

Audition MCF 141

Alexandre Vigny

<https://alexandre-vigny.github.io/mcf141.pdf>

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Alexandre Vigny



- 31 years old, born in 1992.
- PhD defense September 2018 (*5 years ago*).
- Living in Bremen, Germany.

- Master in University Paris Diderot, 2015.
Logique Mathématique et Fondement de l'Informatique (LMFI).
- PhD in University Paris Diderot. (3 years)
With Arnaud Durand & Luc Segoufin
- Post-doc in Warsaw. (1 year)
With Szymon Toruńczyk & Mikołaj Bojańczyk
- Post-doc in Bremen. (3 years)
With Sebastian Siebertz

Summary

Area of Research

- Graph theory
- Logic
- Distributed computing

Highlights

- Publication in J.ACM
- 2 Upcoming journal papers
- 9 Conference papers
- Co-organizer of a workshop

PODC-DARe: Distributed Algorithms on

REalistic network models

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PODC-DARe: Distributed Algorithms on

REalistic network models

Teaching

- In Paris (3 years)
~ 180h
- In Bremen (3 years)
~ 270h

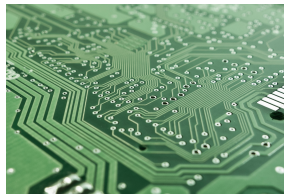
Highlights

- Creation of syllabuses
- Responsible for two courses
- All level bachelor to master
- Both in French and English

Algorithmic graph theory

Given a graph G and a property P : “Does G satisfy P ?”

→ Is G planar?

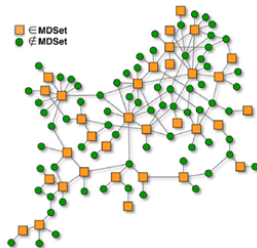


Algorithmic graph theory

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→ Does G have a k -dominating set?



Algorithmic graph theory

Given a graph G and a property P : “Does G satisfy P ?”

→ Is G planar?

→ Does G have a k -dominating set?

→ Is G connected?



Goal: Efficient algorithms ...

... at least for restricted graph classes and/or simple properties.

Logic

First-order (FO) logic

- Can express k -independent set: *There are k vertices, that are not adjacent*
$$\exists x_1 \dots \exists x_k \bigwedge_{i < j} (\neg E(x_i, x_j) \wedge x_i \neq x_j)$$
- Cannot express : connectivity, planarity, 2-colorability, ...

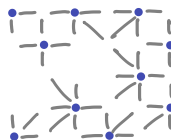
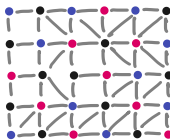
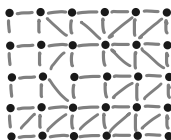
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Monadic second-order (MSO) logic

- More general than FO
- Can express : 3-colorability:
$$\exists X_1 \exists X_2 \exists X_3 (\forall x \bigvee_{i < 3} x \in X_i) \wedge (\forall x \forall y E(x, y) \rightarrow \bigwedge_{i < 3} (x \notin X_i \vee y \notin X_i))$$



Distributed computing

Distributed Computing : Local model

- Different notion of **efficient**
- Time needed VS Information needed

Distributed computing

Distributed Computing : Local model

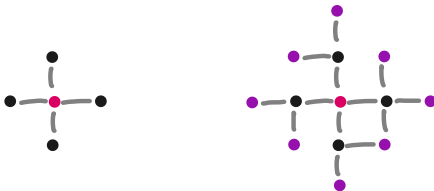
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Distributed computing

Distributed Computing : Local model

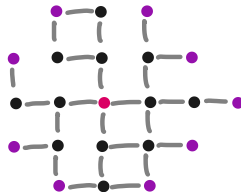
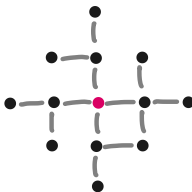
- Different notion of **efficient**
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Distributed computing

Distributed Computing : Local model

- Different notion of **efficient**
- Time needed VS Information needed



- Can you decide locally?

