The Multi-Processor Computing Framework (MPC)



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Introduction

MPC is a thread-based MPI + OpenMP runtime.

Originally developed by CEA

- ▶ Initiated as a PhD work¹
- ▶ Started more than ten years ago as part of a complex 3D AMR code²
- ▶ Rapidly evolved to become an independent MPI
- ▶ ParaTools collaborates with CEA to develop it

Thread-based since the origin

- ► Full MPI_THREAD_MULTIPLE support (no overhead)
- ▶ Provides its own user-level MxN scheduler

Built as a framework



^{1.} Marc Pérache. Contribution à l'élaboration d'environnements de programmation dédiés au calcul scientifique hautes performances. Thèse de Doctorat, spécialité informatique, CEA/DAM Île de France, Université de Bordeaux 1, Domaine Universitaire, 351 Cours de la libération, 33405 Talence Cedex, October 2006. Note: 141 pages.

^{2.} Jourdren, H. (2005). **HERA: a hydrodynamic AMR platform for multi-physics simulations**. In Adaptive Mesh Refinement-Theory and Applications (pp. 283-294). Springer Berlin Heidelberg.

MPC is in practice a modular set of components:

Message Passing Layer

- ▶ Inter-thread, TCP, SHM, IB, Portals (in progress)
- ▶ Highly modular and configurable
- ▶ Very compact (easily extensible) and thread-safe
- Usable as a stand-alone component (library mode)
- ▶ Supports RDMA (to be released), P2P and regular collectives

MPI Layer (Fully 1.3 but practically close to 3.0)

- ▶ Built on top of the Message Passing Layer
- ▶ Full thread-based MPI with MPI_THREAD_MULTIPLE support
- ▶ MPI 3.0 support within one year

User-level MxN thread scheduler

- Avoids busy-waiting (replaced by context-switching)
- ▶ Full support for POSIX Thread API (e.g. TBB porting)
- ▶ Its policy can be configured
- Can schedule a (very) large number of threads (oversubscribing)
- ▶ Global view of all supported runtimes (MPI, OpenMP...)



MPC Framework (2/4)

OpenMP Runtime

- ▶ Support for OpenMP 3.1
- ▶ Compatible with GNU and Intel compilers (C, C++ and Fortran)
- Optimized for hybrid OpenMP + MPI payloads

NUMA-Aware Allocator

- ▶ Per-thread heap to improve data locality and avoid contention
- ▶ Lock-free hierarchical heaps
- Built to limit system calls (memory recycling)

Patched Compiler Chain

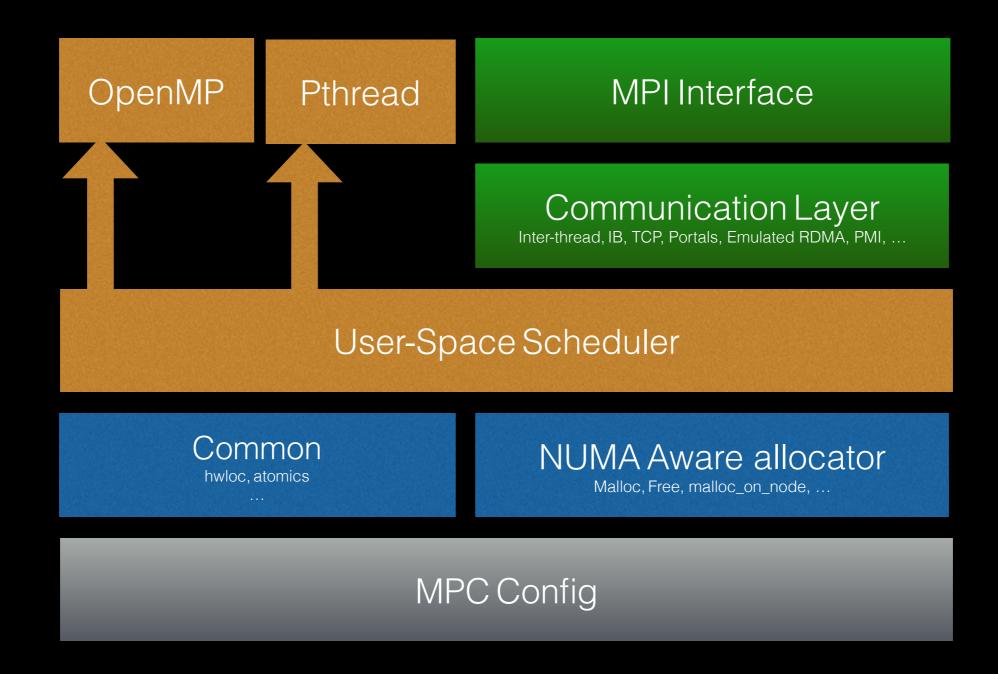
- Automatic privatization of global variables (GNU and Intel compilers)
- ▶ Hierarchical Local Storage (HLS) for GCC

Generic configuration Module

- XML-based configuration
- Easy to extend
- Designed for maximum modularity



MPC Framework (3/4)





