

# *Practical Quantum Computing*

Lecture 12  
Surface Code, Lattice Surgery,  
Decoding, Compiling Circuits

| Week                                 | Tuesday (3h)  |                               |                               | Wednesday (3h)  |                         |     | Deadlines   |  |
|--------------------------------------|---|-------------------------------|-------------------------------|---|-------------------------|-----|-------------|--|
| 1. The Basics                        | <u>Introduction</u>   | Gates                         | Circuit Identities            | Qiskit  | Cirq/Qualtran           | Q&A |             |  |
|                                      | <b>Programming Assignment 1:</b> <u>The basics of a quantum circuit simulator</u> |                               |                               | <b>Programming Assignment 1:</b> The building blocks of a quantum circuit simulator |                         |     |             |  |
| 2. Entanglement and its Applications | Teleportation   | Superdense Coding             | Quantum Key Distribution      | Qualtran/Assignment2  | Terminology of Projects | Q&A |             |  |
|                                      | <b>Programming Assignment 2:</b> The basics of a quantum circuit optimizer        |                               |                               | <b>Programming Assignment 2:</b> The building blocks of a quantum circuit optimizer |                         |     |             |  |
| 3. Computing                         | Phase Kickback and Toffoli  | Distinguishing quantum states | The First Algorithms          | Invited TBA   |                         | Q&A | 11 May 2024 |  |
| 4. Advanced Topics*                  | Arithmetic Circuits*  | Fault-Tolerance*              | Surface QEC*<br>Grover's Alg* | Invited TBA   | Invited TBA             | Q&A | 18 May 2024 |  |

\* not evaluated

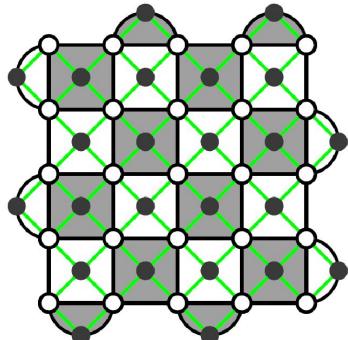
# Learning goals - 12 Surface QECC and Grover's (Advanced)

1. What you have learned by now
  - a. Quantum circuits: mathematics, diagrams and circuit identities
  - b. Entanglement: teleportation, quantum games, QKD
  - c. Superpositions, Phase Kickback and finding hidden strings
  - d. Quantum algorithms and quantum arithmetic
  - e. Fault-Tolerance and Quantum Error Correction
2. The surface QECC
  - a. What it is
  - b. How it is built in hardware
  - c. High-level description of its functionality
3. Lattice Surgery
  - a. Implementing error corrected gates
  - b. Instruction set for universal gate set
4. Decoding
  - a. Finding the most probable error
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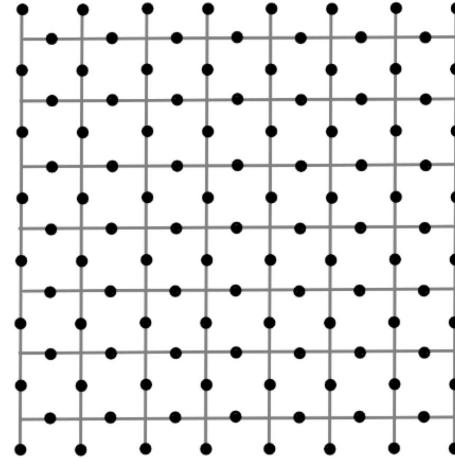
- Deadline for programming Assignment 2
- 18 May 2024

# The Surface Code

- Quantum error correcting codes are defined by the measurements we make
- Let's move beyond the simple  $Z_j Z_{(j+1)}$  of the repetition code
- In the surface code we use a 2D lattice of code qubits, and define observables for plaquettes and vertices



- data qubit
- measure qubit
- coupler
- detect X
- detect Z



$$A_v = \sigma_x^i \sigma_x^j \sigma_x^k$$

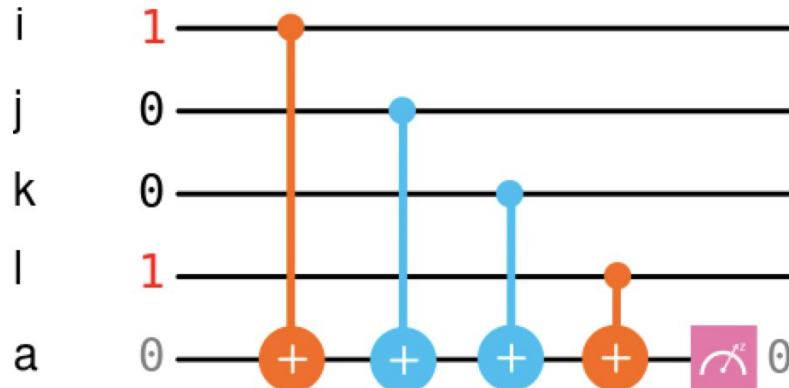
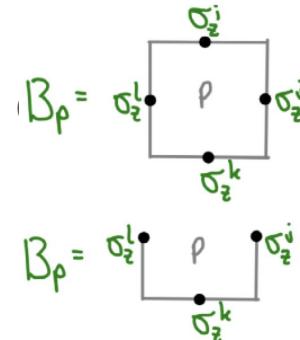
$$A_v = \sigma_x^i \sigma_x^j \sigma_x^k$$

$$\beta_p = \sigma_z^l \sigma_z^m \sigma_z^n$$

$$\beta_p = \sigma_z^l \sigma_z^m \sigma_z^n$$

# Plaquette Syndrome

- First let's focus on the plaquette syndrome
- These are similar to the two qubit measurements in the repetition
- Instead we measure the parity around plaquettes in the lattice
- Can again be done with CX gates and an extra qubit

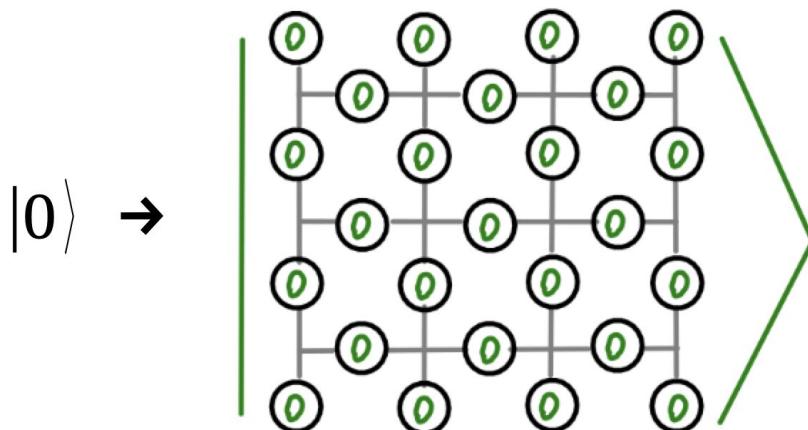


# Plaquette Syndrome

We can define a classical code (storing only a bit) based on the plaquette syndrome alone

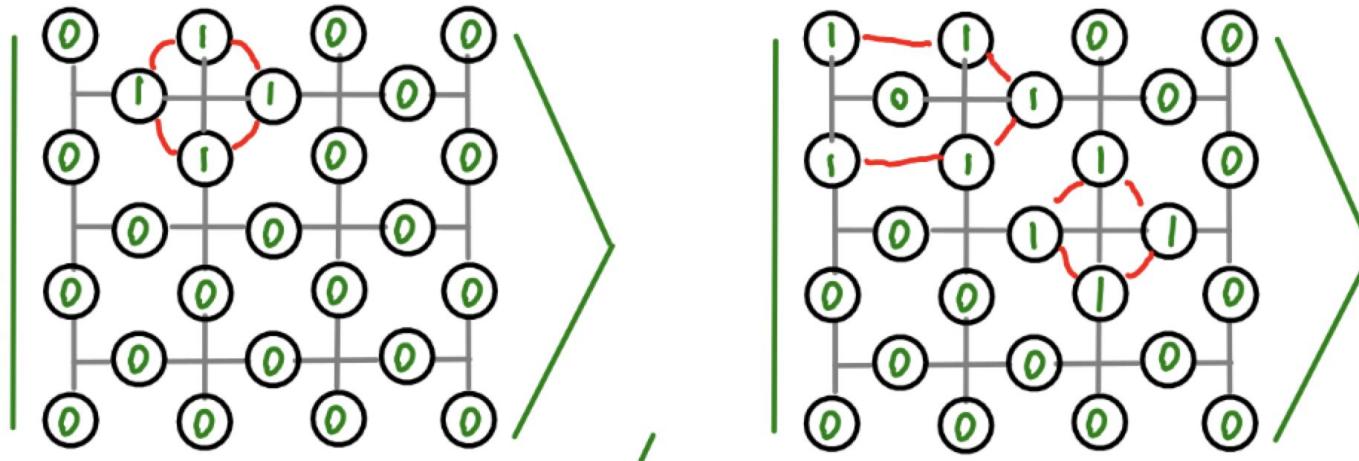
Valid states are those with trivial outcome for all plaquette syndrome measurements:

- Even parity on all plaquettes
- How to store a 0 in this?
- How about the state where every code qubit is  $|0\rangle$ ?



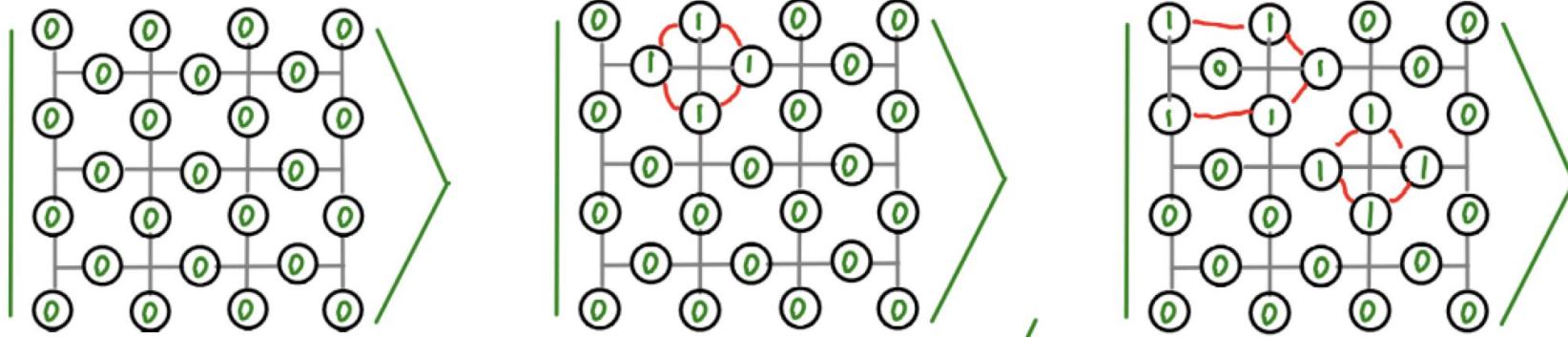
# Plaquette Syndrome

- There are ‘nearby’ states that also have even parity on all plaquettes
- These can’t be a different encoded state: they are only a few bit flips away from our encoded 0 state
- We’ll treat them as alternative ways to store a 0



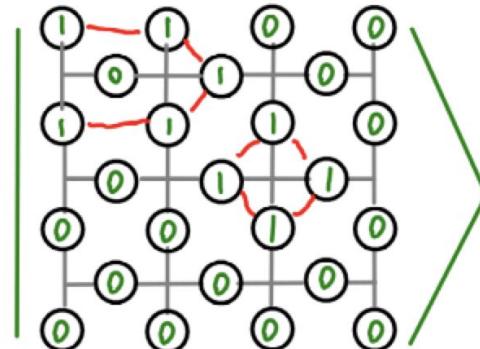
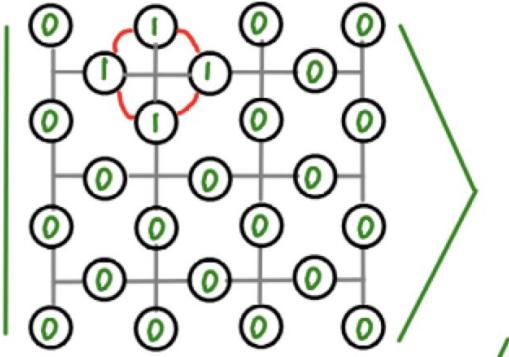
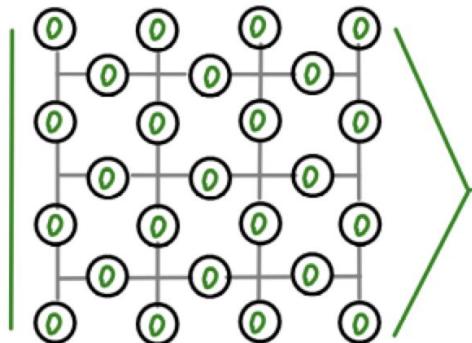
# Plaquette Syndrome

- Given any state for an encoded 0
  - Pick a vertex
  - Apply bit flips around that vertex
- Now you have another valid state for 0
- This generates an exponentially large family



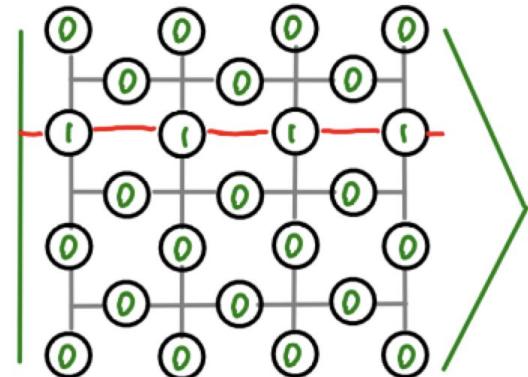
# Plaquette Syndrome

- The states in this family can be very different
- But they all share a common feature
  - Any line from top to bottom (passing along edges) has even parity
- This is how we can identify an encoded 0
- And it gives us a clue about how to encode a 1



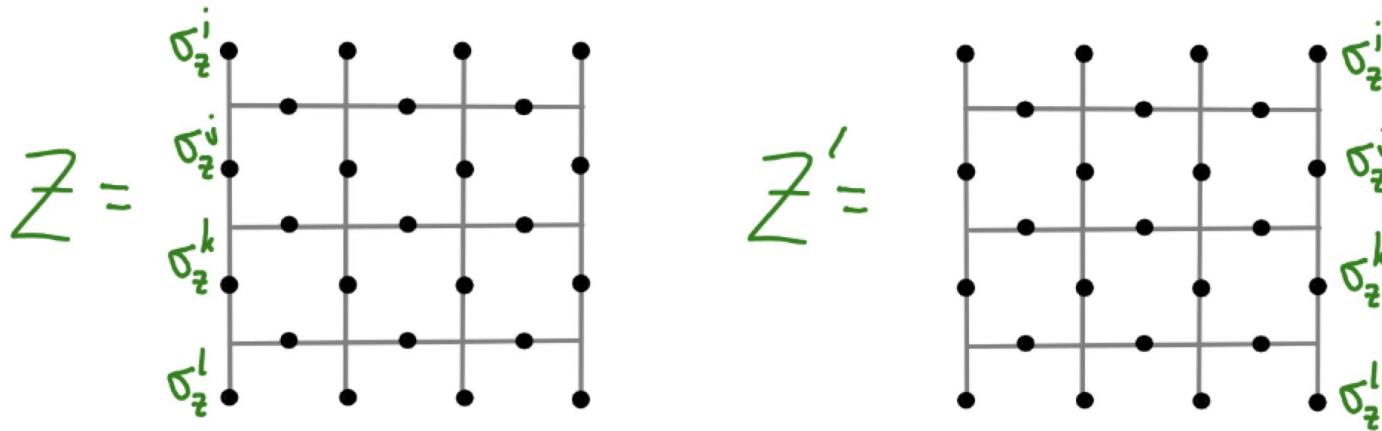
# Plaquette Syndrome

- For our basic encoded 1, we use a bunch of 0s with a line from left to right (passing through plaquettes)
- This also spawns an exponentially large family
- All have *odd* parity for a line from top to bottom
- Unlike the repetition code, distinguishing encoded 0 and 1 requires some effort (which is good!)

 $|1\rangle \rightarrow$ 

# Logical X and Z

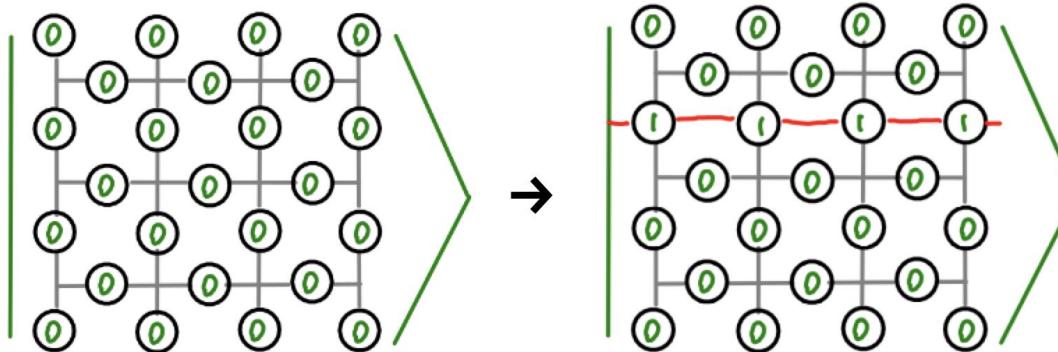
- Distinguishing 0 and 1 corresponds to measuring Z on the physical qubit
- The following observables detect what we need



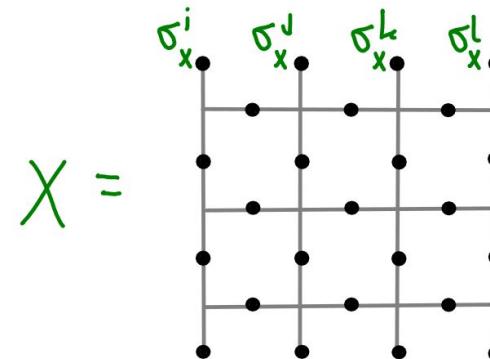
- Or the same on any line from top to bottom
- Uses the edges has a nice advantage: we can think of them as large (unenforced) plaquettes

# Logical X and Z

- To flip between 0 and 1, we can flip a line of qubits

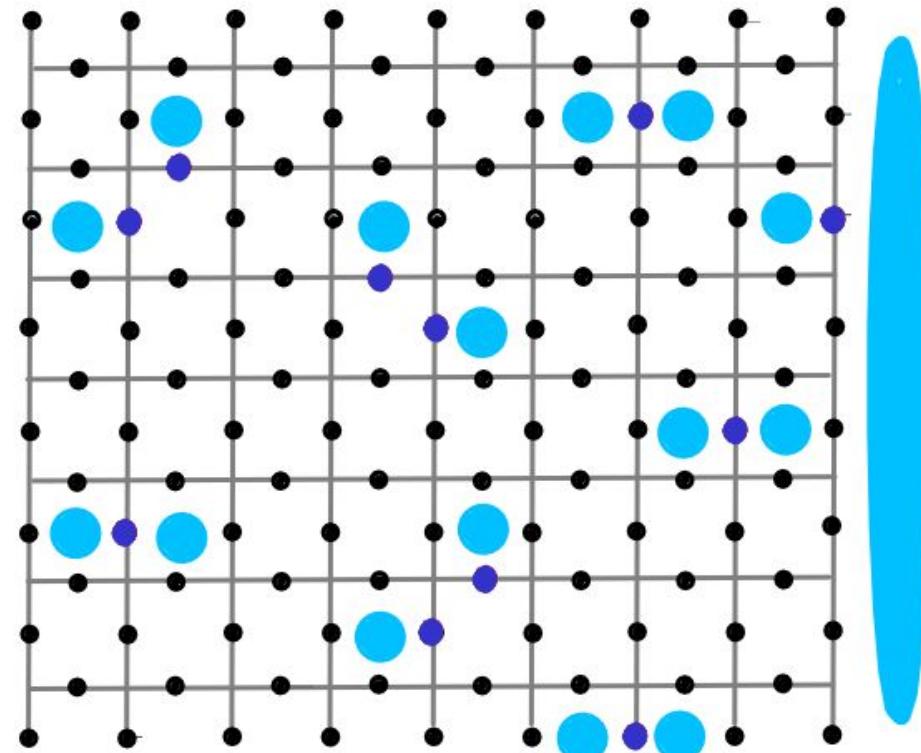


- Such lines of flips act as an X on the logical qubit



# Effects of Errors

- Applying an X to any code qubit changes the parity of its two plaquettes
  - An isolated X creates a pair of defects
  - Further Xs can move a defect, or annihilate pairs of them
  - A logical X requires many errors to stretch across the lattice
- 
- With the plaquette operators, we can encode and protect a *bit*

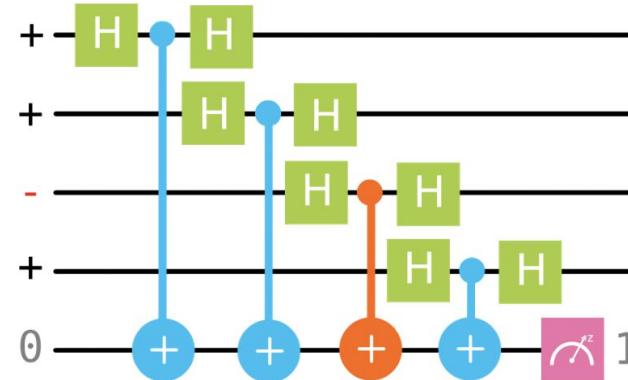
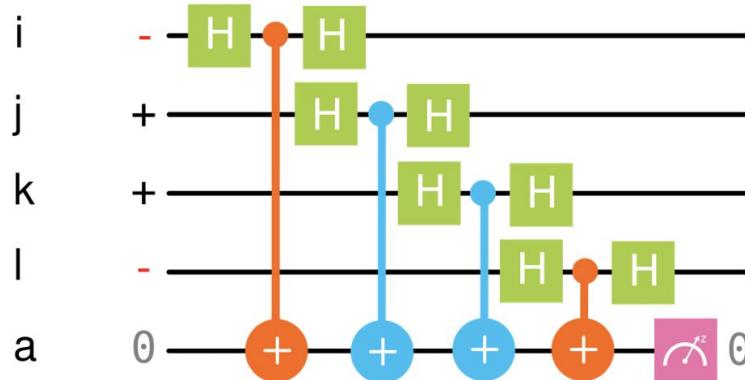


# Vertex Syndrome

- Now forget the plaquettes and focus on vertices
- These observables can also be measured using CX gates and an ancilla
- In this case they look at the  $|+\rangle$  and  $|-\rangle$  states, and count the parity of the number of  $|-\rangle$ s

