

§ Basics of Higher Cts III

§ 1 Recap on Straightening

why we need it:

most important tool in higher Cts. Theory; gives us

- lots of functors : $R \mapsto R\text{Mod}$
- homotopy coherent algebra : sections + prop

$$\begin{array}{ccc} \text{Assoc} & \dashrightarrow & C^\otimes \\ \parallel \approx & \downarrow & \downarrow \text{Assoc} \end{array}$$

§ 2 Kan extensions

"extend" give functors universally

e.g. $Q\text{coh}(-) : \text{PreSch} \longrightarrow \widehat{\text{Cat}}$ (locally)
needed heavily later on: covergent¹ of finite
type prestacks
coconnective prestacks

§ 3 Finality

Simplify computations for
(∞) limits, Kan extensions

S1 Straightening

John

Given $\Phi: E \rightarrow B$ we may consider
its fibers (for $b: B$)

$$\begin{array}{ccc} E_b & \rightarrow & E \\ \downarrow & \lrcorner & \downarrow \\ \{b\} & \rightarrow & B \end{array}$$

and ask: For which $\Phi: E \rightarrow B$

is

$$B \ni b \mapsto E_b \in \text{Cat}$$

a functorial... note that there
are two ways to be functorial: covariantly or
contravariantly

- $\mathcal{F} \dashv : E \rightarrow B$
- a ω Cartesian fibration \Rightarrow $B \rightarrow \text{Cat}$
 $b \mapsto E_b$
 - a Cartesian fibration $\Rightarrow B^{\text{op}} \rightarrow \text{1Cat}$
 $b \mapsto E_b$

Straightening Unstraightening Then says

$$(\text{Colant}/B)_{\text{gleich}} \xrightleftharpoons[\sim]{} \text{Fun}(B, \text{1Cat})$$

$$(\text{Cart}/B)_{\text{flach}} \xrightleftharpoons[\sim]{} \text{Fun}(B^{\text{op}}, \text{1Cat})$$

Let us spell this out again:

df $\Phi: E \rightarrow B$

- A morphism $e \xrightarrow{f} \tilde{e}$ is called Φ -coCartesian

if $\forall \bar{e} \in E :$

action
of Φ
on mapping
spaces

$$\begin{array}{ccc} \text{Map}_E(\tilde{e}, \bar{e}) & \xrightarrow{f^*} & \text{Map}_E(e, \bar{e}) \\ \downarrow & & \downarrow \\ \text{Map}_B(\Phi\tilde{e}, \Phi\bar{e}) & \xrightarrow{(\Phi f)^*} & \text{Map}_B(\Phi e, \Phi\bar{e}) \end{array}$$

is a pullback square.

- A morphism $e \xrightarrow{f} \tilde{e}$ in E is called Φ -Cartesian if

$\forall \bar{e} \in E$

$$\begin{array}{ccc} \text{Map}_E(\bar{e}, e) & \xrightarrow{f^*} & \text{Map}_E(\bar{e}, \tilde{e}) \\ \downarrow & \lrcorner & \downarrow \\ \text{Map}_B(\Phi\bar{e}, \Phi e) & \longrightarrow & \text{Map}_B(\Phi\bar{e}, \Phi\tilde{e}) \end{array}$$

is a pullback square.

Remark: $f: e \rightarrow \tilde{e}$ being \mathbb{E} -Cartesian means:

$$\begin{array}{ccc} e & \xrightarrow{f} & \tilde{e} \\ \text{Hh} \downarrow & \simeq & \downarrow \exists! g_\alpha \\ \underline{e} & & \end{array} \quad \begin{array}{c} \Phi_{g_\alpha} \simeq \alpha \\ \text{unique up to a} \\ \text{contractible} \\ \text{space of} \\ \text{choices.} \end{array}$$

\downarrow

$$\begin{array}{ccc} \mathbb{E}e & \xrightarrow{\mathbb{E}f, \mathbb{E}\tilde{e}} & \mathbb{E}\tilde{e} \\ \mathbb{E}h \downarrow & \simeq & \downarrow H\alpha \\ \mathbb{E}\tilde{e} & & \end{array}$$

