MPI 2.1 at **MPI Forum** Chicago, April 28-30, 2008 Ballot 4

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Höchstleistungsrechenzentrum Stuttgart



First official vote: MPI-21. Ballot 3, Items 1-6, 8, and 10-11 official (at April 28-30, 2008 meeting) Do you

Incorrect use of MPI IN PLACE in description of MPI ALLGATHER and MPI ALLGATHERV

2.a-g Edits to MPI-2 Chapter 4, Miscellany

Interlanguage use of Attributes and Error in Example 4.13 in MPI-2 3. (Use of Attributes in C and Fortran)

What info kevs can be set?

Datatypes and MPI PROBE

MPI PROCNULL and RMA Part 2 Number

(7. Significance of non-root arguments to MPI_COMM_SPAWN → MPI 2.2)

8. Thread safety and collective communication

const in C++ specification of predefined MPI objects → MPI 2.2)

10.a-d Questions about Cart create, and Cart coords

10.e Questions about Cart_create, and Cart_coords (Typo)

11.a-c Questions about Graph create

Abstain:

Official

Do you

accept

and 4

Items

Number

of insti-

tutions:

Yes:

23

No:

Abstain:

Not in the

room:

vote:

1st

vote:

accept

MPI-2.1

Ballot 4

Items

1-6.8.

and 11-12

of insti-

tutions:

Yes:

No:

the

MPI 2.1 Rolf Rabenseifne Slide 2 Höchstleistungsrechenzentrum Stuttgart





1st official vote: Do you accept

MPI-2.1 Ballot 4 Items 12-21, and 22.A

Number of institutions:

Yes:

No:

Abstain:

First official vote: MPI-21. Ballot 3. Items 12-21, and 22.A (at April 28-30, 2008 meeting)

Why no MPI INPLACE for MPI EXSCAN?

MPI GET PROCESSOR NAME, MPI COMM GET NAME and Fortran/C

Interpretation of user defined datarep

15.a-b MPI Abort

MPI Type create f90 real etc.

17. MPI File get info

MPI File set view 18.

MPI IN PLACE for MPI Reduce scatter 19.

Blocklengths of zero in MPI_TYPE_STRUCT and in MPI 20. Datatypes

Which thread is the funneled thread? 21.

22.A Change "INOUT" to "IN" for MPI Handle Parameters in several routines. Proposal A

H L R S

All official votes on Ballot 3 and 4 together

the listed All votes MPI-2.1

 2nd vote on Ballot 3 Ballot 3

1st vote on Ballot 4

are done together if nobody wants to vote on an item separately

Ballot 3

1. MPI COMM PARENT instead of MPI COMM GET PARENT

2. MPI UNPACK EXTERNAL

3. Additional C++ binding errors

4. MPI REQUEST CANCEL used where MPI CANCEL intended

→ MPI 2.2)

7. Error in MPI_Scan Example

8. Missing newline in Fortran binding

9. Misspelled argument in Fortran binding

10. Error in MPI-1. Example 3.12

11. Error in MPI-1, Example 3.34

12. Change MPI-2, page 343, lines 22-23 13. MPI 1.1, strlen in first pt-to-pt example

14. Formatting error on MPI 1.1, page 58

Ballot 4

Incorrect use of MPI IN PLACE in description of MPI ALLGATHER and MPI ALLGATHERV

2.a-g Edits to MPI-2 Chapter 4, Miscellany

Interlanguage use of Attributes and Error in Example 4.13 in MPI-2 (Use of Attributes in C and Fortran)

4. What info keys can be set?

5. Datatypes and MPI PROBE

6. MPI PROCNULL and RMA Part 2

(7.+9. → MPI 2.2)

Thread safety and collective communication

10.a-d Questions about Cart create & Cart coords

Questions a. Cart_create & Cart_coords (Typo)

11.a-c Questions about Graph create

Why no MPI INPLACE for MPI EXSCAN?

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19 20. Blocklengths of zero in MPI TYPE STRUCT

and in MPI Datatypes 21. Which thread is the funneled thread?

Change "INOUT" to "IN" for MPI Handle Parameters in several routines, Proposal A





Following slides: Results and votes from March 2008 Meeting



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MPI 2.1 - Ballot 4 - Goals

Scope of Effort:

- Clarification to the MPI standards document.
- i.e., defining the outcome of MPI function if the current definition is not clear.
- not changing the current definition if it is "bad" (→ this may be task of MPI 2.2-3.0)

Working plan for Ballot 4:

- · This meeting:
 - Monday: Official reading and discussion if necessary (2:00-3:00, 3:15-5:15pm)
 - If additional discussions are needed: 9:00-10:00pm
 - Tuesday: Official reading of such changes
 - Tuesday: Straw votes on each Ballot 4 item (9:00-9:30am).

Not later than Monday evening:

You all must be sure, that you will vote on each item with "yes" only, if it is okay!

If necessary, please discuss it with some members of the Forum. If there is still a problem, please contact me before Monday 9:00pm.



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Ballot 4 - Categories

- · Numbering of items: same as in http://www.cs.uiuc.edu/homes/wgropp/projects/parallel/MPI/mpi-errata/
- Sequence in the meeting:
 - (S) Small and simple (Items 1, 2a-g, 3, 5, 6, 10e, 11c, 12, 14, 15a-b, 18, 19, 21, X)
 - Long, but still simple (Items 4, 8, 13, 17, 20)
- R May be relevant for implementations that have not implemented this decision (Items 4, 6, 7, 9, 10a-d, 11a-b, 13, 16, 17, 18, 20)

Decisions are necessary (Items 7, 10a-d, 11a-b, 16, 22, 9)



Ballot 4, Item 1 – Incorrect use of MPI IN PLACE in description of MPI ALLGATHER and MPI ALLGATHERV

Question: Do you accept this entry?

Yes:

All=32 No:

Abstain:

Mail discussion, proposed by Jeff Squyres, Sep. 15, 2005 http://www.cs.uiuc.edu/homes/wgropp/projects/parallel/MPI/mpi-errata/discuss/gatherinplace/

On MPI-2.0, page 158, lines 25-31, remove the text

Specifically, the outcome of a call to MPI ALLGATHER in the "in place" case is as if all processes executed n calls to

```
MPI_GATHER( MPI_IN_PLACE, 0, MPI_DATATYPE_NULL, recvbuf,
     recvcount, recvtype, root, comm )
for root = 0, \ldots, n-1.
```

On MPI-2.0, page 159, lines 23-28, remove the text

Specifically, the outcome of a call to MPI ALLGATHER in the "in place" case is as if all processes executed n calls to

```
MPI_GATHERV( MPI_IN_PLACE, 0, MPI_DATATYPE_NULL, recvbuf,
recvcount, displs, recvtype, root, comm )
```

for root = $0, \ldots, n-1$.

Reason:

The text is wrong. The text is not needed because the outcome of MPI ALLGATHER and MPI ALLGATHERV is already specified and this analogy is not needed.



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Ballot 4, Item 2.a - Edits to MPI-2 Chapter 4, Miscellany

Question Do you accept this entry?

Yes:

ΑII No:

Abstain:

Mail discussion, proposed by Jesper Larson Traeff, Jan 09, 2008

http://www.cs.uiuc.edu/homes/wgropp/projects/parallel/MPI/mpi-errata/discuss/mpi2misc/

MPI-2.0, Sect. 4.1, page 37, **remove** line 44-46:

This is advice to implementors, rather than a required part of MPI-2. It is not suggested that this be the only way to start MPI programs. If an implementation does provide a command called mpiexec, however, it must be of the form described here.

Rationale for this remove:

It is largely a repetition of lines 34-36 (except the statement "It is not suggested that this be the only way..."): Instead, MPI specifies an mpiexec startup command and recommends but does not require it, as advice to implementors. However, if an implementation does provide a command called mpiexec, it must be of the form described below.



