

Dictionary-based model reduction for state estimation

Alexandre PASCO

Centrale Nantes, Nantes Université, France
In collaboration with Anthony NOUY.



Context

Context

- Parametric equation in a Hilbert space U ,

$$\mathcal{F}(u, \xi) = 0,$$

$u \in U$ the **state to recover** and $\xi \in \mathcal{P} \subset \mathbb{R}^d$ an **unknown** parameter.

Context

- Parametric equation in a Hilbert space U ,

$$\mathcal{F}(u, \xi) = 0,$$

$u \in U$ the **state to recover** and $\xi \in \mathcal{P} \subset \mathbb{R}^d$ an **unknown** parameter.

Observations

of m continuous linear measurements

$\ell_1(u), \dots, \ell_m(u)$, i.e. $w = P_W u$

with $W := \text{span}\{R_U \ell_i\}_{i=1}^m$

Context

- Parametric equation in a Hilbert space U ,

$$\mathcal{F}(u, \xi) = 0,$$

$u \in U$ the **state to recover** and $\xi \in \mathcal{P} \subset \mathbb{R}^d$ an **unknown** parameter.

Observations

+

Prior knowledge

of m continuous linear measurements
 $\ell_1(u), \dots, \ell_m(u)$, i.e. $\mathbf{w} = P_W u$
with $W := \text{span}\{R_U \ell_i\}_{i=1}^m$

on the solution manifold
 $\mathcal{M} := \{u(\xi) : \xi \in \mathcal{P}\}$ based
on model reduction

Context

- Parametric equation in a Hilbert space U ,

$$\mathcal{F}(u, \xi) = 0,$$

$u \in U$ the **state to recover** and $\xi \in \mathcal{P} \subset \mathbb{R}^d$ an **unknown** parameter.

Observations

+

Prior knowledge

of m continuous linear measurements

$\ell_1(u), \dots, \ell_m(u)$, i.e. $w = P_W u$

with $W := \text{span}\{R_U \ell_i\}_{i=1}^m$

on the solution manifold

$\mathcal{M} := \{u(\xi) : \xi \in \mathcal{P}\}$ based

on model reduction



Compute a recovery $A(w) \simeq u$

One space approach

Linear MOR

One space: formulation

- Approximate \mathcal{M} by a linear subspace V ,

$$\text{dist}(V, \mathcal{M}) \leq \varepsilon, \quad \dim(V) = n$$

One space: formulation

- Approximate \mathcal{M} by a linear subspace V ,

$$\text{dist}(V, \mathcal{M}) \leq \varepsilon, \quad \dim(V) = n$$

- Parameterized-Background Data-Weak (PBDW) [Maday et al., 2015]:

$$A_V(\textcolor{blue}{w}) := v^* + \eta^*, \quad v^* := \arg \min_{v \in V} \|P_W(u - v)\|, \quad \eta^* := \textcolor{blue}{w} - P_W v^*$$

One space: formulation

- Approximate \mathcal{M} by a linear subspace V ,

$$\text{dist}(V, \mathcal{M}) \leq \varepsilon, \quad \dim(V) = n$$

- Parameterized-Background Data-Weak (PBDW) [Maday et al., 2015]:

$$A_V(\textcolor{blue}{w}) := v^* + \eta^*, \quad v^* := \arg \min_{v \in V} \|P_W(u - v)\|, \quad \eta^* := \textcolor{blue}{w} - P_W v^*$$