Remark

Note that the definition of a \underline{E} -Galoisian morphism is not circular i.e

can define $f^*: \text{Map}_{\underline{E}}(\tilde{e}, \bar{e}) \rightarrow \text{Map}_{\underline{E}}(e, \bar{e})$
without straightening / unstraightening . . .
Idea: Use "Axiom of ∞ -cats"
 $\text{Fun}([2], E) \xrightarrow{\sim} \underset{12}{\text{Fun}}([1], E) \times \text{Fun}([1], E)$



def $\Phi: E \rightarrow B$ is called

- *coCartesian fib.* if for every $e \in E$ and every morphism $\beta: b \rightarrow \tilde{b}$ w/ $b \simeq \Phi e$,

there exists a Φ -coCartesian lift of β :

$$\text{lift}_0^{\Phi}(e, \beta): e \rightarrow \beta_! e \quad \begin{matrix} \leftarrow \\ \text{notation for} \\ \text{the codomain} \end{matrix}$$

i.e. $\text{lift}_0^{\Phi}(e, \beta): e \rightarrow \beta_! e$ is Φ -coCart. and

$$\Phi(\text{lift}_0^{\Phi}(e, \beta)) \simeq \beta.$$

- Cartenian fibration if for all $\tilde{e} \in E$
and all $\beta: b \rightarrow \tilde{b}$ w/ $\tilde{b} \simeq \Phi \tilde{e}$, there
is a Φ -Cartenian lift

$$\text{lif}_{\tilde{e}}^{\Phi}(\tilde{e}, \beta) : \underbrace{\beta^{*\tilde{e}}}_{\text{domain}}$$

notation for

domain

Rank Many equivalent (model independent) def of
(c) Cartesian fibrations ... see e.g Reehl-Verity
Elements of ∞ -Cats ...

A conceptually very nice equivalent def
 $i : \underline{\Phi} : E \rightarrow B$ is a Colartesian fib

if $\vec{ev}_0^\Phi : E^{[1]} \xrightarrow{\quad} E \overset{\vec{x}}{\times}_{\underline{\Phi}} B$ has a
left adjoint section ... See also
Crossen
Cranshi etc.

Formalization
of Higher Cts

Recall For $B: \text{1Cat}$, one defines

$$\begin{array}{ccc} \text{1Cat}/B & \xrightarrow{\quad} & \text{Fun}([1], \text{1Cat}) \\ \downarrow & \lrcorner & \downarrow \\ * & \longrightarrow & \text{1Cat} \end{array}$$

Objects : functors $E \rightarrow B$

morphisms :

$$\begin{array}{ccc} E & \xrightarrow{\psi} & \tilde{E} \\ & \searrow \approx & \downarrow \\ & & B \end{array}$$

- df
- $\text{Colart}/B \subset 1\text{Cat}/B$ full silicat.
on CoCartesian fibrations on B .
 - $\text{Cart}/B \subset 1\text{Cat}/B$ analogously . . .
 - $(\text{Colart}/B)_{\text{strict 1-full}} \subset \text{Colart}/B$
1-full silicatigally on those morphisms
such that ϕ
sends $\overline{\Phi}$ -CoCartian
morphisms to
 $\widetilde{\Phi}$ -CoCartian
morphisms.

$$\begin{array}{ccc} E & \xrightarrow{\psi} & \widetilde{E} \\ \overline{\Phi} \searrow & \cong & \swarrow \widetilde{\Phi} \\ & B & \end{array}$$

- Analogously : $(\text{Cart}/B)_{\text{strict}} \subset 1\text{Cart}/B$
- $0\text{-CoCart}/B \underset{\text{full}}{\subset} \text{CoCart}/B$ as there coCartesian fibrations $E \xrightarrow{\Phi} B$ such that E_b is a space for every $b \in B$.

Note :

$0\text{-CoCart}/B \cap (\text{CoCart}/B)_{\text{strict}}$
 \Downarrow is an equivalence
 $0\text{CoCart}/B$

- Analogously : $0\text{Cart}/B \underset{\text{full}}{\subset} \text{Cart}/B$

Then

There are equivalences

$$\begin{array}{ccc} \text{Fun}(B, \text{Cat}) & \xrightleftharpoons[\sim]{\text{uf}} & ((\text{Gart}/B)_{\text{strict}} \\ \text{IU} & & \text{IU} \\ \text{Fun}(B, \text{Spc}) & \xrightleftharpoons[\sim]{} & 0-\text{Gart}/B \end{array}$$

Contravariantly natural in B w.r.t/
precomposition on the left and pullback
on the right.