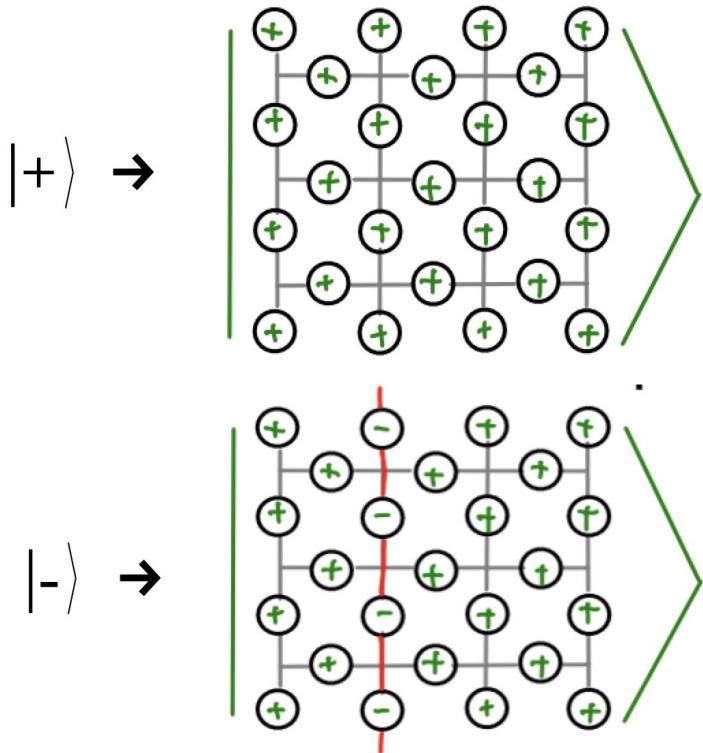
$$A_v = \begin{array}{c} \sigma_x^i \\ \hline \sigma_x^j & v & \sigma_z^k \\ \hline \sigma_x^l \end{array}$$

$$A_v = \begin{array}{c} \sigma_i \\ \hline \sigma_j & v & \sigma_k \\ \hline \sigma_l \end{array}$$



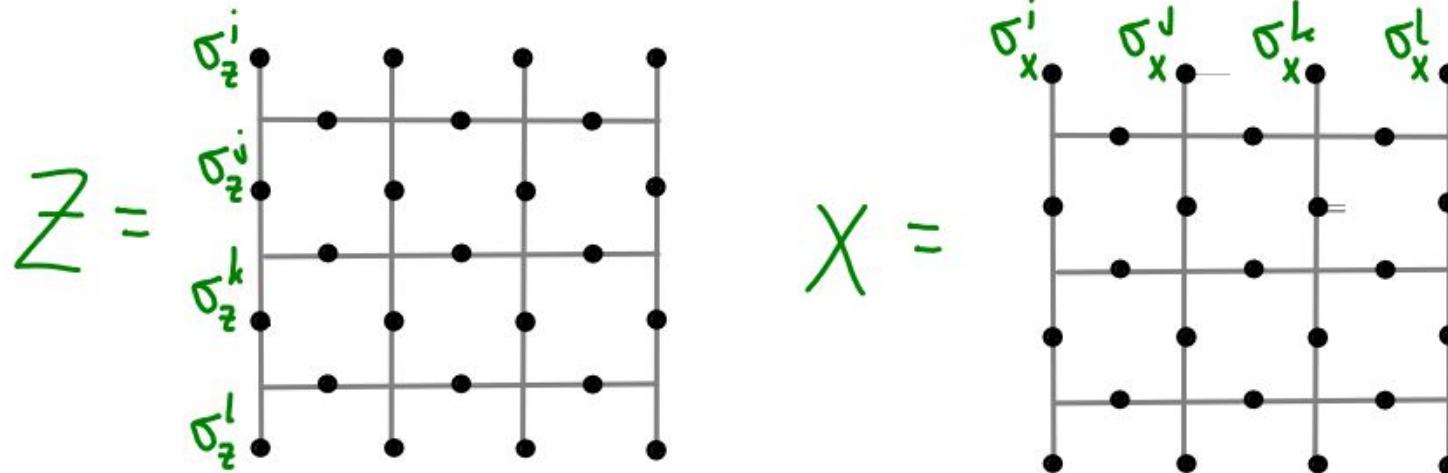
# Vertex Syndrome

- These operators also allow us to encode and protect a bit value
- In this case, let's use + and - to label the two states
- They are encoded using suitable patterns of  $|+\rangle$  and  $|-\rangle$  states for the code qubits
- As with the plaquettes, these also correspond to exponentially large families of states



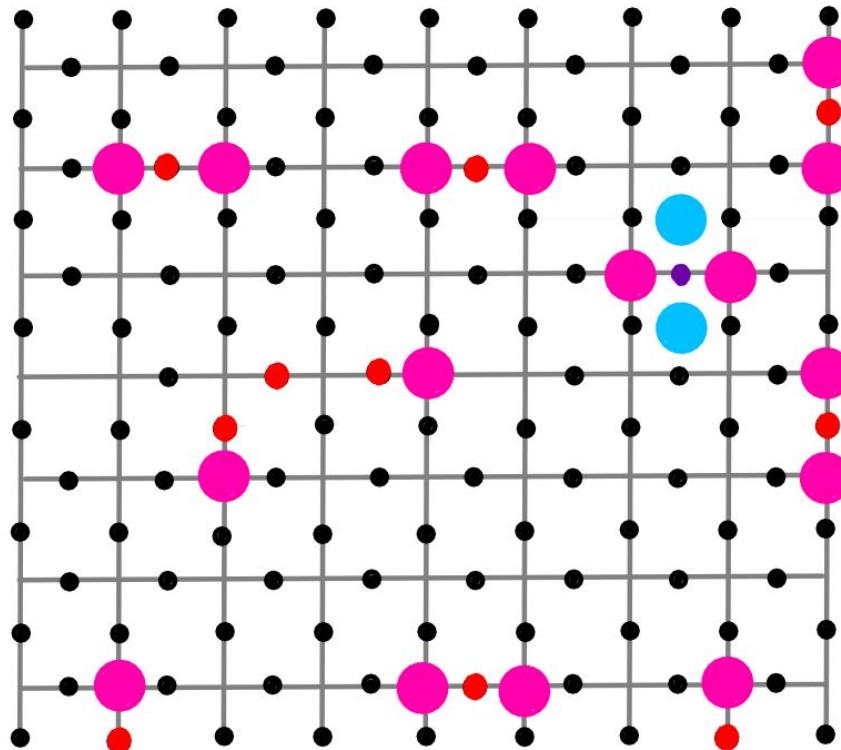
# Logical X and Z

- What is the X operator (distinguish between  $|+\rangle$  and  $|-\rangle$  )?
- What is the Z operator (flip between  $|+\rangle$  and  $|-\rangle$  )?
- Turns out they are exactly the same as before!



# Effects of Errors

- Applying a Z to any code qubit changes the X parity of its two vertices
- An isolated Z creates a pair of defects
- Further Zs can be move a defect, or annihilate pairs of them
- A logical Z requires many errors to stretch across the lattice
- With the vertex operators, we can encode and protect a *bit*

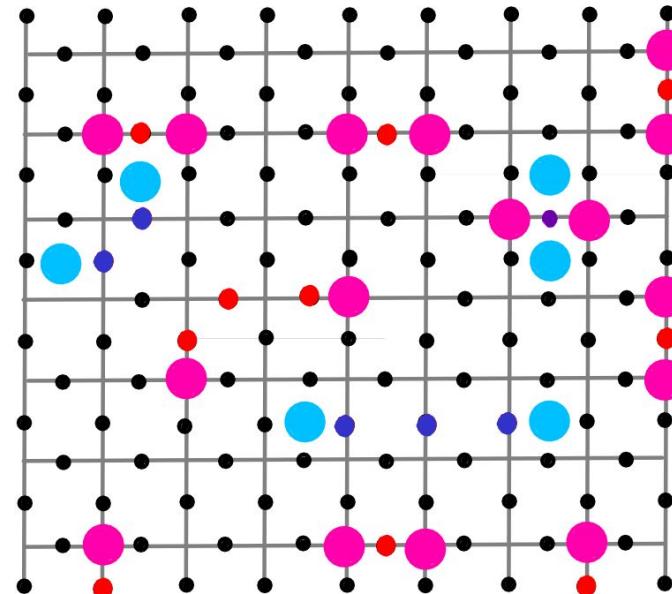


# Putting it all Together

- The plaquette and vertex operators commute
- This allows us to detect both X and Z errors
- Since  $Y \sim XZ$ , we can detect Y errors too

$$A_v = \sigma_x^i \sigma_x^j \sigma_x^k$$

$$A_v = \sigma_x^i \sigma_x^j \sigma_x^k$$

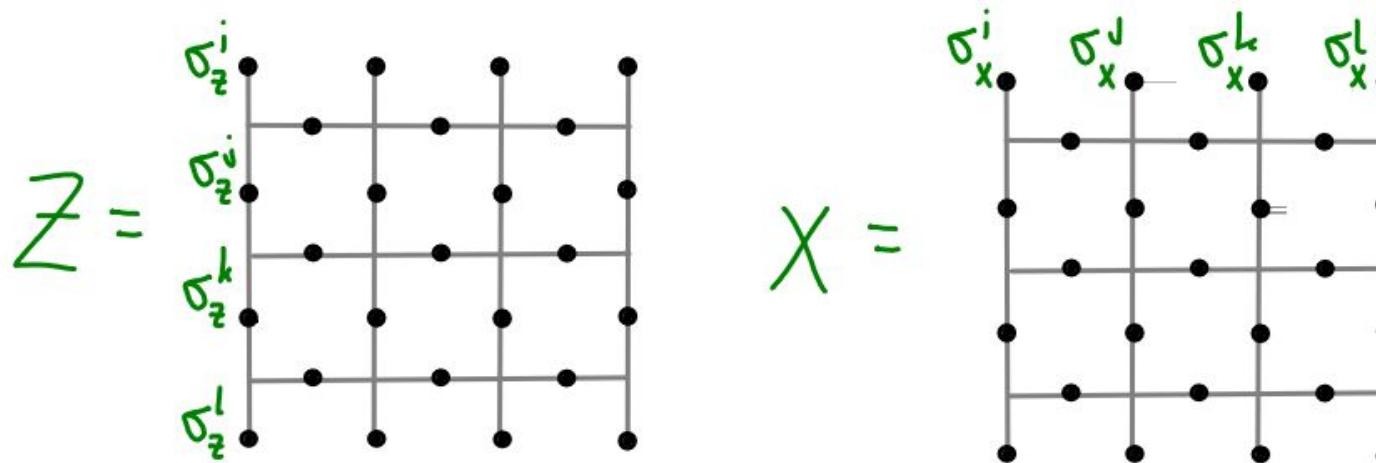


$$B_p = \sigma_z^l \sigma_z^m \sigma_z^n$$

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# Putting it all Together

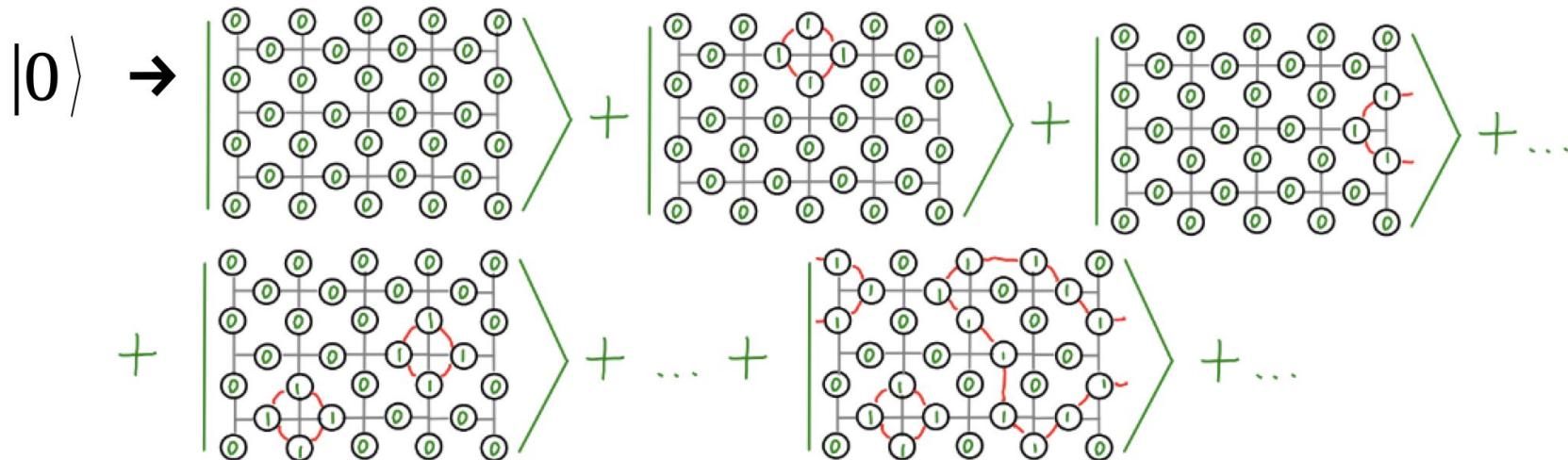
- The Z and X operators on the encoded qubit are exactly the same as before



- These, and the Hadamard, can be performed fault-tolerantly

# Putting it all Together

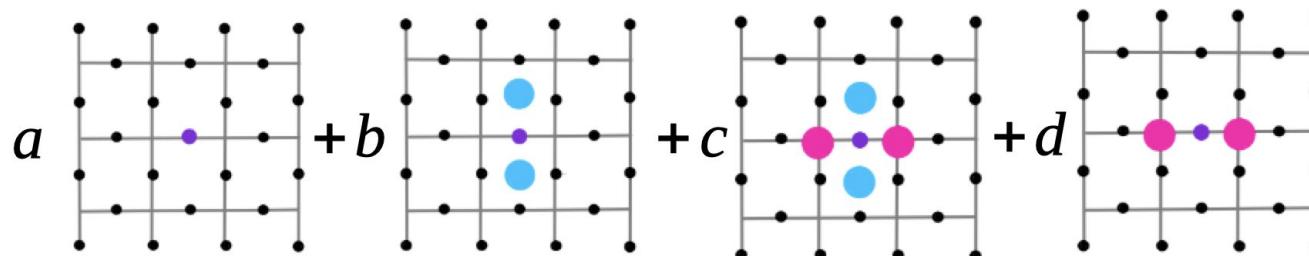
- The states we need are highly entangled quantum states
- They are examples of topologically ordered states



- Though such things can be hard to make, we create and protect them with the syndrome measurements

# Putting it all Together

- We are not just protected against X and Z, but all local errors
- As mentioned earlier,  $Y \sim XZ$
- Everything else can be expressed  
 $E = aI + bX + cY + dZ$
- This creates a superposition of different types of error on the surface code
- Measuring the stabilizers collapses this to a simple X, Y or Z
- Though such things can be hard to make, we create and protect them simply by making the stabilizer measurements

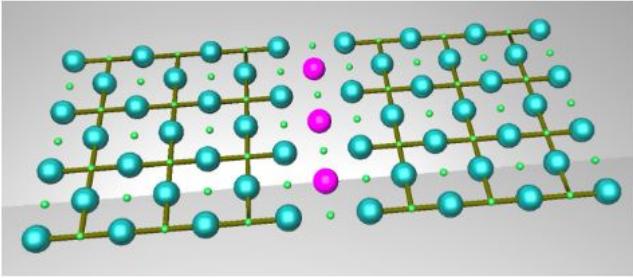


# Learning goals - 12 Surface QECC and Grover's (Advanced)

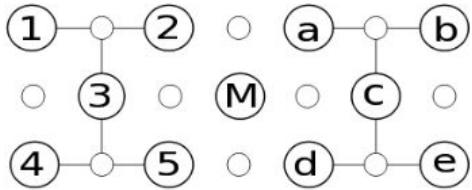
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# Lattice Surgery



**Figure 5.** Arrangements of physical qubits for rough lattice merging. Left and right continuous surfaces encode separate logical qubits. The pink qubits form the intermediate qubit line for the merging operation.



**Figure A1.** Lattice qubits for merging two rough surfaces of distance 2 into a single surface.

|                 | 1   | 2   | 3   | 4   | 5   | $M$ | $a$ | $b$ | $c$ | $d$ | $e$ |
|-----------------|-----|-----|-----|-----|-----|-----|-----|-----|-----|-----|-----|
| $\alpha\alpha'$ | $X$ | $X$ | $X$ |     |     |     |     |     |     |     |     |
|                 |     |     | $X$ | $X$ | $X$ |     |     |     |     |     |     |
|                 | $Z$ |     | $Z$ | $Z$ |     |     |     |     |     |     |     |
|                 |     | $Z$ | $Z$ |     |     | $Z$ |     |     |     |     |     |
|                 | $Z$ | $Z$ |     |     |     |     |     |     |     |     |     |
|                 |     |     |     |     |     |     | $Z$ |     |     |     |     |
|                 |     |     |     |     |     |     |     | $Z$ | $Z$ |     |     |
|                 |     |     |     |     |     |     |     | $Z$ | $Z$ | $Z$ |     |
|                 |     |     |     |     |     |     |     | $Z$ | $Z$ | $Z$ | $Z$ |
|                 |     |     |     |     |     |     | $X$ | $X$ | $X$ |     |     |

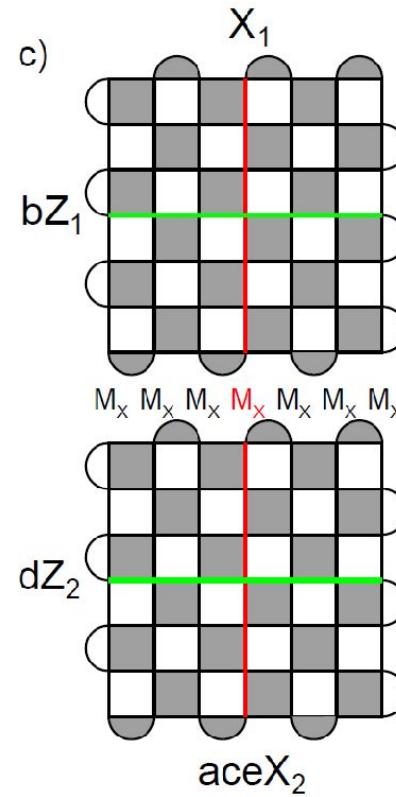
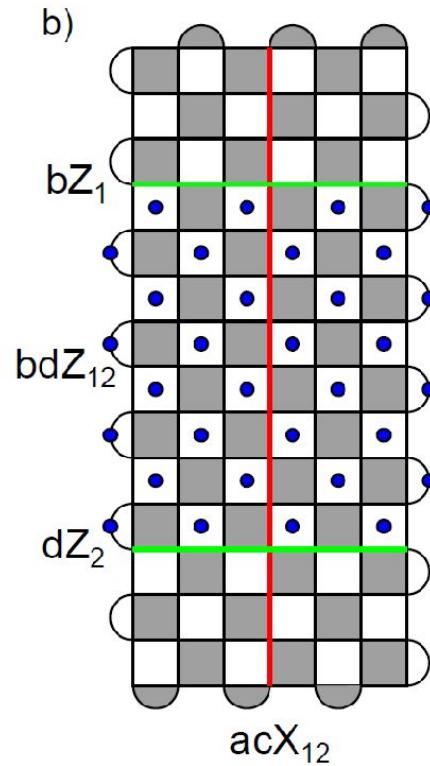
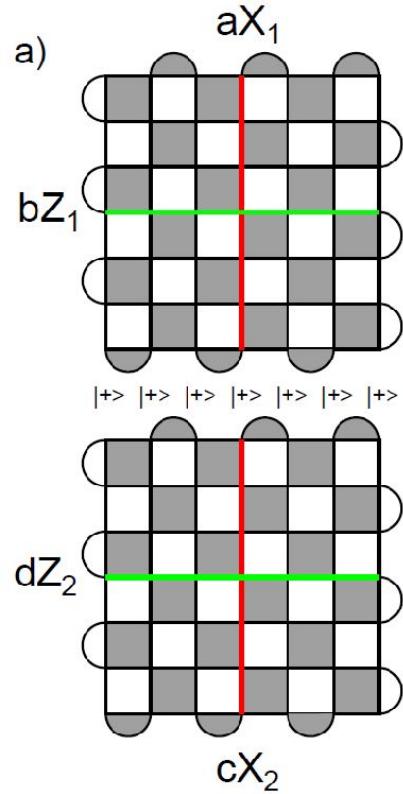
The stabilizers  $X_2X_MX_a$  are measured across the join to merge the surfaces, with

|                 | 1        | 2   | 3   | 4   | 5   | $M$ | $a$ | $b$ | $c$ | $d$ | $e$ |
|-----------------|----------|-----|-----|-----|-----|-----|-----|-----|-----|-----|-----|
| $\alpha\alpha'$ | $(-1)^m$ | $X$ |     |     |     |     | $X$ | $X$ |     |     |     |
|                 |          | $X$ | $X$ | $X$ |     |     |     |     |     |     |     |
|                 |          |     | $X$ | $X$ | $X$ |     |     |     |     |     |     |
|                 |          |     |     |     |     |     | $X$ | $X$ | $X$ |     |     |
|                 |          |     |     |     |     |     |     | $X$ | $X$ | $X$ |     |
|                 |          |     | $Z$ | $Z$ | $Z$ |     | $Z$ | $Z$ |     |     |     |
|                 |          |     | $Z$ | $Z$ | $Z$ |     | $Z$ | $Z$ |     |     |     |
|                 |          |     |     |     |     |     | $Z$ | $Z$ | $Z$ |     |     |
|                 |          |     |     |     |     |     | $Z$ | $Z$ | $Z$ |     |     |
|                 |          |     | $Z$ | $Z$ |     |     | $Z$ | $Z$ | $Z$ |     |     |

The stabilizer  $X_5X_MX_d$  across the join is now measured, with outcome  $m'$ , leaving the state as

|                 | 1        | 2           | 3   | 4   | 5   | $M$ | $a$ | $b$ | $c$ | $d$ | $e$ |
|-----------------|----------|-------------|-----|-----|-----|-----|-----|-----|-----|-----|-----|
| $\alpha\alpha'$ | $(-1)^m$ | $X$         |     |     |     |     | $X$ | $X$ |     |     |     |
|                 |          | $(-1)^{m'}$ |     |     |     |     | $X$ | $X$ |     |     |     |
|                 |          |             | $X$ | $X$ | $X$ |     |     |     |     |     |     |
|                 |          |             |     | $X$ | $X$ | $X$ |     |     |     |     |     |
|                 |          |             |     |     |     |     | $X$ | $X$ | $X$ |     |     |
|                 |          |             |     |     |     |     |     | $X$ | $X$ | $X$ |     |
|                 |          |             | $Z$ | $Z$ | $Z$ |     | $Z$ | $Z$ |     |     |     |
|                 |          |             | $Z$ | $Z$ | $Z$ |     | $Z$ | $Z$ |     |     |     |
|                 |          |             |     |     |     |     | $Z$ | $Z$ | $Z$ |     |     |
|                 |          |             | $Z$ | $Z$ |     |     | $Z$ | $Z$ | $Z$ |     |     |

arXiv:1808.06709, arXiv:1808.02892



# Logical XX measurement

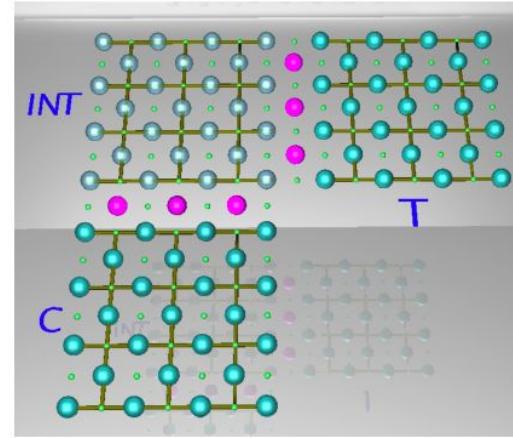
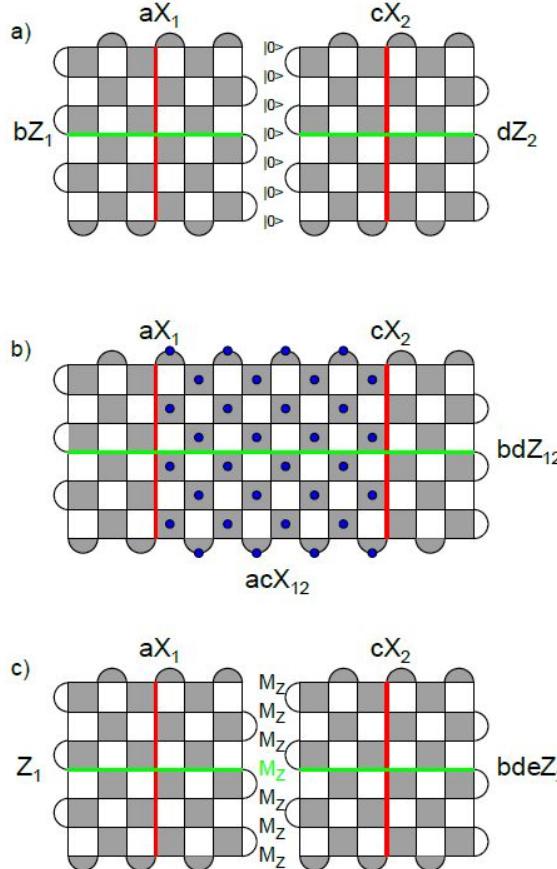


Figure 7. Layout of qubits for a CNOT operation with lattice surgery. Control (C) and target (T) surfaces interact by merging and splitting with the intermediate surface (INT).