MPC Framework (4/4)

Thread-based approach

- ▶ Reduced launch time (spawns only one process per node)
- Memory savings (network buffers, runtime internals, process contexts)
- ► MPI + X fair mixing over the unified scheduler
- ► Fast inter-thread communication (one message => one copy)

Lightweight and extensible MPI + X

- ▶ Relatively small codebase
- Supports a steadily increasing part of the standards (MPI, OpenMP)
- Generic network support (complete rail system)
- ▶ Unified XML configuration system

Compact RDMA-capable communication library

- ▶ Multirail support for IB, TCP, Portals*, SHM*
- ▶ Portable RDMA interface*
- ▶ Compile MPC as an auxiliary communication library*

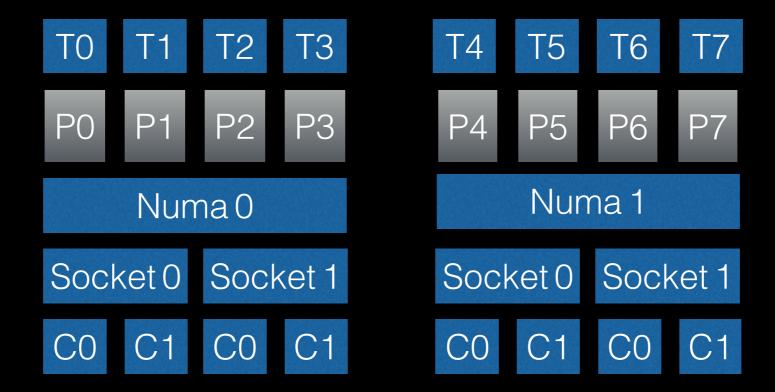
MPC is supported by profiling and debugging tools

- ▶ Patched GDB (user-level thread support)
- ▶ Support for OMPT*, MPIT is a short-term milestone
- Compatible with Allinea's DDT
- Compatible with TAU



^{*} to be released in a near future

Execution Configurations (1/3)



Process-Based MPI

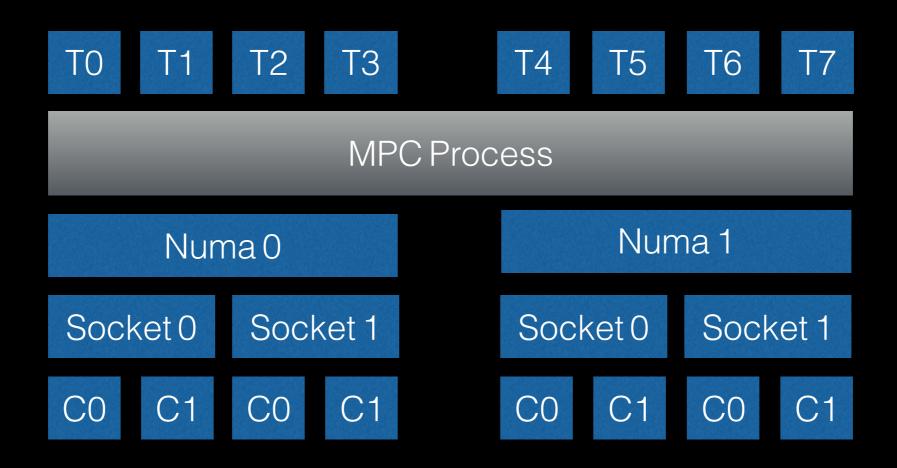
mpcrun -N=1 -p=8 -n=8 ./a.out

Launch is described with four parameters (replacing -np):

- ▶ N: Number of nodes
- p: Number of system processes
- n: Number of MPI tasks
- c: Number of cores per process



Execution Configurations (2/3)



Thread-Based MPI

mpcrun -N=1 -p=1 -n=8
$$./a$$
.out



Execution Configurations (3/3)



Mixed Approach

mpcrun -N=1 -p=2 -n=8 -c=4 ./a.out



Privatization Mechanism (1/6)

Motivating example:

```
#include <mpi.h>
#include <stdio.h>
int rank = -1;
int main( int argc, char ** argv )
{
    MPI Init( &argc, &argv );
    MPI Comm rank ( MPI COMM WORLD, &rank );
    printf("My rank %d", rank );
    MPI Finalize();
    return 0;
                           10
```

Privatization Mechanism (2/6)

It's working with $n == p \longrightarrow Process-based MPI$

```
$ mpcrun -n=4 ./a.out
Running MPC with HYDRA job manager
MPC version 3.0.0_devel C/C++ (4 tasks 1 processes 4 cpus (3.59GHz) ethread_mxn) MPC allocator
Initialization time: 0.1s - Memory used: 11MB
My rank 1
```

Overwritten global variable — Thread-based MPI



Privatization Mechanism (3/6)