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Ballot 4, Item 2.b - Edits to MPI-2 Chapter 4, Miscellany

Question: Do you accept this

entry? Yes:

All-1 No:

Abstain:

Mail discussion, proposed by Jesper Larson Traeff, Jan 09, 2008 http://www.cs.uiuc.edu/homes/wgropp/projects/parallel/MPI/mpi-errata/discuss/mpi2misc/

MPI-2.0, Sect. 4.10, page 43, line 34:

Replace

"It consists of (key, value) pairs"

by

"It stores an unordered set of (key, value) pairs"

MPI-2.0, Sect. 4.10, page 43, line 34: Replace

"A key may have only one value."

by

"A key can have only one value."

Rationale: To emphasize that the info object is a kind of dictionary (data structure)



MPI 2.1 Slide 10

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Ballot 4, Item 2.c - Edits to MPI-2 Chapter 4, Miscellany

Question: Do you accept this entry?

Yes:

All=33 No:

Abstain:

Mail discussion, proposed by Jesper Larson Traeff, Jan 09, 2008 http://www.cs.uiuc.edu/homes/wqropp/projects/parallel/MPI/mpi-errata/discuss/mpi2misc/

MPI-2.0, Sect. 4.11, page 49, **add** after line 21:

/* no memory is allocated */

MPI-2.0, Sect. 4.11, page 49, add after line 22:

/* memory allocated */

Rationale:

To make consistent with Fortran example, add these comments.



Ballot 4, Item 2.d - Edits to MPI-2 Chapter 4, Miscellany

Question: Do you accept this entry?

Yes: all

No:

Abstain:

Mail discussion, proposed by Jesper Larson Traeff, Jan 09, 2008 http://www.cs.uiuc.edu/homes/wgropp/projects/parallel/MPI/mpi-errata/discuss/mpi2misc/

MPI-2.0, Sect. 4.12.2, page 50, line 9: Remove first

"in"

Reason: Typo



















Ballot 4, Item 2.e - Edits to MPI-2 Chapter 4, Miscellany

Question: Do you accept this entry?

Yes: AII-3

No:

Abstain:

Mail discussion, proposed by Jesper Larson Traeff, Jan 09, 2008

http://www.cs.uiuc.edu/homes/wgropp/projects/parallel/MPI/mpi-errata/discuss/mpi2misc/ MPI 2.0, Sect. 4.12.6, Exa. 4.12, page 55, line 21-22 read:

INTEGER TYPE, IERR

INTEGER (KIND=MPI_ADDRESS_KIND) ADDR

but should read:

INTEGER TYPE, IERR, AOBLEN(1), AOTYPE(1) INTEGER (KIND=MPI_ADDRESS_KIND) AODISP(1)

MPI 2.0, Sect. 4.12.6, Exa. 4.12, page 55, line 25-26 read:

CALL MPI GET ADDRESS (R, ADDR, IERR)

CALL MPI_TYPE_CREATE_STRUCT(1, 5, ADDR, MPI_REAL, TYPE, IERR)

but should read:

AOBLEN(1) = 5

CALL MPI_GET_ADDRESS(R, AODISP(1), IERR)

AOTYPE(1) = MPI REAL

CALL MPI TYPE CREATE STRUCT(1, AOBLEN, AODISP, AOTYPE, TYPE, IERR)



(Text changed in March 2008 meeting)

Reason: It was bad Fortran style and hard to read for C programmer.



Ballot 4, Item 2.g - Edits to MPI-2 Chapter 4, Miscellany

Question:

Do you accept this entry?

Yes:

No:

Abstain:

Mail discussion, proposed by Jesper Larson Traeff, Jan 09, 2008

http://www.cs.uiuc.edu/homes/wqropp/projects/parallel/MPI/mpi-errata/discuss/mpi2misc/ MPI-2.0, Sect. 4.12.6, page 56, line 29:

"assciated"

should be

"associated"

MPI-2.0, Sect. 4.14.5, page 74, line 9: Replace

"it erroneous"

by

"it is erroneous"

MPI-2.0, Sect. 5.3.2, page 85, line 25: Replace

"as the as the"

by

"as the"

Reason: Typos



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No:



Ballot 4, Item 2.f - Edits to MPI-2 Chapter 4, Miscellany

Mail discussion, proposed by Jesper Larson Traeff, Jan 09, 2008

Do you accept this

Question:

entry?

Yes:

AII-3

Abstain:

http://www.cs.uiuc.edu/homes/wgropp/projects/parallel/MPI/mpi-errata/discuss/mpi2misc/

MPI 2.0, Sect. 4.12.10, Exa. 4.14, page 60, line 31-32 read:

INTEGER TYPE, IERR, MYRANK

INTEGER (KIND=MPI_ADDRESS_KIND) ADDR

but should read:

INTEGER TYPE, IERR, MYRANK, AOBLEN(1), AOTYPE(1) INTEGER (KIND=MPI ADDRESS KIND) AODISP(1)

MPI 2.0, Sect. 4.12.10, Exa. 4.14, page 55, line 35-36 read:

CALL MPI GET ADDRESS (R, ADDR, IERR)

CALL MPI_TYPE_CREATE_STRUCT(1, 5, ADDR, MPI_REAL, TYPE, IERR)

but should read:

AOBLEN(1) = 5CALL MPI_GET_ADDRESS(R, AODISP(1), IERR)

AOTYPE(1) = MPI REAL

CALL MPI TYPE CREATE STRUCT(1, AOBLEN, AODISP, AOTYPE, TYPE, IERR)



(Text changed in March 2008 meeting)

Reason: It was bad Fortran style and hard to read for C programmer.



Ballot 4, Item 3 – Interlanguage use of Attributes and Error in Example 4.13 in MPI-2 (Use of Attributes in C and Fortran)

Question: Do you accept this entry?

Yes: all

No:

Abstain:

Mail discussion, proposed by Nicholas Nevin (Mar.25, 1999) et al. http://www.cs.uiuc.edu/homes/wgropp/projects/parallel/MPI/mpi-errata/discuss/getattr/

and Mail discussion, Jeff Squyres (June 24,2005) et al. http://www.cs.uiuc.edu/homes/wgropp/projects/parallel/MPI/mpi-errata/discuss/attrcandf/

Change MPI-2.0, Sect.4.12, page 58, line 36 reads:

IF (val.NE.5) THEN CALL ERROR

but should read

IF (val.NE.address of i) THEN CALL ERROR

Rationale for this modification:

MPI-2.0, Sect. 4.12, page 58, lines 12-13 and 16-18 clearly state that if an attribute is set by C, retrieving it in Fortran will obtain the address of the attribute.



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Question:

Do you

accept

entry?

this

Yes:

AII-2

No:

Abstain:

Ballot 4, Item 4 - What info keys can be set?

Mail discussion, proposed by Linda Stanberry and Bill Gropp, Jun 17, 1999 http://www.cs.uiuc.edu/homes/wgropp/projects/parallel/MPI/mpi-errata/discuss/infoset/ MPI 2.0. Sect. 4.10 Info Objects, page 43, line 38-40 read

If a function does not recognize a key, it will ignore it, unless otherwise specified. If an implementation recognizes a key but does not recognize the format of the corresponding value, the result is undefined.

but should read

An implementation must support info objects as caches for arbitrary (key, value) pairs, regardless of whether it recognizes the key. Each function that takes hints in the form of an MPI_Info must be prepared to ignore any key it does not recognize. This description of info objects does not attempt to define how a particular function should react if it recognizes a key but not the associated value, MPI INFO GET NKEYS, MPI INFO GET NTHKEY, MPI INFO GET VALUELEN, and MPI INFO GET must retain all (key, value) pairs so that layered functionality can also use the Info object.

