

# Query enumeration on sparse graphs

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Alexandre Vigny

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# About me

- Alexandre Vigny
- I started my Ph.D in October 2015
- My advisors are :
  - ▶ Arnaud Durand (IMJ-PRG, Paris 7)
  - ▶ Luc Segoufin (LSV, ENS Cachan)

# Introduction

- Query  $q$
- Database  $D$
- Compute  $q(D)$

*small*

*huge*

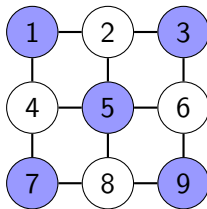
*gigantic*

Examples :

query  $q$

$$q(x, y) := \exists z(B(x) \wedge E(x, z) \wedge \neg E(y, z))$$

database  $D$



solutions  $q(D)$

$$\{(1,2) \ (1,3) \ (1,4) \\ (1,6) \ (1,7) \ \dots \\ (3,1) \ (3,2) \ (3,4) \\ (3,6) \ (3,7) \ \dots \\ \dots \}$$

# Enumeration

Input :  $\|D\| := n$  &  $\|q\| := k$  ( $k \ll n$ )

Goal : output solutions one by one

- STEP 1: Preprocessing

Prepare the enumeration : Database  $D \longrightarrow$  Index  $I$

Preprocessing time :  $f(k) \cdot n \rightsquigarrow O(n)$

- STEP 2 : Enumeration

Enumerate the solutions : Index  $I \longrightarrow \overline{x_1}, \overline{x_2}, \overline{x_3}, \overline{x_4}, \dots$

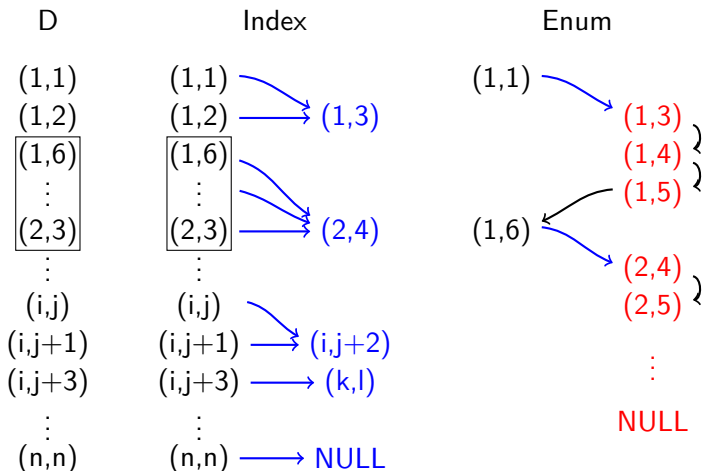
Delay :  $O(f(k)) \rightsquigarrow O(1)$

**Constant delay enumeration after linear preprocessing** ( $CD \circ Lin$ )

## Example

Database  $D := \langle \{1, \dots, n\}; E \rangle$        $\|D\| = |E|$  ( $E \subseteq D \times D$ )

Query  $q(x, y) := \neg E(x, y)$



# Enumeration VS Model-Checking

Model-Checking : Is there a solution ? Yes / No

Constant-delay Enumeration : First solution computed in time  $O(n)$

Constant-delay Enumeration  $\implies$  Linear Model-Checking

Under some complexity hypothesis, the Model-Checking is not doable in polynomial time.



