

assignment06

June 5, 2019

1 Functional Programming SS19

2 Assignment 06 Solutions

2.0.1 Exercise 1 (Fixpoints)

Give the value of $Val\llbracket e \rrbracket_\rho$ for the following simple Haskell expressions $e \in \{e_1, e_2, e_3\}$ for the environment $\rho = \omega + \rho'$, where ω is the initial environment and ρ' is the environment with $\rho'(x) = 2, \rho'(y) = 18, \rho'(z) = 3$, and ρ' is undefined for all other variables.

Describe your computation in detail by applying rules for *Val* step by step. Also, for each higher-order function $f : Dom \rightarrow Dom$, where $lfp\ f$ is needed in the calculation, determine what the function $f^n(\perp), n \in \mathbb{N}$, computes.

```
e1 = let isAnswer = \x -> if x < 17 then 2 else 3
      in isAnswer y
```

```
e2 = let sum = \x -> if x < 4 then x + sum (x+1)
      else 0
      in sum x
```

```
e3 = let pow = \y -> \ z -> if z <= 0 then 1
      else y * pow y (z-1)
      in pow 4 z
```

Hints: * You may switch between infix- and prefix-notation for Haskell operators without any intermediate steps if needed. * You may simplify $Val\llbracket exp_1\ exp_2\ exp_3 \rrbracket_\rho$ to $f\ (Val\llbracket exp_2 \rrbracket_\rho)\ (Val\llbracket exp_3 \rrbracket_\rho)$ where $f = Val\llbracket exp_1 \rrbracket_\rho$ in Functions in Dom in one step if exp_1 represents a function expecting two arguments, e.g. $+, -, *, <, <=$, etc.

- $e1 = \text{let isAnswer} = \backslash x \rightarrow \text{if } x < 17 \text{ then } 2 \text{ else } 3 \text{ in isAnswer } y$

$$\begin{aligned} Val\llbracket e1 \rrbracket &= Val\llbracket \text{let isAnswer} = \backslash x \rightarrow \text{if } x < 17 \text{ then } 2 \text{ else } 3 \text{ in isAnswer } y \rrbracket_\rho \\ &= Val\llbracket \text{isAnswer } y \rrbracket(\rho + \{\text{isAnswer} / lfp\ f\}) \end{aligned}$$

for f in Functions in Dom where

$$f(k) = \text{Val}[\![\backslash x \rightarrow \text{if } x < 17 \text{ then } 2 \text{ else } 3]\!](\rho + \{\text{isAnswer}/k\})$$

Here, we have $\text{Val}[\![\backslash x \rightarrow \text{if } x < 17 \text{ then } 2 \text{ else } 3]\!] = h$ in Functions in Dom where

$$\begin{aligned} h(d) &= \text{Val}[\![\text{if } x < 17 \text{ then } 2 \text{ else } 3]\!](\rho + \{x/d\}) \\ &= \begin{cases} 2 & \text{if } \text{Val}[\![<(x \ 17)]\!](\rho + \{x/d\}) = \text{True in Constructions}_0 \text{ in Dom} \\ 3 & \text{if } \text{Val}[\![<(x \ 17)]\!](\rho + \{x/d\}) = \text{False in Constructions}_0 \text{ in Dom} \\ \perp & \text{otherwise} \end{cases} \end{aligned}$$

such that $\text{Val}[\![<(x \ 17)]\!](\rho + \{x / d\}) = g$ in Functions in Dom.

Since isAnswer is not a recursive function, we have that $\text{lfp } f = h$ and so

$$\begin{aligned} \text{Val}[\![e1]\!] &= \text{Val}[\![\text{isAnswer } y]\!](\rho + \{\text{isAnswer} / h\}) \\ &= \text{Val}[\![\text{if } x < 17 \text{ then } 2 \text{ else } 3]\!](\rho + \{x / y\}) \\ &= \text{Val}[\![\text{if } x < 17 \text{ then } 2 \text{ else } 3]\!](\{x / 18\}) \\ &= \text{Val}[\![\text{if } 18 < 17 \text{ then } 2 \text{ else } 3]\!] \\ &= 3 \end{aligned}$$