Rationale for this clarification: The MPI-2.0 text allowed that also MPI INFO DELETE, MPI_INFO_SET, MPI_INFO_GET, and MPI_INFO_DUP could ignore (key, value) pairs that are not recognized in routines in other chapters that take hints with info arguments. The proposed clarification is necessary when we assume, that layered implementation of parts of the MPI-2 standard should be possible and may use the MPI Info objects for their needs. This was a goal of the MPI-2 Forum and the MPI-2.0 specification.

(Text changed in March 2008 meeting)





Ballot 4, Item 5 - Datatypes and MPI PROBE

Question: Do you accept this

entry?

Yes:

All-1

No:

Abstain:

Mail discussion, proposed by Patrick H. Worley, Nov 24, 1999, and Bill Gropp http://www.cs.uiuc.edu/homes/wgropp/projects/parallel/MPI/mpi-errata/discuss/probedatatype/ MPI 1.1. page 22. line 48 reads

used after a call to MPI_PROBE. (End of rationale.)

but should read

used after a call to MPI_PROBE or MPI_IPROBE. With a status returned from MPI PROBE or MPI IPROBE, the same dataypes are allowed as in a call to MPI RECV to receive this message. (End of rationale.)

Advice to users. The buffer size required for the receive can be affected by data conversions and by the stride of the receive datatype. In most cases, the safest approach is to use the same datatype with MPI GET COUNT and the receive. (End of advice to users.)

Rationale for this clarification: Reason for the first part: The current MPI-1.1 text says "The datatype argument should match the argument provided by the receive call that set the status variable." With MPI_PROBE, there isn't such a receive call.

Reason for the advice to users: It helps to write portable code. Because malloc needs a byte count, users may write wrong programs by using MPI BYTE.

(Text changed in March 2008 meeting)



Question

Do you accept

this entry?

Yes:

No:

Abstain:

Ballot 4, Item 6 – MPI PROCNULL and RMA Part 2

Mail discussion, proposed by Richard Treumann, Mar 8, 2002 http://www.cs.uiuc.edu/homes/wgropp/projects/parallel/MPI/mpi-errata/discuss/procnull2/

MPI-2, page 114, after line 4 (and after the lines added about MPI_PROC_NULL), add

After any RMA operation with rank MPI_PROC_NULL, it is still necessary to finish the RMA epoch with the synchronization method that started the epoch.

Reason: The behavior of one-sided RMA with target MPI PROC NULL was not clear.

(Text changed in March 2008 meeting)





Ballot 4, Item 7 - Significance of non-root arguments to MPI COMM SPAWN

Moved into MPI-2.2

Mail discussion, proposed by Jeff Squyres, Jul 30, 1999, Cra http://www.cs.uiuc.edu/homes/wgropp/projects/parallel/MPI/mpi-errata/discuss/spawn/

Proposal A

Add in MPI-2.0, page 88, after line 24:

Advice to users. If the non-root processes do not use MPI ERRCODES IGNORE, then they have to allocate the appropriate number of entries (see maxproc at root) in the array_of_errcodes although the maxproc argument is unused in non-root processes. It is allowed to use an array of erroodes at some of the calling processes and MPI_ERRCODES_IGNORE at some others. (End of advice to users.)

Rationale for this clarification:

It was not clear that maxproc is significant as input argument only at root while it is needed at all processes to define the length of array_of_errcodes. It was not explicitly forbidden that MPI_ERRCODES_IGNORE is used only at some processes. And there isn't a general rule that all arguments must be the same.





















Question: Which proposal is better?

Abstain:

Question: Do you accept proposal Yes:

Abstain:

Ballot 4, Item 7 - Significance of non-root arguments to MPI COMM SPAWN

Moved into MPI-2.2

Mail discussion, proposed by Jeff Squyres, Jul 30, 1999, et al.

http://www.cs.uiuc.edu/homes/wgropp/projects/parallel/MPI/mpi-errata/discuss/spawn/ Proposal B

MPI 2.0 page 84 line 45 reads (in MPI COMM SPAWN):

OUT array_of_errcodes one code per process (array of integer)

but should read:

OUT array_of_errcodes one code per process (array of integer,

significant only at root)

MPI 2.0 page 89 line 42 reads (in MPI COMM SPAWN MULTIPLE):

OUT array_of_errcodes one code per process (array of integer)

but should read:

OUT array_of_errcodes one code per process (array of integer, significant only at root)

Comment: This proposal modifies the MPI interface. User codes may be broken. Another reason, not to do this modification, is that the non-root processes have no chance (in error-return-mode) to detect an error. And after an error. MPI does not guarantee that MPI communication still works. Only MPI Abort should be guaranteed to work.





Ballot 4. Item 8 -Thread safety and collective communication

Question:

Do you accept this entry?

Yes: AII-4

No:

Abstain:

Mail discussion, proposed by Karl Feind, Nov 16, 1998, et al.

http://www.cs.uiuc.edu/homes/wgropp/projects/parallel/MPI/mpi-errata/discuss/thread-safety/ Add new paragraphs after MPI-2, 8.7.2 page 195 line 9 (the end of the clarification on "Collective calls"):

Advice to users. With three concurrent threads in each MPI process of a communicator comm. it is allowed that thread A in each MPI process calls a collective operation on comm, thread B calls a file operation on an existing filehandle that was formerly opened on comm. and thread C invokes one-sided operations on an existing window handle that was also formerly created on comm. (End of advice to users.)

Rationale. As already specified in MPI_FILE_OPEN and MPI_WIN_CREATE, a file handle and a window handle inherit only the group of processes of the underlying communicator, but not the communicator itself. Accesses to communicators, window handles and file handles cannot affect one another. (End of rationale.)

Advice to implementors. If the implementation of file or window operations internally uses MPI communication then a duplicated communicator may be cached on the file or window object. (End of advice to implementors.)



Rationale for this clarification:

The emails have shown, that the current MPI-2 text can be misunderstood.





Question Do you accept entry? Yes:

No: Abstain: Ballot 4, Item 9 - const in C++ specification of predefined MPI objects (was just datatypes)

Item 22.A: INOUT → Variant A/D: Same as voted for Bal. 3

Mail discussion, by Richard Treumann and Rolf Rabenseifner Jun 13 - Jul 26, 2001 http://www.cs.uiuc.edu/homes/wgropp/projects/paralle Moved into MPI-2.2

- MPI-2, page 345, line 37: **Remove** the const from **const MPI::Op**.
- MPI-2, page 346, line 20: **Remove** the const from **const MPI::Group**.
- MPI-2, page 346, add after line 34:

Advice to implementors: If an implementation does not change the value of predefined handles while execution of MPI Init, the implementation is free to define the predefined operation handles as const MPI::Op and the predefined group handle MPI::GROUP_EMPTY as const MPI::Group. Other predefined handles must not be "const" because they are allowed as INOUT argument in the MPI COMM SET NAME/ATTR and MPI TYPE SET NAME/ATTR routines. (End of advice to implementors.)

- Reason: MPI Init may change the predefined handles, because MPI 1.1, page 10, lines 9-10 says: "Opaque objects accessed by constant handles are defined and do not change value between MPI initialization (MPI_INIT() call) and MPI completion (MPI_FINALIZE() call)." Therefore they must not be defined as const in the MPI standard.
- Alternative A: Keep INOUT for handles with call by value in C/C++ In this case, the text above should be used (as already proposed for Ballot 3)

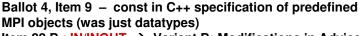


Question: Do you accept this entry?

Yes:

No:

Abstain:



Item 22.B : IN/INOUT → Variant B: Modifications in Advice

Mail discussion, by Richard Treumann and Rolf Rabenseifner Jun 13 - Jul 26, 2001 http://www.cs.uiuc.edu/homes/wgropp/projects/paralle Moved into MPI-2.2

- MPI-2, page 345, line 37: **Remove** the const from **const MPI::Op**.
- MPI-2, page 346, line 20: **Remove** the const from **const MPI::Group**.
- MPI-2, page 346, add after line 34:

Advice to implementors: If an implementation does not change the value of predefined handles while execution of MPI Init, the implementation is free to define the predefined operation handles as const MPI::Op and the predefined group handle MPI::GROUP EMPTY as const MPI::Group. Other predefined handles must not be "const" because they are allowed as INOUT argument inthe MPI COMM SET NAME/ATTR and MPI TYPE SET NAME/ATTR routines. (End of advice to implementors.)