Logic & Meta theorems

Problems can be expressed in **logic**. (FO, MSO,...)

The \mathcal{L} , \mathcal{C} model-checking problem:

Given $\varphi \in \mathcal{L}$ and $G \in \mathcal{C}$, does $G \models \varphi$?

Goal: **fixed parameter tractable** algorithms $O(f(\varphi) \cdot |G|^c)$

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Courcelle's Theorem (1990):

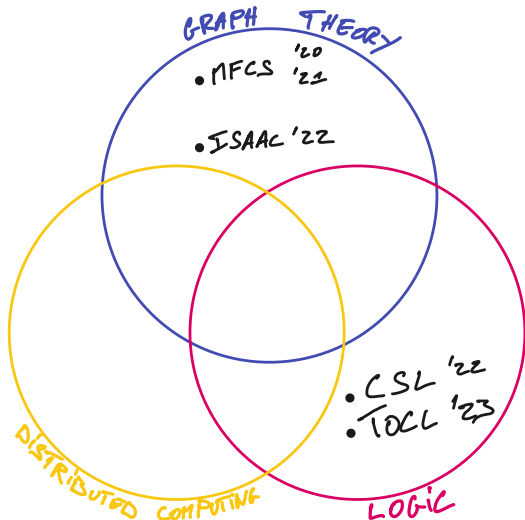
for $\varphi \in \text{MSO}$ and $\text{treewidth}(G) \leq k$, in time $O(f(\varphi, k) \cdot |G|)$

→ Generalize many known results, ex:

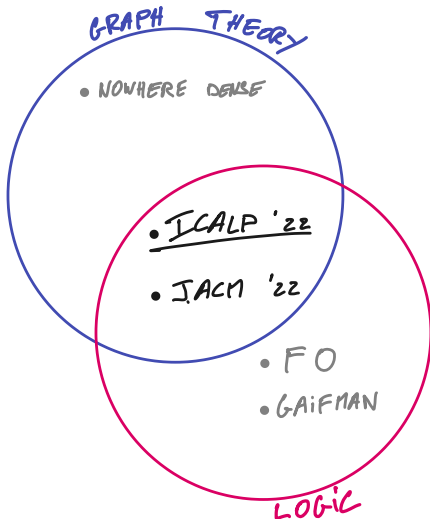
Arnborg, Proskurowski 1989:

independent sets, dominating sets, graph coloring, Hamiltonian, ...
are **linear** on **partial k -tree**.

Result overview



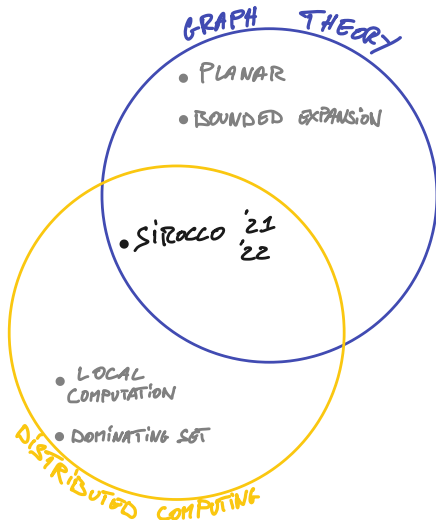
Result overview



• ENUMERATION FO QUERIES

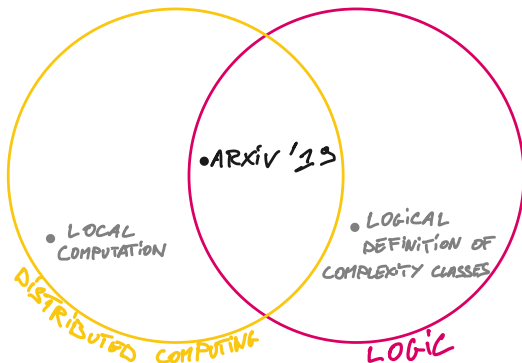
→ PODS 2018
JACM 2022

Result overview



- ENUMERATION FO QUERIES
→ PODS 2018
→ SACH 2022
- DISTRIBUTED DOMINATION ON SPARSE GRAPH CLASSES
→ SIROCCO 2021/2022
→ EUR. J. COMB.
(TO APPEAR)