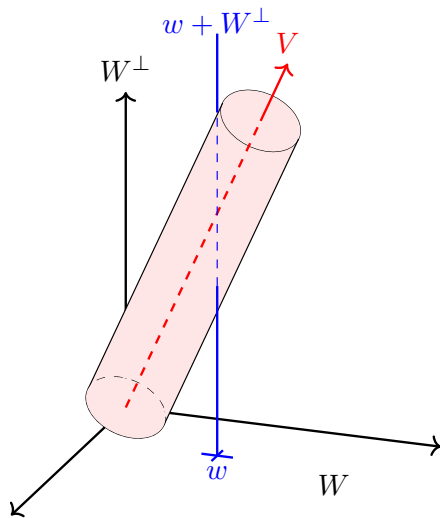
- Sharp error bound [Binev et al., 2017]: $\|u - A_V(\textcolor{blue}{w})\| \leq \varepsilon \mu(V, W)$

$$\mu(V, W) := \left(\inf_{v \in V} \sup_{w \in W} \frac{\langle v, w \rangle}{\|v\| \|w\|} \right)^{-1}$$

One space: Geometry

Linear MOR

$$\text{dist}(V, \mathcal{M}) \leq \varepsilon$$



One space: Pros and cons

Pros

- Online efficiency $\mathcal{O}(n^3)$.
- No need to know ε .
- Optimal (worst case sense)
when \mathcal{M} is a cylinder
centered in V
[Binev et al., 2017].

One space: Pros and cons

Pros

- Online efficiency $\mathcal{O}(n^3)$.
- No need to know ε .
- Optimal (worst case sense) when \mathcal{M} is a cylinder centered in V [Binev et al., 2017].

Cons

- Requires $n \leq m$.
- Trade-off between ε and $\mu(V, W)$.
- Limited by the Kolmogorov m -width, $\varepsilon \geq d_m(\mathcal{M})$ where

$$d_n(\mathcal{M}) := \inf_{\substack{X \subset U \\ \dim X = n}} \max_{u \in \mathcal{M}} \text{dist}(X, u).$$

Multi-space approach

Library-based MOR

Multi-space: library-based MOR

- Consider a library of spaces $\mathcal{L}_n^N := \{V_1, \dots, V_N\}$

$$\text{dist}\left(\bigcup_{k=1}^N V_k, \mathcal{M}\right) \leq \varepsilon, \quad \dim(V_k) \leq n \leq m.$$

Multi-space: library-based MOR

- Consider a library of spaces $\mathcal{L}_n^N := \{V_1, \dots, V_N\}$

$$\text{dist}\left(\bigcup_{k=1}^N V_k, \mathcal{M}\right) \leq \varepsilon, \quad \dim(V_k) \leq n \leq m.$$

- New benchmark: non-linear Kolmogorov (n, N) -width [Temlyakov, 1998]

$$d_n(\mathcal{M}, N) := \inf_{\#\mathcal{L}_n^N = N} \max_{u \in \mathcal{M}} \min_{V \in \mathcal{L}_n^N} \text{dist}(V, u),$$

which is expected to decay much faster than $d_n(\mathcal{M})$.

Multi-space: library-based MOR

- Consider a library of spaces $\mathcal{L}_n^N := \{V_1, \dots, V_N\}$

$$\text{dist}\left(\bigcup_{k=1}^N V_k, \mathcal{M}\right) \leq \varepsilon, \quad \dim(V_k) \leq n \leq m.$$

- New benchmark: non-linear Kolmogorov (n, N) -width [Temlyakov, 1998]

$$d_n(\mathcal{M}, N) := \inf_{\#\mathcal{L}_n^N = N} \max_{u \in \mathcal{M}} \min_{V \in \mathcal{L}_n^N} \text{dist}(V, u),$$

which is expected to decay much faster than $d_n(\mathcal{M})$.

- Aim for low ε with a low n , thus better stability.

Multi-space: How ?

Each space $V \in \mathcal{L}_n^N$ gives a one-space estimate $A_V(w)$...

Multi-space: How ?

Each space $V \in \mathcal{L}_n^N$ gives a one-space estimate $A_V(w)$...



How to select $V^*(w)$
among this library ?