Alternative B: Change into IN/INOUT (i.e. mark handles as IN and referenced opaque object as INOUT) In this case, **remove** the last sentence from the proposal:

"Other predefined handles must not be "const" because they are allowed as INOUT argument in the MPI COMM SET NAME/ATTR and MPI TYPE SET NAME/ATTR routines."



Question Do you accept this entry? Yes:

No: Abstain:

Ballot 4, Item 9 - const in C++ specification of predefined MPI objects (was just datatypes)

Item 22.B: IN → Variant C: Modifications in Advice

Mail discussion, by Richard Treumann and Rolf Rabenseifner Jun 13 - Jul 26, 2001 http://www.cs.uiuc.edu/homes/wgropp/projects/parallel Moved into MPI-2.2

- MPI-2, page 345, line 37: Remove the const from const MPI::Op.
- MPI-2, page 346, line 20: Remove the const from const MPI::Group.
- MPI-2, page 346, add after line 34:

Advice to implementors: If an implementation does not change the value of predefined handles while execution of MPI Init, the implementation is free to define the predefined operation handles as const MPI::Op and the predefined group handle MPI::GROUP EMPTY as const MPI::Group. Other predefined handles must not be "const" because they are allowed as INOUT argument in the MPI COMM SET NAME/ATTR and MPI TYPE SET NAME/ATTR routines. (End of advice to implementors.)

Alternative C: Change into IN (i.e. mark handles as IN and forget about referenced opaque object as) In this case, **remove** the last sentence from the proposal:

"Other predefined handles must not be "const" because they are allowed as INOUT argument in the MPI_COMM_SET_NAME/ATTR and MPI TYPE SET NAME/ATTR routines."



Ballot 4. Item 10.a -Questions about Cart create, and Cart coords

Mail discussion, proposed by Jesper L. Traeff, Mar 4, 2002, and Rolf Rabenseifner http://www.cs.uiuc.edu/homes/wgropp/projects/parallel/MPI/mpi-errata/discuss/topo/

Here is a proposal about handling zero-dimensional Cartesian communicators that are produced with MPI_Cart_sub if all remain_dims are false.

MPI-1.1 Sect.6.5.4, page 187, line 42 (end of definition of MPI_Cart_sub) reads

(This function is closely related to MPI COMM SPLIT.)

but should read

If all entries in remain_dims are false or comm is already associated with a zero-dimensional Cartesian topology then newcomm is associated with a zero-dimensional Cartesian topology. (This function is closely related to MPI COMM SPLIT.)







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Ballot 4. Item 10.b -Questions about Cart create, and Cart coords

MPI-1.1 Sect.6.5.4. page 183. add at end of lines 30 (definition of MPI Cartdim get and MPI Cart get):

If comm is associated with a zero-dimensional Cartesian topology. MPI Cartdim get returns ndims=0 and MPI Cart get will keep all output arguments unchanged.

MPI-1.1 Sect.6.5.4, page 184, add a new paragraph after line 23 (definition of MPI Cart rank):

If comm is associated with a zero-dimensional Cartesian topology, coord is not significant and 0 is returned in rank.

MPI-1.1 Sect.6.5.4, page 184, add a new paragraph after lines 39 (definition of MPI_Cart_coords):

If comm is associated with a zero-dimensional Cartesian topology, coords will be unchanged.



Ballot 4. Item 10.c -

Questions about Cart create, and Cart coords

Question:

Which alternat. is better?

A: all

B:

Abstain:

Alternative A:

MPI-1.1 Sect.6.5.5, page 186, after line 47 (end of definition of MPI_Cart_shift), the following paragraph is added:

It is erroneous to call MPI CART SHIFT with a direction that is either negative or greater than or equal to the number of dimensions in the Cartesian communicator. This implies that it is erroneous to call MPI_CART_SHIFT with a comm that is associated with a zero-dimensional Cartesian topology.

Alternative B:

MPI-1.1 Sect.6.5.5, page 186, after line 47 (end of definition of MPI Cart shift), the following paragraph is added:

If comm is associated with a zero-dimensional Cartesian topology, then the input arguments direction and disp are ignored and always MPI_PROC_NULL is returned in rank_source and rank_dest.

(End of Alternative A and B)

Rationale: A is directly in the spirit of the current definition.

















Question Do you accept total Item 10.a-d? Yes: AII-1 No:

Abstain:

Ballot 4. Item 10.d -Questions about Cart create, and Cart coords

MPI-1.1 Sect.6.5.1, page 179, lines 29-30 (end of definition of MPI_Cart_create) reads

The call is erroneous if it specifies a grid that is larger than the group size. but should read

If ndims is zero then a zero-dimensional Cartesian topology is created. The call is erroneous if it specifies a grid that is larger than the group size or if ndims is negative.

(Text changed in March 2008 meeting)



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Ballot 4, Item 10.a-d - Background information (1)

This report shows significant differences of MPI implementations in the (extreme) case of zero-dimensional topologies!

Therefore clarifications may be necessary.

(Although hopefully never an application may produce zerodimensional Cartesian topologies!)

If there are no objections, I would produce clarifications for MPI-2.1 according to the behavior of mpich2 and MPI from IBM (except the rank -767705708, see below).



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Ballot 4, Item 10.a-d - Background information (2)

MPI-1.1 Sect. 6.5.6, page 187, routine MPI_CART_SUB defines on lines 38-42:

If a cartesian topology has been created with MPI CART CREATE, the function MPI CART SUB can be used to partition the communicator group into subgroups that form lower-dimensional cartesian subgrids, and to build for each subgroup a communicator with the associated subgrid cartesian topology. (This function is closely related to MPI COMM SPLIT.)

The text clearly says, that the new communicator must

- (1) have a Cartesian topology associated
- (2) be lower-dimensional in the case of a subgrid.

There is no restriction on the input array remain dims. Therefore all MPI implementations (that I tested) allow that all entries in remain dims are "false".



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Ballot 4, Item 10.a-d - Background information (3)

I tested several MPI libraries with creating a subgrid with MPI Cart sub(remain_dims=0) from a 1-dim Cartesian topology.

I tested the subgrid communicator with MPI Topo test.

In the case of MPI CART, I used MPI Cartdim get to retrieve ndims, and MPI Cart get for further details:

- mpich2 (1.0.3):
 - MPI CART. ndims=0.
 - MPI_Cart_get works, but keeps all OUT arguments unchanged
- IBM (on SP):
 - MPI CART. ndims=0.
 - MPI Cart get works, but keeps all OUT arguments unchanged
- OpenMPI 1.2.4:
 - MPI CART. ndims=1.
 - MPI Cart get works and returns dims=1, periods=0, coords=0 independent from process or periods in the original comm (may be wrong because (2) is not fulfilled).
- NEC MP/EX: not MPI_CART (may be wrong because (1) is not fulfilled)
- Voltaire: not MPI_CART (mpich1) (may be wrong because (1) is not fulfilled)



Ballot 4, Item 10.a-d - Background information (4)

With the implementations that return a correct zero-dim Cartesian topology. I tested further usage of this zero-dim communicator:

- MPI Comm size returns 1 and MPI Comm rank returns 0 because this communicator is like MPI COMM SELF, but with Cartesian topology associated.
- MPI Cart rank(IN ZeroDimComm, IN coords=0, OUT rank) Rationale. A zero-dim communicator has zero coords, i.e., this routine should not examine the coords input argument.
 - mpich2: rank = 0 is returned (may be because this is the only existing rank in this communicator, this value may make sense, independent of the coord, that should not be analyzed)
 - IBM: rank = -767705708 is returned (strange value, not MPI PROC NULL, not MPI UNDEFINED)



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Ballot 4, Item 10.a-d - Background information (5)

- MPI Cart coords(IN ZeroDimComm, IN rank=0, OUT coords) Rationale. A zero-dim communicator has zero coords, i.e., this routine should not return anything in the coords output argument.
 - mpich2 and IBM: coords is not modified (as expected)
- MPI Cart sub(IN ZeroDimComm, IN remain dims=0, OUT subsubcomm) Rationale. A zero-dim communicator has zero dimensions, i.e., this routine should not examine remain dims and the returned communicator should be again a zerodim Cartesian communicator.
 - mpich2 and IBM: subsubcomm is a zero-dim Cartesian communicator (as
- MPI Cart shift(IN ZeroDimComm, IN direction=0, IN disp=1, OUT src, OUT dest) Rationale. This call is erroneous because in a zero-dim communicator, the direction=0 does not exist.
 - mpich2 and IBM: They detect the error and abort.