Result overview



- ENUMERATION FO QUERIES
→ PODS 2018
SACM 2022
- DISTRIBUTED DOMINATION
ON SPARSE GRAPH CLASSES
→ SIROCCO 2021/2022
→ EUR. J. COMB.
(TO APPEAR)
- PARAMETERIZED DISTRIBUTED
COMPLEXITY THEORY :
A LOGICAL APPROACH
→ ARXIV

Presentation
○○

Past Research
○○○○
○○○○

Future and integration
●○○○○

Teaching
○○○
○○

Proposition pédagogique
○○○○○

Research project

Research project

First (*short term*) goal: new logics

Recently definition of a new logic: $\text{FO} + \text{conn}$

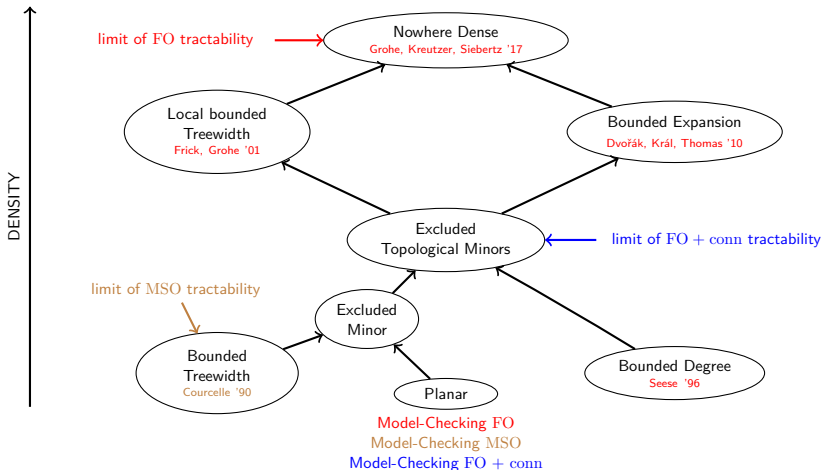
- Pilipczuk, Schirrmacher, Siebertz, Torunczyk, Vigny. ICALP'22
- more expressive than FO ,
- less than MSO .

Beyond $\text{FO} + \text{conn}$

- What can be added?
- What do we want to express?
- Example: a path of even length, using only blue nodes, ...

Keeping in mind algorithmic applications

Characterization of graph classes



Directed graphs

Direction on edges

- More general
- Problems are harder
E.g. Directed Dominating Set is NP-complet on DAGs
- Some problems do not care about orientation

Reconfiguration problems

- No need to find a set
- Here, the orientation matters!

Distributed computing & Certification

Local computing

- Other notion of **efficient**
- Still looking for meta theorems

Complexity classes

- Hard problems may not be equally hard
- Define complexity classes through logic

Compact certification

- Feuilloley, Bousquet, Pierron
*What Can Be Certified Compactly? Compact local certification of **MSO** properties in **tree-like** graphs.* In PODC'22
- Fraigniaud, Montealegre, Rapaport, Todinca
A Meta-Theorem for Distributed Certification. In SIROCCO'22

Integration

LINKS (Linking Dynamic Data)

- S. Salvati, C. Paperman, I. Boneva, F. Capelli, M. Monet
- Tree-like structures & MSO (S. Salvati)
- Connections computation-logic (S. Salvati and C. Paperman)
- Graph databases (I. Boneva)
- Knowledge compilation & tree-like structures (F. Capelli, M. Monet)
- Distributed computing is new
- Adding a solid background on graph theory