Passing to opposite Cts:

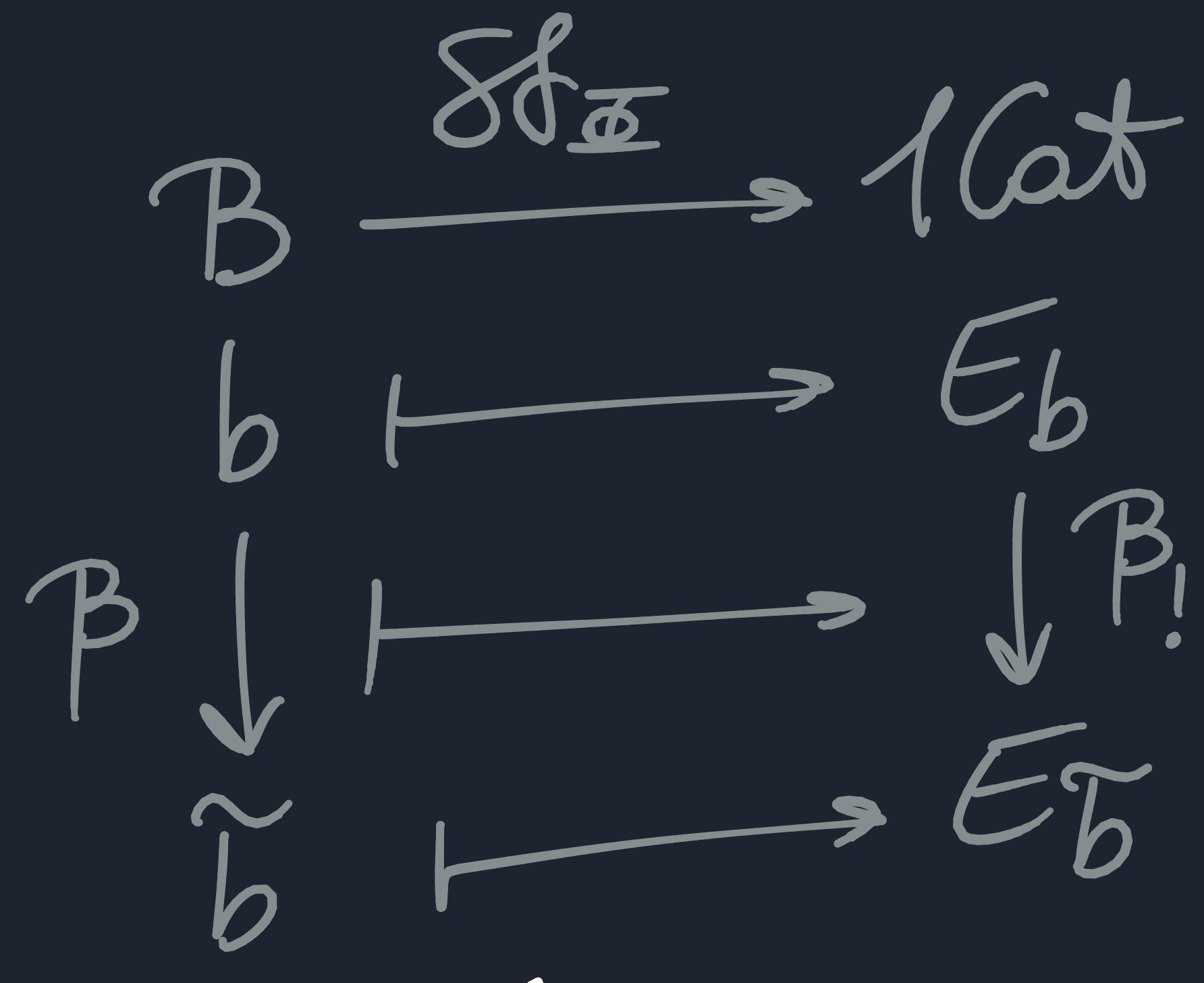
$$\begin{array}{ccc} \text{Fun}(B^{\text{op}}, \text{1Cat}) & \xrightleftharpoons[\sim]{\text{uf}} & (\text{Cat}/B)_{\text{strict}} \\ \text{IU} & & \text{IU} \\ \text{Fun}(B^{\text{op}}, \text{Spc}) & \xrightleftharpoons[\sim]{} & \text{0-Cat}/B \end{array}$$

More explicitly :

Given
obtain

$$\underline{\Phi}: E \rightarrow B$$

coCartesian fibration, we



action on morphism spaces is induced by
 $\underline{\Phi}$ -coCartesian lifts.