(In OpenMPI, all these calls work as expected on a 1-dimensional topology on a "MPI COMM SELF")



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Ballot 4, Item 10.a-d - Background information (6)

The last test is not addressed by MPI-1.1:

Is it possible to build a zero-dim Cartesian topology directly by calling MPI Cart create:

- MPI Cart create(IN MPI COMM SELF, IN ndims=0, IN dims=1, IN Periods, IN reorder, OUT ZeroDimComm)
 - Results: All tested MPI implementations return an error and abort. (Same on MPI COMM WORLD)



Ballot 4. Item 10.e -Questions about Cart create, and Cart coords

Question: Do you accept this typo correction Item 10.e? Yes:

all

No:

Abstain:

Typo correction:

MPI-1.1 Sect.6.5.4. page 184. lines 30 (definition of MPI Cart coords) reads

length of vector coord in the calling program (integer) IN maxdims but should read (missing "s" at coords)

Mail discussion, proposed by Jesper L. Traeff, Mar 4, 2002, and Rolf Rabenseifner

http://www.cs.uiuc.edu/homes/wgropp/projects/parallel/MPI/mpi-errata/discuss/topo/

maxdims length of vector coords in the calling program (integer)



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Question

Do you

accept

entry?

this

Yes:

AII-1

No:

Abstain:

Ballot 4, Item 11.a - Questions about Graph create

Mail discussion, proposed by Jesper L. Traeff, Mar 4, 2002, and Rolf Rabenseifner http://www.cs.uiuc.edu/homes/wgropp/projects/parallel/MPI/mpi-errata/discuss/topo/

MPI-1.1, Sect. 6.5.3, page 181, line 1-3 read:

If the size, nnodes, of the graph is smaller than the size of the group of comm. then some processes are returned MPI COMM NULL, in analogy to MPI CART CREATE and MPI COMM SPLIT.

but should read

If the size, nnodes, of the graph is smaller than the size of the group of comm, then some processes are returned MPI_COMM_NULL, in analogy to MPI_CART_CREATE and MPI_COMM_SPLIT. If the graph is empty, i.e., nnodes == 0, then MPI_COMM_NULL is returned in all processes.

Rationale for this clarification:

As in MPI_COMM_CREATE, empty groups are allowed, but empty groups are described here in a different way, and should be mentioned explicitly therefore.

Empty graphs have been allowed in the past



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Ballot 4, Item 11.b - Questions about Graph create

Question:

Do you accept this entry?

Yes: AII-2=30

No:

Abstain:

After MPI-1.1, Sect. 6.5.3, page 181, line 35, the following paragraph should be added:

A single process is allowed to be defined multiple times in the list of neighbors of a process (i.e., there may be multiple edges between two processes). A process is also allowed to be a neighbor to itself (i.e., a self loop in the graph). The adjacency matrix is allowed to be non-symmetric.

Advice to users. Performance implications of using multiple edges or a nonsymmetric adjacency matrix are not defined. The definition of a nodeneighbor edge does not imply a direction of the communication. (End of advice to users.)

Rationale for this clarification:

The Example 6.3, MPI-1.1, page 15, line 29 - page 186, line 13, clearly shows multiple edges between nodes and self loops: the two (multiple) self-loops of node 0 and of node 7. It is nowhere forbidden, that the graph has edges only in one

(Text modified at March 2008 meeting)



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S

Ballot 4, Item 11.c - Questions about Graph create

Question: Do you accept

this

entry?

Yes:

All=32

Abstain:

No:

After MPI-1.1, Sect. 6.4, page 178, end of the sentence on lines 6-7, the following sentence should be added:

All input arguments must have identical values on all processes of the group of comm_old.

Rationale for this clarification:

This statement is missing.



Ballot 4. Item 12 -Why no MPI INPLACE for MPI EXSCAN?

Question: Do you accept this entry?

AII-1=30

No:

Yes:

Abstain:

Mail discussion, proposed by Falk Zimmermann and Tony Skjellum, Mar 19, 1999 http://www.cs.uiuc.edu/homes/wgropp/projects/parallel/MPI/mpi-errata/discuss/exscan/

MPI-2, Sect. 7.3.6, page 167, lines 6-8 read:

The reason that MPI-1 chose the inclusive scan is that the definition of behavior on processes zero and one was thought to offer too many complexities in definition, particularly for user-defined operations. (End of rationale.)

but should read:

No in-place version is specified for MPI_EXSCAN because it is not clear what this means for the process of rank zero.

The reason that MPI-1 chose the inclusive scan is that the definition of behavior on processes zero and one was thought to offer too many complexities in definition, particularly for user-defined operations. (End of rationale.)













Question Do you accept

entry? Yes:

this

AII-1=32 No:

Abstain:

Ballot 4. Item 13 -MPI GET PROCESSOR NAME and Fortran

Mail discussion, proposed by William Gropp et al., Dec 09, 1999

http://www.cs.uiuc.edu/homes/wgropp/projects/parallel/MPI/mpi-errata/discuss/procname/ Add the following clarification to the current interface definitions of MPI GET PROCESSOR NAME and MPI_COMM_GET_NAME.

MPI 1.1, Sect. 7.1, routine MPI GET PROCESSOR NAME, page 193, add after line 20:

In C, a null character is additionally stored at name [resultlen]. resultlen cannot be larger then MPI_MAX_PROCESSOR_NAME-1. In Fortran, name is padded on the right with blank characters. resultlen cannot be larger then MPI MAX PROCESSOR NAME.

MPI-1.1, Sect. 7.1, page 193, beginning of line 29 reads

examine the ouput argument

but should read (additional t in output)

(Text modified at March 2008 meeting)

examine the output argument

MPI 2.0, Sect. 8.4, routine MPI_COMM_GET_NAME, page 178, add after line 48:

In C, a null character is additionally stored at name [resultlen]. resultlen cannot be larger then MPI MAX OBJECT NAME-1. In Fortran, name is padded on the right with blank characters. resultlen cannot be larger then MPI MAX OBJECT NAME.



There is additional information, including results of checking the behavior of many MPI implementations, in the mail discussion. The test programs and some results are also available here:

- mpi comm get name test protocol.txt, mpi comm get name test.c, mpi comm get name test.f
- mpi get processor name test.c, mpi get processor name test.f, mpi get processor name test protoco





Ballot 4, Item 14 – Interpretation of user defined datarep

Question: Do you

accept this entry? Yes:

AII=33

No:

Abstain:

Mail discussion, proposed by Hubert Ritzdorf et al., Feb 02, 1999 http://www.cs.uiuc.edu/homes/wgropp/projects/parallel/MPI/mpi-errata/discuss/datarep/

MPI-2.0 Sect. 9.5.3 User-defined Data Representations, page 254, lines 13-15 read:

Then in subsequent calls to the conversion function, MPI will increment the value in position by the count of items converted in the previous call.

but should read:

Then in subsequent calls to the conversion function, MPI will increment the value in position by the count of items converted in the previous call, and the userbuf pointer will be unchanged.