Global variables should be duplicated (or **privatized**) One idea is to rely on TLS:

```
_{\underline{}}thread int rank = -1;
```

```
$ mpcrun -net=tcp -p=1 -n=8 ./a.out
Running MPC with HYDRA job manager
MPC version 3.0.0_devel C/C++ (8 tasks 1 processes 4 cpus (3.59GHz) ethread_mxn) MPC allocator
Initialization time: 0.1s - Memory used: 9MB
My rank 1
My rank 1
My rank 7
My rank 7
My rank 7
My rank 5
My rank 5
My rank 3
My rank 3
My rank 3
TLS solution not working in oversubscribing
Context (more MPI tasks than cores)
My rank 3
```

Conclusion: need to design a mechanism able to duplicate global variables and to deal with the context switching logic



Privatization Mechanism (4/6)

MPC solution: a compile-time approach for automatic global variable privatization

- Simplifies porting of existing codes and libraries
- Allows MPI applications to be run in shared-memory

The loader was also patched to optimize context-switching using the GS register on x64 targets.

```
$ mpc_cc t.c
(Front-end C) Automatic privatization to MPC task (variable rank in file t.c line 6)
```

- Promotes global variables to TLS level
- Adds the context-swiching logic to the generated code
- Implementation in GNU compilers (C/C++ and Fortran)
 - Uses the -fmpc-privatize flag (enabled by default)



Privatization Mechanism (5/6)

CEA collaborated with Intel to have this same option implemented in the Intel C/C++ and Fortran compilers.

Excerpt from the Intel ICC compiler manual (starting with 15.0):

-fmpc-privatize (L*X only)

-fno-mpc-privatize (L*X only)

Enables or disables privatization of all static data for the MultiProcessor Computing environment (MPC) unified parallel runtime.

Architecture Restrictions: Only available on Intel(R) 64 architecture

Arguments:

None

Default:

-fno-mpc-privatize

The privatization of all static data for the MPC unified parallel runtime is disabled.

Description:

This option enables or disables privatization of all static data for the MultiProcessor Computing environment (MPC) unified parallel runtime.

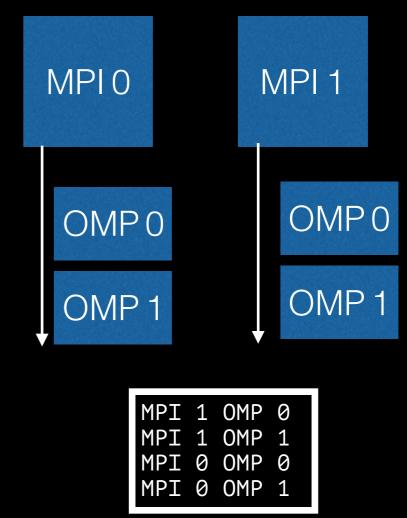
Option -fmpc-privatize causes calls to extended thread-local-storage (TLS) resolution, run-time routines that are not supported on standard Linux* OS distributions.

This option requires installation of another product. For more information, see Feature Requirements



Privatization in Hybrid Context (6/6)

```
#include <mpi.h>
int rank = -1;
int omprank = -1;
#pragma omp threadprivate(omprank)
int main( int argc, char ** argv )
   MPI_Init( &argc, &argv );
   MPI_Comm_rank( MPI_COMM_WORLD, &rank );
   #pragma omp parallel
       omprank = omp get thread num();
       printf("MPI %d OMP %d\n", rank, omprank );
  MPI Finalize();
   return 0;
```



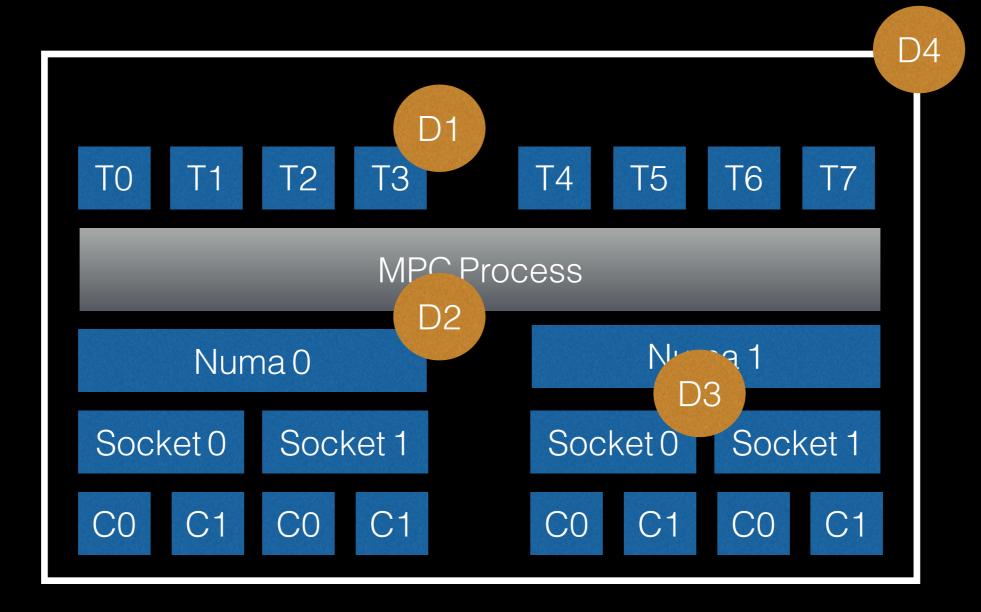
2 Processes on 4 cores

Hybrid MPI+OpenMP mode requires a specific handling of thread-private OpenMP variables provided by the Extended TLS mechanism