- $e2 = \text{let sum} = \backslash x \rightarrow \text{if } x < 4 \text{ then } x + \text{sum } (x+1) \text{ else } 0 \text{ in sum } x$

$$\begin{aligned} \text{Val}[\![e2]\!] &= \text{Val}[\![\text{let sum} = \backslash x \rightarrow \text{if } x < 4 \text{ then } x + \text{sum } (x+1) \text{ else } 0 \text{ in sum } x]\!]_{\rho} \\ &= \text{Val}[\![\text{sum } x]\!](\rho + \{\text{sum} / \text{lfp } f\}) \end{aligned}$$

for f in Functions in Dom where

$$f(k) = \text{Val}[\![\backslash x \rightarrow \text{if } x < 4 \text{ then } x + \text{sum } (x+1) \text{ else } 0]\!](\rho + \{\text{sum}/k\})$$

As above, we have $\text{Val}[\![\backslash x \rightarrow \text{if } x < 4 \text{ then } x + \text{sum } (x+1) \text{ else } 0]\!] = h$ in Functions in Dom where

$$\begin{aligned} h(d) &= \text{Val}[\![\text{if } x < 4 \text{ then } x + \text{sum } (x+1) \text{ else } 0]\!](\rho + \{\text{sum} / \text{lfp } f, x/d\}) \\ &= \begin{cases} x + \text{sum}(x+1) & \text{if } \text{Val}[\![<(x \ 4)]\!](\rho + \{\text{sum} / \text{lfp } f, x/d\}) = \text{True in Constructions}_0 \text{ in Dom} \\ 0 & \text{if } \text{Val}[\![<(x \ 4)]\!](\rho + \{\text{sum} / \text{lfp } f, x/d\}) = \text{False in Constructions}_0 \text{ in Dom} \\ \perp & \text{otherwise} \end{cases} \end{aligned}$$

such that $\text{Val}[\![<(x \ 4)]\!](\rho + \{\text{sum} / \text{lfp } f, x / d\}) = g$ in Functions in Dom.

We are now dealing with a recursive function, so let us consider $f^n(\perp), n \in \mathbb{N}$:

$$\begin{aligned}
g_0 &= f^0(\perp) = \perp \\
g_1 &= f^1(\perp) = \begin{cases} x + \perp(x+1) & \text{if } x < 4 \\ 0 & \text{if } x \geq 4 \\ \perp & \text{otherwise} \end{cases} = \begin{cases} \perp & \text{if } x < 4 \\ 0 & \text{if } x \geq 4 \\ \perp & \text{otherwise} \end{cases} \\
g_2 &= f^2(\perp) = f(g_1) = \begin{cases} x + g_1(x+1) & \text{if } x < 4 \\ 0 & \text{if } x \geq 4 \\ \perp & \text{otherwise} \end{cases} = \begin{cases} \perp & \text{if } x < 3 \\ 3 & \text{if } x = 3 \\ 0 & \text{if } x > 4 \\ \perp & \text{otherwise} \end{cases} \\
g_3 &= f^3(\perp) = f(g_2) = \begin{cases} x + g_2(x+1) & \text{if } x < 4 \\ 0 & \text{if } x \geq 4 \\ \perp & \text{otherwise} \end{cases} = \begin{cases} \perp & \text{if } x < 3 \\ 5 & \text{if } x = 2 \\ 3 & \text{if } x = 3 \\ 0 & \text{if } x > 4 \\ \perp & \text{otherwise} \end{cases}
\end{aligned}$$

This is thus the function that computes the sum of all number from x up to 3 and returns 0 otherwise; thus

$$\text{lfp } f = \begin{cases} \sum_{i=x}^3 i & \text{if } x < 4 \\ 0 & \text{if } x \geq 4 \\ \perp & \text{otherwise} \end{cases}$$

We thus have that

$$\begin{aligned}
\text{Val}[\![e2]\!] &= \text{Val}[\![\text{sum } x]\!](\rho + \{\text{sum} / \text{lfp } f\}) \\
&= \text{Val}[\![\text{lfp } f \ x]\!](\rho) \\
&= \text{Val}[\![\text{lfp } f \ 2]\!] \\
&= \sum_{i=2}^3 i = 5
\end{aligned}$$