Restricted databases or/and queries

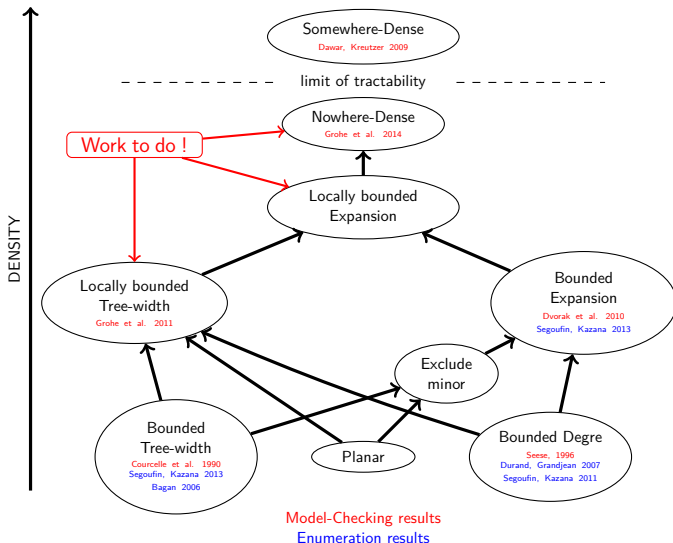
Bounded degree, planar  $\dots$

MSO, quantifier free  $\dots$

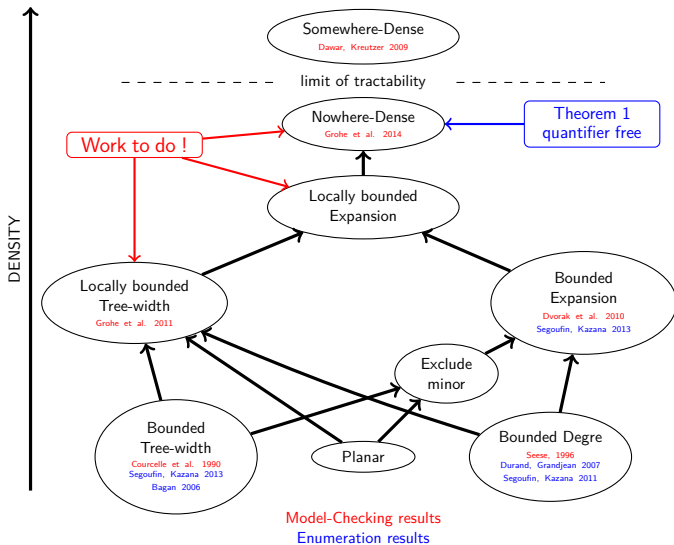
**Nowhere-dense**

**First Order**

# Classes of graphs closed under taking sub-graphs

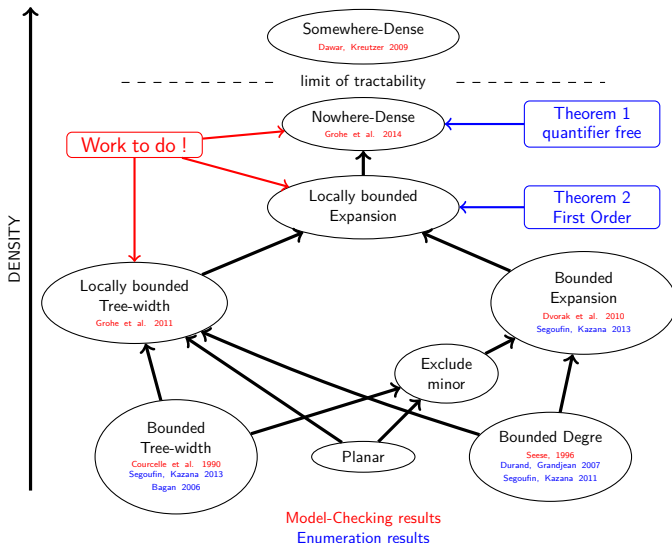


# Classes of graphs closed under taking sub-graphs





# Classes of graphs closed under taking sub-graphs



# First results

## Theorem 1

The enumeration of **quantifier free** first-order queries over **nowhere dense** class of graphs is in  $CD \circ Lin$ .

## Theorem 2

The enumeration of **first-order** queries over class of graphs with **locally bounded expansion** is in  $CD \circ Lin$ .

## Future work

- Generalize the Nowhere-Dense case.
- Enumeration with update.

Thank you !