Rationale for this clarification:

It was not clear, whether the userbuf pointer must also be moved in the subsequent calls. This clarification was already done in 1999 and should already be implemented in existing implementations of user-defined data representations.

(Text changed in March 2008 meeting)



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Question:

Do you

accept

entry?

this

Yes:

No:

Abstain:

Ballot 4, Item 15.a - MPI Abort

Mail discussion, proposed by William Gropp et al., Jan 31, 2001 http://www.cs.uiuc.edu/homes/wgropp/projects/parallel/MPI/mpi-errata/discuss/abort/

MPI-1.1 Sect. 7.5, MPI Abort, page 200, lines 23-26 read:

This routine makes a "best attempt" to abort all tasks in the group of comm. This function does not require that the invoking environment take any action with the error code. However, a Unix or POSIX environment should handle this as a return errorcode from the main program or an abort (errorcode).

but should read (" or an abort(errorcode)" removed):

This routine makes a "best attempt" to abort all tasks in the group of comm. This function does not require that the invoking environment take any action with the error code. However, a Unix or POSIX environment should handle this as a return errorcode from the main program.

Rationale for this clarification:

POSIX defines void abort(void). The routine void exit(int status) may be used to implement "handle this as a return errorcode from the main program". abort(errorcode) was not substituted by exit(errorcode) because this is technically not enough, if the MPI implementation wants to return it also from mpiexec, see next proposal.



Ballot 4, Item 15.b - MPI Abort

Question: Do you accept this entry?

Yes: all

No:

Abstain:

Mail discussion, proposed by William Gropp et al., Jan 31, 2001 http://www.cs.uiuc.edu/homes/wgropp/projects/parallel/MPI/mpi-errata/discuss/abort/

MPI-1.1 Sect. 7.5, MPI Abort, page 200, add after line 34 (end of rationale):

Advice to users. Whether the errorcode is returned from the executable or from the MPI process startup mechanism (e.g., mpiexec), is an aspect of quality of the MPI library but not mandatory. (End of advice to users.)

Advice to implementors. Where possible, a high quality implementation will try to return the errorcode from the MPI process startup mechanism (e.g. mpiexec or singleton init). (End of advice to implementors.)

Rationale for this clarification:

The intent of word "should" in "should handle this as a return errorcode from the main program" is only a quality of implementation aspect and not a must. This was not clear and could be misinterpreted.













Question:

Do you

accept

entry?

this

Yes:

AII-5

No:

Abstain:

Ballot 4, Item 16 - MPI Type create f90 real etc.

Mail discussion, proposed by Nicholas Nevin et al., Feb 09, 2001

http://www.cs.uiuc.edu/homes/wgropp/projects/parallel/MPI/mpi-errata/discuss/typef90real/ Problem:

An application may repeatedly call (probably with same (p,r) combination) the MPI TYPE CREATE F90 xxxx routines.

Proposal: Add after MPI-2.0 Sect. 10.2.5. MPI TYPE CREATE F90 xxxx. page 295. line 47 (End of advice to users.):

Advice to implementors. An application may often repeat a call to MPI_TYPE_CREATE_F90_xxxx with the same combination of (xxxx,p,r). The application is not allowed to free the returned predefined, unnamed datatype handles. To prevent the creation of a potentially huge amount of handles, the MPI implementation should return the same datatype handle for the same (REAL/COMPLEX/INTEGER, p, r) combination. Checking for the combination (p,r) in the preceding call to MPI_TYPE_CREATE_F90_xxxx and using a hash-table to find formerly generated handles should limit the overhead of finding a previously generated datatype with same combination of (xxxx,p,r). (End of advice to implementors.)

Rationale for this clarification: (Text changed in March 2008 meeting)

Currently most MPI implementations are handling the MPI_TYPE_CREATE_F90_xxxx functions wrong or not with the requested quality.





Question:

this

entry?

All-1=31

Abstain:

Yes:

No:

Ballot 4, Item 17 - MPI File get info

Mail discussion, proposed by Rolf Rabenseifner, Jan. 31, 2008 http://www.cs.uiuc.edu/homes/wgropp/projects/parallel/MPI/mpi-errata/discuss/filegetinfonull/

Do you Add text stating that the return value can be null: accept

MPI-2.0 Sect. 9.2.8, File Info, page 219, lines 11-13 read:

MPI FILE GET INFO returns a new info object containing the hints of the file associated with fh. The current setting of all hints actually used by the system related to this open file is returned in info_used.

The user is responsible for freeing info used via MPI INFO FREE.

MPI FILE GET INFO returns a new info object containing the hints of the file associated with fh. The current setting of all hints actually used by the system related to this open file is returned in info_used.

If no such hints exist, a handle to a newly created info object is returned that contains no key/value pair.

The user is responsible for freeing info_used via MPI_INFO_FREE.

Rationale for this clarification:

(Text modified at March 2008 meeting)

This text was missing. It was not clear, whether a MPI Info handle would be returned that would return nkeys=0 from MPI INFO GET NKEYS. From user's point of view, this behavior might have been expected without this clarification. For most implementations, this clarification is irrelevant because they always return several default hints, e.g., the filename.





Question:

Do you

accept

entry?

AII-1=32

Abstain:

this

Yes:

No:

Ballot 4, Item 18 - MPI File set view

Mail discussion, proposed by Rolf Rabenseifner, Jan. 31, 2008

http://www.cs.uiuc.edu/homes/wgropp/projects/parallel/MPI/mpi-errata/discuss/filesetview/

Does the info passed to MPI File set view replace all of the previous info keys? (The answer given in this clarification is "no".)

Proposal: Add in MPI-2.0 Sect. 9.2.8, File Info, page 218, after line 18 the following sentences:

When an info object that specifies a subset of valid hints is passed to MPI_FILE_SET_VIEW or MPI_FILE_SET_INFO, there will be no effect on previously set or defaulted hints that the info does not specify.

Rationale for this clarification:

This text was missing. It was not clear, whether an info object in MPI_FILE_SET_VIEW and MPI FILE SET INFO was intended to replace only the mentioned hints or was intended to substitute a complete new set of hints for the prior set.









Question:

Do you

accept

entry?

Yes:

AII=33

Abstain:

No:

this

Ballot 4, Item 19 - MPI IN PLACE for MPI Reduce scatter

Mail discussion, proposed by Rajeev Thakur, Oct. 10, 2002

http://www.cs.uiuc.edu/homes/wgropp/projects/parallel/MPI/mpi-errata/discuss/redscat/

MPI-2.0. Sect. 7.3.3. routine MPI REDUCE SCATTER, page 163, delete the sentence on line 19-20:

Note that the area occupied by the input data may be either longer or shorter than the data filled by the output data.

Rationale for this clarification:

The sentence makes no sense because the input data can never be shorter than the output data. The output, determined by recycounts[il, is a subset of the input.











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Ballot 4, Item 20 - Blocklengths of zero in MPI TYPE STRUCT and in MPI Datatypes

Mail discussion, proposed by Jesper Larsson Traeff, April 23, 2007 http://www.cs.uiuc.edu/homes/wgropp/projects/parallel/MPI/mpi-errata/discuss/blklenzero/ And Mail discussion, proposed by Peter Ganster and Bill Gropp, Jul. 25 2001 http://www.cs.uiuc.edu/homes/wgropp/projects/parallel/MPI/mpi-errata/discuss/structblocklen/

Problem:

Blocklengths of zero are allowed. Do we need to add a statement to this effect? The mail threads contain a deeper discussion of the implications of zero-length blocks in MPI datatypes.

Rationale for the clarification and modification on the following slides:

The outcome of zero-count entries in the type map was not defined. For this, a clarification was needed. The interfaces of MPI_TYPE_CREATE_HINDEXED and MPI_TYPE_CREATE_STRUCT was inconsistent to the rest derived datatype routines. This was probably due to editing errors. A meaning of negative values was never defined not intended. Therefore, portable applications could not use negative values. These editing errors are fixed by this proposal.



