

Introduction to Audio Content Analysis

module 10.1: alignment — dynamic time warping

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introduction overview

Georgia Center for Music Tech || Technology College of Design

corresponding textbook section

section 10.1

lecture content

- Dynamic Time Warping (DTW): synchronization of two sequences with similar content
- learning objectives
 - explain the standard DTW algorithm
 - discuss disadvantages of and modifications to the standard DTW algorithm
 - implement DTW



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lecture content

- Dynamic Time Warping (DTW): synchronization of two sequences with similar content
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synchronize two sequences

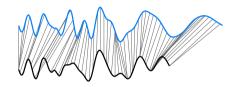
- similar musical content
- different tempo and timing

$$\textit{A}(\textit{n}_{\mathrm{A}}) \quad \textit{n}_{\mathrm{A}} \in [0; \mathcal{N}_{\mathrm{A}} - 1]$$

$$\textit{B(n}_{\mathrm{B}}) \quad \textit{n}_{\mathrm{B}} \in [0; \mathcal{N}_{\mathrm{B}} - 1]$$



- minimizing pairwise distance between sequences
- covering whole sequence
- moving only forward in time



dynamic time warping overview

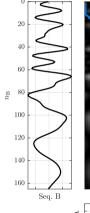


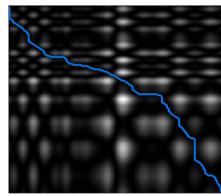
- dynamic programming technique
- time is warped non-linearly to match sequences
- finds optimal match between two sequences given a cost function
- the overall cost indicates the overall distance between the sequences

dynamic time warping processing steps

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- extract suitable features ⇒ two series of feature vectors
- 2 compute distance matrix $D_{AB}(n_A, n_B)$
- 3 compute alignment path $p(n_{\rm P})$ with $n_{\rm P} \in [0; \mathcal{N}_{\rm P} 1]$ \Rightarrow minimal overall distance
- 4 (align sequences using dynamic time stretching)





60



120

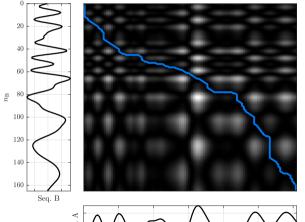
100

dynamic time warping distance matrix computation

■ given 2 sequences of vectors,

- lacktriangle compute distance matrix $oldsymbol{D}_{
 m AB}(n_{
 m A},n_{
 m B})$
 - example $D_{AB}(1, n_B)$ is the distance of the first vector in Seq. A to all vectors in Seq. B

compute the distance between all pairs of observations



60

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120

100

dynamic time warping path properties 1/2

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boundaries: covers both *A*, *B* from beginning to end

$$m{
ho}(0) = [0,0] \ m{
ho}(\mathcal{N}_{
m P}-1) = [\mathcal{N}_{
m A}-1,\mathcal{N}_{
m B}-1]$$

causality: only forward movement

$$n_{\mathrm{A}}|_{\boldsymbol{p}(n_{\mathrm{P}})} \leq n_{\mathrm{A}}|_{\boldsymbol{p}(n_{\mathrm{P}}+1)}$$
 $n_{\mathrm{B}}|_{\boldsymbol{p}(n_{\mathrm{P}})} \leq n_{\mathrm{B}}|_{\boldsymbol{p}(n_{\mathrm{P}}+1)}$

continuity: no jumps

$$n_{\rm A}\big|_{m{p}(n_{
m P}+1)} \le (n_{
m A}+1)\big|_{m{p}(n_{
m P}+1)}$$
 $n_{
m B}\big|_{m{p}(n_{
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dynamic time warping path properties 1/2

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dynamic time warping path properties 1/2

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boundaries: covers both A, B from beginning to end

$$m{
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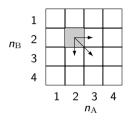
$$n_{\mathrm{B}}\big|_{\boldsymbol{\rho}(n_{\mathrm{P}})} \leq n_{\mathrm{B}}\big|_{\boldsymbol{\rho}(n_{\mathrm{P}}+1)}$$

continuity: no jumps

$$egin{aligned} n_{\mathrm{A}}ig|_{oldsymbol{
ho}(n_{\mathrm{P}}+1)} &\leq (n_{\mathrm{A}}+1)ig|_{oldsymbol{
ho}(n_{\mathrm{P}})} \ n_{\mathrm{B}}ig|_{oldsymbol{
ho}(n_{\mathrm{P}}+1)} &\leq (n_{\mathrm{B}}+1)ig|_{oldsymbol{
ho}(n_{\mathrm{P}})} \end{aligned}$$