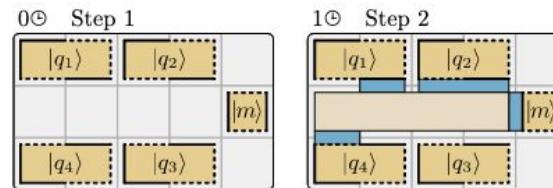


Figure 8: Example of a  $Z_{|q_1\rangle} \otimes Y_{|q_2\rangle} \otimes X_{|q_4\rangle} \otimes Z_{|m\rangle}$  measurement to implement a  $(Z \otimes Y \otimes \mathbb{1} \otimes X)_{\pi/8}$  gate.

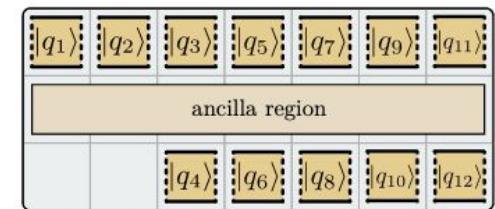
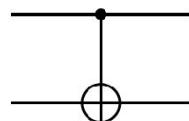


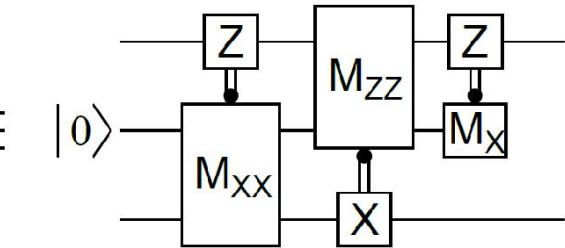
Figure 9: A compact block stores  $n$  data qubits in  $1.5n + 3$  tiles. The consumption of a magic state can take up to  $9\odot$ .

### Logical CNOT

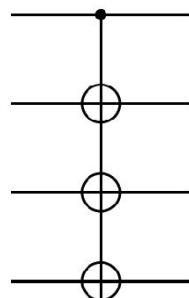
a)



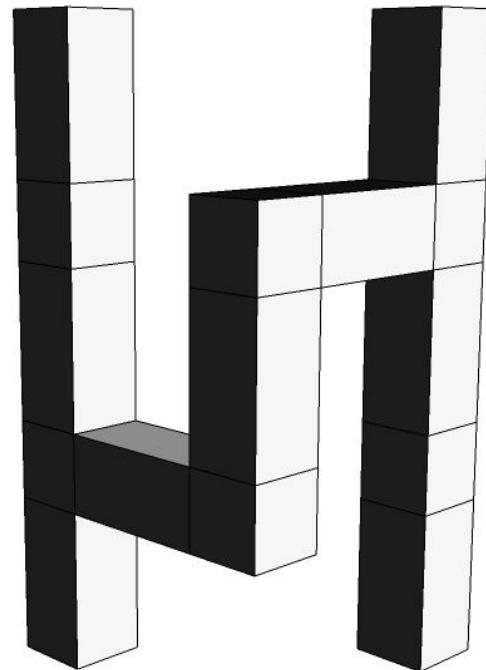
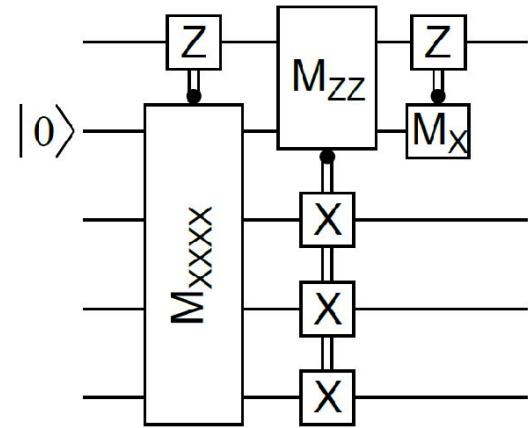
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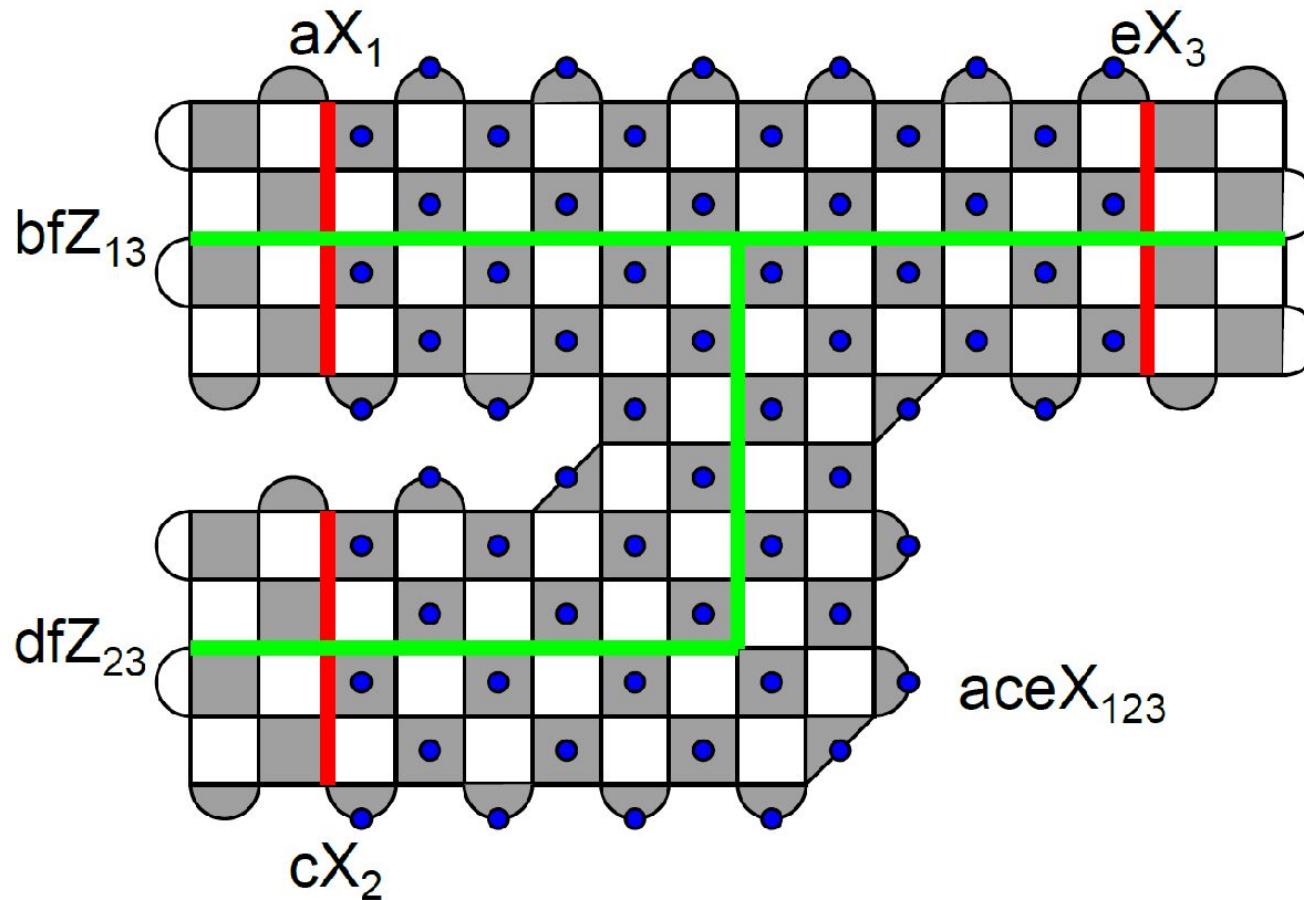
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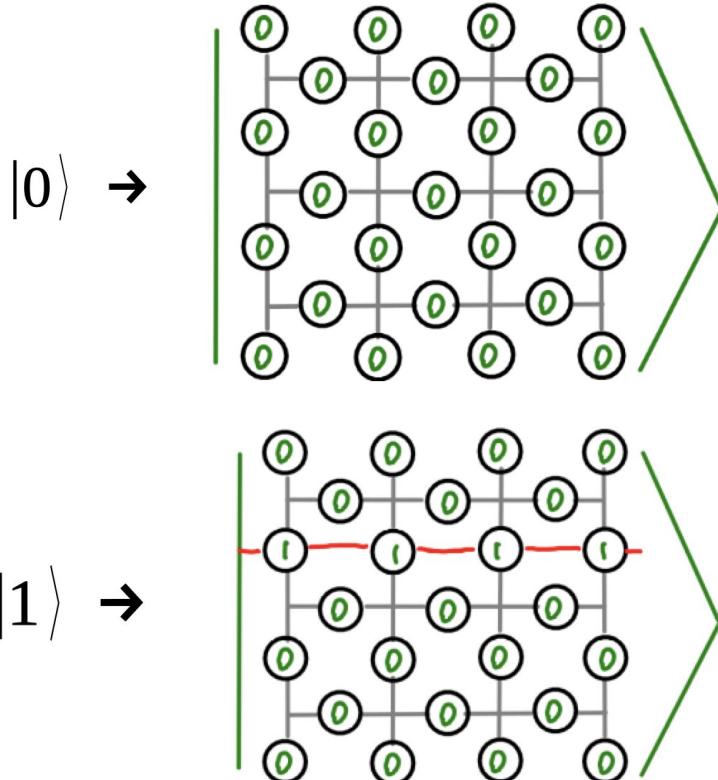


### Multi-body logical X measurement



# More Logical Gates

- We've seen how to do logical X and Z
- A logical CX can be done without much trouble
- A logical H requires the lattice to be rotated, but that can be done
- Other logical Clifford gates can be done with some crazy tricks
- But that's all! No other logical operations can be done fault-tolerantly.
- A solution is *magic state distillation*, using the logical gates we have to clean up the one we don't



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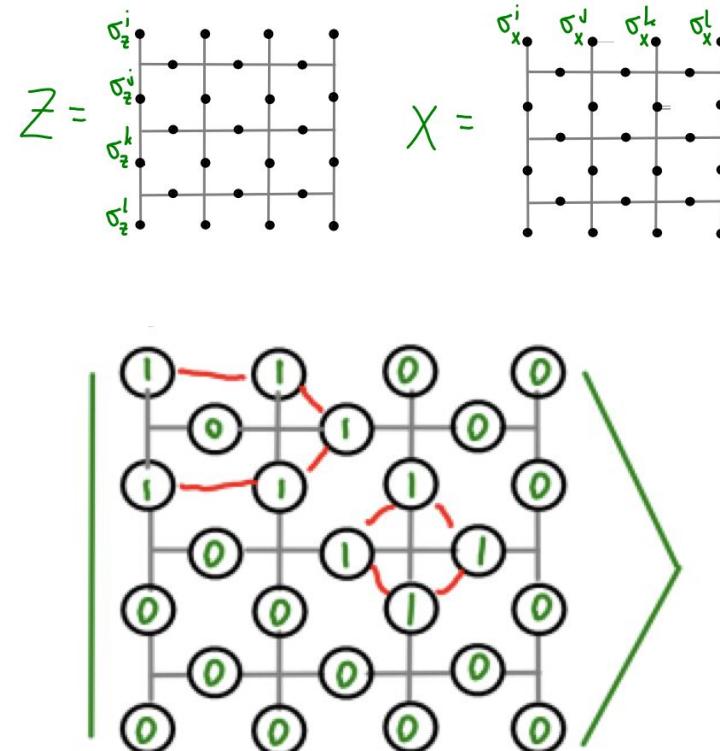
## 4. Decoding

- a. Finding the most probable error
- b. Applying the correction

- Deadline for programming Assignment 2
- 18 May 2024

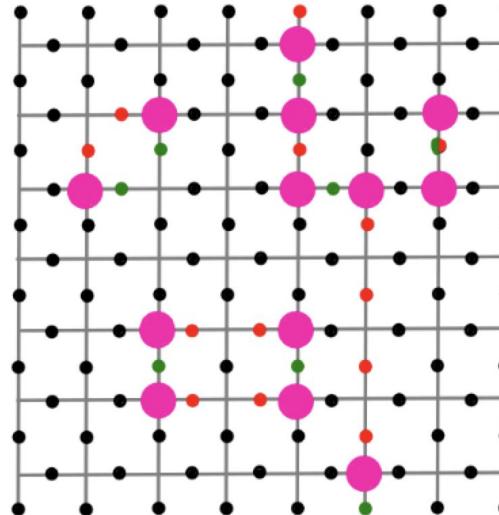
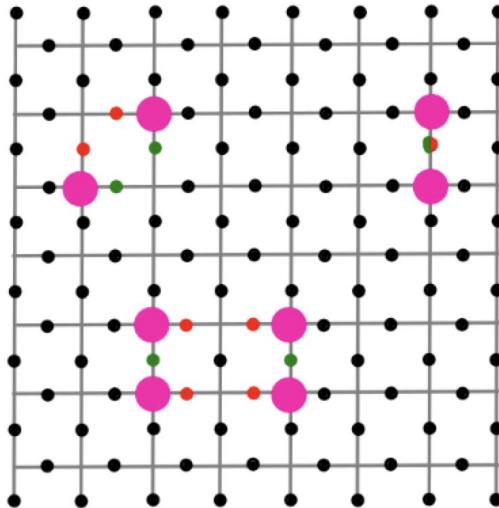
# Final Readout

- The logical operators are many-body observables
- So how do we read them out fault-tolerantly
- When you decide on a basis for final measurement, you stop caring about some errors
- You can then measurement in a product basis
- Final readout and final stabilizer measurement can be constructed from the result
- Measurement errors are effectively the same as errors before measurement



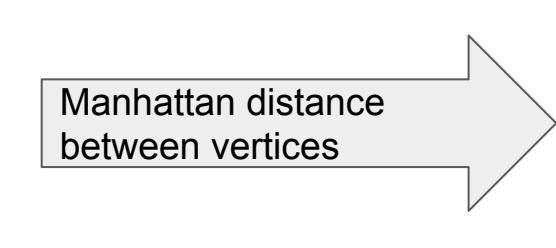
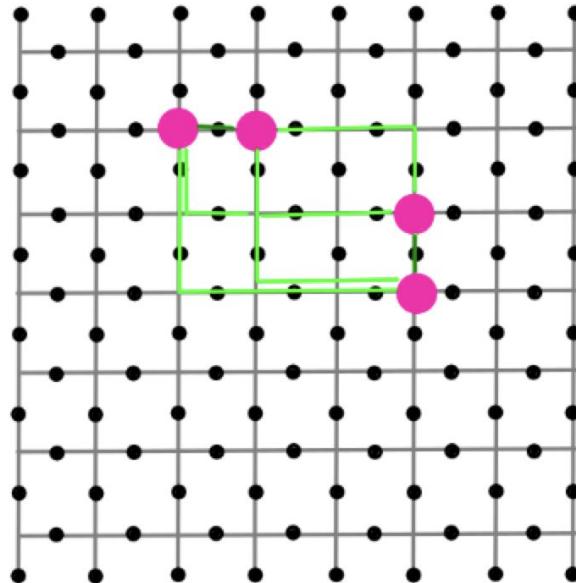
# Decoding

- Given the measurement results, we need to work out what errors happened
- More specifically, the ‘equivalence class’ of errors
- This is the job of the decoding algorithm



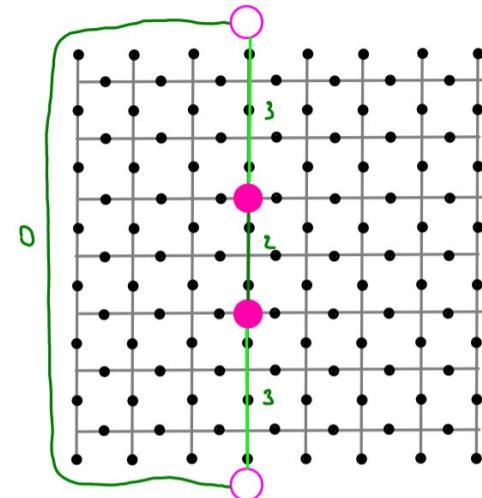
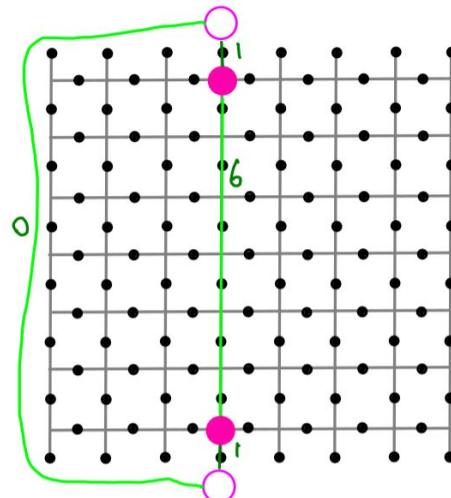
# Decoding with MWPM

- A good option is Minimum Weight Perfect Matching
- We start with the simple and unrealistic case: errors only between measurement
- Each round can be decoded separately, corresponding to MWPM on a 2D graph
- Decoding for X and Z errors can also be done independently



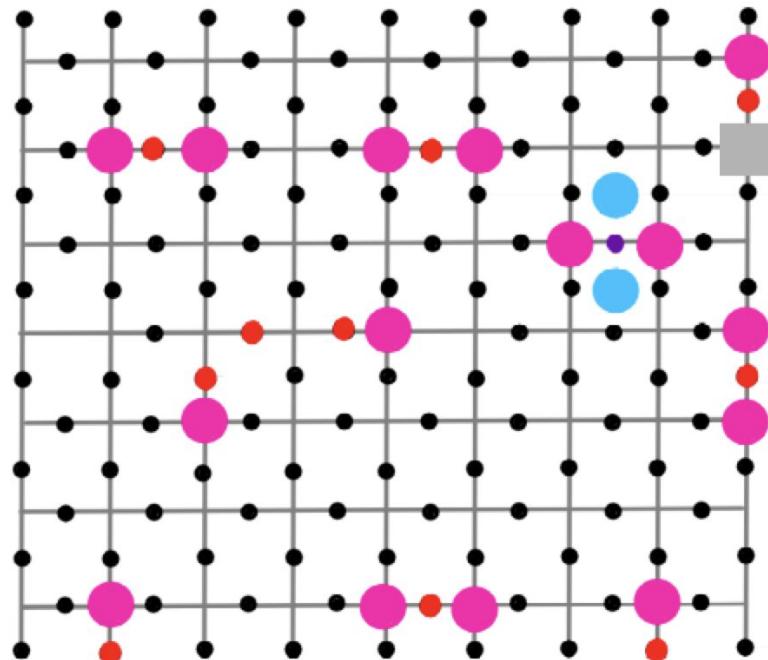
# Decoding

- We need to be careful to account for the effects of the edges
- This is done by introducing extra ‘virtual nodes’

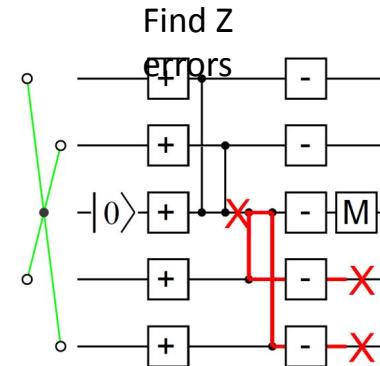
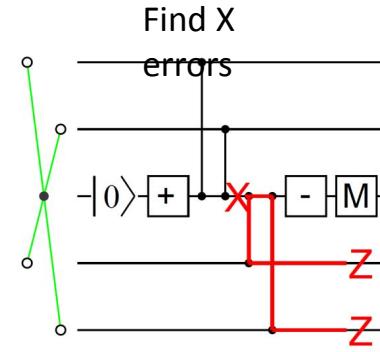
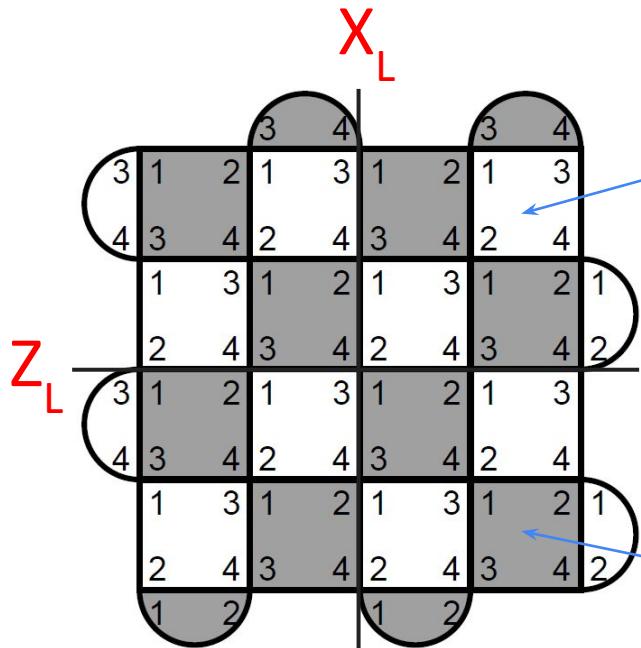