Carribault, P., Pérache, M., & Jourdren, H. (2011). **Thread-local storage extension to support thread-based MPI/openMP applications**. In *OpenMP in the Petascale Era* (pp. 80-93). Springer Berlin Heidelberg.



Hierarchical Local Storage (HLS) (1/2)



Extends the TLS mechanism to store "global" arrays (e.g., huge material data tables) at a given level of the memory hierarchy (Node, NUMA, Cache, Core).



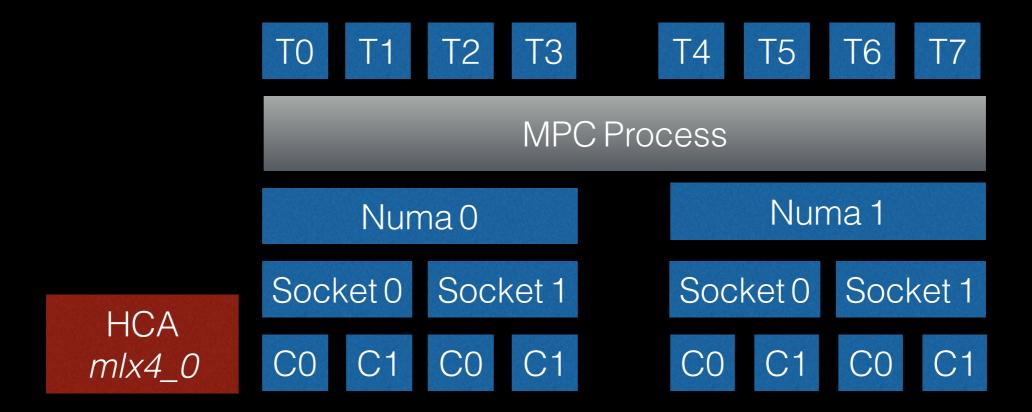
Hierarchical Local Storage (HLS) (2/2)

```
int a,b
#pragma hls node(a)
#pragma hls numa(b)
void f()
    #pragma hls single(a)
        // executed by one instruction flow per node
        a = 4;
    #pragma hls single(b)
        // executed by one instruction flow per NUMA node
        b = 2 :
```

Compiler-enabled (GCC) data sharing between MPI tasks with OpenMP-like directives.



Collaborative Polling



The HCA being a shared ressource, waiting tasks can help others to progress by polling the network or copying buffers (also true for inter-thread comm).

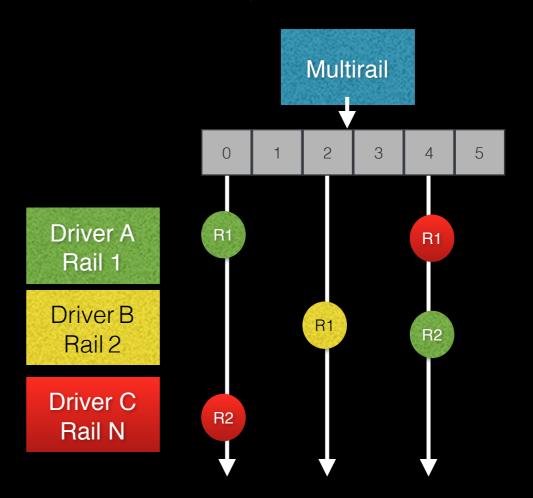
Didelot, S., Carribault, P., Pérache, M., & Jalby, W. (2014). **Improving MPI communication overlap with collaborative polling**. *Computing*, *96*(4), 263-278.



Network Rail Support and Configuration (1/4)

```
<config>
   <name>tcp_config_mpi</name>
   <driver>
           <tcp/>
   </driver>
</config>
<rail>
   <name>tcp mpi</name>
   <priority>1</priority>
   <topology>ring</topology>
   <config>tcp config mpi</config>
</rail>
<rail>
   <name>tcp_large</name>
   <priority>10</priority>
   <topology>none</topology>
   <config>tcp_config_mpi</config>
   <qates>
       <gate>
          <minsize>
              <value>32KB</value>
          </minsize>
       </gate>
   </gates>
</rail>
<cli>option>
   <name>multirail tcp</name>
   <rails>
       <rail>tcp large</rail>
       <rail>tcp mpi</rail>
   </rails>
</cli option>
```

Endpoint-based Rail-election through priorities and gates Multi-Driver, Multi-Device



Launch command:

mpcrun -net=**multirail_tcp** -p=8 -n=16 ./a.out



Network Rail Support and Configuration (2/4)