Multi-space: Selection

Idea [Cohen et al., 2022]: select the “closest” to \mathcal{M}

Multi-space: Selection

Idea [Cohen et al., 2022]: select the “closest” to \mathcal{M}

- Assume we have \mathcal{S} such that for any $v \in U$,

$$c \operatorname{dist}(v, \mathcal{M}) \leq \mathcal{S}(v, \mathcal{M}) \leq C \operatorname{dist}(v, \mathcal{M})$$

Multi-space: Selection

Idea [Cohen et al., 2022]: select the “closest” to \mathcal{M}

- Assume we have \mathcal{S} such that for any $v \in U$,

$$c \operatorname{dist}(v, \mathcal{M}) \leq \mathcal{S}(v, \mathcal{M}) \leq C \operatorname{dist}(v, \mathcal{M})$$

- Select $V^*(w)$ as

$$V^*(w) \in \operatorname{argmin}_{V \in \mathcal{L}_n^N} \mathcal{S}(A_V(w), \mathcal{M}).$$

Multi-space: Selection

Proposition (Near optimal selection [Cohen et al., 2022])

Assuming that P_W is injective on \mathcal{M} and that $\mu(\mathcal{M}, W) < \infty$,

$$\|u - A_{V^*}(w)\| \leq 2 \frac{C}{c} \mu(\mathcal{M}, W) \min_{V \in \mathcal{L}_n^N} \|u - A_V(w)\|,$$

$\mu(\mathcal{M}, W)$ reflects how well \mathcal{M} and W are aligned.

Dictionary approach

Dictionary-based MOR

Dictionary approach: Library

- Dictionary of K vectors (or snapshots),

$$\mathcal{D}_K = \{v^{(1)}, \dots, v^{(K)}\}$$

Dictionary approach: Library

- Dictionary of K vectors (or snapshots),

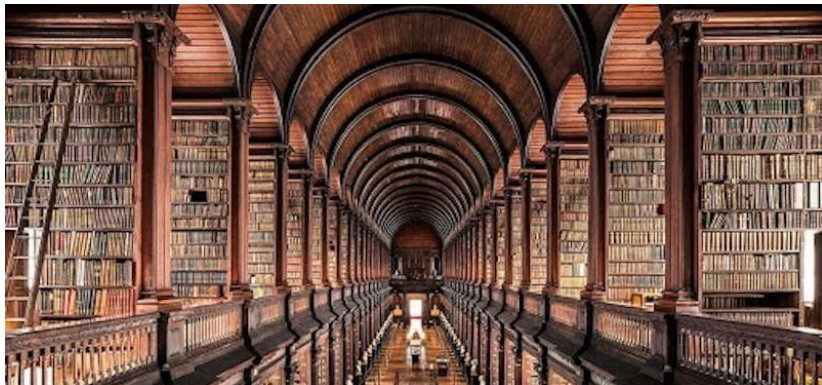
$$\mathcal{D}_K = \{v^{(1)}, \dots, v^{(K)}\}$$

- Take as library $\mathcal{L}_m^N = \mathcal{L}_m(\mathcal{D}_K)$, containing all the subspaces spanned by at most m vectors from \mathcal{D}_K ,

$$\mathcal{L}_m(\mathcal{D}_K) := \left\{ \sum_{k=1}^K x_k v^{(k)} : x \in \mathbb{R}^K, \|x\|_0 \leq m \right\},$$

Dictionary approach: Library

$\mathcal{L}_m(\mathcal{D}_K)$ is **large** \rightarrow low ε but not fully explorable.



Dictionary approach: Compressive sensing

Sparse approximation given a
few linear measurements

Dictionary approach: Compressive sensing

Sparse approximation given a
few linear measurements



Compressive Sensing

Dictionary approach: Compressive sensing

Sparse approximation given a few linear measurements \rightarrow Compressive Sensing

- Consider the Basis Pursuit Denoising problem for $\alpha > 0$

$$\mathbf{x}_\alpha(\mathbf{w}) := \operatorname{argmin}_{\mathbf{x} \in \mathbb{R}^K} \frac{1}{2} \|\mathbf{C}\mathbf{x} - \mathbf{w}\|_2^2 + \alpha \|\mathbf{x}\|_1,$$

whose solution is unique and m -sparse under some assumptions.