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Ballot 4, Item 20 - Blocklengths of zero in MPI TYPE STRUCT and in MPI Datatypes

Proposal:

Add the following paragraph in MPI 1.1, Sect. 3.12, page 62, after line 2 (i.e., after ... "of the types defined by Typesig."):

Most datatype constructors have replication count or block length arguments. Allowed values are nonnegative integers. If the value is zero, no elements are generated in the type map and there is no effect on datatype bounds or extent.

MPI 1.1, Sect 3.12.1, MPI_TYPE_HINDEXED, page 67, line 22-24 read:

IN count number of blocks - also number of entries in

array of displacements and array of blocklengths

(integer)

but should read:

IN count number of blocks - also number of entries

array of displacements and array of blocklengths

(nonnegative integer)



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Ballot 4, Item 20 - Blocklengths of zero in MPI TYPE STRUCT and in MPI Datatypes

MPI 1.1, Sect 3.12.1, MPI TYPE STRUCT, page 68, line 19-22 read:

IN count number of blocks (integer) - also

number of entries in arrays array_of_types,

array_of_displacements and array of blocklengths

IN array_of_blocklength number of elements in each block

(array of integer)

but should read:

IN count number of blocks (nonnegative integer) - also

number of entries in arrays array of types,

array_of_displacements and array_of_blocklengths

IN array_of_blocklength number of elements in each block

(array of nonnegative integer)







Rolf Rabenseifner



Ballot 4, Item 20 - Blocklengths of zero in MPI TYPE STRUCT and in MPI Datatypes

MPI 2.0, Sect 4.14.1, MPI TYPE CREATE HINDEXED, page 66, line 36-38 read:

number of blocks - also number of entries in IN count

array of displacements and array of blocklengths

(integer)

but should read:

IN count number of blocks - also number of entries

array of displacements and array of blocklengths

(nonnegative integer)



Ballot 4, Item 20 - Blocklengths of zero in MPI TYPE STRUCT and in MPI Datatypes

Question

Do you accept this Item 20?

Yes:

all No:

Abstain:

MPI 2.0, Sect 4.14.1, MPI TYPE CREATE STRUCT, page 67, line 14-18 read:

number of blocks (integer) - also

number of entries in arrays array_of_types,

array_of_displacements and array_of_blocklengths

IN array_of_blocklength number of elements in each block

(array of integer)

but should read:

number of blocks (nonnegative integer) - also IN count

number of entries in arrays array_of_types, array of displacements and

array_of_blocklengths

IN array_of_blocklength number of elements in each block

(array of nonnegative integer)



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Ballot 4, Item 21 - Which thread is the funneled thread?

MPI THREAD FUNNELED The process may be multi-threaded.

http://www.cs.uiuc.edu/homes/wgropp/projects/parallel/MPI/mpi-errata/discuss/funneled/

Mail discussion, proposed by Grea Lindahlet al., Nov. 30, 2007

but only the main thread will make MPI calls

(all MPI calls are "funneled" to the main thread).

MPI-2.0 Sect. 8.7.3. MPI Init thread, page 196, lines 25-26 read:

Question:

Do you accept this entry?

Yes: AII=33

No:

Abstain:

but should read:

MPI THREAD FUNNELED The process may be multi-threaded. but the application must ensure that only the main thread makes MPI calls (for the definition of main thread, see MPI_IS_THREAD_MAIN).

Rationale for this clarification:

The existing document doesn't make it clear that the MPI user has to funnel the calls to the main thread; it's not the job of the MPI library. I have seen multiple MPI users confused by this issue, and when I first read this section, I was confused by it, too.



MPI 2.1 Slide 54

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Ballot 4, Item 22 - Change "INOUT" to "IN" for MPI Handle Parameters in several routines (1)

Mail discussion, proposed by Jeff Squyres and Rolf Rabenseifner, Jan. 21, 2008 http://www.cs.uiuc.edu/homes/wgropp/projects/parallel/MPI/mpi-errata/discuss/inout/ Problem:

There is a inconsistency between the INOUT description for handle arguments, and their usage in the language independent definitions of MPI-1.1 and MPI-2.0. There are 3 possibilities to solve this. The MPI Forum should decide, which possibility is the best.



Ballot 4, Item 22 - Change "INOUT" to "IN" for MPI Handle Parameters in several routines (2)

Background:

MPI 2.0 Sect. 2.3 Procedure Specification, page 6 lines 30-34 read:

There is one special case - if an argument is a handle to an opaque object (these terms are defined in Section 2.5.1), and the **object is updated** by the procedure call, then the argument is marked **OUT**. It is marked this way **even though the** handle itself is not modified - we use the OUT attribute to denote that what the handle references is updated. Thus, in C++, IN arguments are either references or pointers to const objects.

Example: MPI 1.1, page 171, lines 26-34:

MPI ATTR PUT(comm, keyval, attribute val)

IN communicator to which attribute will be attached (handle) IN kevval key value, as returned by MPI_KEYVAL_CREATE (integer) attribute val attribute value

int MPI_Attr_put(MPI_Comm comm, int keyval, void* attribute_val)

Example: MPI 2.0, page 44, lines 36-43:

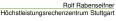
MPI INFO_SET(info, key, value)

info object (handle) INOUT info IN key (string)

call by value (IN) The object behind is modified (INOUT)

Handle with

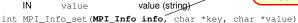














Ballot 4, Item 22 - Change "INOUT" to "IN" for MPI Handle Parameters in several routines (3)

List of routines with IN handles and INOUT objects:

```
MPI_ATTR_PUT, MPI_ATTR_DELETE,
 MPI_ERRHANDLER_SET
MPT 2.0:
 MPI INFO SET,
                             MPI INFO DELETE,
 MPI_COMM_SET_ERRHANDLER,
                             MPI_TYPE_SET_ERRHANDLER,
 MPI_WIN_SET_ERRHANDLER,
 MPI_GREQUEST_COMPLETE,
 MPI_COMM_SET_NAME,
                        MPI_TYPE_SET_NAME,
                                               MPI_WIN_SET_NAME,
 MPI_COMM_SET_ATTR,
                        MPI_TYPE_SET_ATTR,
                                               MPI_WIN_SET_ATTR,
 MPI_COMM_DELETE_ATTR, MPI_TYPE_DELETE_ATTR, MPI_WIN_DELETE_ATTR,
  MPI FILE SET SIZE,
                        MPI FILE PREALLOCATE,
 MPI_FILE_SET_INFO,
                        MPI_FILE_SET_VIEW,
 MPI FILE WRITE AT.
                        MPI_FILE_WRITE_AT_ALL, MPI_FILE_IWRITE_AT,
 MPI_FILE_READ,
                        MPI_FILE_READ_ALL,
 MPI FILE WRITE,
                        MPI_FILE_WRITE_ALL,
 MPI_FILE_IREAD,
                        MPI_FILE_IWRITE,
 MPI_FILE_SEEK,
 MPI FILE READ SHARED,
                                MPI FILE WRITE SHARED,
 MPI_FILE_IREAD_SHARED,
                                MPI_FILE_IWRITE_SHARED,
 MPI_FILE_READ_ORDERED,
                                MPI_FILE_WRITE_ORDERED,
  MPI_FILE_SEEK_SHARED,
                                MPI_FILE_WRITE_AT_ALL_END,
  MPI_FILE_WRITE_AT_ALL_BEGIN,
  MPI FILE READ ALL BEGIN,
                                MPI FILE READ ALL END,
 MPI_FILE_WRITE_ALL_BEGIN,
                                MPI_FILE_WRITE_ALL_END,
  MPI FILE READ ORDERED BEGIN, MPI FILE READ ORDERED END,
 MPI_FILE_WRITE_ORDERED_BEGIN, MPI_FILE_WRITE_ORDERED_END,
```





Ballot 4, Item 22 - Change "INOUT" to "IN" for MPI Handle Parameters in several routines (4)

Question: Do you

accept this entry?