# Imperfect Measurements

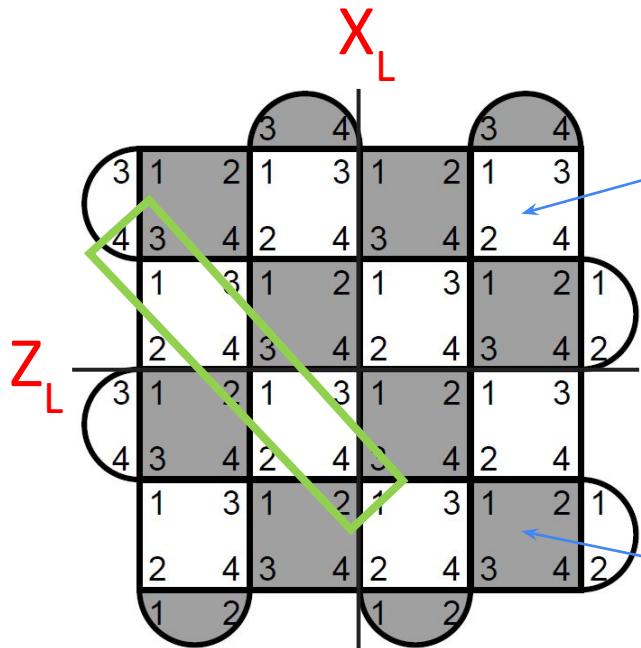
- We have the problem of imperfect measurements
  - The measurements might lie
  - Errors on the additional qubit
  - Errors in the CX gates
- We base the decoding using syndrome changes
- This leads to a 3D MWPM problem (2D space + time)



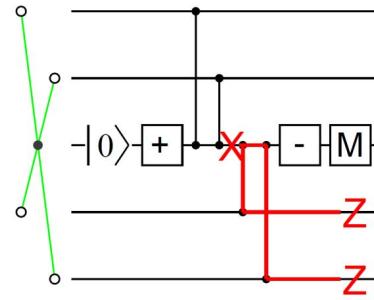
# Gate sequence:



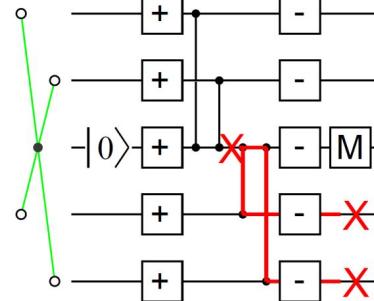
# Gate sequence:



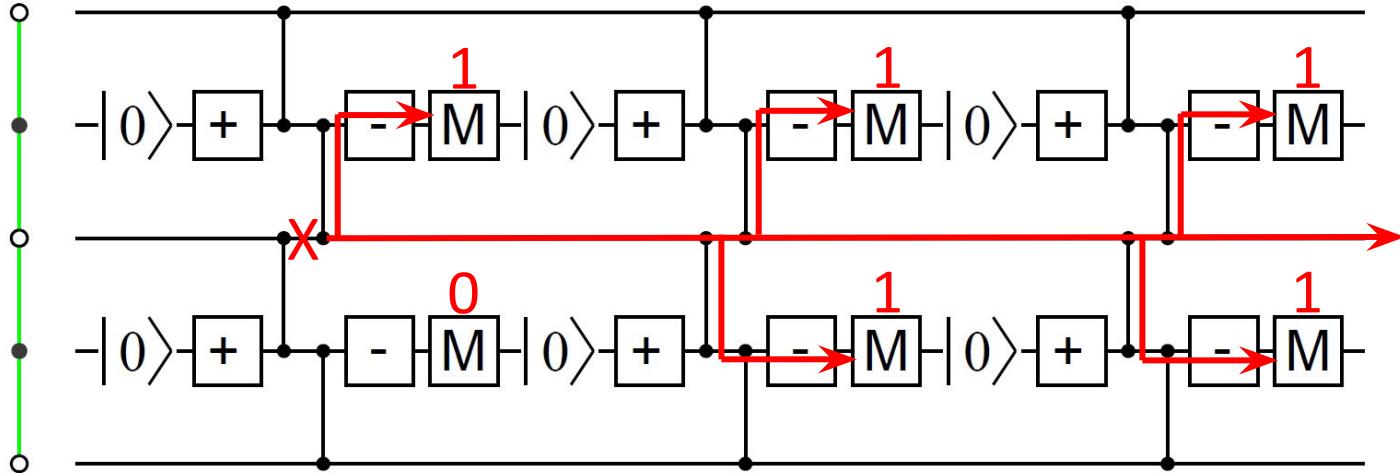
Find X errors



Find Z errors

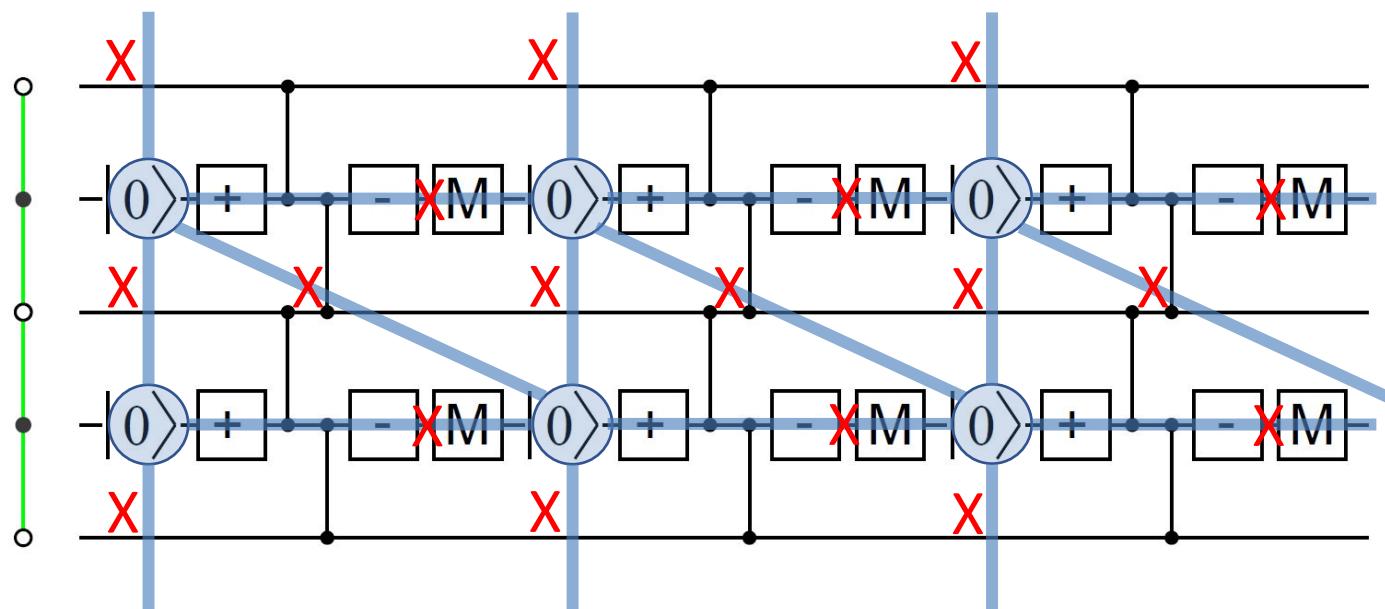


# How to do memory:



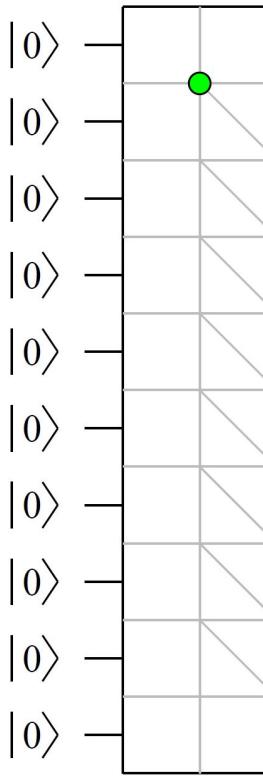
Measurement value change = detection event

# How to do memory:



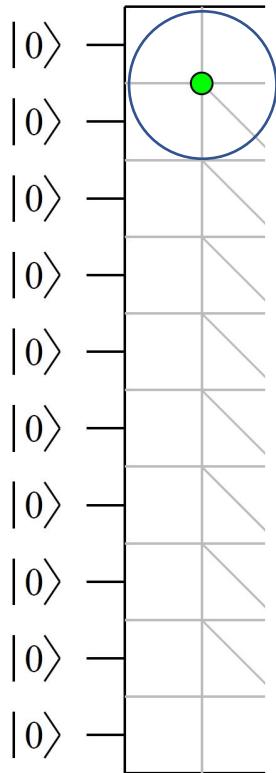
Build graph of all possible detection events

# Classical processing



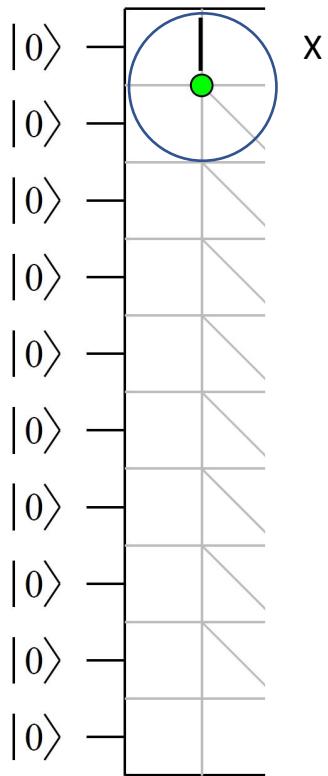
- 10 data qubits
- One detection event

# Classical processing



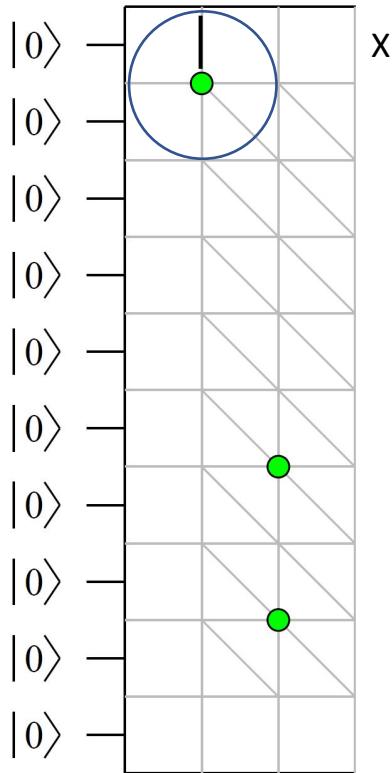
- 10 data qubits
- One detection event
- Explore uniformly,  
boundary found

# Classical processing



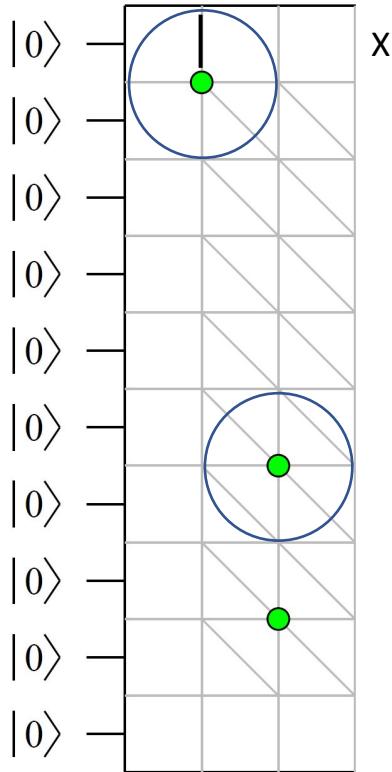
- 10 data qubits
- One detection event
- Explore uniformly,  
boundary found
- Match detection event to  
boundary, record belief that  
X error present

# Classical processing



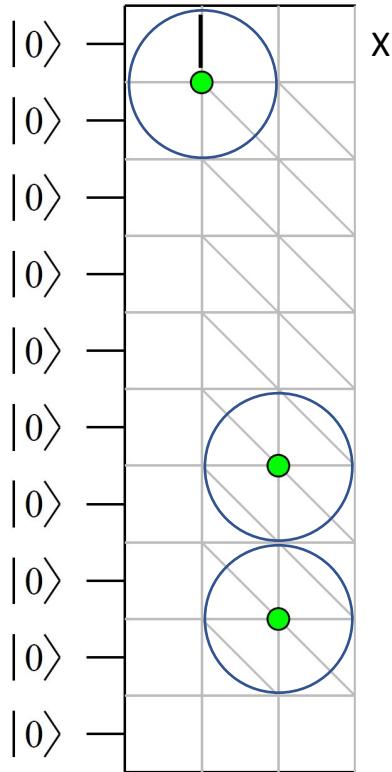
- Two more detection events

# Classical processing



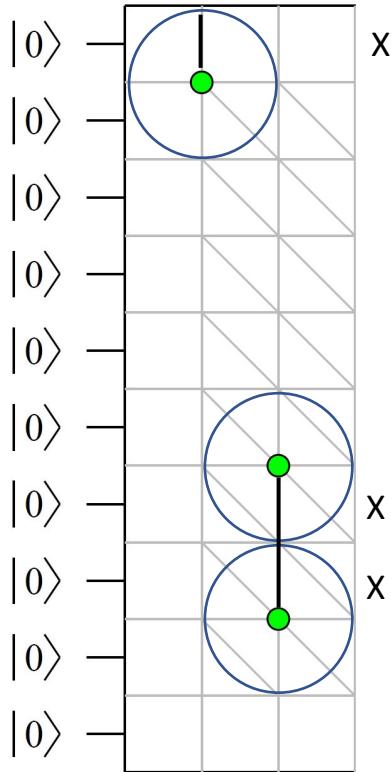
- Two more detection events
- Pick one, explore, current time boundary encountered

# Classical processing



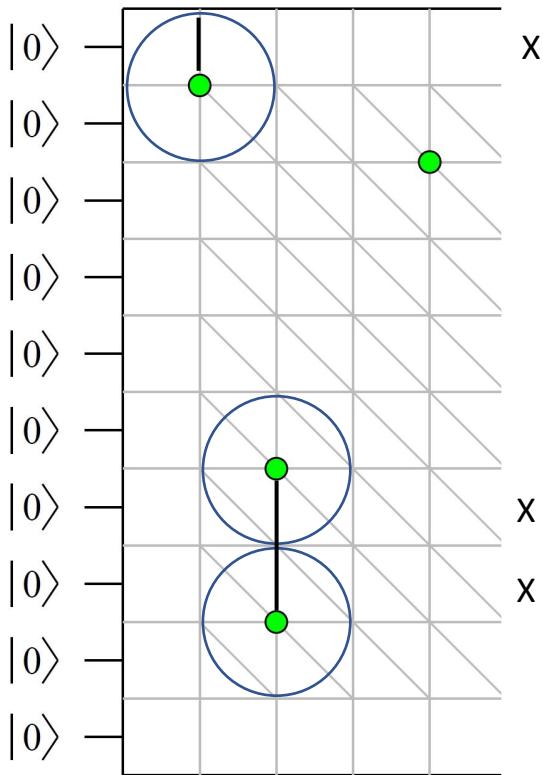
- Two more detection events
- Pick one, explore, current time boundary encountered
- Explore around other, exploratory regions touch

# Classical processing



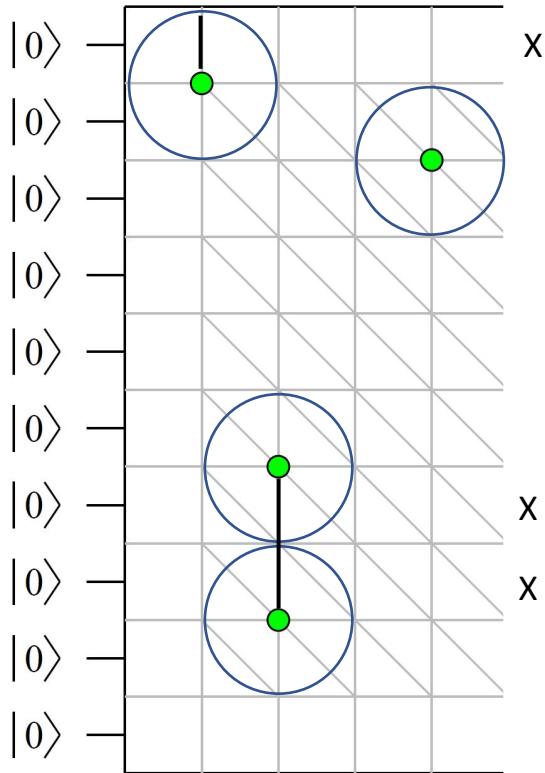
- Two more detection events
- Pick one, explore, current time boundary encountered
- Explore around other, exploratory regions touch
- Match, record belief that two more X errors present

# Classical processing



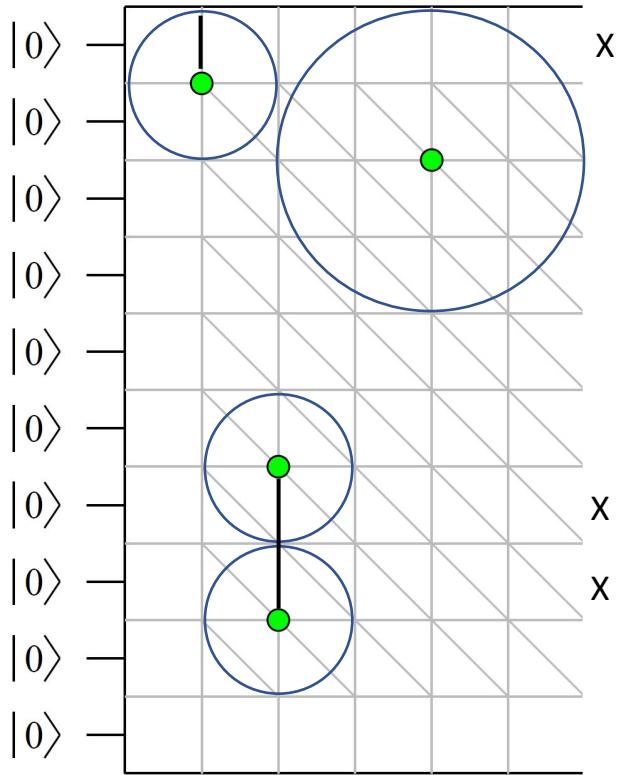
- One more detection event

# Classical processing



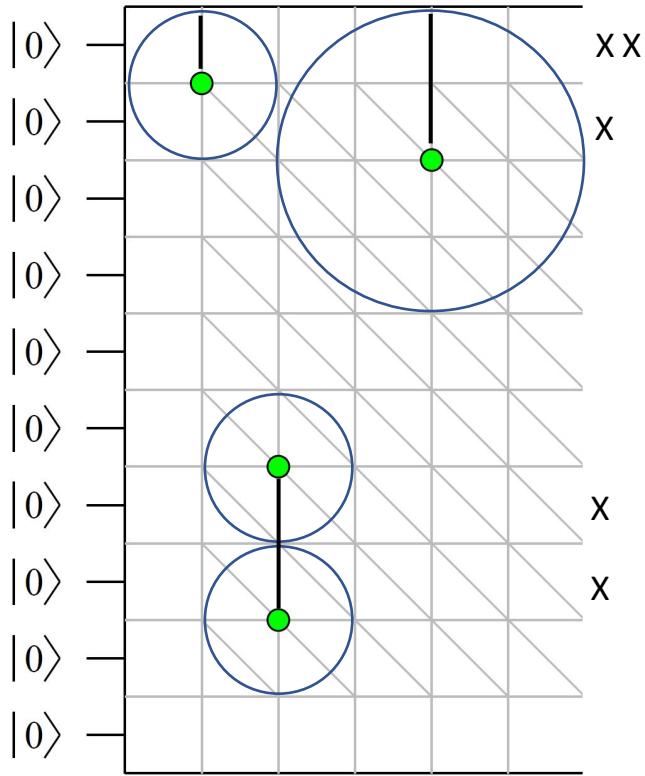
- One more detection event
- Explore, current time boundary encountered, must wait for more data

# Classical processing



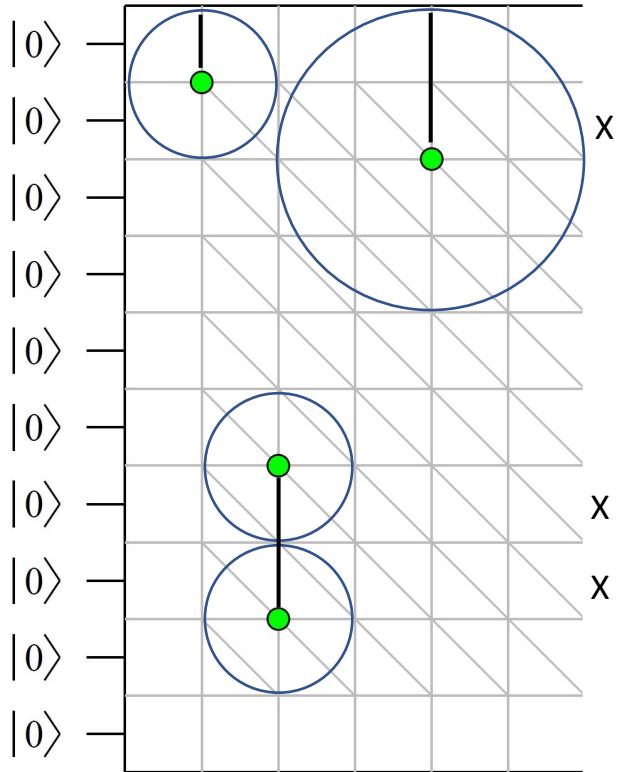
- One more detection event
- Explore, current time boundary encountered, must wait for more data
- Explore further, boundary encountered

# Classical processing



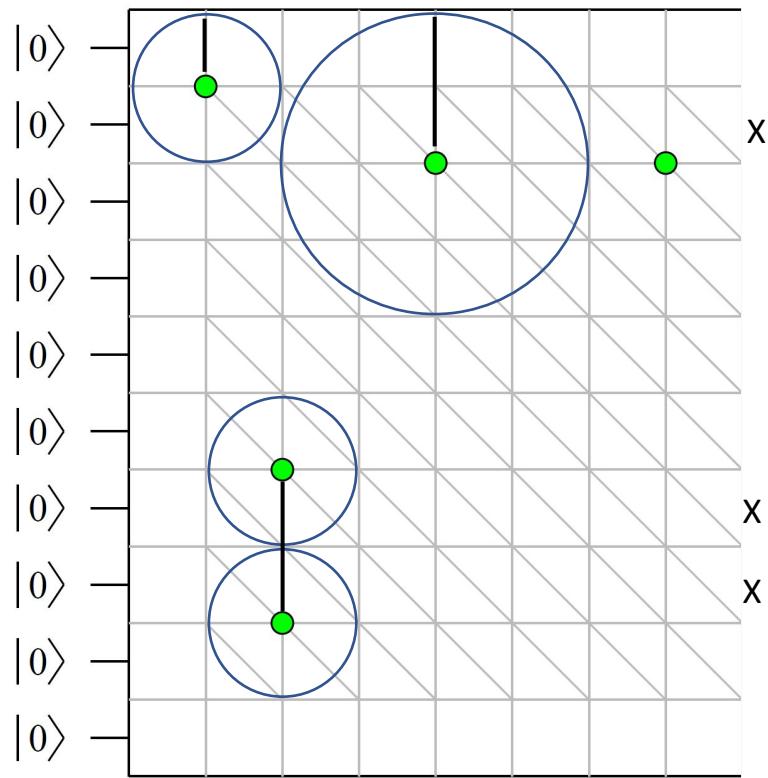
- One more detection event
- Explore, current time boundary encountered, must wait for more data
- Explore further, boundary encountered
- Match, record belief that two more X errors present