Dictionary approach: Compressive sensing

Sparse approximation given a few linear measurements \rightarrow Compressive Sensing

- Consider the Basis Pursuit Denoising problem for $\alpha > 0$

$$\mathbf{x}_\alpha(\mathbf{w}) := \operatorname{argmin}_{\mathbf{x} \in \mathbb{R}^K} \frac{1}{2} \|\mathbf{C}\mathbf{x} - \mathbf{w}\|_2^2 + \alpha \|\mathbf{x}\|_1,$$

whose solution is unique and m -sparse under some assumptions.

- Focus on the space spanned by the support,

$$V_\alpha(\mathbf{w}) := \operatorname{span}\{v^{(i)} : \mathbf{x}_\alpha(\mathbf{w})_i \neq 0\} \in \mathcal{L}_m(\mathcal{D}_K)$$

Dictionary approach: Selection

- Use \mathcal{S} to select $V_{\mathcal{S}}(\boldsymbol{w})$ among the **smaller** library $\mathcal{L}(\boldsymbol{w}) := \{V_{\alpha}(\boldsymbol{w}) : \alpha > 0\}$ and define

$$A_{\mathcal{S}}^{\text{dic}}(\boldsymbol{w}) := A_{V_{\mathcal{S}}(\boldsymbol{w})}(\boldsymbol{w})$$

Dictionary approach: Selection

- Use \mathcal{S} to select $V_{\mathcal{S}}(w)$ among the **smaller** library $\mathcal{L}(w) := \{V_{\alpha}(w) : \alpha > 0\}$ and define

$$A_{\mathcal{S}}^{\text{dic}}(w) := A_{V_{\mathcal{S}}(w)}(w)$$

- In practice for efficiency and numerical stability reasons, we only compute $\hat{\mathcal{L}}(w) \subset \mathcal{L}(w)$ with $\#\hat{\mathcal{L}}(w) \leq \tau K$.

Dictionary approach: Selection

- Use \mathcal{S} to select $V_{\mathcal{S}}(\boldsymbol{w})$ among the **smaller** library $\mathcal{L}(\boldsymbol{w}) := \{V_{\alpha}(\boldsymbol{w}) : \alpha > 0\}$ and define

$$A_{\mathcal{S}}^{\text{dic}}(\boldsymbol{w}) := A_{V_{\mathcal{S}}(\boldsymbol{w})}(\boldsymbol{w})$$

- In practice for efficiency and numerical stability reasons, we only compute $\hat{\mathcal{L}}(\boldsymbol{w}) \subset \mathcal{L}(\boldsymbol{w})$ with $\#\hat{\mathcal{L}}(\boldsymbol{w}) \leq \tau K$.

Proposition (Near optimal selection)

Assuming that P_W is injective on \mathcal{M} and that $\mu(\mathcal{M}, W) < \infty$,

$$\|u - A_{\mathcal{S}}^{\text{dic}}(\boldsymbol{w})\| \leq 2 \frac{C}{c} \mu(\mathcal{M}, W) \min_{V \in \hat{\mathcal{L}}(\boldsymbol{w})} \|u - A_V(\boldsymbol{w})\|,$$

Parameterized PDEs

Offline-online decomposition for dictionary-based multi-space

Parameterized PDEs: Offline-online decomposition

- Discrete framework $U = \mathbb{R}^{\mathcal{N}}$ with large \mathcal{N} .
- **Offline:** Heavy pre-computations independently on w .
- **Online:** Fast computation of $A_S^{\text{dic}}(w)$.