Conversely: A functor

$$F: \mathcal{B} \rightarrow \text{Cat}$$

$$\begin{array}{ccc} b & \mapsto & Fb \\ \beta \downarrow & \longmapsto & \downarrow F_\beta \\ \tilde{b} & \mapsto & F\tilde{b} \end{array}$$

gives rise to

a Cartesian fib.

$$\text{Un}_F: \mathcal{B}^f \longrightarrow \mathcal{B}$$

where

\mathcal{B}^f is informally described by:

- objects: $(b \in \mathcal{B}, x: Fb) \iff \begin{array}{c} * \\ \downarrow x \\ Fb \end{array}$

- morphisms:

$$\begin{array}{ccc} b & & \begin{array}{c} * \\ \swarrow x \quad \searrow \tilde{x} \\ Fb \xrightarrow{\quad F_\beta \quad} F\tilde{b} \end{array} \\ \beta \downarrow & & \\ \tilde{b} & & \end{array} \iff \begin{array}{c} F_\beta(x) \xrightarrow{\cong} \tilde{x} \\ \uparrow (\text{Un}_F\text{-Cart} \iff \cong \text{iso}) \end{array}$$

The functor

$$\text{Un}_F: \mathcal{B}^f \longrightarrow \mathcal{B} \quad \text{does the obvious thing.}$$
$$(b, x) \mapsto b$$

Rmk There exists a universal cCartesian fib.

1Cat.



1Cat

such that

$$\begin{array}{ccc} \mathcal{B}^{\mathbb{I}} & \xrightarrow{\quad} & 1\text{Cat.} \\ \downarrow \text{un}_F & \lrcorner & \downarrow \pi^{\text{univ}} \\ \mathcal{B} & \xrightarrow[\quad F \quad]{} & 1\text{Cat} \end{array}$$

"Any cCartesian fib. is obtained as a
pullback from π^{univ} ".

Examples:

- C cat, $c \in C$
- $C_{c/} \rightarrow C$ colCartesian fib
 \downarrow st
 $C(c, -) : C \rightarrow \text{Spc}$
- monoidal cats will be certain colCartesian fib
- C monoidal category, C^+ a left C -module category.
One defines a Cartesian fibration
 $C^+ \rightarrow \Delta^{\text{op}}$
 $\text{Alg}_{+Mod}(C, C^+) \longrightarrow \text{Alg}(C)$
which straightens to $\begin{array}{ccc} \text{Alg}(C)^{\text{op}} & \xrightarrow{\quad} & \hat{\text{Cat}} \\ A & \longmapsto & {}_A\text{Mod}(C^+) \end{array}$

Example (1.5 Yoneida) ← will not
in the book come this
in my talk

Note that both

$\text{Fun}([1], \mathcal{C})$

$\text{ev}_0 \downarrow$

\mathcal{C}

and

$\mathcal{C} \times \mathcal{C}$

$\downarrow \text{proj}_1$

\mathcal{C}

are cartesian fibrations.

An ev_0 -cartesian morphism in $\text{Fun}([1], \mathcal{C})$ is a
a morphism $d: f \rightarrow f'$ s.t.

\Downarrow

$$\begin{array}{ccc} c_0 & \xrightarrow{d_0} & c'_0 \\ f \downarrow & \simeq & \downarrow f' \\ c_1 & \xrightarrow{d_1} & c'_1 \end{array}$$

$$\text{ev}_1(d) = d_1$$

is an isomorphism.

Analogously: proj_1 -cartesian morphism
is a morphism $(d_0, d_1): (c_0, c_1) \rightarrow (c'_0, c'_1)$
such that $\text{proj}(d_0, d_1) = d_1$ is an iso.

In particular, the commutative triangle

$\text{Fun}([1], \mathcal{C}) \xrightarrow{(\text{ev}_0, \text{ev}_1)} \mathcal{C} \times \mathcal{C}$ is a

$$\begin{array}{ccc} \text{ev}_0 & \simeq & \text{proj}_1 \\ \searrow & & \swarrow \end{array}$$

morphism

in

$(\text{Cart}/\mathcal{C})_{\text{strict}}$

$\simeq \text{Fun}(\mathcal{C}^{\text{op}}, \text{1Cat})$

$(\text{ev}_0, \text{ev}_1)$ charly sends

ev_0 -cartesian morphisms to proj_1 -cartesian morphisms

Straightening this morphism yields a natural transformation

$$\otimes (c \mapsto c_C) \Rightarrow \subseteq$$

Components are

$$C_C \xrightarrow{ev_1} C$$

by def of straightening

constant diagram functor

$$C^{\text{op}} \rightarrow \text{1Cat}$$

$$x \mapsto C$$

will not speak about this

Note that :

