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# Semester Learning Portfolio

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September 19, 2025

## ABSTRACT

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# 1 Authentication + Link Layer Security

## 1.1 Network Security Assignment part 1

### Assignment 1.1.

**Objective:** Research and write a concise paragraph about techniques used to mitigate ARP spoofing and Spanning Tree Protocol (STP) attacks (Layer 2 attacks). Please write details on how the chosen technique detects and prevents the attack, and any potential limitations they may have in a network environment. Please also mention your opinion about the complexity of the techniques you found.

**Answer 1.2.** (ARP Spoofing Mitigation):

- (i) **Static ARP entries:** Using static entries in the ARP table means the IP–MAC mapping cannot be altered by ARP spoofing. The limitation is that when a new device joins the network, its IP–MAC pair must be manually added to the ARP tables of the relevant devices. its nice that i can setup this in a static mac adreses but let say that i ahve to do this for a capnut and maitnign this so alle vesties are update date and
- (ii) **Dynamic ARP Inspection (DAI):**Is a technique where the switches are configured to map each device in the network to a specific IP–MAC pair. If an ARP spoofing attack occurs, then the switch detects that there is an unauthorized ARP request. The limitation of this method is that the switch must be set up with DAI and must be a supported type of switch. This is a better solution than the static assigning since there is a dynamic system in the switches that can help manage the ARP spoofing attacks instead of manually setting each device.
- (iii) **XArp:** Is an anti-spoofing software that can detect if an ARP spoofing attack is being performed on a target system that has installed the XArp on the system, and this is the limitation—that I have to install the XArp and make sure that it's up to date and has no vulnerabilities in this program.

**Answer 1.3.** (STP Attacks Mitigation):

- (i) **BPDU Guard** is a security feature that automatically puts a PortFast-configured access port into an error-disabled state when it receives any BPDU, protecting the STP domain from rogue switches or misconfiguration
- (ii) **Root Guard** is a security feature that prevents non-root ports from becoming root ports by placing them into a root-inconsistent state if they receive superior BPDUs, ensuring the STP topology remains stable and protecting the network from rogue root bridge elections.

## 1.2 Network Security Assignment part 2

### Assignment 1.4.

**Objective:** In this assignment we are going to emulate a Man-in-the-Middle (MITM) attack using this network topology.

As an attacker we should connect to the switch to be able to communicate with the target/victim hosts. From now on, we refer to our two targets hosts as victims.

**Answer 1.5.** (Experiencing Layer 2 attacks):

- (i) The setup I have is two lightweight Lubuntu systems and a Kali Linux where the Man-in-the-Middle attack will be performed. The network is connected to a NAT network through my local machine.

```

root@vboxClone:~# ip a
1: lo: <LOOPBACK,UP,LOWER_UP> mtu 65536 qdisc noqueue state UNKNOWN group default qlen 1000
    link/loopback 00:00:00:00:00:00 brd 00:00:00:00:00:00
    inet 127.0.0.1/8 scope host lo
        valid_lft forever preferred_lft forever
    inet ::1/128 scope host 
        valid_lft forever preferred_lft forever
2: enpos3: <BROADCAST,MULTICAST,UP,LOWER_UP> mtu 1500 adfisc fd_codel state UP group default qlen 1000
    link/ether 08:00:27:57:e3:0b brd ff:ff:ff:ff:ff:ff
    inet 10.0.2.4/24 brd 10.0.2.255 scope global dynamic enp0s3
        valid_lft 30sec preferred_lft 30sec
    inet6 fe80::a00:27ff:fe57:e30b/64 scope link
        valid_lft forever preferred_lft forever
root@vboxClone:~# arp -oC
root@vboxClone:~# ping -C
root@vboxClone:~# pi C
root@vboxClone:~# arp a
a: Host name lookup failure
root@vboxClone:~# arp -a
? (10.0.2.15) at 08:00:27:cb:61:3b [ether] on enpos3
gateway (10.0.2.1) at 52:54:00:12:35:00 [ether] on enpos3
? (10.0.2.3) at 08:00:27:7b:a2:b6 [ether] on enpos3
? (10.0.2.5) at 08:00:27:33:75:72 [ether] on enpos3
? (10.0.2.6) at 08:00:27:d1:f8:5d [ether] on enpos3
root@vboxClone:~# arp -a
? (10.0.2.15) at 08:00:27:cb:61:3b [ether] on enpos3
gateway (10.0.2.1) at 52:54:00:12:35:00 [ether] on enpos3
? (10.0.2.3) at 08:00:27:7b:a2:b6 [ether] on enpos3
? (10.0.2.5) at 08:00:27:33:75:72 [ether] on enpos3
? (10.0.2.6) at 08:00:27:d1:f8:5d [ether] on enpos3
root@vboxClone:~# "C
root@vboxClone:~# "C

```

Figure 1: The two Lubuntu machines

In [Figure 1](#) there are the two lightweight Lubuntu machines. The right machine is performing a ping to the other machine (on the left), and the left machine is running the `arp -a` command to show the devices that are running on this NAT network.