Parameterized PDEs: Framework

- Consider that u solves the operator equation

$$B(\xi)u(\xi) = f(\xi),$$

where the singular values of $B(\xi)$ are **uniformly bounded**,

$$0 < c \leq \min_{v \in U} \frac{\|B(\xi)v\|}{\|v\|} \leq \max_{v \in U} \frac{\|B(\xi)v\|}{\|v\|} \leq C < \infty,$$

Parameterized PDEs: Framework

- Consider that u solves the operator equation

$$B(\xi)u(\xi) = f(\xi),$$

where the singular values of $B(\xi)$ are **uniformly bounded**,

$$0 < c \leq \min_{v \in U} \frac{\|B(\xi)v\|}{\|v\|} \leq \max_{v \in U} \frac{\|B(\xi)v\|}{\|v\|} \leq C < \infty,$$

- Take \mathcal{S} as a **residual norm** as in [Cohen et al., 2022],

$$\mathcal{S}(v, \mathcal{M}) := \min_{\xi \in \mathcal{P}} \|B(\xi)v - f(\xi)\|, \quad \forall v \in U,$$

which satisfies $c \operatorname{dist}(v, \mathcal{M}) \leq \mathcal{S}(v, \mathcal{M}) \leq C \operatorname{dist}(v, \mathcal{M})$.

Parameterized PDEs: Affine decomposition

- Assume the **affine decompositions**

$$B(\xi) = B_0 + \sum_{q=1}^d \theta_q^{(B)}(\xi) B_q \quad \text{and} \quad f(\xi) = f_0 + \sum_{q=1}^{m_f} \theta_q^{(f)}(\xi) f_q,$$

with $B_q : U \longrightarrow U'$ linear operators and $f_q \in U'$.

Parameterized PDEs: Affine decomposition

- Assume the **affine decompositions**

$$B(\xi) = B_0 + \sum_{q=1}^d \theta_q^{(B)}(\xi) B_q \quad \text{and} \quad f(\xi) = f_0 + \sum_{q=1}^{m_f} \theta_q^{(f)}(\xi) f_q,$$

with $B_q : U \longrightarrow U'$ linear operators and $f_q \in U'$.

- Then computing \mathcal{S} requires solving the l.s. system

$$\mathcal{S}(v, \mathcal{M}) := \min_{\xi \in \mathcal{P}} \|G(v)\theta(\xi) - g(v)\|, \quad v \in U,$$

where $\theta(\xi) \in \mathbb{R}^{m_B+m_f}$, $G(v) : \mathbb{R}^{m_B+m_f} \longrightarrow U'$ and $g(v) \in U'$.

Parameterized PDEs: Offline-online decomposition

- **Problem:** precomputing the normal equation for \mathcal{S} costs

$$\mathcal{O}((m_B K + m_f)^2 \mathcal{N})$$

Parameterized PDEs: Offline-online decomposition

- **Problem:** precomputing the normal equation for \mathcal{S} costs

$$\mathcal{O}((m_B K + m_f)^2 \mathcal{N})$$

- **Instead,** we use a random embedding $\Theta \in \mathbb{R}^{k \times \mathcal{N}}$ and consider

$$\mathcal{S}^\Theta(v, \mathcal{M}) := \min_{\xi \in \mathcal{P}} \|\Theta(B(\xi)v - f(\xi))\| = \min_{\xi \in \mathcal{P}} \|G^\Theta(v)\theta(\xi) - g^\Theta(v)\|$$

where $G^\Theta(v) \in \mathbb{R}^{k \times (m_B + m_f)}$ and $g^\Theta(v) \in \mathbb{R}^k$.