# Classical processing

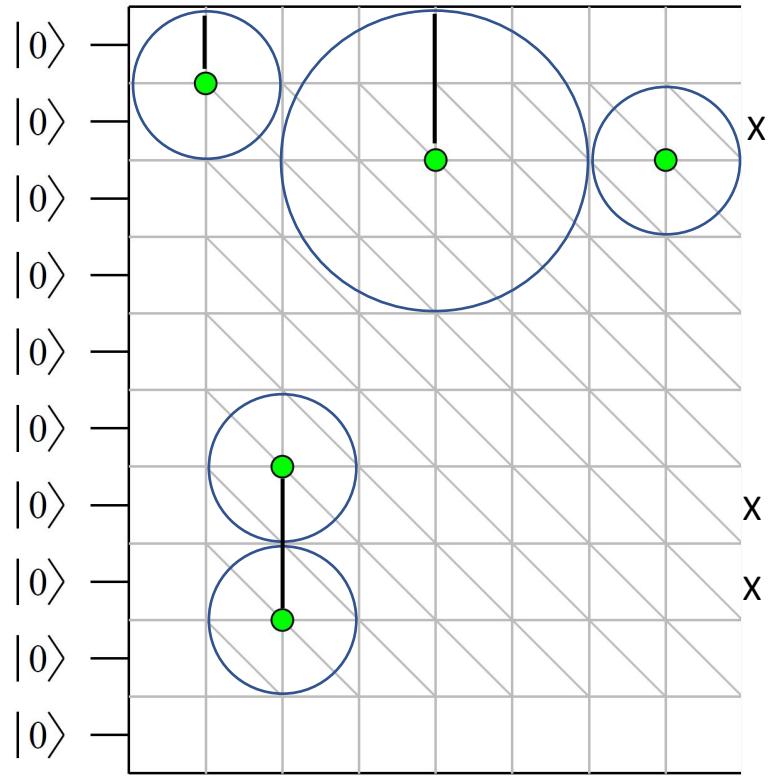


- One more detection event
- Explore, current time boundary encountered, must wait for more data
- Explore further, boundary encountered
- Match, record belief that two more X errors present
- Cancel double error
- **Don't apply physical corrections**