Topology-based rail selection between tasks using rail stacking and a regular expression to match devices.

```
Topological Rail
<rail>
   <name>ib mpi all devices<name>
   <device>!mlx4 *</device>
   <config>ib config mpi</config>
                                                                         Subrails
<rail>
Rail(0) [Topological RAIL (ring) (!mlx4_[0,3])]
                                      4/4:mlx4_0 - 4x QDR (32 Gb/s) - 0] (ring) (mlx4_0)]
    Sub-Rail(1) [IB-MT (v2.0) MPI
    Sub-Rail(2) [IB-MT (v2.0) MPI
                                      1/4:mlx4 3 - 4x QDR (32 Gb/s)
Rail(0) [Topological RAIL (ring) (!mlx4_[0-1])]
    Sub-Rail(1) [IB-MT (v2.0) MPI 4/4:mlx4_0 - 4x QDR (32 Gb/s) - 0] (ring) (mlx4_0)]
    Sub-Rail(2) [IB-MT (v2.0) MPI
                                      3/4:mlx4\ 1 - 4x\ QDR\ (32\ Gb/s)
Rail(0) [Topological RAIL (ring) (!mlx4_*)]
    Sub-Rail(1) [IB-MT (v2.0) MPI
                                      4/4:mlx4 0 - 4x QDR (32 Gb/s) - 0] (ring) (mlx4 0)]
    Sub-Rail(2) [IB-MT (v2.0) MPI
                                      3/4:mlx4\ 1 - 4x\ QDR\ (32\ Gb/s) - 0\ (ring)\ (mlx4\ 1)\ ]
                                      2/4:mlx4_2 - 4x QDR (32 Gb/s) - 0] (ring) (mlx4_2)]
    Sub-Rail(3) [IB-MT (v2.0) MPI
                                      1/4:mlx4_3 - 4x QDR (32 Gb/s) - 0] (ring) (mlx4_3)]
    Sub-Rail(4) [IB-MT (v2.0) MPI
```

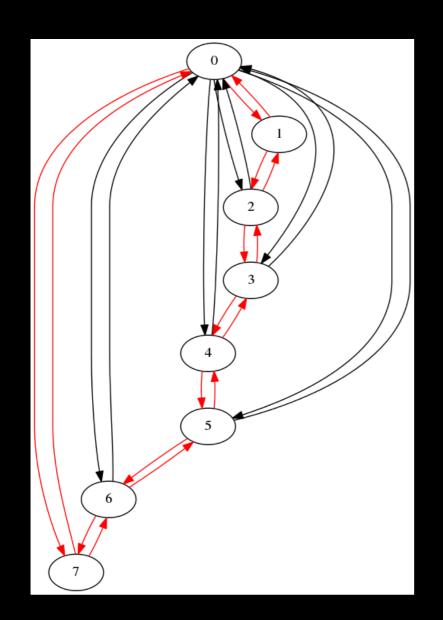
Bull BCS nodes $(32 \times 4 = 128 \text{ cores})$



Network Rail Support and Configuration (3/4)

The bootstrapping process relies on topological promotion:

- First a ring is connected using the PMI (can be on a single rail)
- Then a topology is built using this initial connectivity
- On-demand connections are routed (distance) on this topology

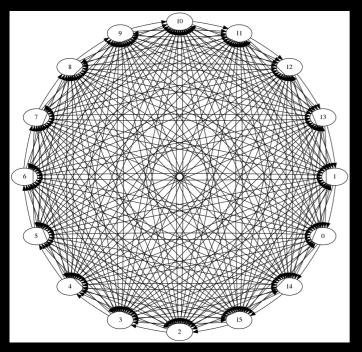


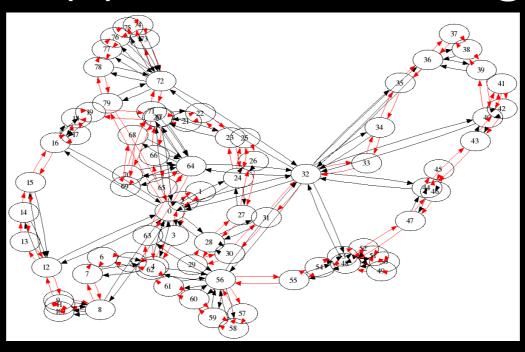
Connectivity after a barrier (8-tree) static routes are in RED, on-demand routes are in black.

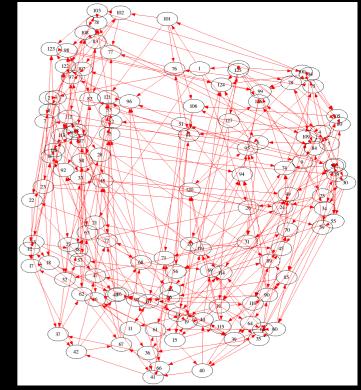
The RED links ensure that any connection can be initiated thanks to routing.