Parameterized PDEs: Offline-online decomposition

- **Problem:** precomputing the normal equation for \mathcal{S} costs

$$\mathcal{O}((m_B K + m_f)^2 \mathcal{N})$$

- **Instead,** we use a random embedding $\Theta \in \mathbb{R}^{k \times \mathcal{N}}$ and consider

$$\mathcal{S}^\Theta(v, \mathcal{M}) := \min_{\xi \in \mathcal{P}} \|\Theta(B(\xi)v - f(\xi))\| = \min_{\xi \in \mathcal{P}} \|G^\Theta(v)\theta(\xi) - g^\Theta(v)\|$$

where $G^\Theta(v) \in \mathbb{R}^{k \times (m_B + m_f)}$ and $g^\Theta(v) \in \mathbb{R}^k$.

Proposition

With $k = \mathcal{O}\left(\epsilon^{-2} (m_B + m_f + \log(\delta^{-1}))\right)$, for any $v \in U$, with probability at least $1 - \delta$ we have

$$\sqrt{1 - \epsilon} \mathcal{S}(v, \mathcal{P}) \leq \mathcal{S}^\Theta(v, \mathcal{P}) \leq \sqrt{1 + \epsilon} \mathcal{S}(v, \mathcal{P}).$$

Parameterized PDEs: Computational aspects

- **Offline cost:** Using a structured embedding Θ (e.g. SRHT),

$$\mathcal{O}\left(\underbrace{mK\mathcal{N}}_{\text{compute } \mathbf{C}} + \underbrace{(m_B K + m_f)\mathcal{N} \log(k)}_{\text{pre-compute } G^\Theta(v) \text{ and } g^\Theta(v)} \right)$$

- **Online cost:** Considering $k = \mathcal{O}(m_B + m_f)$,

$$\mathcal{O}\left(\underbrace{m^2 K}_{\text{LARS}} + \underbrace{k(m m_B + m_f) K}_{\text{prepare l.s.}} + \underbrace{C_{ls} K}_{\text{solve l.s.}} \right)$$

- **Numerical stability** : few affine terms, thus robust to round-off errors.

Parameterized PDEs: Computational aspects

Example when $\theta(\xi) = \xi$ and $f(\xi) = f_0$.

- **Offline cost:**

$$\begin{array}{ccc} \text{With random sketching} & \text{VS} & \text{With normal equation} \\ \mathcal{O}(mKN + dKN \log(k)) & & \mathcal{O}(mKN + d^2 K^2 N) \end{array}$$

- **Online cost:** Considering $k = \mathcal{O}(d)$,

$$\begin{array}{ccc} \text{With random sketching} & \text{VS} & \text{With normal equation} \\ \mathcal{O}(m^2 K + md^2 K + d^3 K) & & \mathcal{O}(m^2 K + m^2 d^2 K + d^3 K) \end{array}$$

Numerical Example

Finite Element space $U \subset \mathcal{H}_{\Gamma_D}^1(\Omega)$ endowed with norm $\|\nabla \cdot\|_{\mathcal{L}^2(\Omega)}$

Numerical: Thermal block

$$\mathcal{N} \sim 8\,000 \quad \text{and} \quad \begin{cases} -\nabla \cdot (\kappa \nabla u) = 1 & \text{in } \Omega, \\ u = 0 & \text{on } \Gamma_D, \\ \kappa = \xi_i & \text{in } \Omega_i, \, 1 \leq i \leq 9, \end{cases}$$

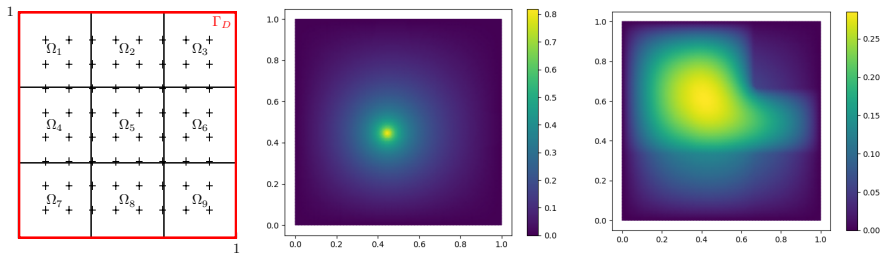


Figure: Left: geometry, with sensors locations (crosses). Middle: Riesz representative of a sensor. Right: a snapshot.