$$\text{Fun}(C^{\text{op}}, \text{1Cat})/\bar{C} \xrightarrow{\cong} \text{Fun}(C^{\text{op}}, \text{1Cat}/C)$$

and hence

\otimes gives us a functor

$$\text{** } C^{\text{op}} \rightarrow \text{1Cat}/C$$

$$c \mapsto (C_C \xrightarrow{ev_1} C)$$

o-Cartesian fib
w/ fiber $\text{Map}_C(c, x) \in \text{Spc}$
at x

and therefore (**) factors as

$$C^{\text{op}} \rightarrow \text{1Cat}/C$$

$$\xrightarrow{\sim} \text{ind}$$

$$o\text{-Cart}/C$$

$$\cong$$

$$\text{Fun}(C, \text{Spc})$$

So by carrying over
get the Hom functor $C^{\text{op}} \times C \rightarrow \text{Spc}$.

$$\text{Map}_C$$

1.6
in the
book

gives an enhanced Straightening/
Unstraightening equivalence



is a
Cartesian
fibration.

Analogously

Cart

\downarrow

1Cat

\Downarrow

$B \mapsto \text{Fun}(B^{\text{op}}, \text{1Cat})$

12

(CoCat/B) strict

* classifies $\text{1Cat}^{\text{op}} \rightarrow \hat{\text{1Cat}}$

$B \mapsto \text{Fun}(B, \text{1Cat}) \simeq (\text{CoCat}/B)^{\text{strict}}$

makes precise that Straightening/Unstraightening
is contravariantly natural in B ...

1.7 Adjoint functors

$F: C_0 \rightarrow C_1$ functor

$\Leftrightarrow [1] \xrightarrow{F} \text{1Cat}$

unstraighten

$\Leftrightarrow \tilde{C} \xrightarrow{\tilde{F}} [1]$ Cartesian fibration

of $F: C_0 \rightarrow C_1$ admits a right adjoint
if $\tilde{C} \xrightarrow{\tilde{F}} [1]$ is also a Cartesian
fibration (i.e. if is biCartesian).

Assume $F: C_0 \rightarrow C_1$ admits a RA. Then we may apply Cartesian straightening to

$$\tilde{C} \xrightarrow{\tilde{F}} [1] \text{ to obtain}$$

$$\begin{array}{ccc} [1]^{\text{op}} & \xrightarrow{\text{forget}} & \text{hence we get a} \\ 0 & \mapsto & C_0 \\ \uparrow & \mapsto & \uparrow F^R \\ 1 & \mapsto & C_1 \end{array} \quad \text{functor} \quad C_1 \xrightarrow{F^R} C_0.$$

Exercise: Show that for all $x \in C_0$ $y \in C_1$ one has

$$\text{Map}_{C_1}(Fx, y) \simeq \text{Map}_{C_0}(x, F^Ry)$$

1.7.3 Partial Adjoints

$\tilde{C} \xrightarrow{\tilde{F}} [1]$ as before without assuming
 \tilde{F} is a Cartesian fibration.

Let $C'_1 \subset C_1$ full subcat. on those
 $e_1, e \in C_1$ such that there
exists an \tilde{F} -Cartesian morphism
 $c_0 \rightarrow c_1$ in \tilde{C} carrying $0 \rightarrow 1$ in $[1]$.

Let $\tilde{C}' \subset \tilde{C}$ full subcat so that
 $\tilde{C}'_0 = C_0$
 $\tilde{C}'_1 = C'_1$

Then :

$\tilde{C}' \rightarrow [1]$ is a Cartesian fibration
⇒ partially defined right adjoint

$$C_1^I \xrightarrow{F'^R} C_0$$

$$\text{Map}_{C_1}(F_{C_0}, c_1) \simeq \text{Map}_{C_0}(c_0, F'^R c_1)$$

$$\text{For all } c_0 \in C_0 \\ c_1 \in C_1^I$$

2. Basic Operations on ∞ -Cts

Q1 How to produce new ∞ -Cts from old ones? (e.g. direct Cts)

Q2 How to construct new functors from old ones?

dt1 (co)limits e.g. localizations
→ direct ∞ -Ct

dt2 Kan extensions e.g. construction of $\mathcal{L}\text{ab}(-)$

2.1.1

$$F: \mathcal{D} \rightarrow \mathcal{C}$$

For any category \mathcal{E} we have

$$\text{Fun}(\mathcal{C}, \mathcal{E}) \xrightarrow{F^*} \text{Fun}(\mathcal{D}, \mathcal{E})$$

precomposition functor

2.1.2

def

Partially defined left (resp. right)

adjoint of F^* is called

left (resp. right) Kan extension

along F and denoted by

LKE_F (resp. RKE_F)