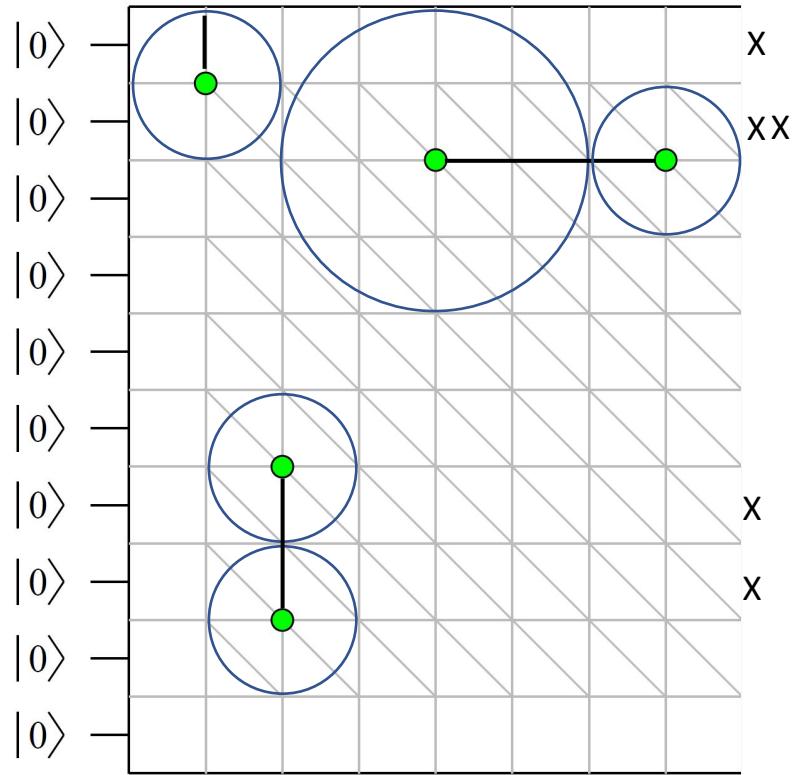
# Classical processing



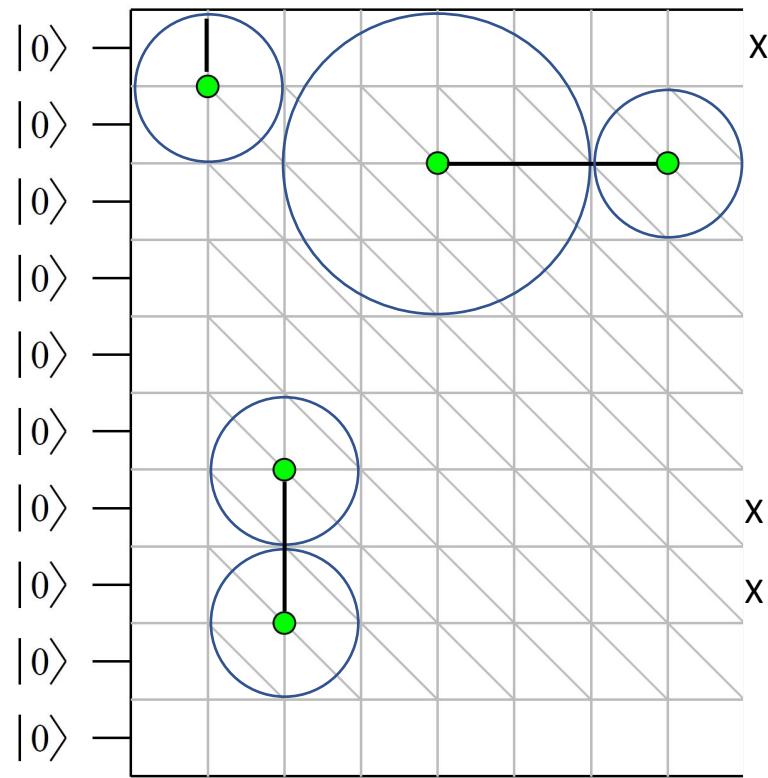
# Classical processing



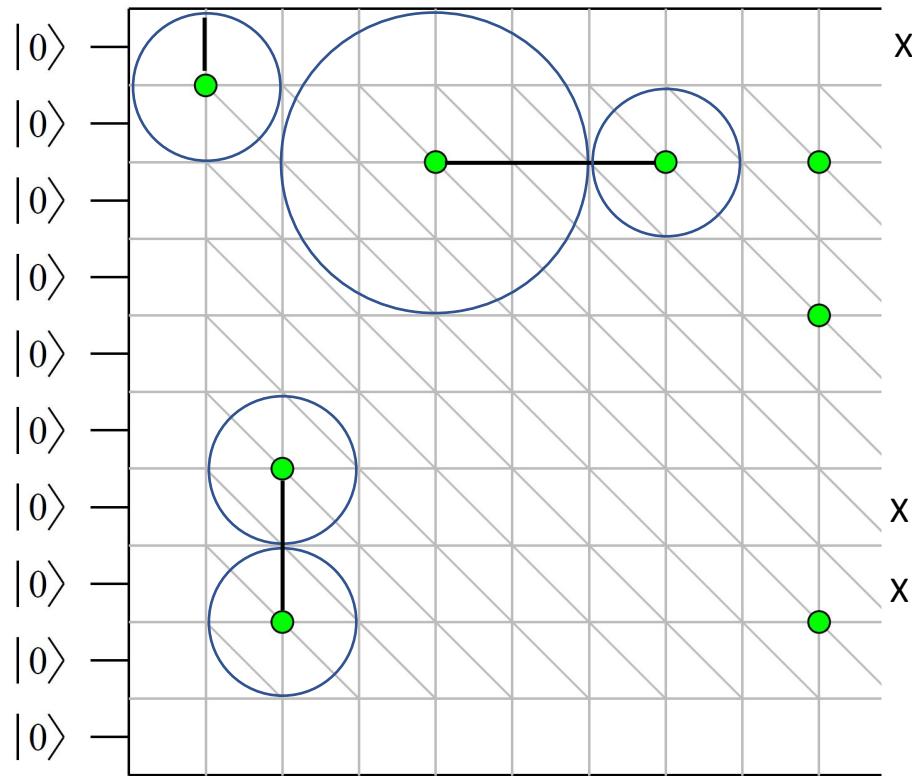
# Classical processing



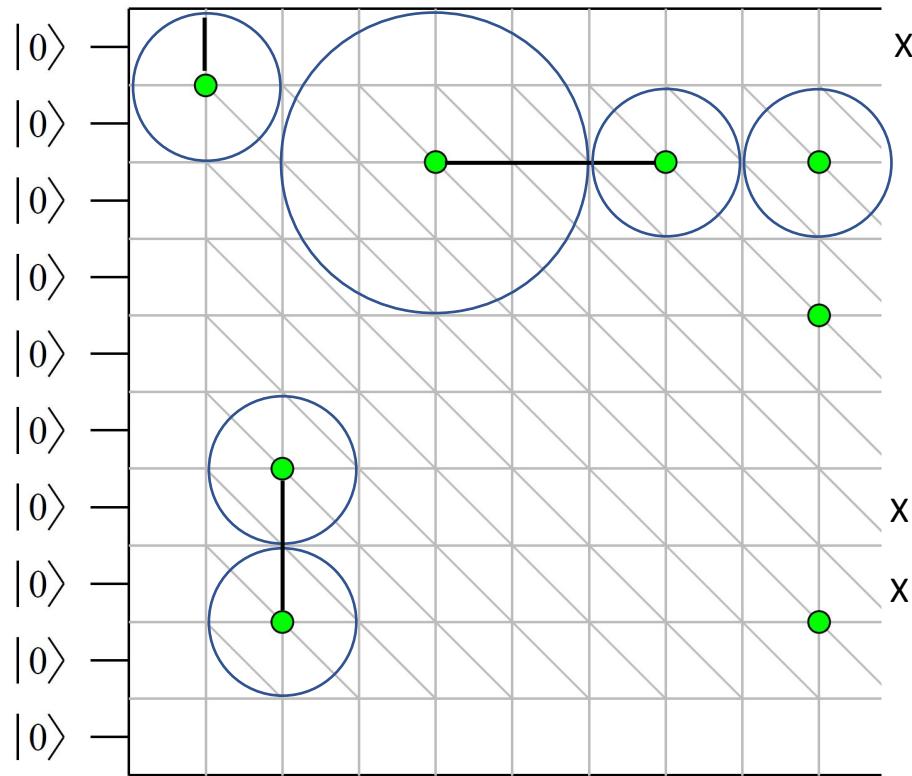
# Classical processing



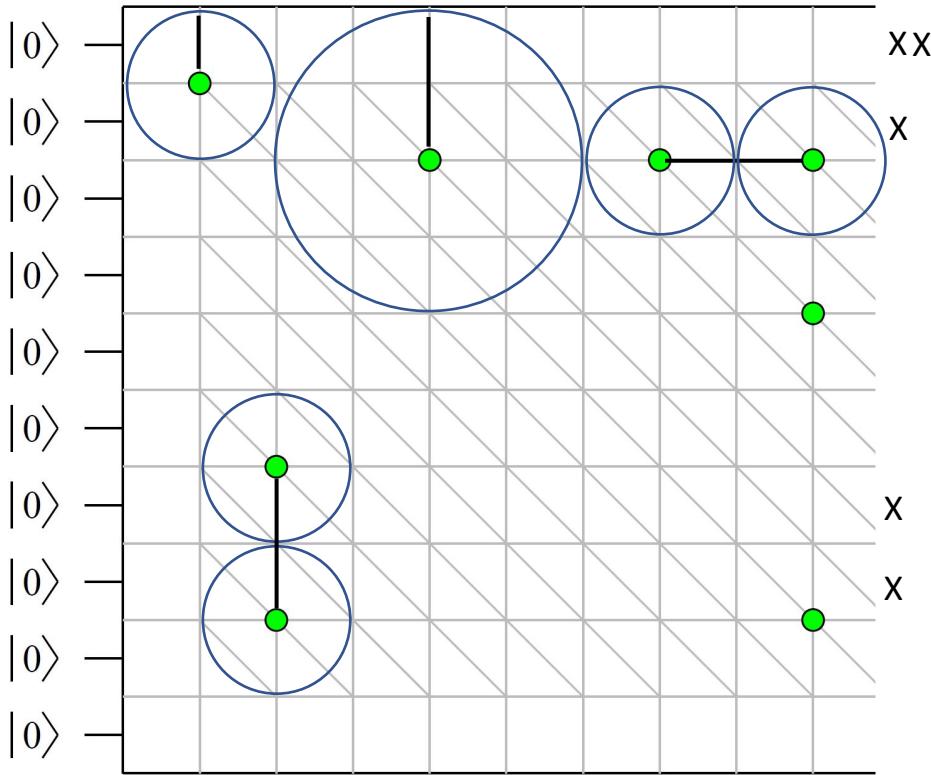
# Classical processing



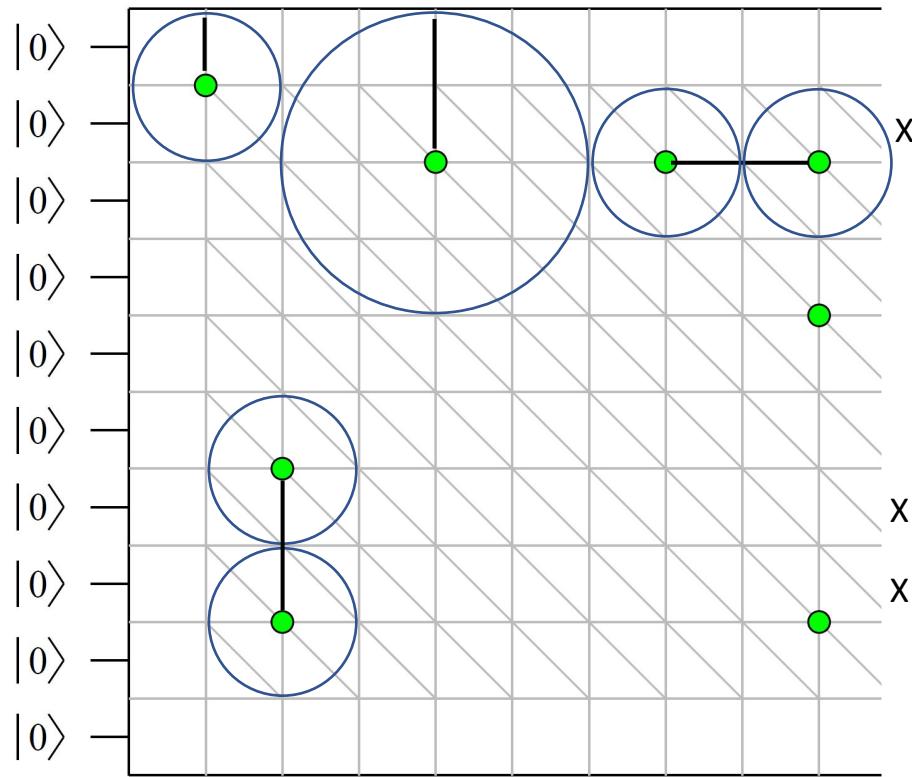
# Classical processing



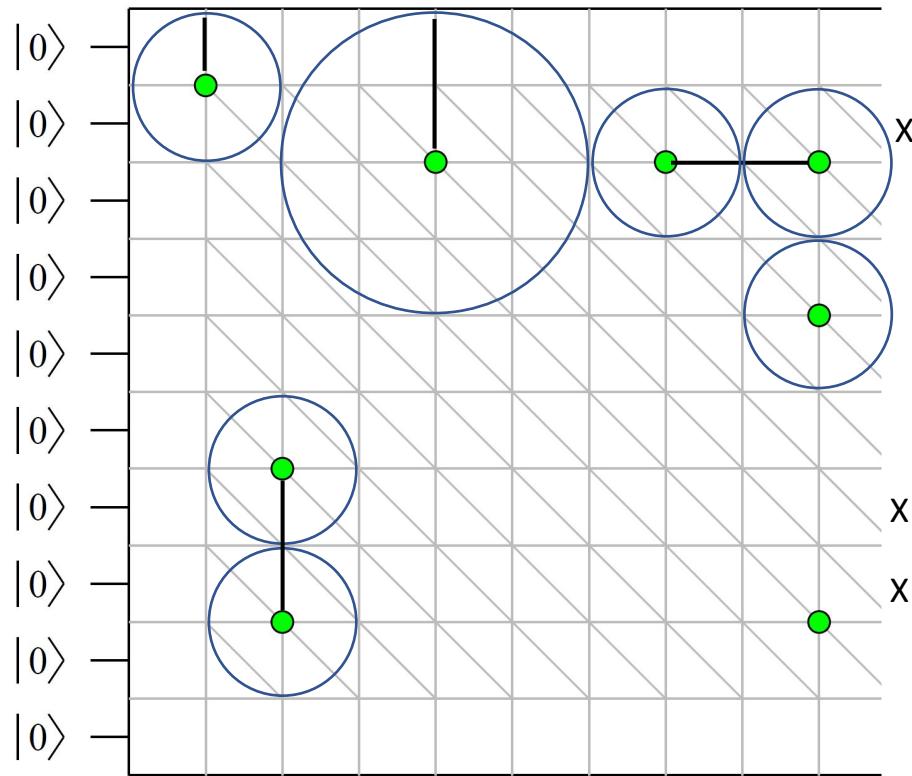
# Classical processing



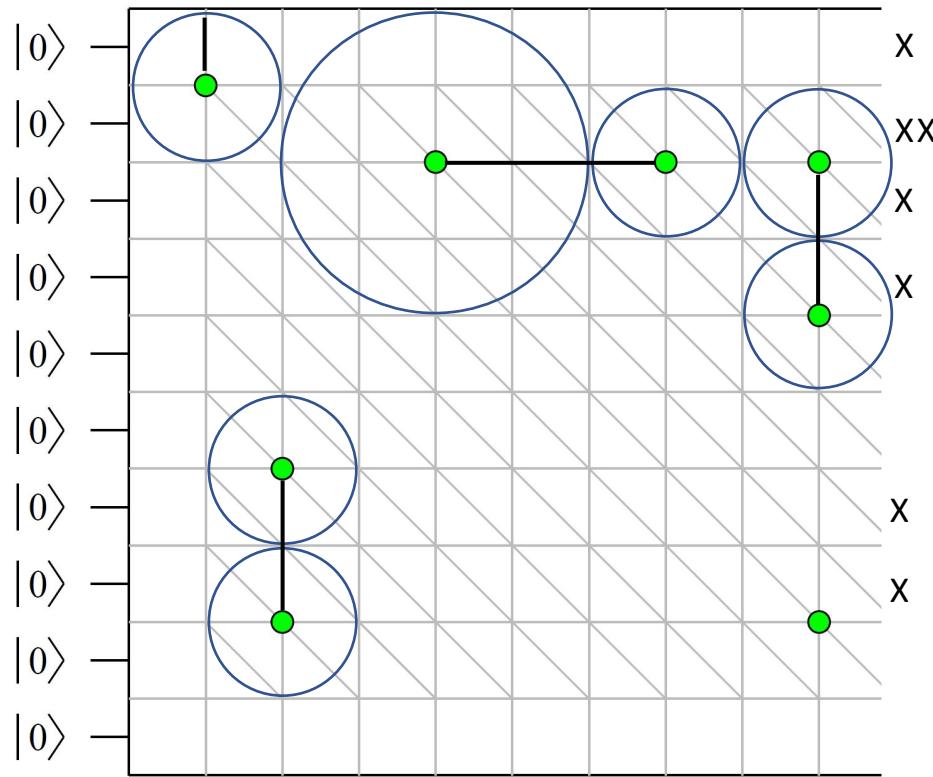
# Classical processing



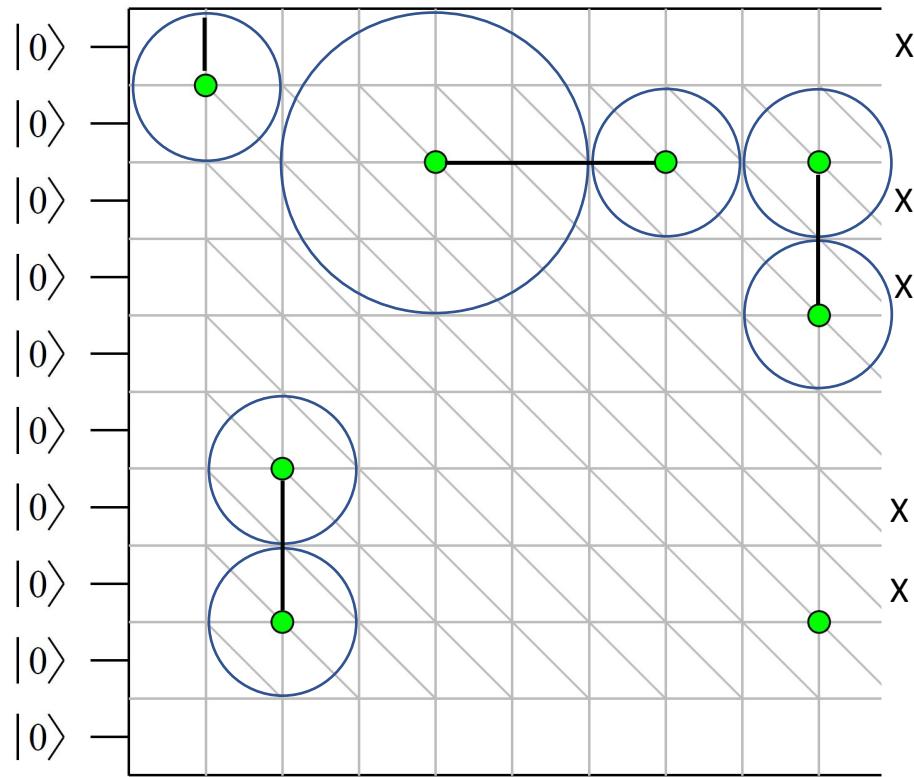
# Classical processing



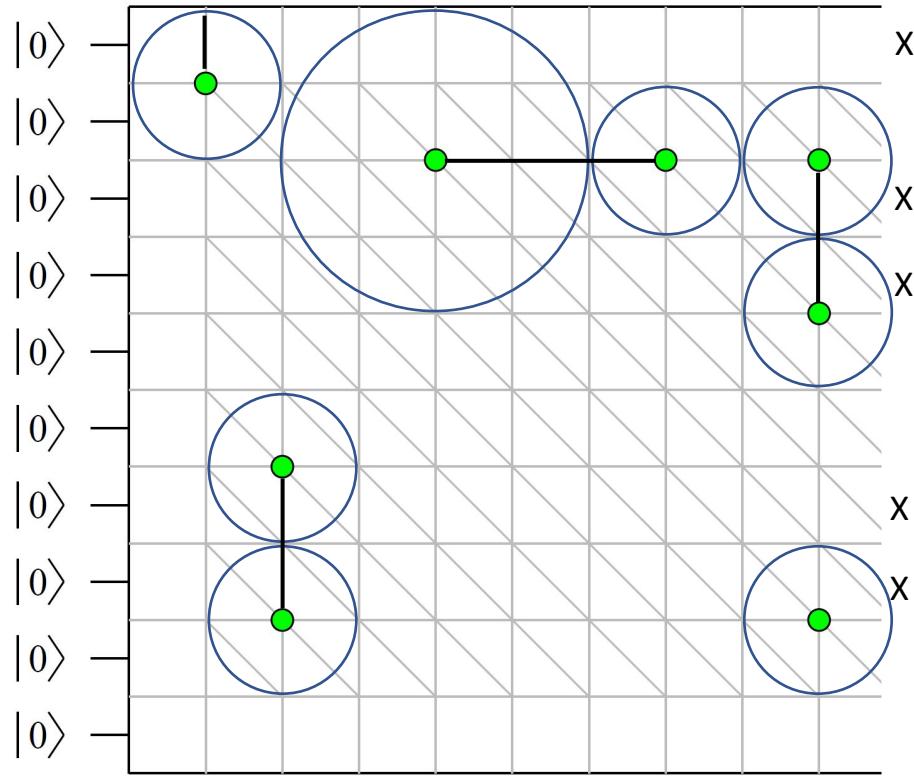
# Classical processing



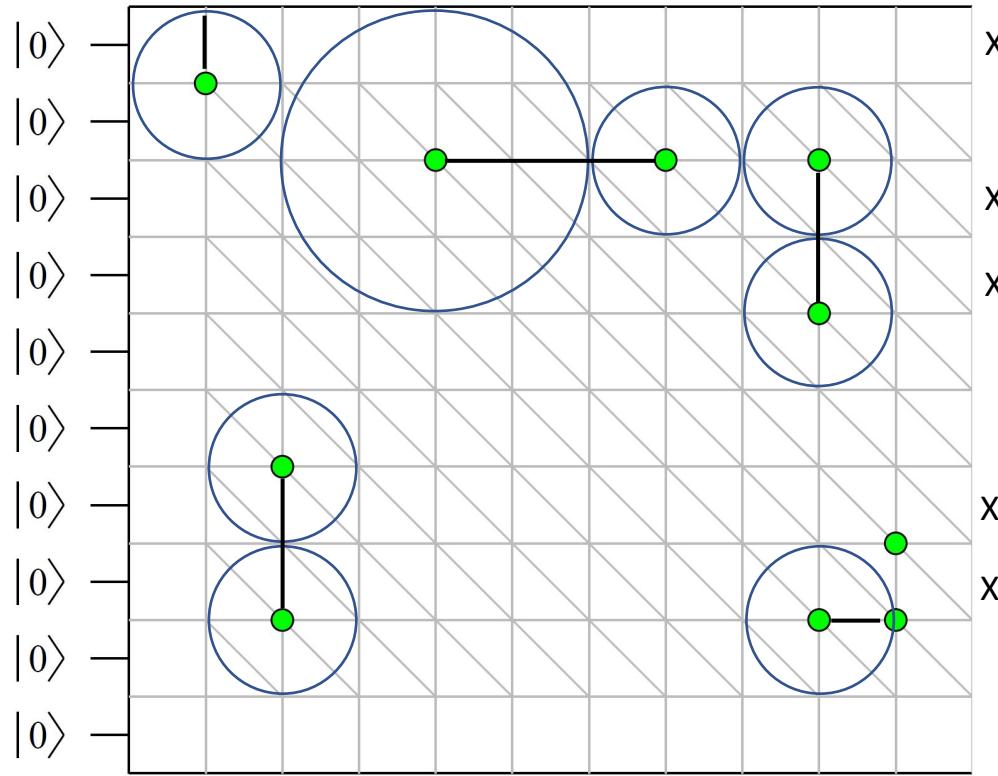
# Classical processing



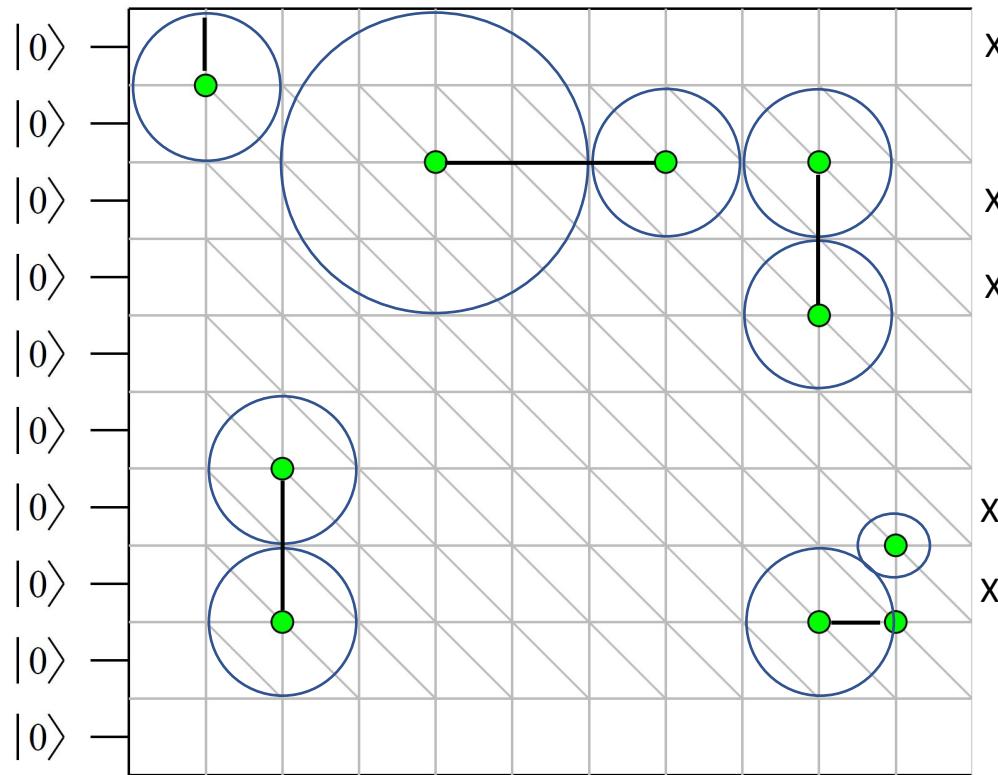
# Classical processing



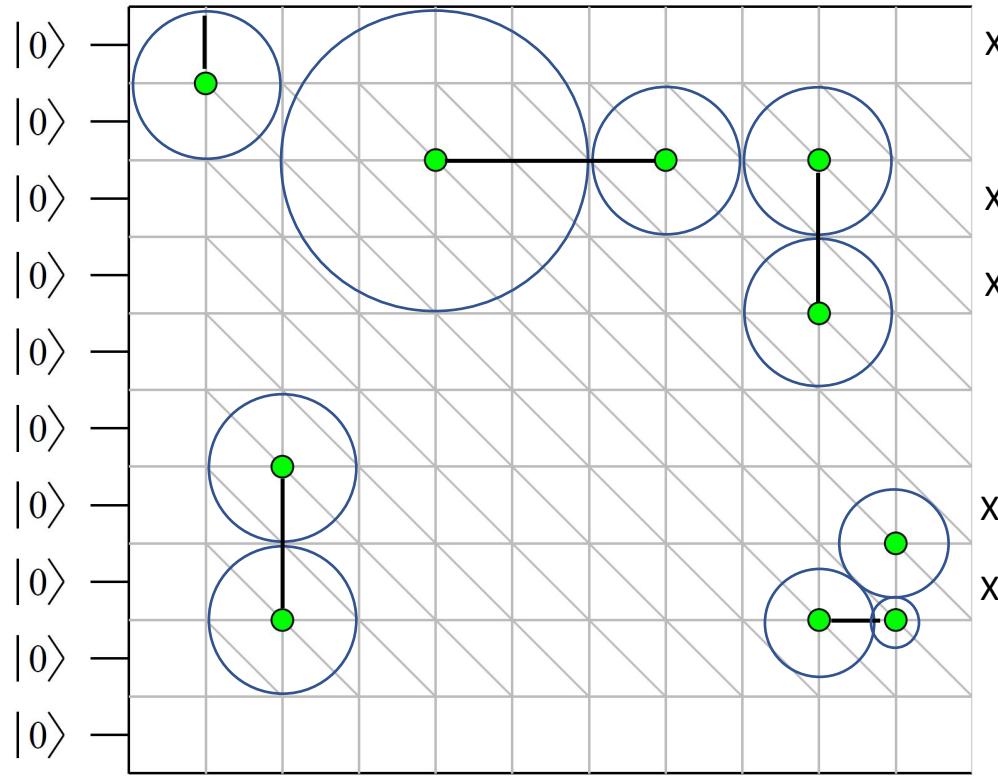
# Classical processing



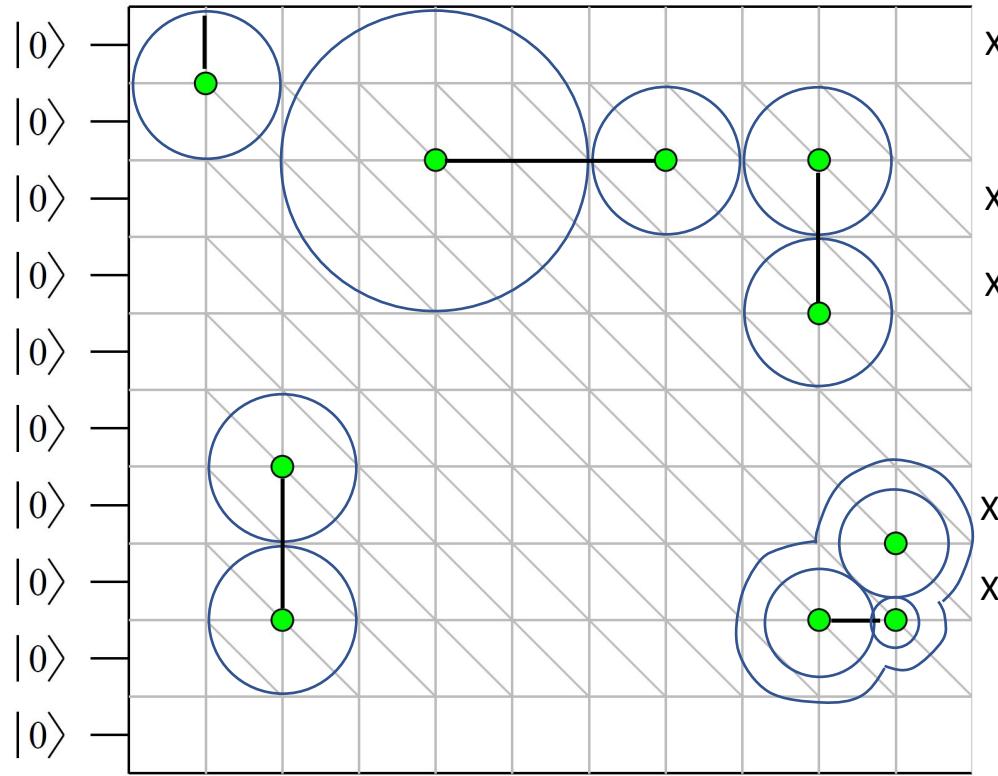
# Classical processing



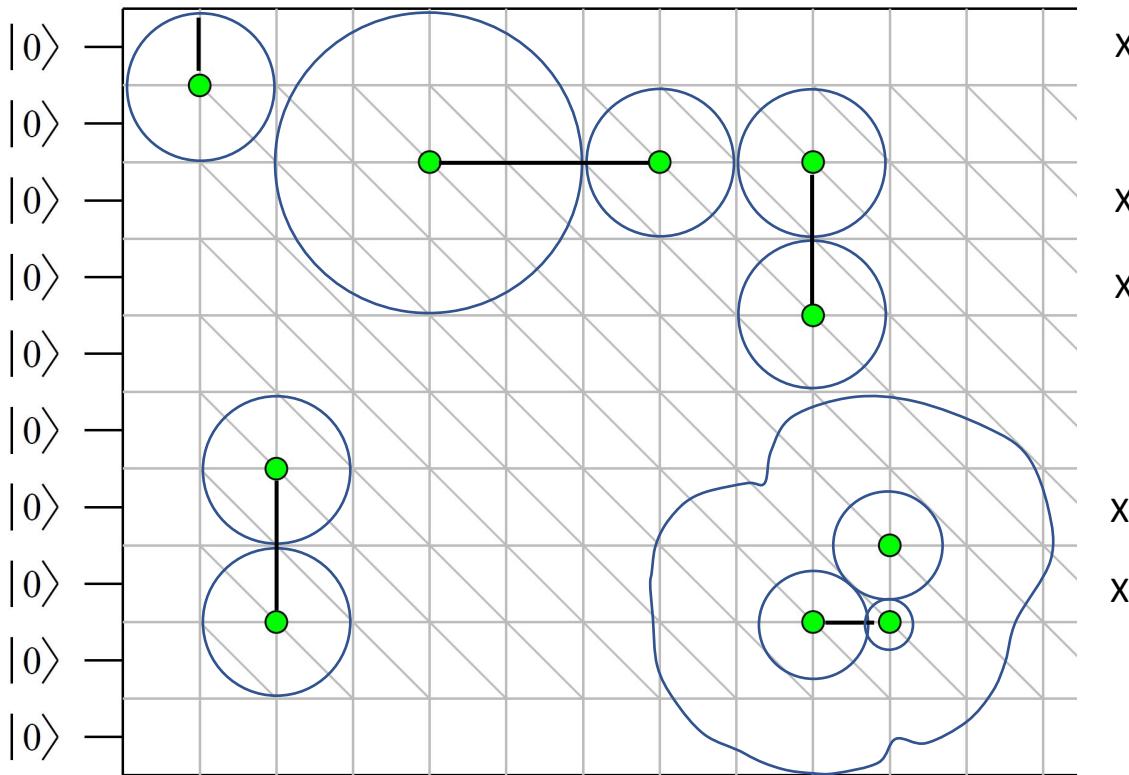
# Classical processing



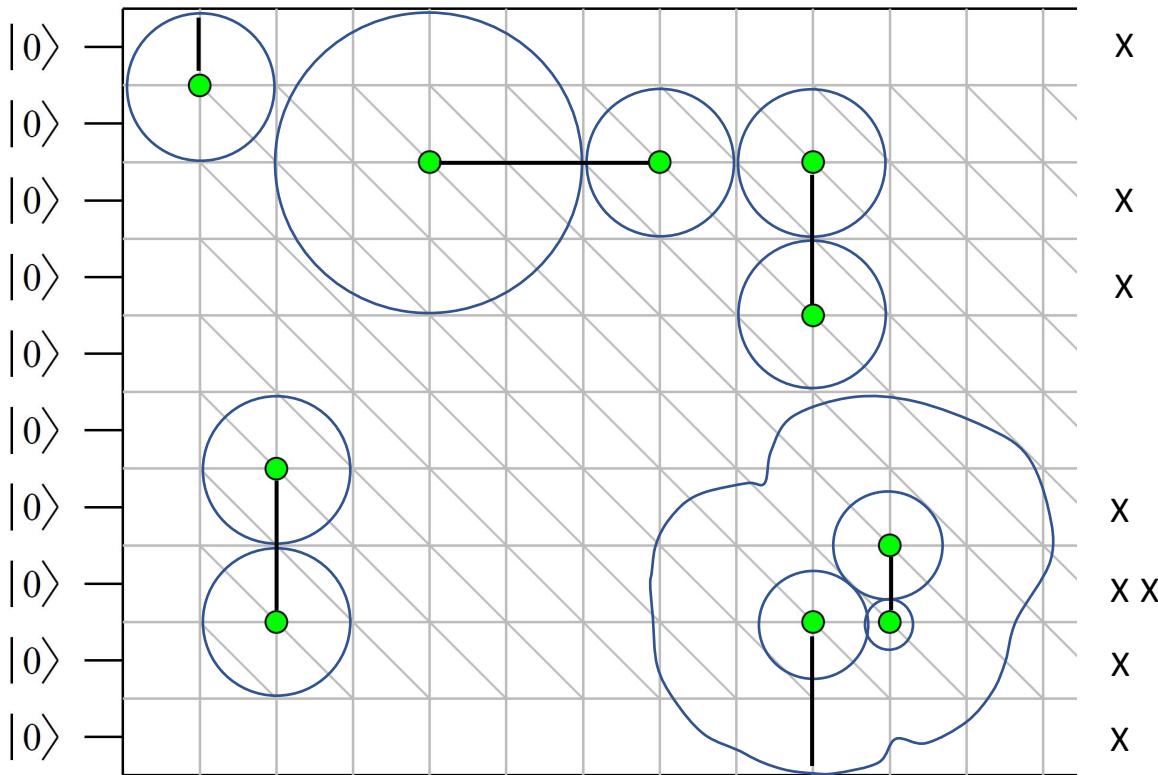
# Classical processing



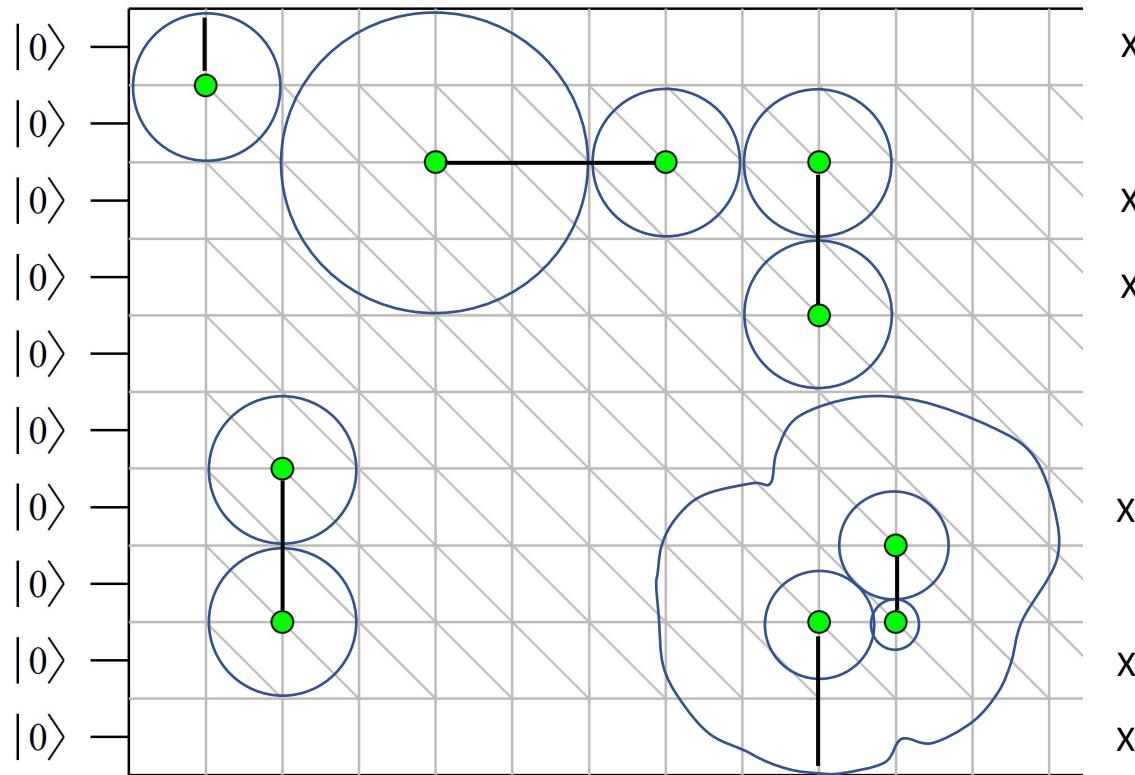
# Classical processing



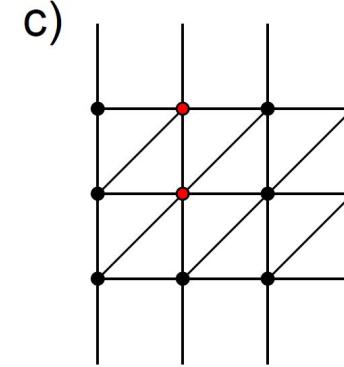
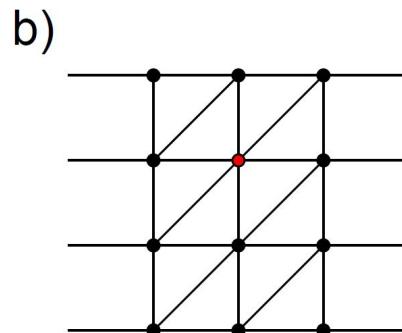
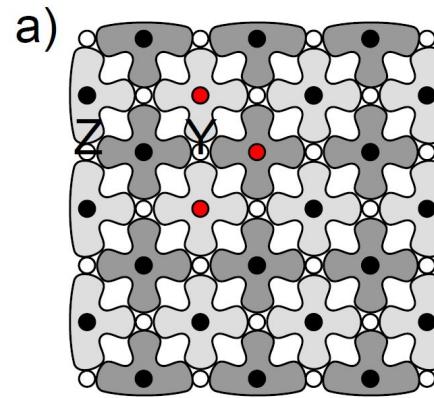
# Classical processing



# Classical processing

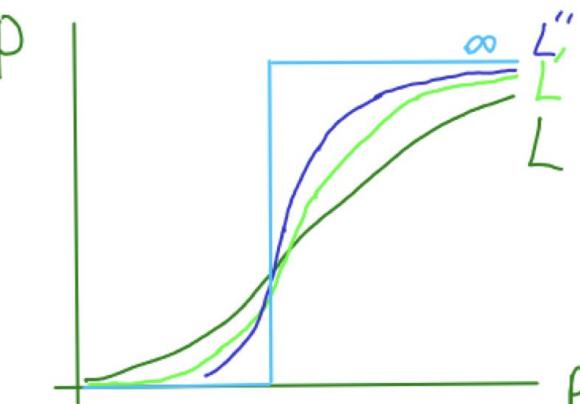
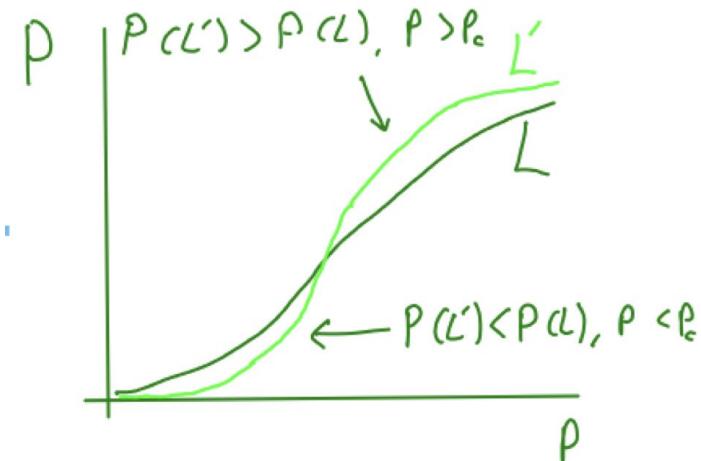