alignment
path properties 2/2



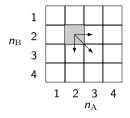


what is the minimum/maximum path length



alignment path properties 2/2





what is the minimum/maximum path length

$$\begin{split} \mathcal{N}_{\mathrm{P,min}} &= \mathsf{max}(\mathcal{N}_{\mathrm{A}}, \mathcal{N}_{\mathrm{B}}) \\ \mathcal{N}_{\mathrm{P,max}} &= \mathcal{N}_{\mathrm{A}} + \mathcal{N}_{\mathrm{B}} - 2 \end{split}$$

$$\mathcal{N}_{\mathrm{P.max}} = \mathcal{N}_{\mathrm{A}} + \mathcal{N}_{\mathrm{B}} - 2$$



alignment DTW: overall cost

every path has an overall cost

$$\mathfrak{C}_{\mathrm{AB}}(j) = \sum_{n_{\mathrm{P}}=0}^{\mathcal{N}_{\mathrm{P}}-1} oldsymbol{D}ig(oldsymbol{p}_{j}(n_{\mathrm{P}})ig)$$

optimal path minimizes the overall cost

$$egin{array}{lll} \mathfrak{C}_{{
m AB}, min} &=& \displaystyle \min_{orall j} \left(\mathfrak{C}_{{
m AB}}(j)
ight) \ j_{
m opt} &=& \displaystyle \operatorname{argmin} \left(\mathfrak{C}_{{
m AB}}(j)
ight) \end{array}$$

⇒ stay in the 'valleys' of distance matrix

how to determine the optimal path



alignment DTW: overall cost



every path has an overall cost

$$\mathfrak{C}_{\mathrm{AB}}(j) = \sum_{n_{\mathrm{P}}=0}^{\mathcal{N}_{\mathrm{P}}-1} oldsymbol{D}ig(oldsymbol{p}_{j}(n_{\mathrm{P}})ig)$$

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⇒ stay in the 'valleys' of distance matrix

how to determine the optimal path



alignment DTW: accumulated cost 1/2

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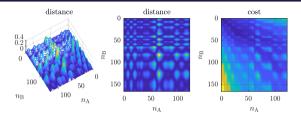
accumulated cost: cost matrix

$$oldsymbol{C}_{ ext{AB}}(n_{ ext{A}},n_{ ext{B}}) = oldsymbol{D}_{ ext{AB}}(n_{ ext{A}},n_{ ext{B}}) + \min \left\{ egin{array}{l} oldsymbol{C}_{ ext{AB}}(n_{ ext{A}}-1,n_{ ext{B}}-1) \ oldsymbol{C}_{ ext{AB}}(n_{ ext{A}}-1,n_{ ext{B}}) \ oldsymbol{C}_{ ext{AB}}(n_{ ext{A}},n_{ ext{B}}-1) \end{array}
ight.$$

initialization

alignment DTW: accumulated cost 2/2





alignment DTW: algorithm description 1/2



■ initialization:

$$oldsymbol{\mathcal{C}}_{\mathrm{AB}}(0,0) = oldsymbol{\mathcal{D}}_{\mathrm{AB}}(0,0), oldsymbol{\mathcal{C}}_{\mathrm{AB}}(n_{\mathrm{A}},-1) = \infty, oldsymbol{\mathcal{C}}_{\mathrm{AB}}(-1,n_{\mathrm{B}}) = \infty$$

recursion:

$$C_{AB}(n_{A}, n_{B}) = D_{AB}(n_{A}, n_{B}) + \min \left\{ egin{array}{l} C_{AB}(n_{A} - 1, n_{B} - 1) \\ C_{AB}(n_{A} - 1, n_{B}) \\ C_{AB}(n_{A}, n_{B} - 1) \end{array}
ight.$$
 $j = \operatorname{argmin} \left\{ egin{array}{l} C_{AB}(n_{A} - 1, n_{B} - 1) \\ C_{AB}(n_{A} - 1, n_{B}) \\ C_{AB}(n_{A}, n_{B} - 1) \end{array}
ight.$
 $C_{AB}(n_{A}, n_{B} - 1) = \left\{ egin{array}{l} [-1, -1] & \text{if } j = 0 \\ [-1, 0] & \text{if } j = 1 \\ [0, -1] & \text{if } j = 2 \end{array}
ight.$

initialization:

$$oldsymbol{\mathcal{C}}_{\mathrm{AB}}(0,0) = oldsymbol{\mathcal{D}}_{\mathrm{AB}}(0,0), oldsymbol{\mathcal{C}}_{\mathrm{AB}}(n_{\mathrm{A}},-1) = \infty, oldsymbol{\mathcal{C}}_{\mathrm{AB}}(-1,n_{\mathrm{B}}) = \infty$$

recursion:

alignment DTW: algorithm description 2/2



■ termination:

$$n_{
m A}=\mathcal{N}_{
m A}-1\wedge n_{
m B}=\mathcal{N}_{
m B}-1$$

path backtracking:

$$p(n_{\rm P}) = p(n_{\rm P}+1) + \Delta p(p(n_{\rm P}+1)), \ n_{\rm P} = \mathcal{N}_{\rm P} - 2, \mathcal{N}_{\rm P} - 3, \dots, 0$$

alignment DTW: algorithm description 2/2



■ termination:

$$n_{
m A}=\mathcal{N}_{
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path backtracking:

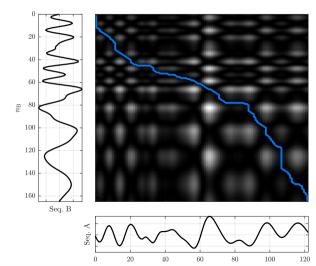
$$p(n_{\rm P}) = p(n_{\rm P}+1) + \Delta p(p(n_{\rm P}+1)), \ n_{\rm P} = \mathcal{N}_{\rm P} - 2, \mathcal{N}_{\rm P} - 3, \dots, 0$$

erview intro DTW distance path cost **example** variants summar 0 0 0 0 0 0

dynamic time warping DTW: example

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dynamic time warping example



$$A = [1, 2, 3, 0],$$

 $B = [1, 0, 2, 3, 1],$



dynamic time warping example



$$A = [1, 2, 3, 0],$$

 $B = [1, 0, 2, 3, 1],$

$$m{D}_{\mathrm{AB}} = \left[egin{array}{cccc} 0 & 1 & 2 & 1 \ 1 & 2 & 3 & 0 \ 1 & 0 & 1 & 2 \ 2 & 1 & 0 & 3 \ 0 & 1 & 2 & 1 \ \end{array}
ight]$$



dynamic time warping example



$$A = [1, 2, 3, 0],$$

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$$m{D}_{\mathrm{AB}} = \left[egin{array}{cccc} 0 & 1 & 2 & 1 \ 1 & 2 & 3 & 0 \ 1 & 0 & 1 & 2 \ 2 & 1 & 0 & 3 \ 0 & 1 & 2 & 1 \end{array}
ight]$$

$$m{D}_{\mathrm{AB}} = \left[egin{array}{cccc} 0 & 1 & 2 & 1 \ 1 & 2 & 3 & 0 \ 1 & 0 & 1 & 2 \ 2 & 1 & 0 & 3 \ 0 & 1 & 2 & 1 \end{array}
ight] \qquad m{C}_{\mathrm{AB}} = \left[egin{array}{ccccc} 0 & \leftarrow 1 & \leftarrow 3 & \leftarrow 4 \ \uparrow 1 & \nwarrow 2 & \nwarrow 4 & \nwarrow 3 \ \uparrow 2 & \nwarrow 1 & \leftarrow 2 & \leftarrow 4 \ \uparrow 4 & \uparrow 2 & \nwarrow 1 & \leftarrow 4 \ \uparrow 4 & \uparrow 3 & \uparrow 3 & \nwarrow 2 \end{array}
ight]$$



dynamic time warping example



$$A = [1, 2, 3, 0],$$

 $B = [1, 0, 2, 3, 1],$

$$\mathbf{D}_{AB} = \begin{bmatrix} 0 & 1 & 2 & 1 \\ 1 & 2 & 3 & 0 \\ 1 & 0 & 1 & 2 \\ 2 & 1 & 0 & 3 \\ 0 & 1 & 2 & 1 \end{bmatrix}$$

