nu(ABC) = 1.$$

Since there’re only 3 players, to get Shapley’s payoffs we can make a table indicating the value brought to a coalition by each player on the way to formation of the gran coalition:

Permutation	Player A	Player B	Player C
ABC	0	0	1
ACB	0	0	1
BAC	0	0	1
BCA	0	0	1
CAB	1	0	0
CBA	0	1	0
TOTAL VALUE	1	1	4

Since in the derivation of the Shapley value it is assumed that the $3! = 6$ permutations/arrival sequences are all *equally likely*, the average value of each biotech company is simply:

$$\psi(A) = \frac{1}{6}, \quad \psi(B) = \frac{1}{6}, \quad \psi(C) = \frac{4}{6} = \frac{2}{3}.$$

So **Ciccio-Pharma**, the drug discoverer, will get two-thirds of the billion, and the big companies split the remaining third.

Example 3. (*Shapley in Simple Games*)

A game $G = (\mathcal{P}, \nu)$ is said to be **simple** if it is *monotone* (i.e., $\nu(C_1) \leq \nu(C_2)$ if $C_1 \subseteq C_2$) and its characteristic function only takes values 0 and 1. The game in Example 1 is clearly *simple* as soon as we rescale the payoffs so that $\nu(\mathcal{P}) = 1$.

In simple games, the Shapley (a.k.a. *Shapley-Shubik power index* in this context) have a particularly attractive interpretation: it measures the **power of a player**, i.e., the probability that she can influence the outcome of the game.

Indeed, the Shapley value of a player j in a simple game $G = (\mathcal{P}, \nu)$ with $\#\mathcal{P} = p$ can be rewritten as follows:

$$\psi^G(j) = \{\text{proportion of permutations where } j \text{ is } \mathbf{pivotal}\} = \frac{\#\{\pi \in \Pi_{\mathcal{P}} \text{ such that } \nu(S_\pi(j)) = 0 \text{ and } \nu(S_\pi(j) \cup \{j\}) = 1\}}{p!} \quad (3)$$

In other words, if agents join the coalition in a random order, $\psi^G(j)$ is exactly the **probability that player j turns a losing coalition into a winning one**.

Example 4. (*Shapley for Variable Importance*)

Suppose we are dealing with a generic (supervised) learning problem where we need to predict a response Y based on a bunch of covariates $\mathbf{X} = (X_1, \dots, X_p)$. The goal here is to measure the **importance** of X_j in this task.

We can cast this problem into a suitable characteristic function game having the p covariates as players $\mathcal{P} = \{1, \dots, p\}$, and some measure of fit as characteristic function ν . To be more specific, for any coalition (of covariates) C , let $\mathbf{X}_C = (X_j : j \in C)$ and $\hat{Y}(C) = \mathbb{E}(Y | \mathbf{X}_C)$, the **ideal** optimal predictor (under a squared risk) based on the covariates in \mathbf{X}_C . Now, if we take

$$\nu(C) = -\mathbb{E}(Y - \hat{Y}(C))^2 \rightsquigarrow \psi(j) = \frac{1}{p!} \sum_{\pi} \mathbb{E}[\hat{Y}(S_\pi(j)) - \hat{Y}(S_\pi(j) \cup \{j\})]^2 \rightsquigarrow \text{population quantity to be estimated!} \quad (4)$$

This is just an averaged version (over all possible submodels) of LOCO, another, **recently introduced** var-importance measure. It is instructive to try to rephrase in this context the previous 4 properties that characterize the Shapley value.

Remark: this is all nice and good but, of course, Shapley for variable importance is not perfect. For example, it defines variable importance with respect to all submodels, but most of those submodels are not of interest and, in addition, it is strongly influenced by the correlation between covariates.

2. The Exercise: Let's get together and feel all right...

↪ Your job ↩

1. Introductory

To check your understanding of Shapley, let's start easy by playing around with a tiny game having only 3 players. So, imagine (!) there're three students curiously named: Antwohnette (A), BadellPadel (B) and Chumbawamba (C). Exactly one of them needs to be working (not necessarily all day long) on this HW to complete it. Here's their working-hours:

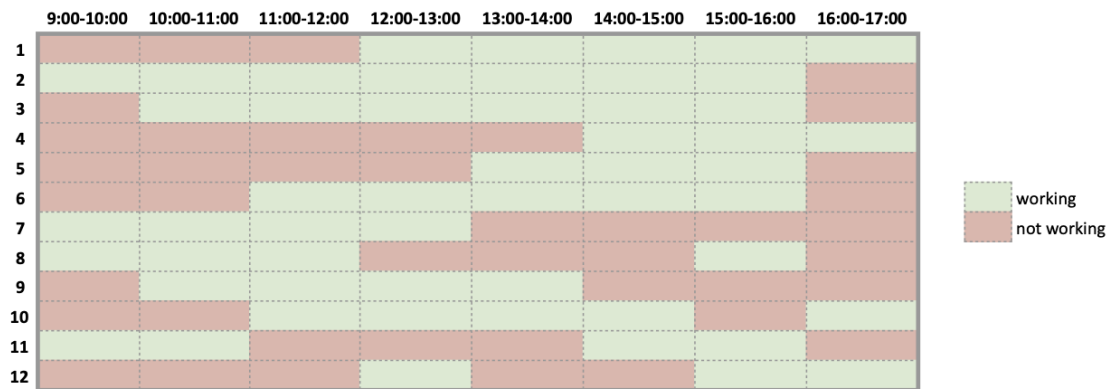
A	Antwohnette	14:00-17:00
B	BadellPadel	11:00-16:00
C	Chumbawamba	9:00-13:00

A coalition C is an agreement by one or more students as to the times they will be *really* working². Its values will be given by:
 $\nu(C) = \{\# \text{ of hours potentially saved by well organized coalition}\}$

Clearly $\nu(A) = \nu(B) = \nu(C) = 0$ and $\nu(ABC) = 4$. **Complete the definition of ν and find Shapley** (see Example 2).

2. Probabilistic

Now let's get serious. Imagine you are in a real, larger team with 12 people having the following working schedule:



Now, I might ask you to find the characteristic function as before... but here you're dealing with $2^{12} = 4096$ coalitions \rightsquigarrow **I built it for you**³ (**Bonus:** write an R function, **the shorter the better**, to get it for a general game of this type). I might also ask you to calculate the Shapley payoffs for each team member... but now there're $12! = 479,001,600$ orderings to consider, not exactly trivial as in Example 2...

Nevertheless, look again at Equation 2: the Shapley is "just" an **expectation** (w.r.t. a *very* specific distribution), so:

- Review our notes on **MonteCarlo Approximation and Error Control**.
- Select 3 players of the team.
- Fix a suitable⁴ (ϵ, α) pair of tolerance and confidence, and design an Hoeffding-based simulation to approximate the Shapley payoffs of those 3 players (HINT: if you look closely, all you need is essentially `sample()`)
- Produce separate Hoeffding-based confidence intervals for these payoffs and briefly comment all your results.

3. Statistical

...ehmmm... actually this morning while writing the solutions, I realized the original version of this point was way too hard and I'm simplifying it... stay tuned...

²This is an example of cooperative game theory applied to a resource allocation problem.

³The characteristic function is stored as a list `char_fun` with $p = 12$ slots so that, for example, `char_fun[[2]]` contains the $\binom{12}{2} = 66$ values associated to coalitions formed by 2 players.

⁴First and foremost in terms of computational time!