# Correlated errors



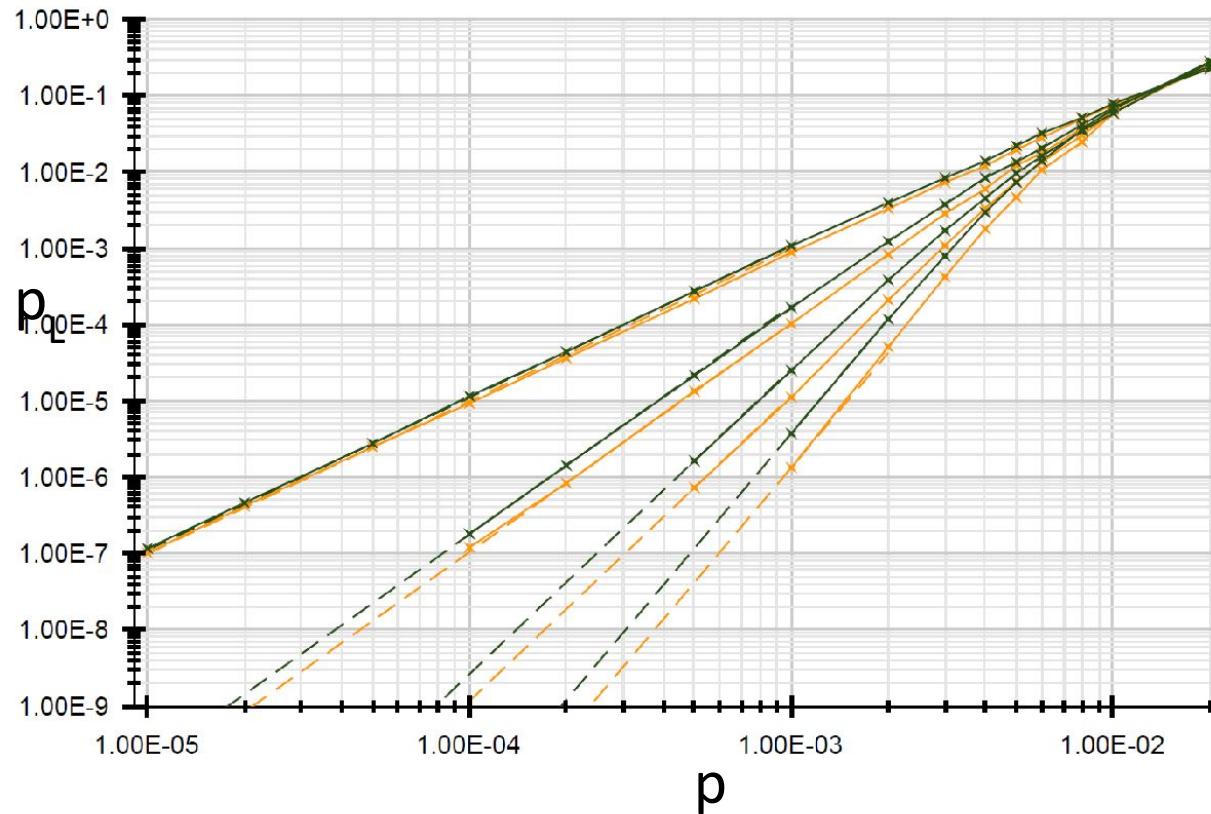
# Threshold

- Correcting according to the right class removes the effects of errors
- Correcting according to the wrong class causes an operation on the encoded qubit (without our knowing)
- What is the probability of such an error,  $P$ , given the probability on the qubits of the code,  $p$ ?
- We find a phase transition as  $L$  is increased (for an  $L \times L$  grid)



# Simulated performance

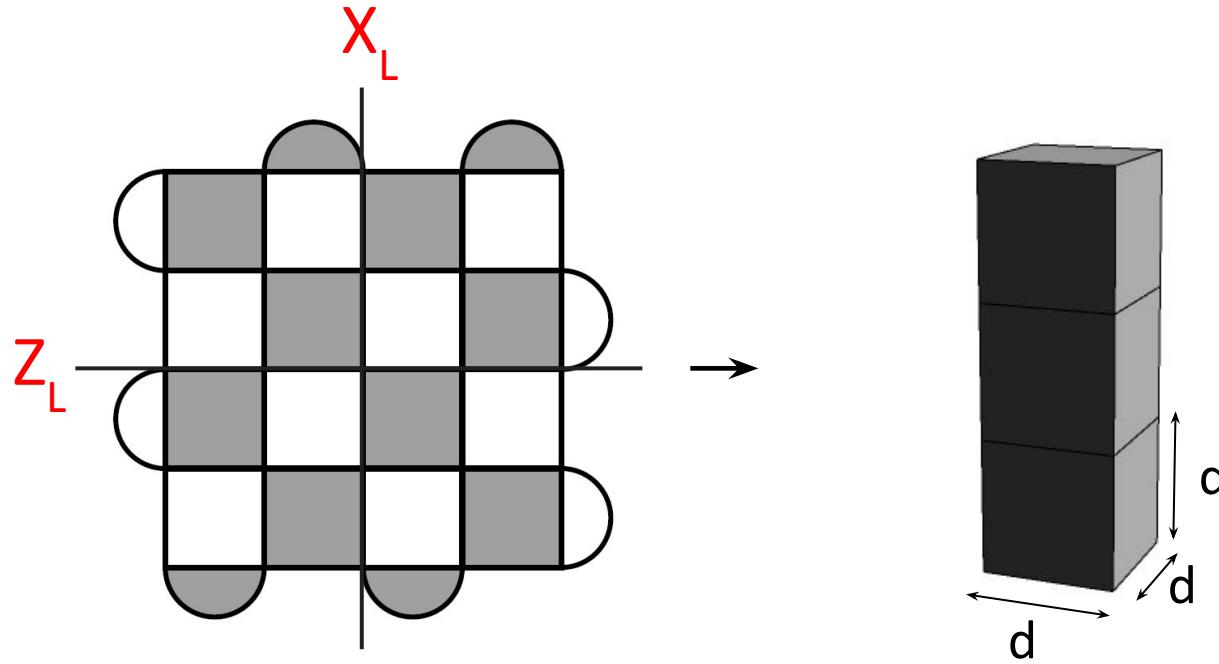
Rotated distance 3, 5, 7, 9 uncorrelated (green) and correlated (orange)



- $p_L = 0.1(100p)^{(d+1)/2}$
- O(1) parallel algorithm
- Low latency

# **Appendix: Logical Gates and Experiments**

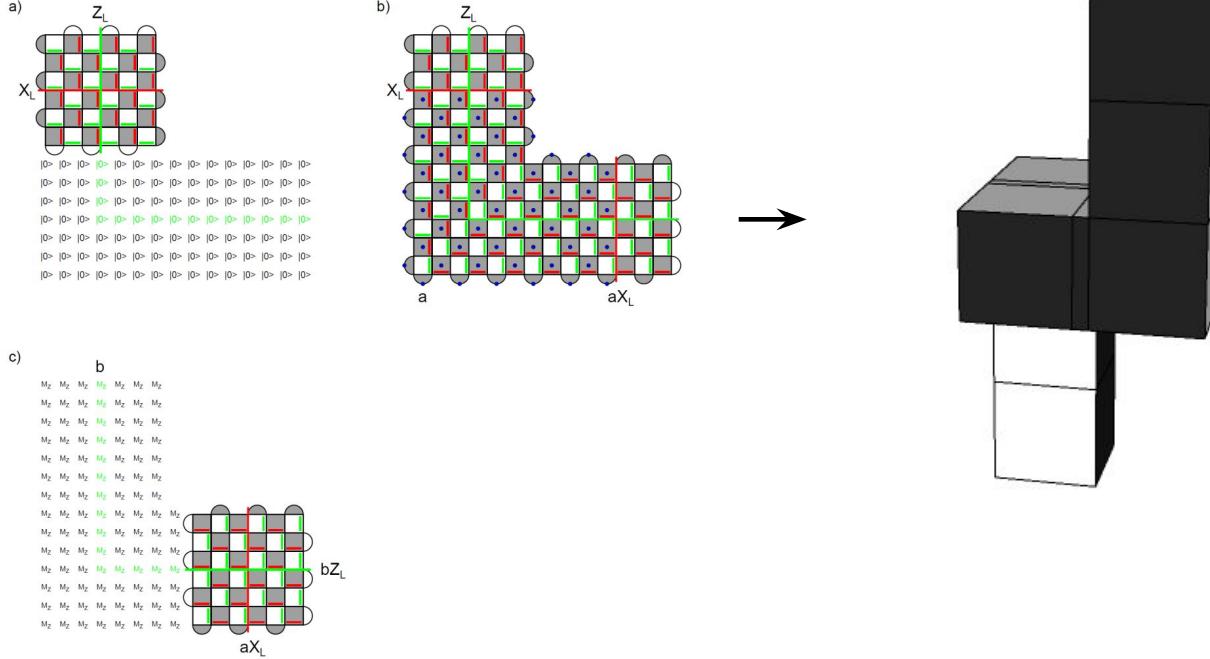
# Logical identity



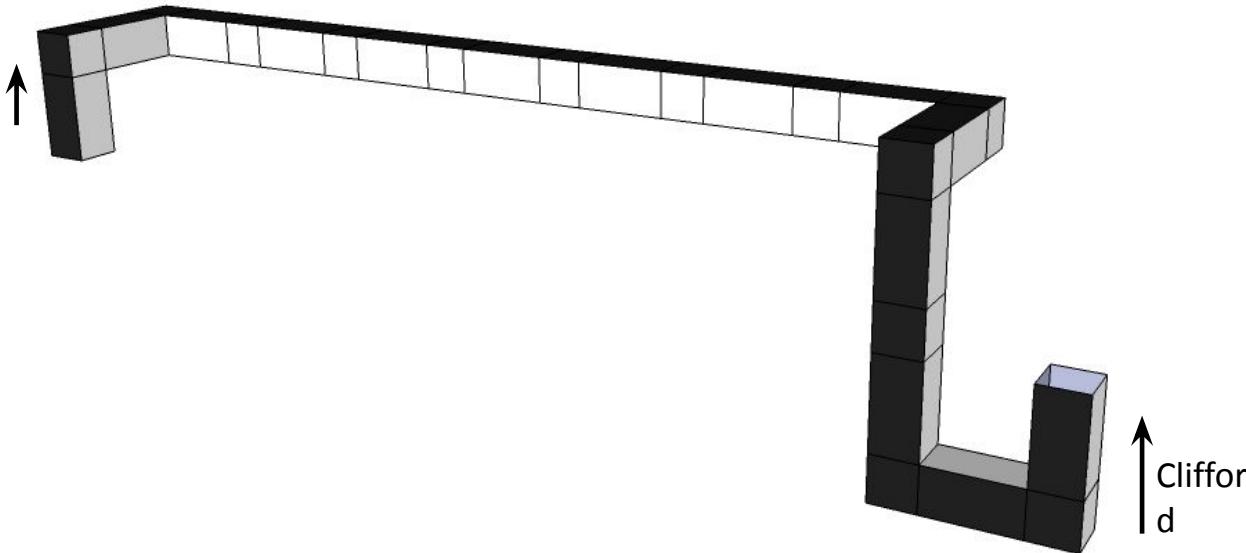
detect X  
 detect Z

$$p_L = 0.1(100p)^{(d+1)/2}$$

# Logical move

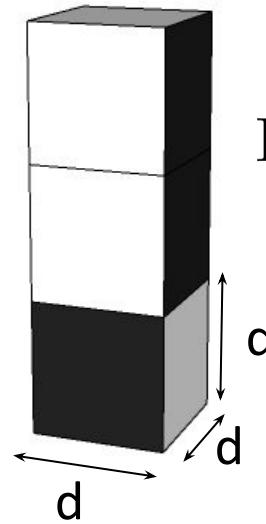
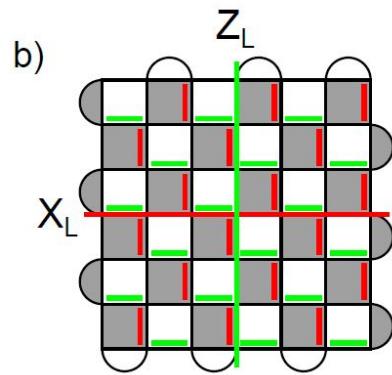
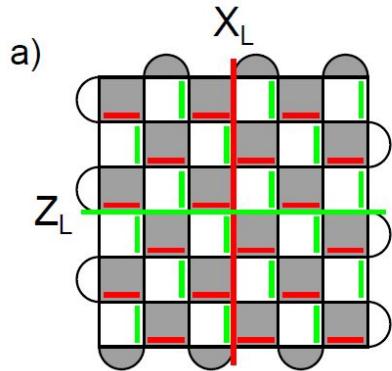


# Logical move



Can move anywhere, even back in time, up to Pauli operators.

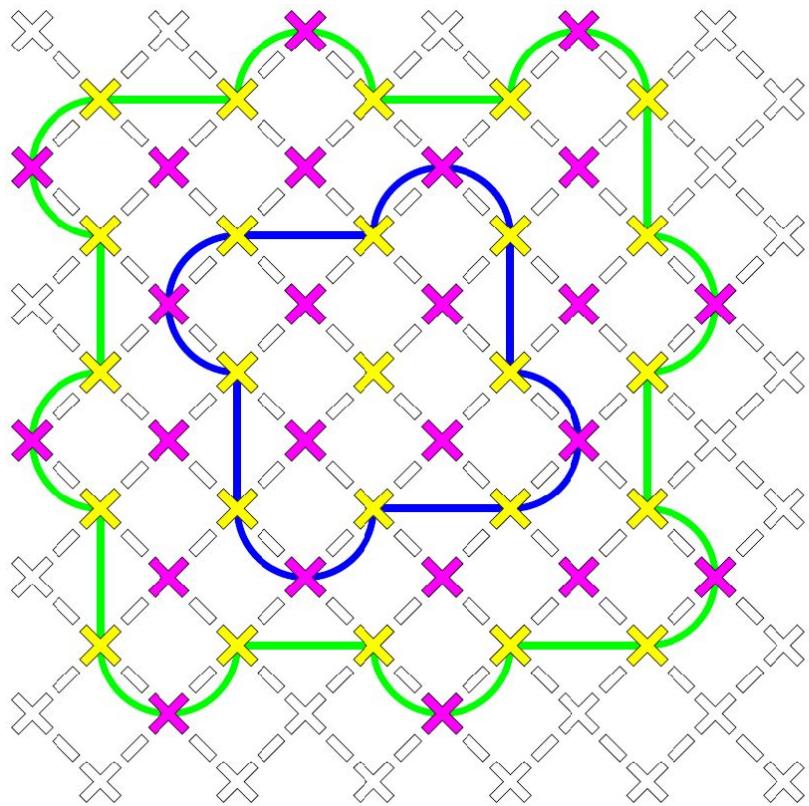
# Logical Hadamard



$$H = \frac{1}{\sqrt{2}} \begin{bmatrix} 1 & 1 \\ 1 & -1 \end{bmatrix}$$

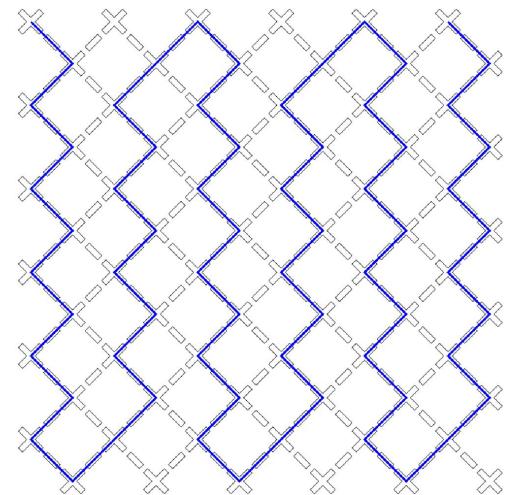
$$p_L = 0.1(100p)^{(d+1)/2}$$

## Device and experiment

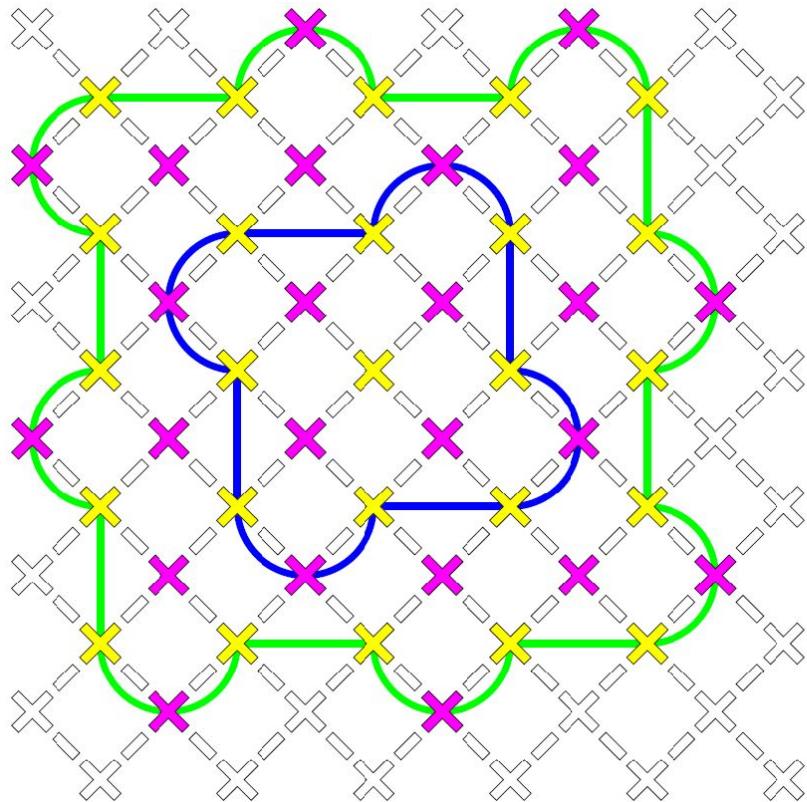


- Show  $d=5$  better than  $d=3$
- Continuous running
- Real-time decoding
- $d=34$  extreme exponential suppression  
( $d=11$  arXiv:2102.06132)

Not used  
 Data qubit  
 Measure qubit



## Challenges



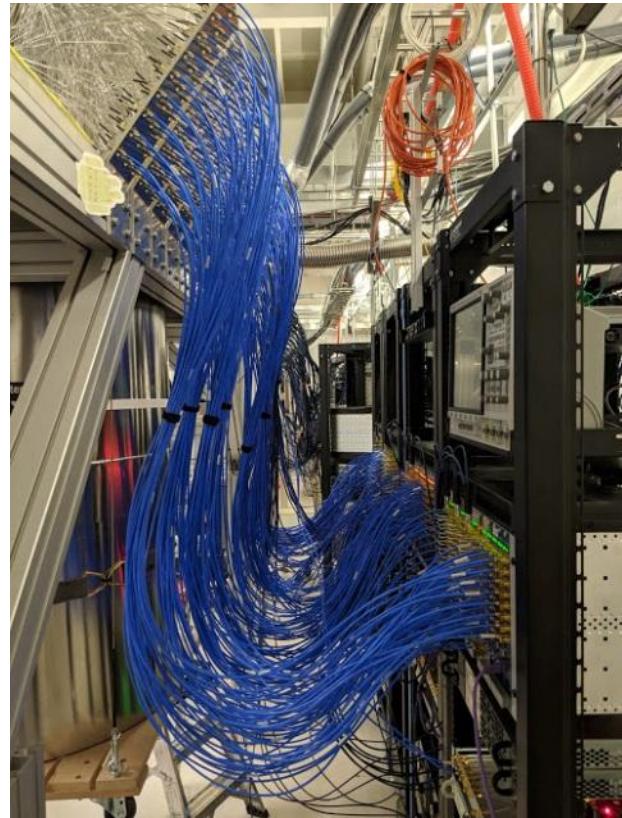
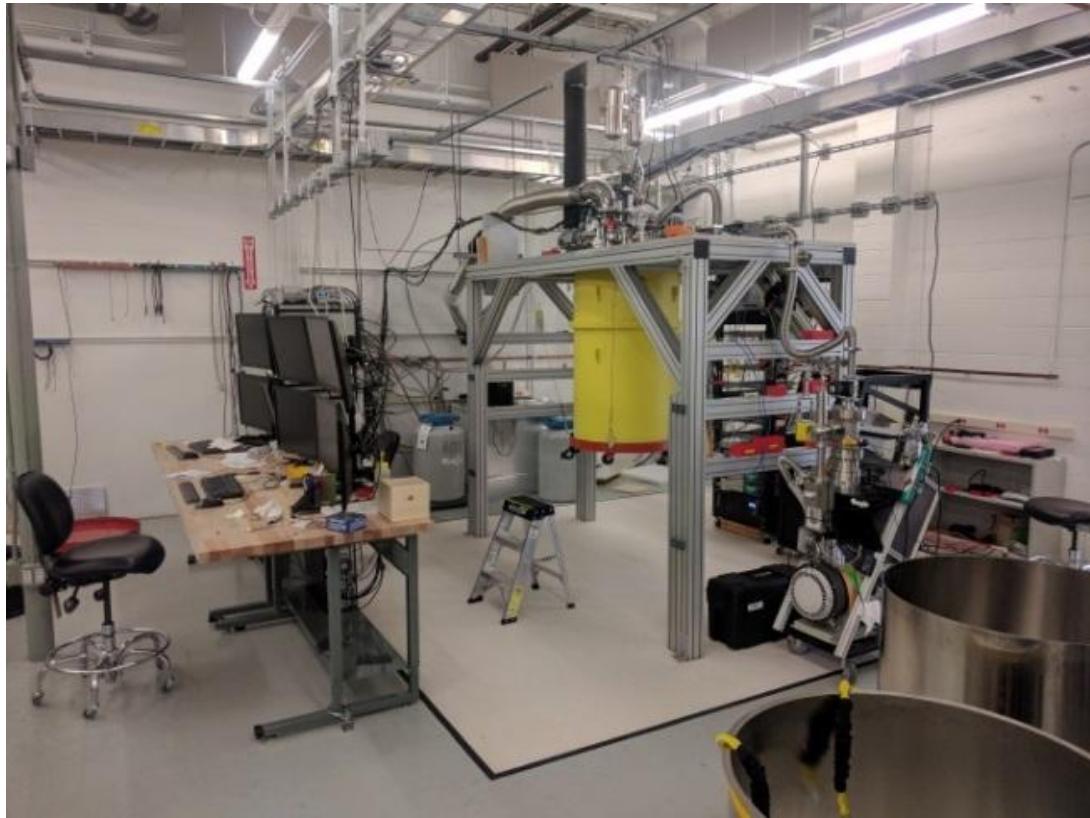
- Yield < 100%
- Coherence ~20us
- Readout and reset ~1us  
(arXiv:2102.06131)
- Calibration 0.5% (arXiv:1907.02510)
- Cosmic rays (arXiv:2104.05219)
- Decoding (arXiv:1202.5602)
- Target 1.5 to 10x suppression