Numerical: Thermal block

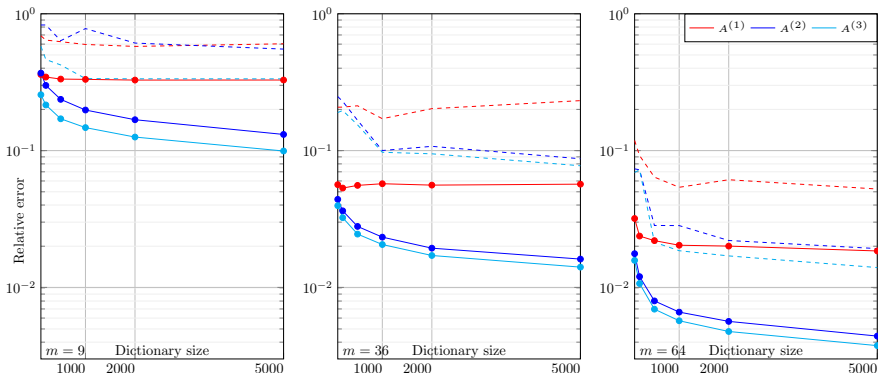
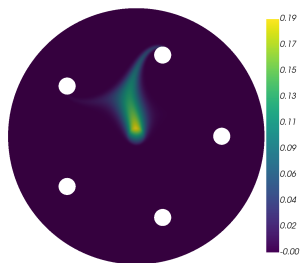
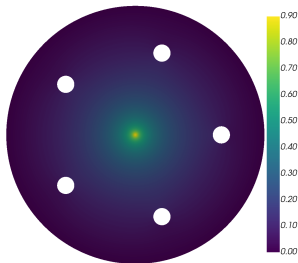
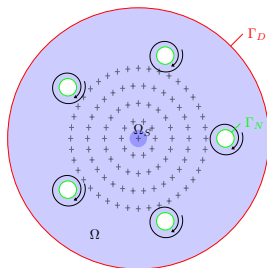


Figure: Evolution of the recovery errors in U -norm, on 500 test snapshots, with growing dictionary sizes K , for $m \in \{9, 36, 64\}$.

Numerical: Advection diffusion

$$\mathcal{N} \sim 150\,000 \quad \text{and} \quad \begin{cases} -0.01\Delta u + \mathcal{V}(\xi) \cdot \nabla u = \frac{100}{\pi} \mathbb{1}_{\Omega_S} & \text{in } \Omega, \\ u = 0 & \text{on } \Gamma_D, \\ n \cdot \nabla u = 0 & \text{on } \Gamma_N, \end{cases}$$

$$\mathcal{V}(\xi) = \sum_{i=1}^5 \frac{1}{\|x - x^{(i)}\|} \left(\xi_i e_r(x^{(i)}) + \xi_{i+5} e_\theta(x^{(i)}) \right)$$



Numerical: Advection diffusion

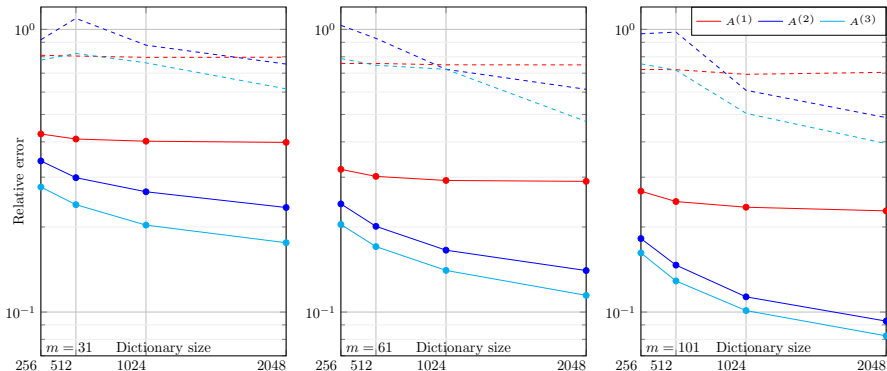


Figure: Evolution of the recovery errors in U -norm, on 500 test snapshots, with growing dictionary sizes K , for $m \in \{31, 61, 101\}$.