- (ii) The next step is to perform the ARP spoofing attack on the two targets. To do that, on the Kali machine I use the program Ettercap to scan for the two targets and select them as victims, where it will then perform the spoofing attack.

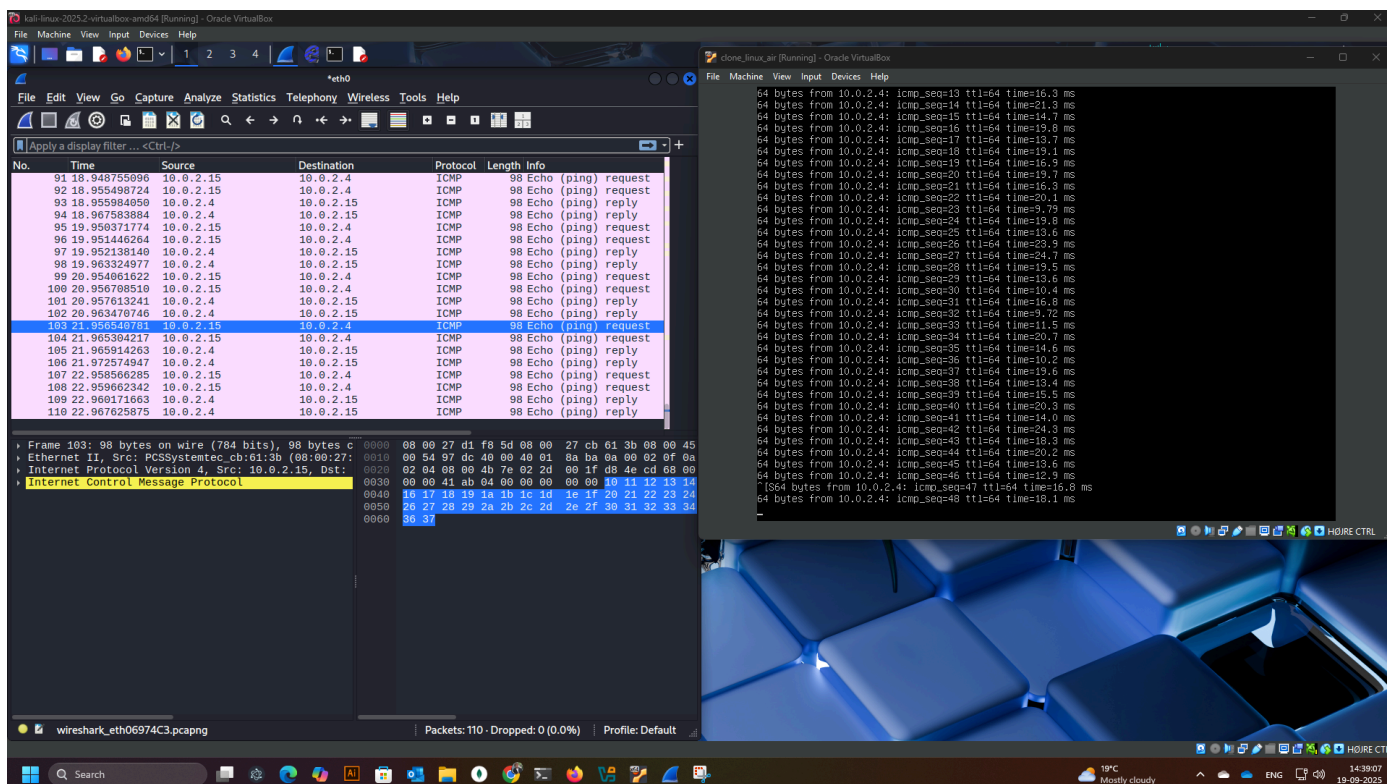


Figure 2: The two Lubuntu machines

In [Figure 2](#), it shows how the attack is under execution, where on the left is the Lubuntu machine that performs a ping to the other Lubuntu machine (on the right). But since we have created a Man-in-the-Middle between the two targets, the traffic can now be seen on the Kali machine, as shown in the image. In this, Wireshark is capturing the traffic between the two machines.