

1 Question 1

Here we consider a graph G with 2 connected components, and we want to know the number of edges and triangles:

- The first component is a complete graph of $n = 100$ vertices. The number of edges is thus:

$$\binom{n}{2} = \frac{n!}{2!(n-2)!} = \frac{n \cdot (n-1)}{2} = 4\,950 \quad (1)$$

since we can reach $(n-1)$ edges from every of the n edges, without repetition. Since every triangle is closed in a complete graph, the number of triangles is:

$$\frac{1}{3} * \binom{n}{3} = \frac{161\,700}{3} = 53\,900 \quad (2)$$

- The second component is a bipartite graph of $n = 50$ vertices. The number of edges simplifies to $n * n = 2500$. Since it's a bipartite graph, there are no closed triangles, so the count of triangles here is 0.

To sum up the number of edges is $4\,950 + 2\,500 = 7\,450$ and the number of triangles is $53\,900$.

2 Question 2

The global clustering coefficient is:

$$C = \frac{\text{number of closed triplets}}{\text{number of total triplets (open and closed)}} \quad (3)$$

So we can see that the maximum value for C is 1, and this happens only when the number of open triplet is 0: this is the case in a **complete graph**.

3 Question 3

Here we consider a connected graph with a single connected component. Since the matrix $L = D - W$ is symmetric, positive and semi-definite, its smallest trivial eigenvalue is **0 associated to the unit eigenvector 1**.

With the **Rayleigh-Ritz Theorem**, we know that the eigenvector corresponding to smallest eigenvalue (0) offers no useful information when solving the Two-Way Cut from the Laplacian, and so removing it before applying spectral clustering doesn't affect the results.

4 Question 4

Let's see the algorithm for spectral clustering:

Here we can see clearly that step 1, 2 and 3 are purely deterministic since they are mathematical computation with no approximation or random process involved. However, when applying the K-means algorithm, we are doing a stochastic step. Even if the results of the K-means step and thus clustering algorithm should attain a local optimum every time we run it, running k-means with different random seeds could give different solutions to the minimization problem. At that level, the Spectral clustering algorithm's output is **stochastic**.

5 Question 5

Here is the formula for calculating the modularity of a graph clustering:

$$Q = \sum_{c=1}^{n_c} \left[\frac{l_c}{m} - \left(\frac{d_c}{2m} \right)^2 \right] \quad (4)$$

Now let's compute the modularities for each of the 2 scenarios.

Algorithm 1 Spectral Clustering

Input: Graph $G = (V, E)$ and parameter k **Output:** Clusters C_1, C_2, \dots, C_k (i.e., cluster assignments of each node of the graph)

- 1: Let A be the adjacency matrix of the graph
 - 2: Compute the Laplacian matrix $L_{rw} = I - D^{-1}A$. Matrix D corresponds to the diagonal degree matrix of graph G (i.e., degree of each node v (= number of neighbors) in the main diagonal)
 - 3: Apply eigenvalue decomposition to the Laplacian matrix L_{rw} and compute the eigenvectors that correspond to d smallest eigenvalues. Let $U = [u_1 | u_2 | \dots | u_d] \in \mathbb{R}^{m \times d}$ be the matrix containing these eigenvectors as columns
 - 4: For $i = 1, \dots, m$, let $y_i \in \mathbb{R}^d$ be the vector corresponding to the i -th row of U . Apply k -means to the points $(y_i)_{i=1, \dots, m}$ (i.e., the rows of U) and find clusters C_1, C_2, \dots, C_k
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Figure 1: Algorithm for Spectral Clustering

- Scenario a): For both green and blue community, we have $l_c = 6$ and $d_c = 13$ and $m = 13$, so:

$$Q = 2 * \left(\frac{6}{13} - \left(\frac{13}{26} \right)^2 \right) \approx 0.423 \quad (5)$$

- Scenario b): For the green community we have $l_c = 2$, $d_c = 11$ and $m = 13$. For the blue community, we have $l_c = 4$ and $d_c = 15$, so we get:

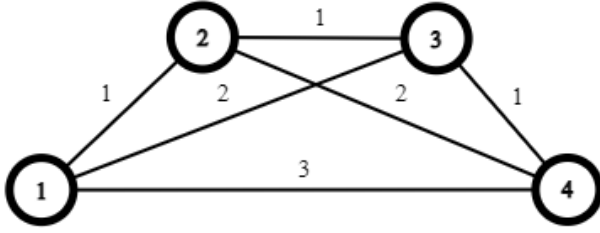
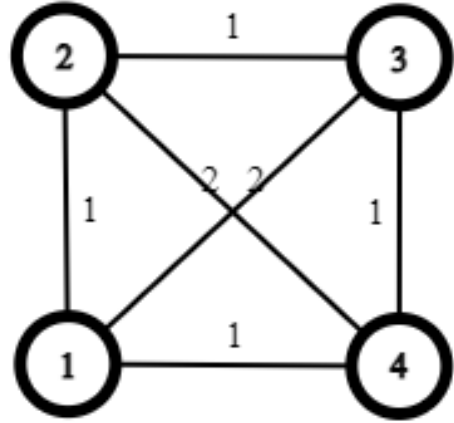
$$Q = \left(\frac{2}{13} - \left(\frac{11}{26} \right)^2 \right) + \left(\frac{4}{13} - \left(\frac{15}{26} \right)^2 \right) \approx -0.050 \quad (6)$$

6 Question 6

Let's recall the formula for calculating the shortest path kernel between two Floyd-transformed graphs G_1 and G_2 :

$$k(G_1, G_2) = \sum_{e_1 \in E_1} \sum_{e_2 \in E_2} k_{edge}(e_1, e_2) \quad (7)$$

So let's see the two graphs we have, P_n which is a path of n vertices and C_n , a cycle of n vertices (here $n = 4$):

Figure 2: A P_n graph representationFigure 3: A C_n graph representation

We can see that the row vector associated with P_n is $p = (3, 2, 1)$, and the one associated with C_n is $c = (4, 2, 0)$. If we denote Sk_1 the shortest path kernel for (C_4, C_4) , Sk_2 the shortest path kernel for (C_4, P_4) and Sk_3 the shortest path kernel for (P_4, P_4) , we finally get:

$$Sk_1 = c \cdot c = 20 \quad (8)$$

$$Sk_2 = c \cdot p = 16 \quad (9)$$

$$Sk_3 = p \cdot p = 14 \quad (10)$$