$$\mathbf{D}_{AB} = \begin{bmatrix} 0 & 1 & 2 & 1 \\ 1 & 2 & 3 & 0 \\ 1 & 0 & 1 & 2 \\ 2 & 1 & 0 & 3 \\ 0 & 1 & 2 & 1 \end{bmatrix} \qquad \mathbf{C}_{AB} = \begin{bmatrix} 0 & \leftarrow 1 & \leftarrow 3 & \leftarrow 4 \\ \uparrow 1 & \nwarrow 2 & \nwarrow 4 & \nwarrow 3 \\ \uparrow 2 & \nwarrow 1 & \leftarrow 2 & \leftarrow 4 \\ \uparrow 4 & \uparrow 2 & \nwarrow 1 & \leftarrow 4 \\ \uparrow 4 & \uparrow 3 & \uparrow 3 & \nwarrow 2 \end{bmatrix}$$



dynamic time warping variants



transition weights: favor specific path directions

$$oldsymbol{\mathcal{C}}_{\mathrm{AB}}(n_{\mathrm{A}},n_{\mathrm{B}}) = \min \left\{ egin{array}{lll} oldsymbol{\mathcal{C}}_{\mathrm{AB}}(n_{\mathrm{A}}-1,n_{\mathrm{B}}-1) & + & \lambda_{\mathrm{d}} \cdot oldsymbol{\mathcal{D}}_{\mathrm{AB}}(n_{\mathrm{A}},n_{\mathrm{B}}) \ oldsymbol{\mathcal{C}}_{\mathrm{AB}}(n_{\mathrm{A}}-1,n_{\mathrm{B}}) & + & \lambda_{\mathrm{v}} \cdot oldsymbol{\mathcal{D}}_{\mathrm{AB}}(n_{\mathrm{A}},n_{\mathrm{B}}) \ oldsymbol{\mathcal{C}}_{\mathrm{AB}}(n_{\mathrm{A}},n_{\mathrm{B}}-1) & + & \lambda_{\mathrm{h}} \cdot oldsymbol{\mathcal{D}}_{\mathrm{AB}}(n_{\mathrm{A}},n_{\mathrm{B}}) \end{array}
ight.$$

step types





dynamic time warping variants



transition weights: favor specific path directions

$$m{C}_{
m AB}(n_{
m A}, n_{
m B}) = \min \left\{ egin{array}{lll} m{C}_{
m AB}(n_{
m A} - 1, n_{
m B} - 1) & + & \lambda_{
m d} \cdot m{D}_{
m AB}(n_{
m A}, n_{
m B}) \\ m{C}_{
m AB}(n_{
m A} - 1, n_{
m B}) & + & \lambda_{
m v} \cdot m{D}_{
m AB}(n_{
m A}, n_{
m B}) \\ m{C}_{
m AB}(n_{
m A}, n_{
m B} - 1) & + & \lambda_{
m h} \cdot m{D}_{
m AB}(n_{
m A}, n_{
m B}) \end{array}
ight.$$

step types







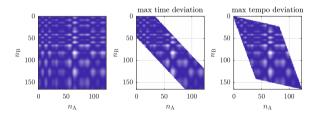
- **challenge**: distance matrix dimensions $\mathcal{N}_A \cdot \mathcal{N}_B$
- ⇒ DTW *inefficient* for long sequences
 - high memory requirements
 - large number of operations

- maximum time and tempo deviation
- 2 sliding window
- 3 multi-scale DTW (severa downsampled iterations)



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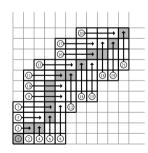
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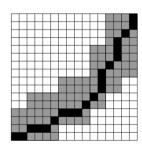


¹S. Dixon and G. Widmer, "MATCH: A Music Alignment Tool Chest," in *Proceedings of the 6th International Conference on Music Information Retrieval (ISMIR)*, London, Sep. 2005.



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- 3 multi-scale DTW (several downsampled iterations)



¹M. Müller, H. Mattes, and F. Kurth, "An Efficient Multiscale Approach to Audio Synchronization," in *Proceedings of the International Society for Music Information Retrieval Conference (ISMIR)*, Victoria, 2006.

dynamic time warping DTW vs. Viterbi

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similarities and differences of DTW and the Viterbi algorithm



dynamic time warping DTW vs. Viterbi

similarities and differences of DTW and the Viterbi algorithm



commonalities

- find path through matrix
- maximizes overall probability/minimizes overall cost
- based on dynamic programming principles

differences

- DTW has more constraints: start/end in corner, move only to neighbor
- DTW is not usually parametrized by training data (transition probs, construction of distance/emission prob matrix)
- Viterbi path length is predefined, DTW path length is not

summary lecture content



- dynamic time warping
 - find globally optimal alignment path between two sequences
- processing steps
 - 1 compute distance matrix
 - 2 compute cost matrix
 - 3 back-track path