- $e3 = \text{let } \text{pow} = \backslash y \rightarrow \backslash z \rightarrow \text{if } z \leq 0 \text{ then } 1 \text{ else } y * \text{pow } y \ (z-1) \text{ in } \text{pow } 4 \ z$

$$\begin{aligned}
\text{Val}[\![e3]\!] &= \text{Val}[\![\text{let } \text{pow} = \backslash y \rightarrow \backslash z \rightarrow \text{if } z \leq 0 \text{ then } 1 \text{ else } y * \text{pow } y \ (z-1) \text{ in } \text{pow } 4 \ z]\!]_{\rho} \\
&= \text{Val}[\![\text{pow } 4 \ z]\!](\rho + \{\text{pow} / \text{lfp } f\})
\end{aligned}$$

for f in Functions in Dom where

$$f(k) = \text{Val}[\![\backslash y \rightarrow \backslash z \rightarrow \text{if } z \leq 0 \text{ then } 1 \text{ else } y * \text{pow } y \ (z-1) \text{ in } \text{pow } 4 \ z]\!](\rho + \{\text{pow} / k\})$$

We have $\text{Val}[\![\backslash y \rightarrow \backslash z \rightarrow \text{if } z \leq 0 \text{ then } 1 \text{ else } y * \text{pow } y \ (z-1)]\!] = h$ in Functions in Dom where

$$h(d_1) = \text{Val}[\![\backslash z \rightarrow \text{if } z \leq 0 \text{ then } 1 \text{ else } y * \text{pow } y \ (z-1)]\!](\rho + \{\text{pow} / \text{lfp } f, y / d_1\})$$

Since we have another lambda expression, we find that $Val[\![z \rightarrow \text{if } z \leq 0 \text{ then } 1 \text{ else } y * \text{pow } y (z-1)]\!] = l$ in Functions in Dom and

$$l(d_2) = Val[\![\text{if } z \leq 0 \text{ then } 1 \text{ else } y * \text{pow } y (z-1)]\!](\rho + \{\text{pow} / \text{lfp } f, y/d_1, z/d_2\}) \\ = \begin{cases} 1 & \text{if } Val[\![z \leq 0]\!](\rho + \{\text{pow} / \text{lfp } f, y/d_1, z/d_2\}) = \text{True in Constructions}_0 \text{ in Dom} \\ y * \text{pow } y (z-1) & \text{if } Val[\![z \leq 0]\!](\rho + \{\text{pow} / \text{lfp } f, y/d_1, z/d_2\}) = \text{False in Constructions}_0 \text{ in Dom} \\ \perp & \text{otherwise} \end{cases}$$

such that $Val[\![z \leq 0]\!](\rho + \{\text{pow} / \text{lfp } f, y/d_1, z/d_2\}) = g$ in Functions in Dom.

Let us now consider $f^n(\perp), n \in \mathbb{N}$:

$$g_0 = f^0(\perp) = \perp \\ g_1 = f^1(\perp) = \begin{cases} 1 & \text{if } z \leq 0 \\ y * \perp(z-1) & \text{if } z > 0 \\ \perp & \text{otherwise} \end{cases} = \begin{cases} 1 & \text{if } z \leq 0 \\ \perp & \text{if } z > 0 \\ \perp & \text{otherwise} \end{cases} \\ g_2 = f^2(\perp) = f(g_1) = \begin{cases} 1 & \text{if } z \leq 0 \\ y * g_1(z-1) & \text{if } z > 0 \\ \perp & \text{otherwise} \end{cases} = \begin{cases} 1 & \text{if } z \leq 0 \\ y & \text{if } z = 1 \\ \perp & \text{if } z > 1 \\ \perp & \text{otherwise} \end{cases} \\ g_3 = f^3(\perp) = f(g_2) = \begin{cases} 1 & \text{if } z \leq 0 \\ y * g_2(z-1) & \text{if } z > 0 \\ \perp & \text{otherwise} \end{cases} = \begin{cases} 1 & \text{if } z \leq 0 \\ y & \text{if } z = 1 \\ y^2 & \text{if } z = 2 \\ \perp & \text{if } z > 2 \\ \perp & \text{otherwise} \end{cases}$$

which is the function that computes the power z -th of a positive integer y ; thus

$$\text{lfp } f(x) = \begin{cases} y^z & \text{if } z \leq 0 \\ \perp & \text{otherwise} \end{cases}$$

We thus have that

$$Val[\![e3]\!] = Val[\![\text{pow } 4 \ z]\!](\rho + \{\text{pow} / \text{lfp } f\}) \\ = Val[\![\text{lfp } f \ 4 \ z]\!](\rho) \\ = Val[\![\text{lfp } f \ 4 \ 3]\!] \\ = 4^3 = 64$$