Example we will see many examples of
Kan extensions in future sessions;

Easy examples are given by

(co) limits :

If $C = *$, $F = (\mathcal{D} \xrightarrow{*} *)$

Then left and right Kan extension
give

$\underset{\mathcal{D}}{\text{Colim}} : \text{Fun}(\mathcal{D}, E) \rightarrow E$

$\underset{\mathcal{D}}{\text{lim}} : \text{Fun}(\mathcal{D}, E) \rightarrow E$

(partially defined)

2.1.3 Formulas for Kan extensions

Lst

$$\begin{array}{ccc} \mathcal{D} & \xrightarrow{\Phi} & E \\ F \downarrow & & \\ C & & \end{array}$$

Consider for $c \in C$

$$\begin{array}{ccc} F\downarrow c & \simeq & \mathcal{D} \times_{C/c} C/c \longrightarrow C/c \\ \pi_D \downarrow & & \downarrow \\ \mathcal{D} & \xrightarrow{F} & C \end{array}$$

and suppose that for every $c \in C$
 notation $\{\underset{\mathcal{D} \times_{C/c}}{\operatorname{Colim}} \Phi := \operatorname{Colim} (\Phi \circ \pi_D)\}_{c \in C}$ exists.

objects : $Fd \xrightarrow{\alpha} c$
 morphisms : $\begin{array}{ccc} & d & \\ \lambda & \downarrow & \tilde{d} \\ Fd & \xrightarrow{F\lambda} & F\tilde{d} \\ \alpha & \searrow \tilde{\alpha} & \downarrow c \end{array}$

Then : $LKE_F \Phi : C \rightarrow E$ exists and
an object is given by

$$c \mapsto \underset{\mathcal{D} \times C/c}{\operatorname{colim}} \Phi \simeq LKE_F \Phi(c)$$

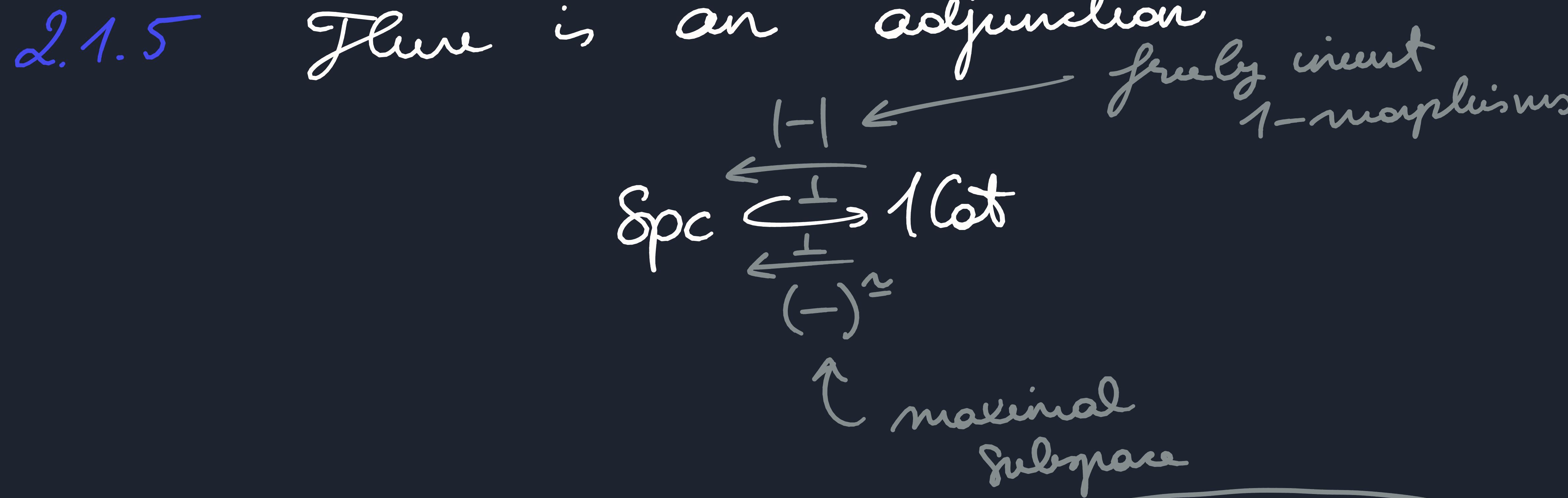
Similarly :

$$RKE_F \Phi(c) \simeq \underset{\mathcal{D} \times C/c}{\operatorname{lim}} \Phi$$