Network Rail Support and Configuration (4/4)







Fully Connected

Ring + Init Barrier

Torus

Promoting the routing topology reduces the network diameter, speeding up **control messages** (internal RPCs). Il also opens the way for "out-of-band" messages (loosely synchronized, one-sided).

- Used internally for emulated RMA, MPI Windows, On-Demand
- User-defined RPCs (remote control, in-situ visualization) "Server Mode"
- ► IDEA: out-of-band messages (load-balancing, profiling)
- ► IDEA: contention management (route instead of connecting N-1 avoidance) on-demand connection hiding (for small payloads)

MPC Development Tracks (1/4)

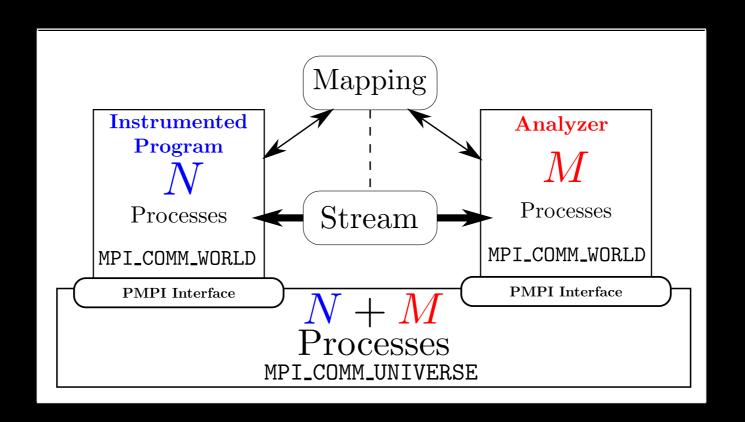
Some features we are working on:

- Better support of Process-Based MPI:
 - ▶ Ease the transition to thread-based / MPI + X
 - ▶ Provide an *mpirun* interface
 - ▶ Implement a state-of-the-art SHM (Linux Cross Memory Attach)
- Enhance the MPI standard coverage:
 - ▶ We recently merged ROMIO for MPI-IO and NBC
 - ▶ Data-types are almost 3.0
 - ▶ Low-level RDMAs are implemented, one sided are on the way
 - Dynamic processes are planned for next year
- Enhance reliability and portability



MPC Development Tracks (2/4)

"Constellation of Services"



Scheme developed during my PhD thesis to manage profiling data using the PNMPI¹ virtualization idea: we would like to apply this approach more widely.



^{1.} Schulz, M., & De Supinski, B. R. (2007, November). **PN MPI tools: A whole lot greater than the sum of their parts**. In *Proceedings of the 2007 ACM/IEEE conference on Supercomputing* (p. 30). ACM.

MPC Development Tracks (3/4)

"Constellation of Services"

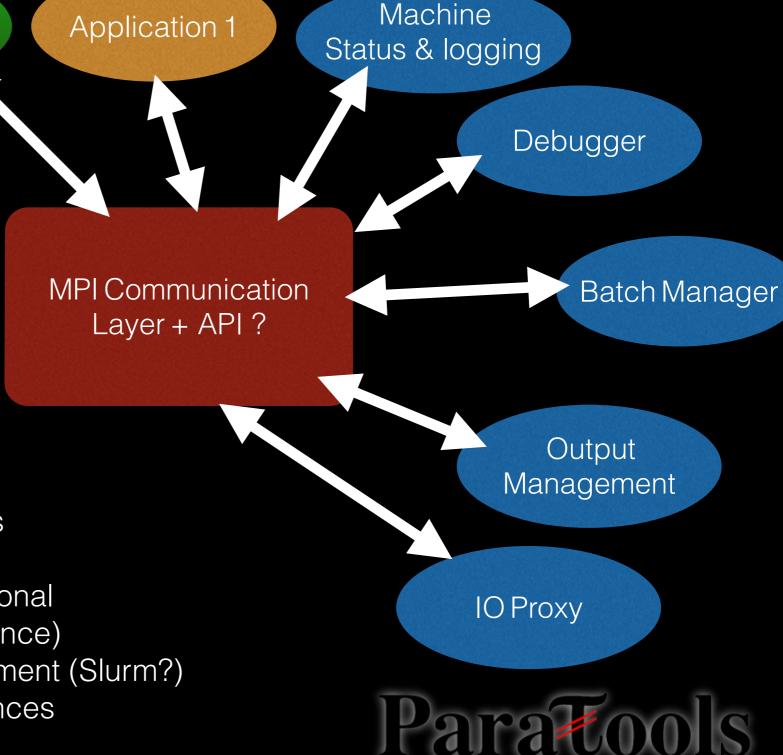
Application 2

Dynamic processes provide a basic block for this idea.

Data management might have to be enhanced (query-based) instead of send-recv — "server" mode?