Yes: all

No:

Abstain:

Alternative A:

Keep the argument definition for handling the opaque objects (INOUT) and add the argument definition for the handles as IN.

Proposal: MPI 2.0 Sect. 2.3 Procedure Specification, page 6 lines 30-34 read:

There is one special case - if an argument is a handle to an opaque object (these terms are defined in Section 2.5.1), and the object is updated by the procedure call, then the argument is marked OUT. It is marked this way even though the handle itself is not modified - we use the OUT attribute to denote that what the handle references is updated. Thus, in C++, IN arguments are either references or pointers to const objects.

but should read:

There is one special case - if an argument is a handle to an opaque object (these terms are defined in Section 2.5.1), and the object is updated by the procedure call, then the argument is marked INOUT or OUT. It is marked this way even though the handle itself is not modified - we use the INOUT or OUT attribute to denote that what the handle references is updated. Thus, in C++, IN arguments are either references or pointers to const objects.



Change the three inconsistent interface definitions from IN to INOUT in MPI-1.1 see list of MPI 1.1 routines below. (Text modified at March 2008 meeting)



Ballot 4, Item 22 - Change "INOUT" to "IN" for MPI Handle Parameters in several routines (5)

MPI_FILE_SYNC

Rationale for this proposal A:

MPI_FILE_SET_ATOMICITY,

This is the minimal change to remove the existing inconsistency. Only the Fortran interfaces of three deprecated MPI-1.1 routines is modified from IN to INOUT. Due to Fortran call by reference, this has no impact for the applications. In the C interfaces, the handle argument is call by value.



Ballot 4, Item 22 - Change "INOUT" to "IN" for MPI Handle Parameters in several routines (6) Moved into MPI-2.2

Alternative B:

Keep the argument definition for handling the opaque objects (INOUT) and add the argument definition for the handles as IN.

Proposal: MPI 2.0 Sect. 2.3 Procedure Specification, page 6 lines 30-34 read:

There is one special case - if an argument is a handle to an opaque object (these terms are defined in Section 2.5.1), and the object is updated by the procedure call, then the argument is marked OUT. It is marked this way even though the handle itself is not modified - we use the OUT attribute to denote that what the handle references is updated. Thus, in C++, IN arguments are either references or pointers to const objects.

but should read:

There is one special case - if an argument is a handle to an opaque object (these terms are defined in Section 2.5.1), and the object is updated by the procedure call but the handle itself is not modified, then the argument is marked IN/INOUT. We use the first part (IN) to specify the use of the handle and the second part (INOUT) to specify the use of the opaque object. Thus, in C++, IN arguments are either references or pointers to const objects, IN/INOUT arguments are references to const handles to non-const objects.



In the listed routines, the **INOUT** handle declaration (in MPI-2.0) and the **IN** handle declaration (in MPI-1.1) is **modified** into a **IN/INOUT** handle declaration.











Ballot 4, Item 22 - Change "INOUT" to "IN" for MPI Handle Parameters in several routines (7)

Moved into MPI-2.2

Rationale for this proposal B:

I have checked the total MPI 1.1 and 2.0 standard to find all routines with an argument specification according to the following declaration pattern:

Language independent interface: INOUT handle C interface MPI_handletype handle

All these routines keep the handle itself unchanged, but the opaque object is modified in a way, that with other MPI routines this change can be detected. For example, an attribute is cached or changed, a file pointer is moved, the content of a file was modified.

The current specification with IN (in MPI 1.1) or INOUT (in MPI 2.0) is inadequate and led to misinterpretation in the const declarations of the C++ interface. It is not explicitly mentioned that IN/IN is abbreviated with IN, and OUT/OUT with OUT. (Therefore no change in all routines with pure IN and pure OUT handles/opaque objects.



This proposal changes the Fortran interface, because the handles itself are now declared as IN. The MPI-2.0 did not decide whether they are IN or INOUT. Only C/C++ interfaces specified call by value for the handles itself. This hasn't any impact for applications. It is not expected that it has any impact on any MPI implementation.



Ballot 4, Item 22 - Change "INOUT" to "IN" for MPI Handle Parameters in several routines (8) Moved into MPI-2.2

Alternative C:

Substitute the argument definition for handling the opaque objects (INOUT) and by the argument definition for the handles (IN).

Proposal:

MPI 2.0 Sect. 2.3 Procedure Specification, page 6 lines 30-34 read:

There is one special case - if an argument is a handle to an opaque object (these terms are defined in Section 2.5.1), and the object is updated by the procedure call, then the argument is marked OUT. It is marked this way even though the handle itself is not modified - we use the OUT attribute to denote that what the handle references is updated. Thus, in C++, IN arguments are either references or pointers to const objects.

but should read:

There is one special case - if an argument is a handle to an opaque object (these terms are defined in Section 2.5.1), and the object is updated by the procedure call but the handle itself is not modified, then the argument is marked IN. Thus, in C++. IN arguments are either references or pointers to const objects, or references to const handles to non-const objects.



all MPI-2.0 routines from the list above, the **INOUT** handle declaration is modified into a IN handle declaration.



Ballot 4, Item 22 - Change "INOUT" to "IN" for MPI Handle Parameters in several routines (9) Moved into MPI-2.2

Rationale for this proposal C:

This proposal is easier, but looses the INOUT information on the opaque object itself. As at Alternative B, this proposal changes the Fortran interface, because the handles itself are now declared as IN. The MPI-2.0 did not decide whether they are IN or INOUT. Only C/C++ interfaces specified call by value for the handles itself. This hasn't any impact for applications. It is not expected that it has any impact on any MPI implementation.



Ballot 4, Item 22 - Change "INOUT" to "IN" for MPI Handle Parameters in several routines (10) Moved into MPI-2.2

Proposal A: Defining only the outcome for the

→ e.a.. INOUT filehandle in write routines.

→ Text in "Terms and Conventions" is **not** changed.

→ Changes at 3 MPI 1.1 routines (IN → INOUT)

Proposal B: Explicitly defining the outcome for handle and opaque object

→ e.g., IN/INOUT filehandle in write routines.

→ Text in "Terms and Conventions" is changed.

→ Changes at 3 MPI 1.1 (IN → IN/INOUT) and 48 MPI 2.0 routines (OUT → IN/INOUT)

Proposal C: Defining only the outcome for the handle itself

→ e.g., IN filehandle in write routines.

- → Text in "Terms and Conventions" is changed.
- → Changes at 48 MPI 2.0 routines (INOUT → IN)

Question 1: Do you (additionally) want the specification of the outcome for handles? i.e., Proposal B or C instead of A

Yes:

No:

Abstain:

Only if we have a positive decision in the first question:

Question 2: Do you want to keep the specification of outcome for the opaque objects? i.e., Proposal B instead of C

Yes:

Abstain:













Do you

accept

entry?

Yes:

No:

Abstain:

 Λ

Decision

Proposal

this



Ballot 4, Item X.1-6 - A clarification is not needed

The following Ballot 4 items do not need a clarification.

Therefore, they are closed and will be removed from the MPI-2 Errata webpage when Ballot 4 is finished.

→ They will go to the same history web-page as all Ballot 4 items.

- **X.1** Description of the send and receive count arguments to MPI Alltoally
- X.2 Shared File Pointers
- **X.3** Overlapping buffers in collective communications
- X.4 MPI_MODE_RDONLY, MPI_MODE_RDWR, MPI_MODE_WRONLY
- **X.5** MPI_FINALIZE in MPI-2 (with spawn)
- **X.6** Reporting invalid handles provided to handle conversion functions

Any reasons for a veto?



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MPI 2.1 - Ballot 4

Thank you very much Ballot 4 is done for now!



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