Exercise :

$$\underset{\mathcal{D}}{\operatorname{colim}} \Phi \simeq \underset{C}{\operatorname{colim}} LKE_F(\Phi)$$

$$\underset{\mathcal{D}}{\operatorname{lim}} \Phi \simeq \underset{C}{\operatorname{lim}} RKE(\Phi)$$



Note that

$$|C| \simeq \operatorname{colim}_C *$$

↑
constant
diagram
functor

$$C \rightarrow \text{Spc}$$

$$c \mapsto *$$

For any $S \in \text{Spc}$ we have

$$|S| \simeq S.$$

Indeed, this may be proven by straightening - unstraightening:

$$\begin{aligned} & \operatorname{Map}_{\text{Spc}}(\operatorname{colim}_C *, S) \\ & \simeq \lim_C \operatorname{Map}_{\text{Spc}}(*, S) \\ & \simeq \lim_C \subseteq \simeq \Gamma(C \times S \rightarrow C) \\ & \simeq \operatorname{Fun}(C, S) \\ & \simeq \operatorname{Map}_{\text{Spc}}(C, S) \end{aligned}$$

as wanted

will not
come up in
lecture

df $C:\mathbf{Cat}$ is called Contractible if

$$|C| \simeq *$$

2.1.6 If $\underline{\Phi} : C \rightarrow \mathbf{Spc}$, then

$$\operatorname{Colim}_C \underline{\Phi} \simeq |\tilde{C}_{\underline{\Phi}}| \text{ where}$$

$\tilde{C}_{\underline{\Phi}} \rightarrow C$ is the Cartesian fib. assoc.
to $\underline{\Phi}$.

Lemma

$$C \in 1\text{Cat}, \quad PC := \text{Fun}(C^{\text{op}}, \text{Spc})$$

Then restriction and lift Kan extension
along the Yoneda embedding

$$C \xrightarrow{\mathcal{L}} PC$$

induce for all cocomplete D an
equivalence

$$\text{Fun}^{\text{colim}}(PC, D) \xrightleftharpoons[\sim]{\mathcal{L}^*, \text{LKE}_{\mathcal{L}}} \text{Fun}(C, D)$$

exhibiting PC as the free cocompletion of C .

In particular: $\text{Fun}^{\text{colim}}(\text{Spc}, D) \simeq D$.

§3 Cfinality

2.2.1

def $F: D \rightarrow C$ is called cfinal if for
any $c: C$, the category
 $c \downarrow F \simeq D \times_C C_{c/}$ is contractible i.e.

$$|D \times_C C_{c/}| \simeq *$$

(and cinitial if $F^{\text{op}}: D^{\text{op}} \rightarrow C^{\text{op}}$ is cfinal)
i.e $|F \downarrow c| \simeq *$

2.2.2

Lemma TFAE :

(i) $F: D \rightarrow C$ is cofinal.

(ii) For any $\Phi: C \rightarrow E$, the natural map

$$\text{colim}_D (\Phi \circ F) \xrightarrow{\sim} \text{colim}_C \Phi$$

is an isomorphism.

(iii) For any $\Phi: C^{\text{op}} \rightarrow E$, the natural map

$$\lim_{C^{\text{op}}} \Phi \xrightarrow{\sim} \lim_{D^{\text{op}}} (\Phi \circ F^{\text{op}})$$

is an iso.

In the book there is also : (ii'), (ii''), (iii'), (iv)

2.2.3

Example

- Let C be a category w/ a terminal object $t \in C$. Then the functor

$* \xrightarrow{t} C$ is cofinal i.e

for $\overline{\Phi}: C \rightarrow E$ we have

$$\operatorname{Colim}_C \overline{\Phi} \simeq \operatorname{Colim}_* \overline{\Phi}(t) \simeq \overline{\Phi}(t)$$