Conclusion

Conclusion

- Dictionary-based multi-space approach for state estimation
- Efficient offline-online decomposition using randomized linear algebra.
- More details in [Nouy and Pasco, 2023] available at <https://arxiv.org/abs/2303.10771>
- Python repo available at <https://github.com/alexandre-pasco/rla4mor/inverse-problems>

Thank you !

Appendix

Appendix: Offline stage

- Write $A_\alpha(w) = \mathbf{U}\mathbf{a}$ with $\mathbf{a} = \mathbf{a}(w) \in \mathbb{R}^{m+K}$ and $\mathbf{U} = (\mathbf{W} \mid \mathbf{V})$.
- Write the l.s. terms to compute \mathcal{S}^Θ as

$$G^\Theta(A_\alpha(w)) := \left(\Theta B^{(1)} \mathbf{U}\mathbf{a} \mid \dots \mid \Theta B^{(d)} \mathbf{U}\mathbf{a} \right)$$

$$g^\Theta(A_\alpha(w)) := \Theta f - \Theta B^{(0)} \mathbf{U}\mathbf{a}.$$

- **Offline:** compute $\underbrace{\mathbf{C}}_{\mathcal{O}(mKN)} \in \mathbb{R}^{m \times K}$ and $\underbrace{\Theta B^{(q)} \mathbf{U}}_{\mathcal{O}(KN \log(\|))} \in \mathbb{R}^{k \times (m+K)}$
- Total offline cost (without snapshot computation) is





$$\mathcal{O}\left((d \log(\mathcal{N}) + m)KN\right) \quad \text{VS} \quad \mathcal{O}\left(d^2 K^2 \mathcal{N}\right)$$

Appendix: Online stage

- **Step 0:** Observe $w = P_w u$.
- **Step 1:** Run LARS to generate $\mathcal{O}(K)$ subspaces, costing $\mathcal{O}(m^2 K)$.
- **Step 2:** Prepare the l.s. systems, each costing $\mathcal{O}(kmd)$.
- **Step 3:** Solve them, each costing $\mathcal{O}(kd^2)$.
- Total online cost with $k = \mathcal{O}(d)$

$$\mathcal{O}\left(\underbrace{m^2 K}_{\text{LARS}} + \underbrace{md^2 K}_{\text{prepare l.s.}} + \underbrace{d^3 K}_{\text{solve l.s.}}\right) = \mathcal{O}((m^2 + md^2 + d^3)K).$$

References I

-  Binev, P., Cohen, A., Dahmen, W., DeVore, R., Petrova, G., and Wojtaszczyk, P. (2017).
Data Assimilation in Reduced Modeling.
SIAM/ASA J. Uncertainty Quantification, 5(1):1–29.
-  Cohen, A., Dahmen, W., Mula, O., and Nichols, J. (2022).
Nonlinear Reduced Models for State and Parameter Estimation.
SIAM/ASA J. Uncertainty Quantification, 10(1):227–267.
-  Maday, Y., Patera, A. T., Penn, J. D., and Yano, M. (2015).
A parameterized-background data-weak approach to variational data assimilation: Formulation, analysis, and application to acoustics.
Int. J. Numer. Meth. Engng, 102(5):933–965.
-  Nouy, A. and Pasco, A. (2023).
Dictionary-based model reduction for state estimation.

References II



Temlyakov, V. N. (1998).
Nonlinear Kolmogorov widths.
Math Notes, 63(6):785–795.