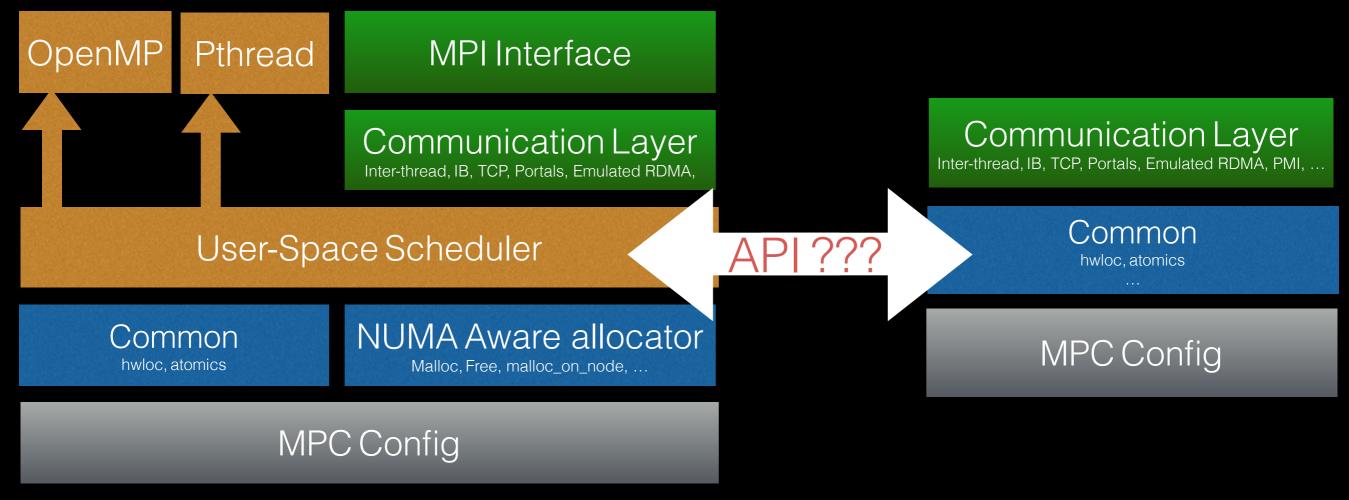
(ANY_TAG, ANY_SOURCE) -> No context Why not RPCs, Streams,

- Attach and detach MPI applications to provide per-user services
- Make data-exchanges more orthogonal
- Support failing services (fault-tolerance)
- Provide a connection/auth management (Slurm?)
- Stream data between parallel instances



MPC Development Tracks (4/4)

"Constellation of Services"



Complete MPC

« Lib-mode » MPC

We were able to link the "lib-mode" MPC inside OpenMPI to provide an auxiliary communication library. Our purpose is to interface MPI instances with an existing service infrastructure taking advantage of RMDAs — trying to extract a standard interface.

MPC Validation Suite (1/3)

To reach its production milestone a test-suite engine has been developed for MPC then forked as **JCHRONOSS** to be used independently.

Released under CECILL-C (fully LGPL compatible) http://jchronoss.hpcframework.com

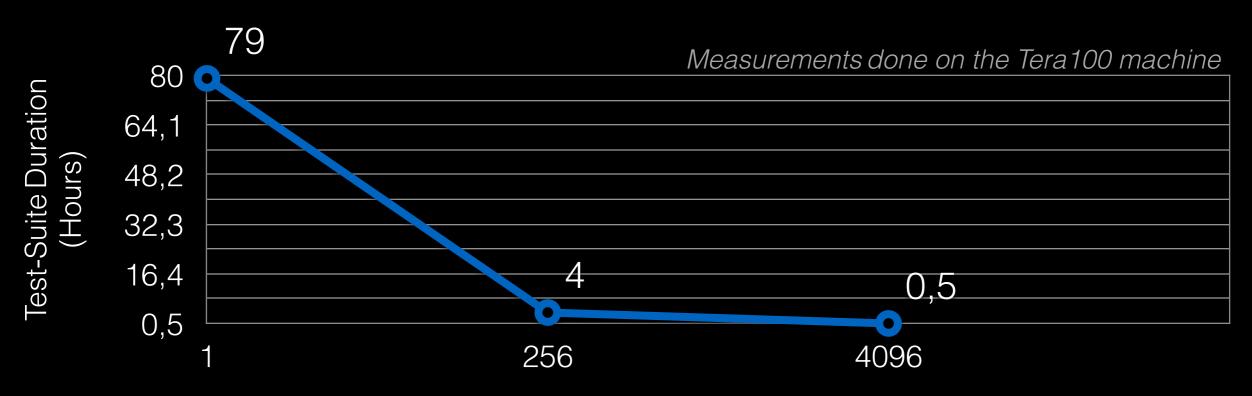


Scheduling of parallel tests in parallel on a supercomputer

- Uses the batch manager (slurm, PBS...) with a surface-based scheduling policy
- Fault-tolerance support (resubmission)
- Support for dependencies (e.g., compilation -> run)
- Can be launched by individual developers (locally)
- Provides a lightweight html GUI for in-situ validation
- Stores the output of failed tests



MPC Validation Suite (2/3)



Maximum number of cores to run the test-suite

Reduces the validation time for the 40,307 tests of the MPC test-suite from 3.2 days to less than an hour (45 minutes), taking advantage of supercomputer parallel resources to validate our parallel runtime in production environment.



