

# High-Dimensional Random Walks

Walking on Matroids

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Alex Albors Juez

Mentor: Linhang Huang

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WDRP, University of Washington

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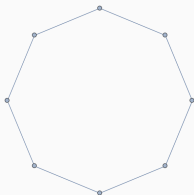
# Motivation

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# Random forests of a graph

Given an undirected graph  $G = (V, E)$ , want to count how many spanning forests there are in a certain graph  $G$ .

1. A **cycle** is a path that starts from a given vertex and ends at the same vertex.
2. A **tree** is a connected graph with no cycles.
3. A **forest** is a graph that has no cycles. Alternatively, it is a union of trees.



(a) Example of a cycle



(b) Example of a tree

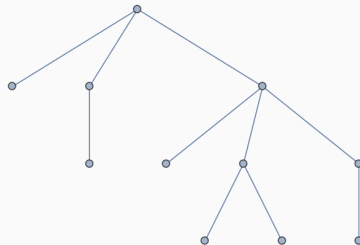


(c) Example of a forest

# Random graphs



(a) Example  $V = 11, E = 17$



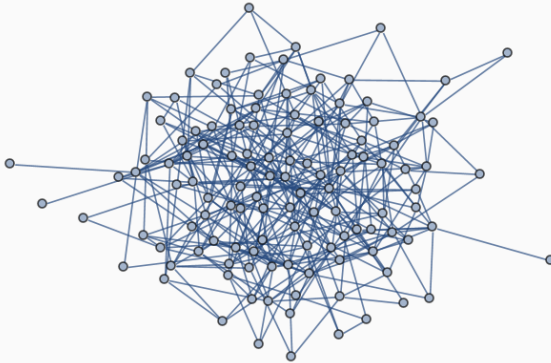
(b) Corresponding spanning tree

**Figure 2:** Possible forest (in this case just one tree)



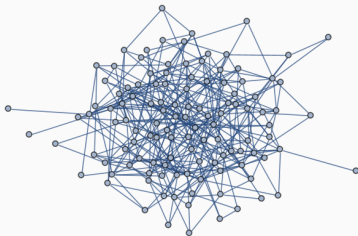
**Figure 3:** Random graph with  $V = 11$ ,  $E = 17$

# Random graphs

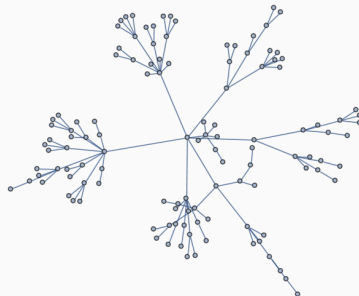


**Figure 4:** Random graph with  $V = 120$ ,  $E = 350$

# Random graphs



(a) Example  $V = 120$ ,  $E = 350$



(b) Corresponding spanning tree

**Figure 5:** Possible forest (in this case just one tree)



# Preliminaries

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A *matroid*  $M$  is a pair  $M = (X, \mathcal{I})$  where  $X$  is a finite set and  $\mathcal{I} \subseteq 2^X$  so that the following holds:

- (i) *Non-emptiness*:  $\emptyset \in \mathcal{I}$
- (ii) *Monotonicity*: If  $Y \in \mathcal{I}$  and  $Z \subseteq Y$  then  $Z \in \mathcal{I}$ .
- (iii) *Exchange property*: If  $Y, Z \in \mathcal{I}$  and  $|Y| < |Z|$ , then for some  $x \in Z \setminus Y$  we have  $Y \cup \{x\} \in \mathcal{I}$

## **Definition (basis)**

Let  $M = (X, \mathcal{I})$  be a matroid. A maximal independent set  $B \in \mathcal{I}$  is called a *basis* of  $X$ .

We can think of a graph as a *graphic matroid*. For a graph  $G = (V, E)$ , we let  $X = E$  and  $\mathcal{I} = \{F \subseteq E \mid F \text{ acyclic}\}$ . Then  $M = (X, \mathcal{I}) = (E, \mathcal{I})$ .

# Random walks

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# Random walk on the bases of M

**High-level idea:** Graph  $G \rightarrow$  Graphic Matroid  $\rightarrow$  Random Walk

**Procedure:**

1. Start with graph  $G = (V, E)$  with  $|E| = n$ .
2. Pick any maximal forest of size  $r$  in  $G$ . Forest is acyclic and of size  $1 \leq r \leq n$ .
3. This corresponds to basis element  $B$  of the graphic matroid of  $G$  (dimension  $r$ ).
4. Drop a random element  $i$  from  $B$ . Pick  $j$  uniformly at random from  $\{1, \dots, n\}$ , and try adding it to  $B \setminus \{i\}$ . Do it until we can.
5. Repeat step 4.

# Random walk on the bases of $M$

**High-level idea:** Graph  $G \rightarrow$  Graphic Matroid  $\rightarrow$  Random Walk

**Procedure:**

1. Start with graph  $G = (V, E)$  with  $|E| = n$ .
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4. Drop a random element  $i$  from  $B$ . Try adding an element  $j \in \{1, \dots, n\}$  to  $B \setminus \{i\}$  uniformly at random. If we can't, try again.
5. Repeat step 4.

## **Theorem (Convergence of algorithm)**

*Given an error  $0 < \epsilon < 1$  and error probability  $0 < \delta < 1$ , after  $r \cdot \log(n^r/\epsilon) \cdot \text{poly}(\log(1/\delta))$  steps, the algorithm will have seen all basis elements up to an error factor of  $1 \pm \epsilon$  with probability at least  $1 - \delta$ .*

# Random walks on bases of $M$

- The efficiency of the algorithm is due to [1, AKOV19]
- Ensuring we can go from sampling to counting is due to [2, JVV86]
- So we can sample an (almost) uniformly random basis element of  $M$  (a spanning forest of  $G$ ) by running this Random Walk long enough.

## Theorem (Convergence of algorithm)

*Given an error  $0 < \epsilon < 1$  and error probability  $0 < \delta < 1$ , after at most  $r \cdot \log(n^r/\epsilon) \cdot \text{poly}(\log(1/\delta))$  steps, the algorithm will have seen all basis elements up to an error factor of  $1 \pm \epsilon$  in the total number of basis elements with probability at least  $1 - \delta$ .*



[1, AKOV19] and [2, JVV86]

- The  $n^r$  comes from the fact there are at most  $\binom{n}{r} \leq n^r$  bases in an  $r$ -dimensional Matroid on the set  $[n]$ .

## References

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-  Anari, N., Liu, K., Gharan, S. O., Vintzant, C. (2019, June). *Log-concave polynomials II: High-dimensional walks and an FPRAS for counting bases of a matroid*. In Proceedings of the 51st Annual ACM SIGACT Symposium on Theory of Computing (pp. 1-12).
-  Mark R. Jerrum, Leslie G. Valiant, Vijay V. Vazirani, *Random generation of combinatorial structures from a uniform distribution*, Theoretical Computer Science, Volume 43, 1986 (pp. 169-188)