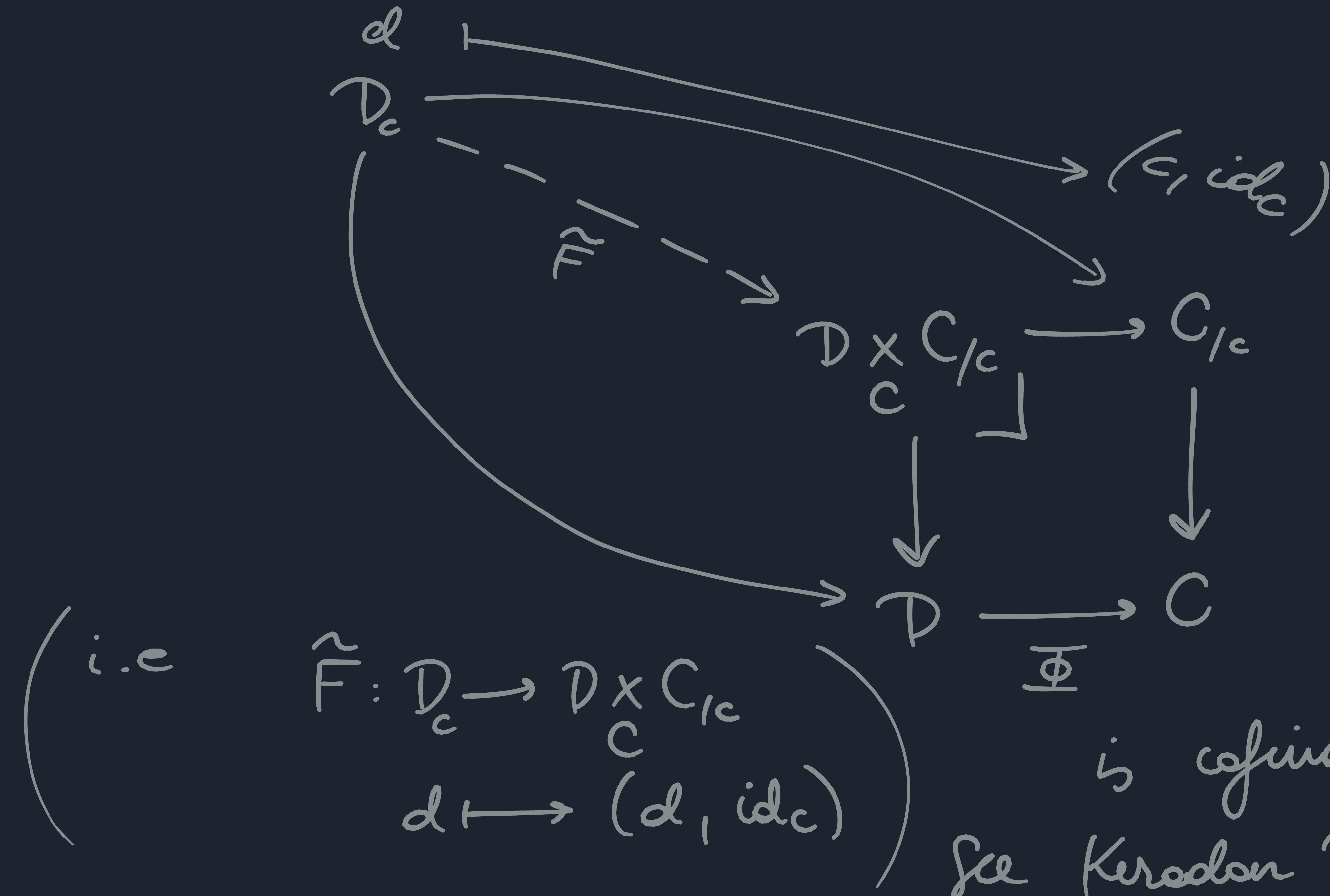
- More generally: Any functor $F: D \rightarrow C$ that admits a left adjoint (i.e F is a right adjoint) is cofinal.

Indeed: $D \times_C C_{c/} \simeq c \downarrow F$ has an initial object for all c and hence $|D \times_C C_{c/}| \simeq *$.

Example / Lemma

$$\mathcal{D} \xrightarrow{F} C \quad \text{Cartesian fibration}$$

Then the canonical functor



is cofinal.
See Kerodon Prop 7.3.4.1

Hence we get:

$$LKE_F \bar{\Phi}(c) \simeq \underset{\mathcal{D} \times C/c}{\operatorname{Colim}} \bar{\Phi} \simeq \underset{\mathcal{D}_c}{\operatorname{Colim}} \bar{\Phi}$$

Similarly, if $\mathcal{D} \xrightarrow{F} C$ is a Cartesian fib.

$$RKE_F \bar{\Phi}(c) \simeq \underset{\mathcal{D}_c}{\operatorname{lim}} \bar{\Phi}$$

2.3 The book also says something on
controllable functors ...
we will skip this for now