MPC Validation Suite (3/3)

```
ID / SZ
                              | LEFT |
                                            TIME
                                                     test 17
                     1/2
                                     18)
                                                     compile 01
PASSED
                     2/2
                                     18)
                                                     compile 04
PASSED
                     1/1
                                     16)
PASSED
                     1/1
                                     15)
                                                     test 01
PASSED
                     1/1
                                     11)
                                             0.00
                                                     test 06
                                     13)
FAILED
                     1/1
                                             0.00
                     1/1
                                     12)
                                                     test 03
                     1/1
                                      7)
                                      8)
                                     10)
                     1/1
                                     6)
                                      4)
                     1/1
                                      5)
                     1/1
                                     14)
                     1/1
                                                     test 10
    --> Fails
     --> Successes : 14
```

Console Output

Test-suite result field visualization



Summary									
Status Nombre Rapport									
success 36988								96 %	
failed 1323								3 %	
error 0								0 %	
skipped	20							0 %	
total	38331	failed						100 %	
Test groups									
Directory		Errors	Failures	Skipped	Success A	Total	% failed	Status	
Threads-interfaces		0	354	16	2489	2859	12	Failed	
MPI-ANL_error-grp_ctxt_comm		0	0	0	1826	1826	0	Success	
MPI-MPICH		0	520	0	1613	2133	24	Failed	
MPI-ANL_error-datatype		0	0	0	1504	1504	0	Success	
MPI-ANL_functional-datatype		0	0	0	1475	1475	0	Success	
MPI-ANL_error-collective		0	0	0	1453	1453	0	Success	
MPI-ANL_error-topo		0	0	0	1325	1325	0	Success	
MPI-Intel_ANL-types		0	0	0	1150	1150	0	Success	
fortran-MPI_Intel_ANL-types		0	13	0	1112	1125	1	Failed	
MPI-Intel_ANL-sync		0	7	0	1055	1062	1	Failed	
MPI-ANL_error-persist_request		0	0	0	964	964	0	Success	
fortran-MPI_Intel_ANL-sync		0	0	0	959	959	0	Success	
MPI-ANL_error-blocking		0	1	0	937	938	0	Failed	
MPI-ANL functional-grp_ctxt_comm		0	32	0	918	950	3	Failed	
MPI-ANL_error-nonblocking		0	24	0	863	887	3	Failed	
MPI-MPICH_TEST_SUITE-pt2pt		0	6	0	845	851	1	Failed	
MPI-ANL_functional-nonblocking		0	2	0	782	784	0	Failed	
MPI-Intel_ANL-collectives		0	0	0	750	750	0	Success	
fortran-MPI_Intel_ANL-communicators		0	11	0	714	725	2	Failed	
MPI-Intel_ANL-async		0	4	0	703	707	1	Failed	
MPI-ANL_functional-persist_request		0	1	0	680	681	0	Failed	
fortran-MPI Intel ANL-collectives		0	0	0	678	678	0	Success	
fortran-MPI Intel ANL-async		0	5	0	645	650	1	Failed	
MPI-Intel_ANL-communicators		0	32	0	645	677	5	Failed	
MPI-ANL error-sendrecv		0	0	0	604	604	0	Success	
MPI-ANL functional-blocking		0	0	0	531	531	0	Success	
MPI Intol AML others		0	25	0	520	554	5	Failed	

Static HTML Interface



Conclusion

We introduced the Multi-Processor Computing framework

Key features

- ▶ MPI & OpenMP implementations (based on unified thread scheduler)
- Automatic privatization (GNU and Intel compilers)
- HLS directives (GNU compilers)
- Collaborative Polling
- Generic rail configuration and topology management
- ▶ Control messages

Development tracks for MPC

- Process-based mode (for transition)
- Constellation of services (low-level back-end)

Test-suite engine

- Validating on a daily basis a complete test-suite (w/ 30k MPI tests)
- ▶ Taking advantage of a parallel machine



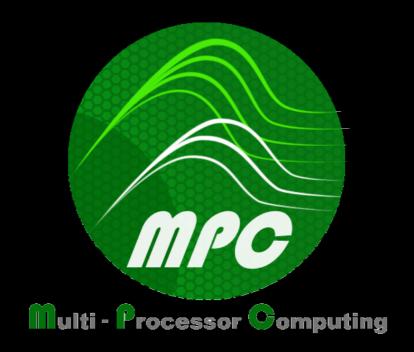
Try MPC

Fully LGPL compatible (French CECILL-C licence) Is an interesting prototyping/research platform

http://paratools.fr/mpc



The Multi-Processor Computing Framework (MPC)



MPI Forum September 24-25, Bordeaux Jean-Baptiste BESNARD jbbesnard@paratools.fr