Past teaching in Paris

Around 180h

- Mainly Bachelor level
- Similar to topics in Faculté des Sciences et Technologies

Subject	Years	Level	Activity
Initiation à la Programmation	2015-2016	L1	TP
Langages et Automates	2015-2016	L2	TD
Éléments d'Algorithmique	2016-2017	L2	TD
Base de données	2016-2017 2017-2018	L3 M1	TP & TD
Concepts Informatique	2017-2018	L1	TD

Past teaching in Bremen

Around 250h

- Master level
- Creation of syllabuses
- Fully in charge of a lecture

Subject	Years	Level	Activity
Finite Model Theory	2019-2020	Master	TD
Sparsity	2019-2020	Master	TD
Parametrized Complexity	2019-2020	Master	TD
Set and Model Theory	2020-2021	Master	TD
Databases, Graphs, Algorithms	2020-2021 2021-2022	Master	Cours & TD
Set and Model Theory	2021-2022	Master	Cours & TD

Highlight

Various settings

- In French, in English
- All levels
- Physical and remote

Various responsibilities

- Creation of the content
- Fully in charge of a lecture

Takeaway

- Polls
- Online white boards

Future 1

L1:

- Outils pour l'informaticien·ne
- Initiation à la programmation
- Algorithmes et programmation

L2:

- Algorithmique et les structures de données
- Programmation orientée objet
- Programmation en C

L3:

- Programmation objet
- Bases de données
- Algorithmique et structures de données

Future 2

Master

- Parcours Machine Learning
 - Algorithmique avancée
 - Bases de données avancées

Responsibilities

- Responsable d'année
- Stages
- ...

Algorithmes & Graphes I, Introduction

→ Notions Élémentaires

- Modélisation par graphes
- Définitions
- Dirigés / non-dirigés
- Connexité / composante connexes
- Arbres et acyclicité

→ Représentation et implémentation

- listes d'adjacence
- matrice d'adjacence

Algorithmes & Graphes II, Dans le vif du sujet

→ Parcours de graphes

- Profondeur (Depth-First Search)
- Largeur (Breadth-First Search)
- Plus court chemin

→ Arbres couvrants

- Algorithme de Jarník-Prim-Dijkstra
- Algorithme de Kruskal

→ Flow

- Directed graphs
- Flow and Cut
- Ford-Fulkerson (max-flow min-cut))

Algorithmes & Graphes III, Ouverture

- Graphes particuliers : Graphes planaires
 - Notion de mineur
 - Formule d'Euler

- Problèmes difficiles
 - 3-coloriage
 - Ensemble dominant
 - Notion de reduction
 - NP-completude

Algorithmes & Graphes IV, Conclusion

- Niveau : L2 / L3
- Prérequis : Algorithmes et programmation
 - Boucle
 - Recursion
 - Liste
 - Tableau
- Conflit avec Algorithmique et structure de données ?
 - Champ d'application de ces structures
 - Introduction de notions de complexité
 - $O(\log(n))$, $O(n)$, $O(n^2)$
 - Polynomial vs NP-complet

Thank you!

- 2019-2023: Postdoc, University of Bremen. With Sebastian Siebertz
- 2018-2019: Postdoc, University of Warsaw. With Szymon Toruńczyk
& Mikołaj Bojańczyk
- 2015-2018: Thesis, University Paris Diderot. With Arnaud Durand
& Luc Segoufin

Info:

- 1 Journal: J.ACM (TOCL & Eur. J. Comb. to appear)
- 9 Conferences: ICDT, PODS, MFCSx2, SIROCCOx2, ISSAC, CSL, ICALP.
- 1 Workshop (co-organizer): <https://podc-dare.github.io/>.
- 1 Popularization: La gazette du GDR-IM.

tinyurl.com/short-polls

◀ back