Not used

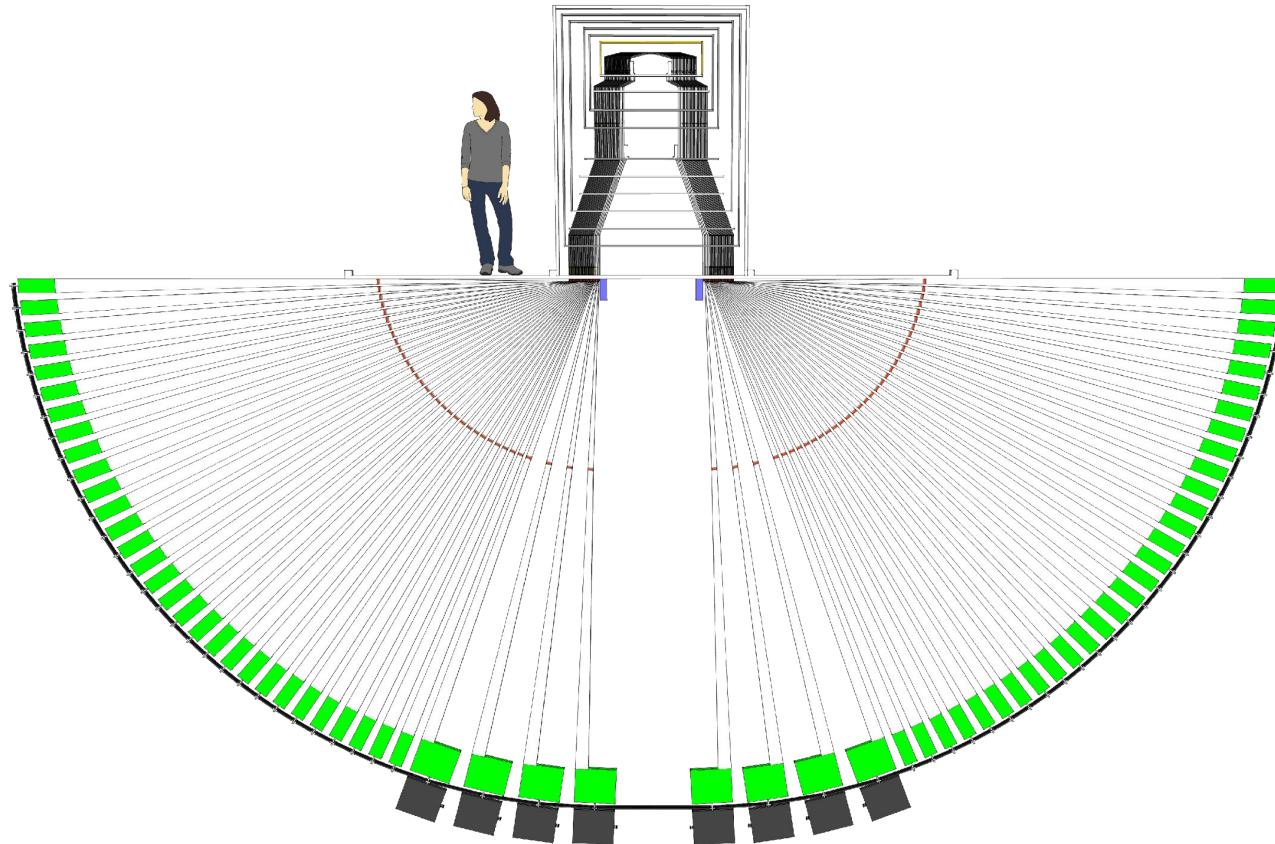
Data qubit

Measure qubit

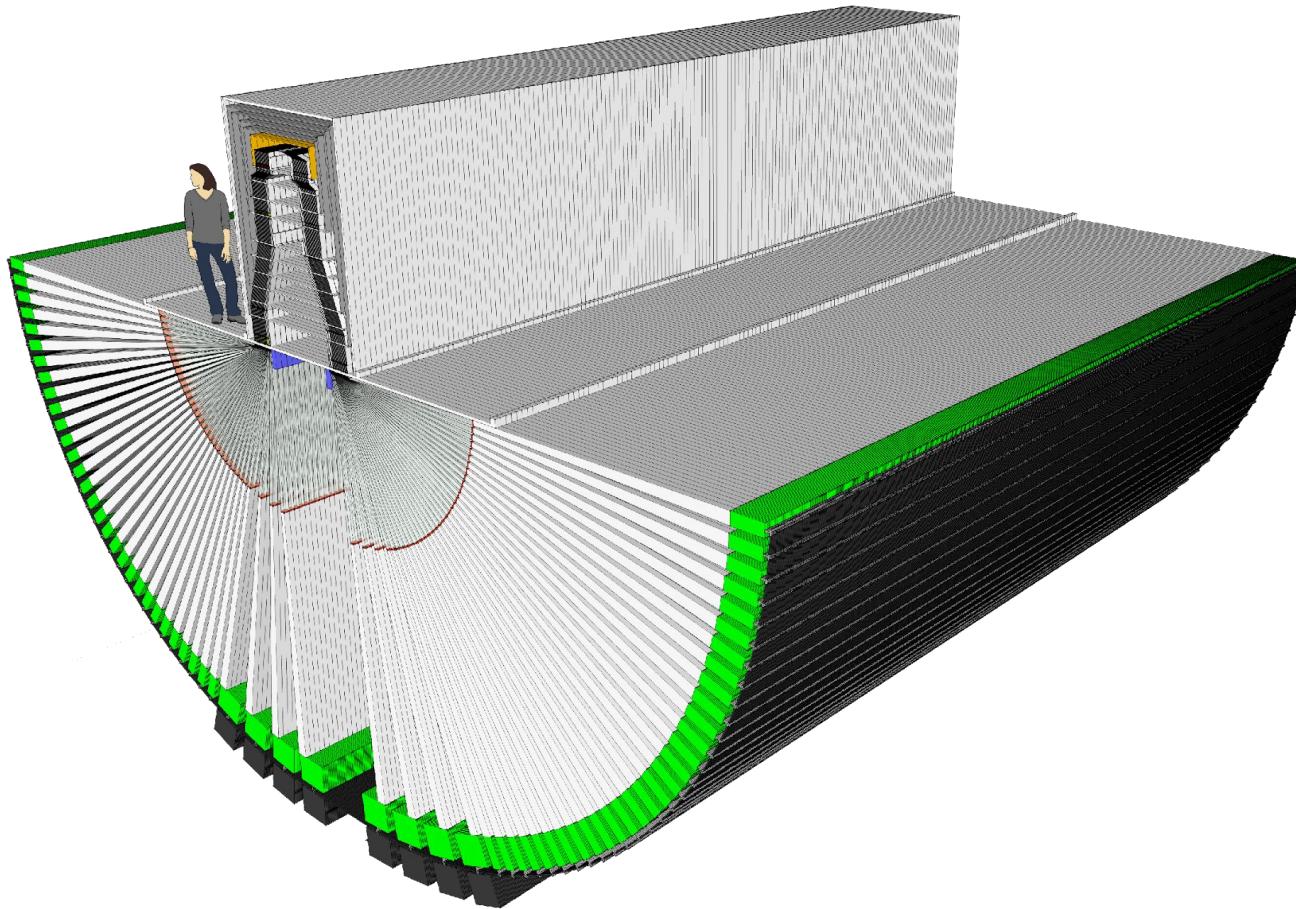
**Scalable?**



**10k scalable qubits**



**1M qubits**



# **Appendix: Compiling**

# Graph states

$$H|0\rangle = |+\rangle$$

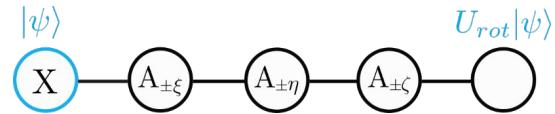
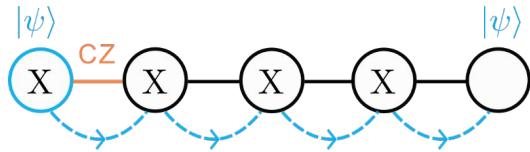
The diagram illustrates the equivalence between a quantum circuit and a graph state. On the left, a circuit with three horizontal lines (qubits) is shown. The top line has a vertical line segment connecting its midpoint to a dot. The middle line has a square box labeled 'Z' followed by a vertical line segment connecting its midpoint to a dot. The bottom line has a square box labeled 'Z'. These three dots are connected by a single vertical line. This circuit is followed by an equals sign. To the right of the equals sign is a graph with three circular nodes labeled 1, 2, and 3. Node 1 is at the top, node 2 is at the bottom left, and node 3 is at the bottom right. Edges connect node 2 to node 1 and node 1 to node 3. Another equals sign follows. To the right of the graph are three terms stacked vertically:  $X Z Z$ ,  $Z X I$ , and  $Z I X$ . These terms represent the local Clifford operations corresponding to the Z gates in the circuit.

- Graph state with nodes  $i$  and edge neighbourhood  $n_i$

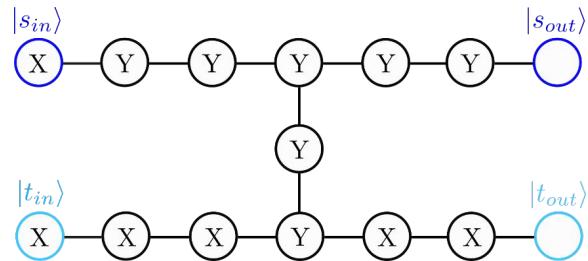
$$s_i = X_i \prod_{j \in n_i} Z_j$$

- Every stabilizer state can be converted to a graph state using only local Clifford operations

# Teleportation based QC



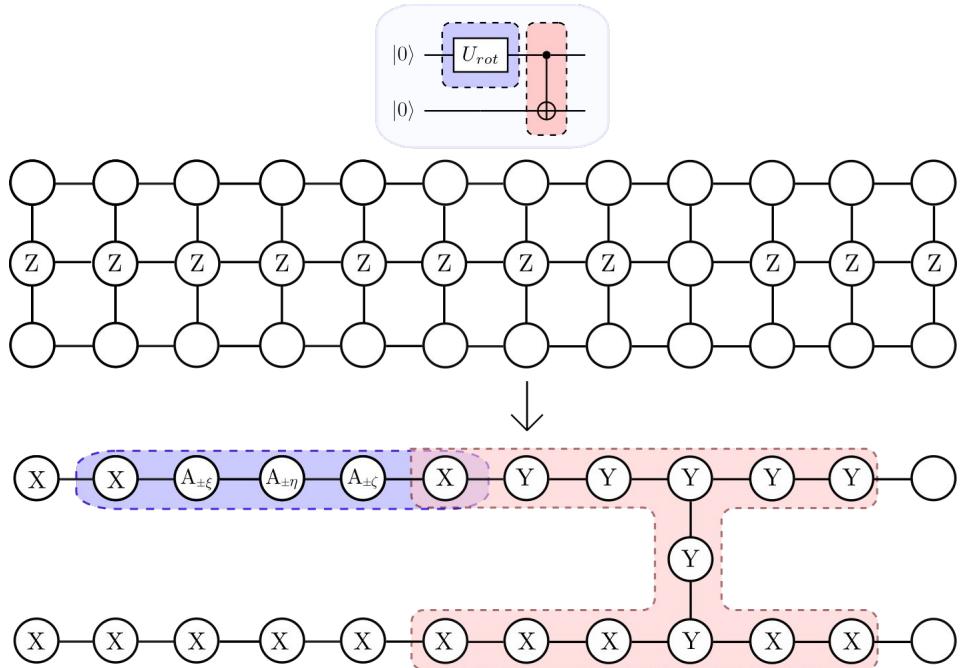
$$U_{rot} = e^{-i\frac{\xi}{2}Z} e^{-i\frac{\eta}{2}X} e^{-i\frac{\zeta}{2}Z}$$



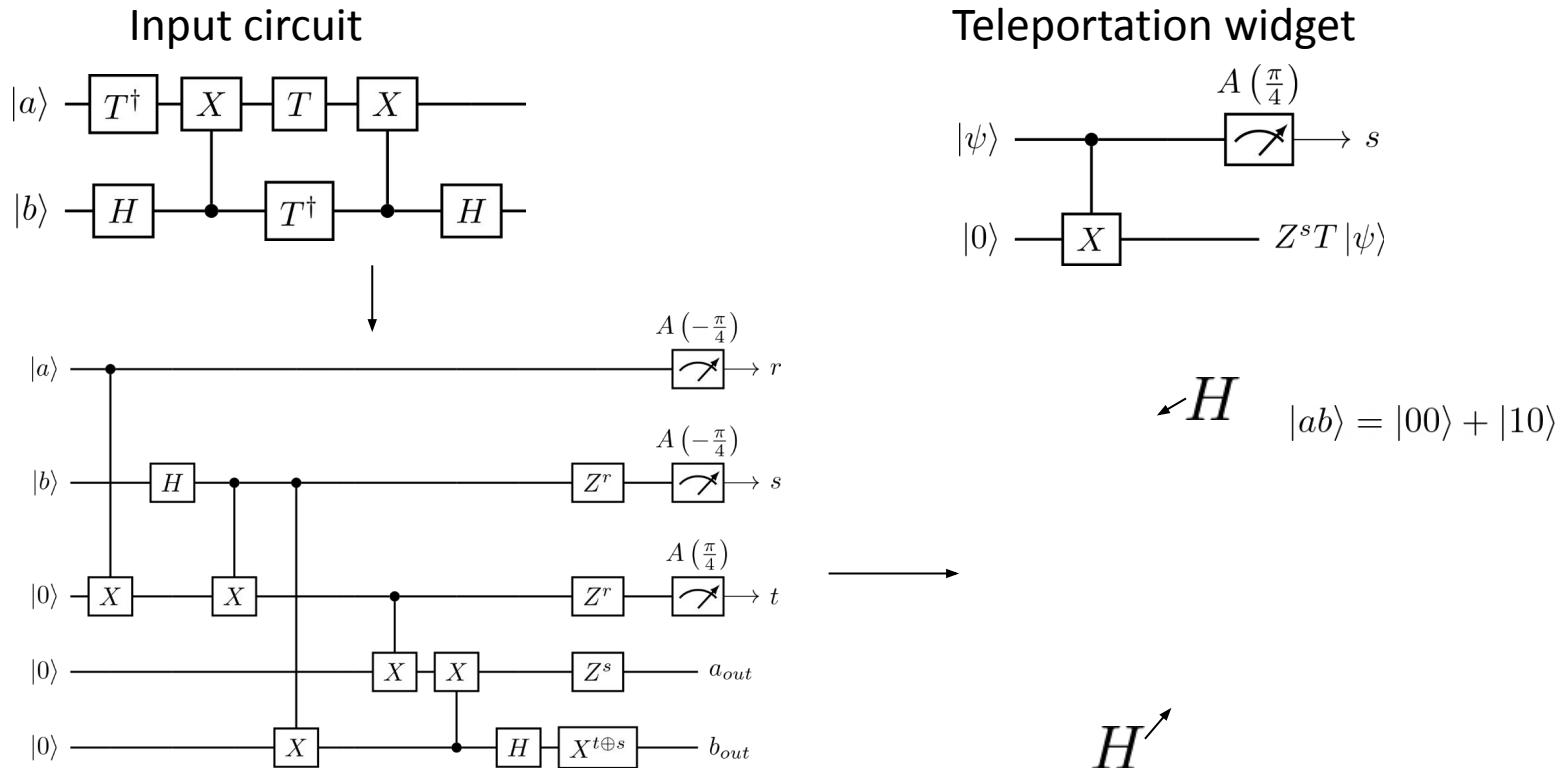
$$|s_{out}\rangle|t_{out}\rangle = CNOT|s_{in}\rangle|t_{in}\rangle$$

# Traditional circuit compilation

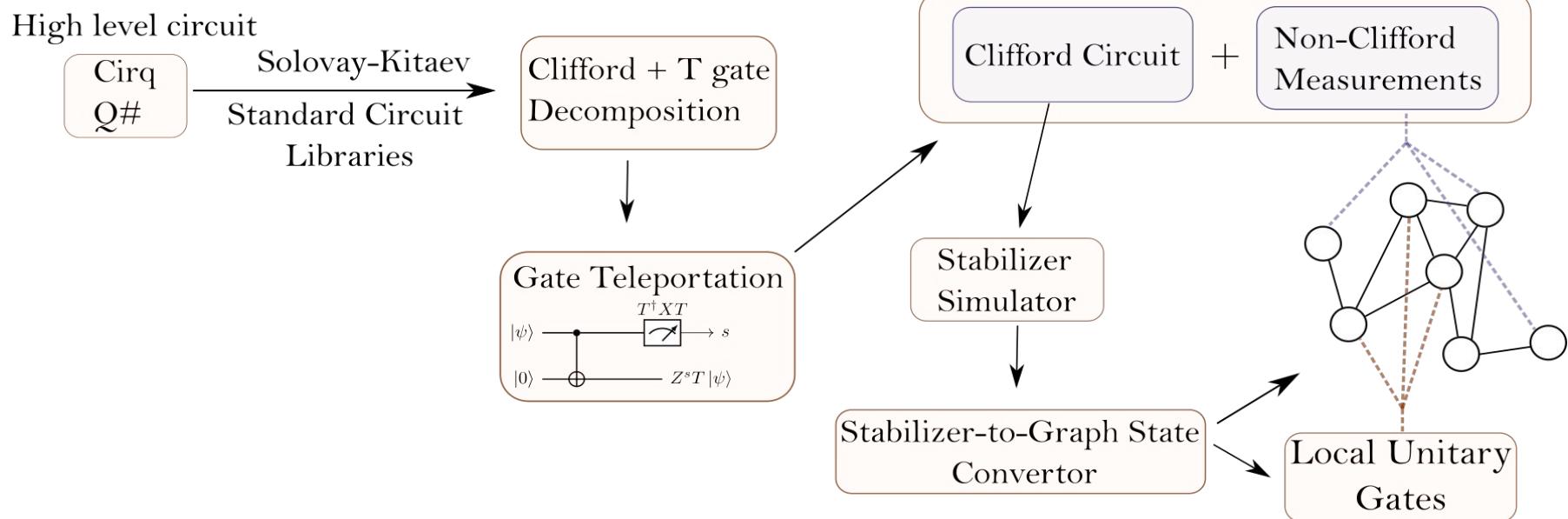
- **Shortcomings of traditional approach**
- Performs Clifford operations on a stabilizer state which can be simulated efficiently on a classical device
- No room for optimisation other than at the circuit level.



# Example of Algorithm specific compilation

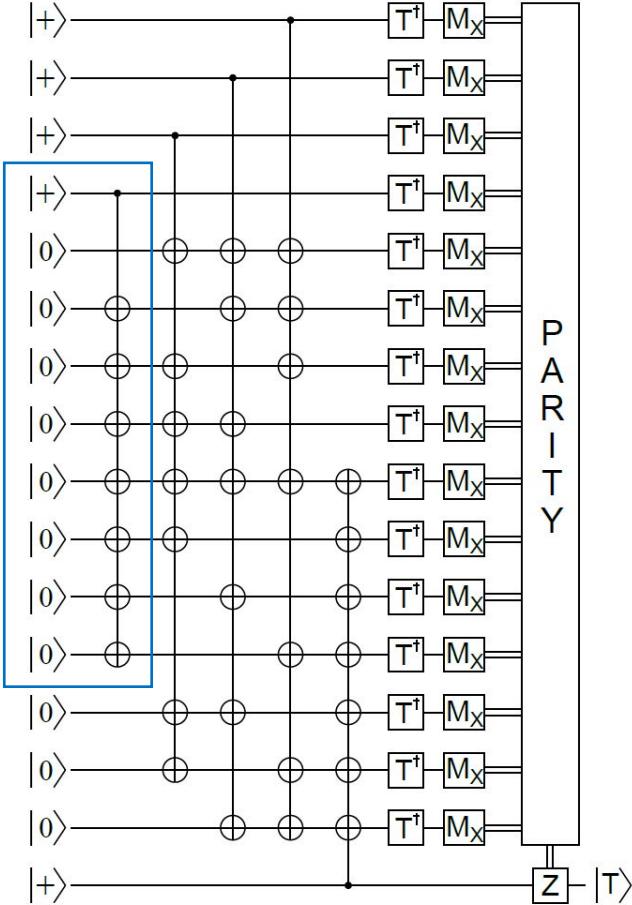


# Algorithm specific graph compiler



Jabalizer – A Julia based graph compiler <https://github.com/QSI-BAQS/Jabalizer.jl>

# State distillation



$$|\psi\rangle = \frac{1}{\sqrt{2}}(|00\dots0\rangle + |11\dots1\rangle)$$

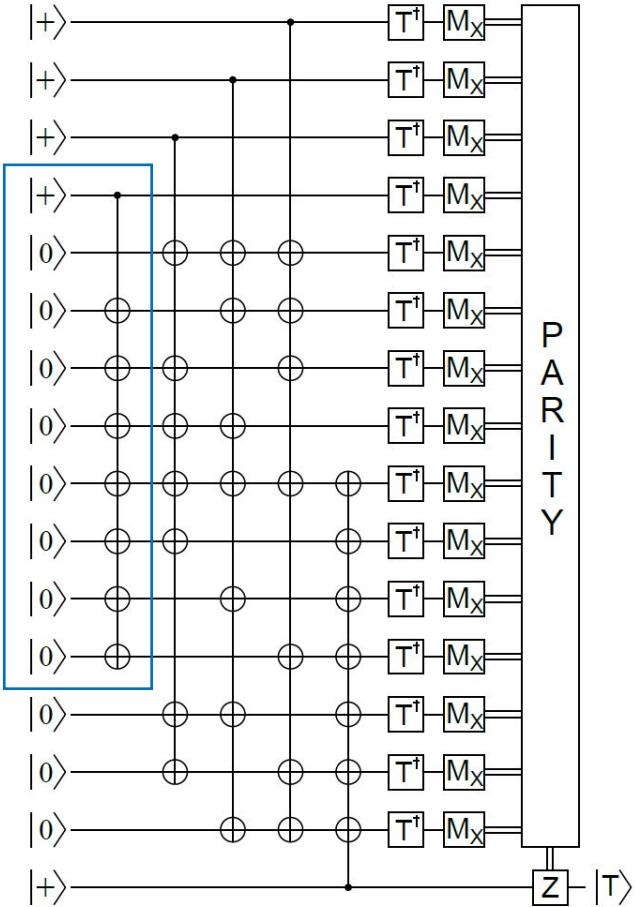
$$XX\dots X|\psi\rangle = |\psi\rangle$$

$$|\phi\rangle = \frac{1}{\sqrt{2}}(|00\dots0\rangle - |11\dots1\rangle)$$

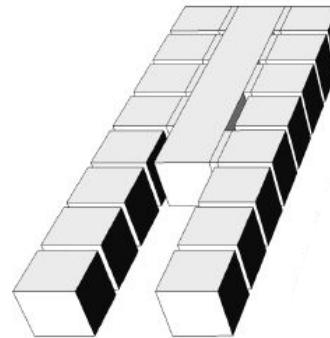
$$XX\dots X|\phi\rangle = -|\phi\rangle$$

$$|00\dots0\rangle = \frac{1}{\sqrt{2}}(|\psi\rangle + |\phi\rangle)$$

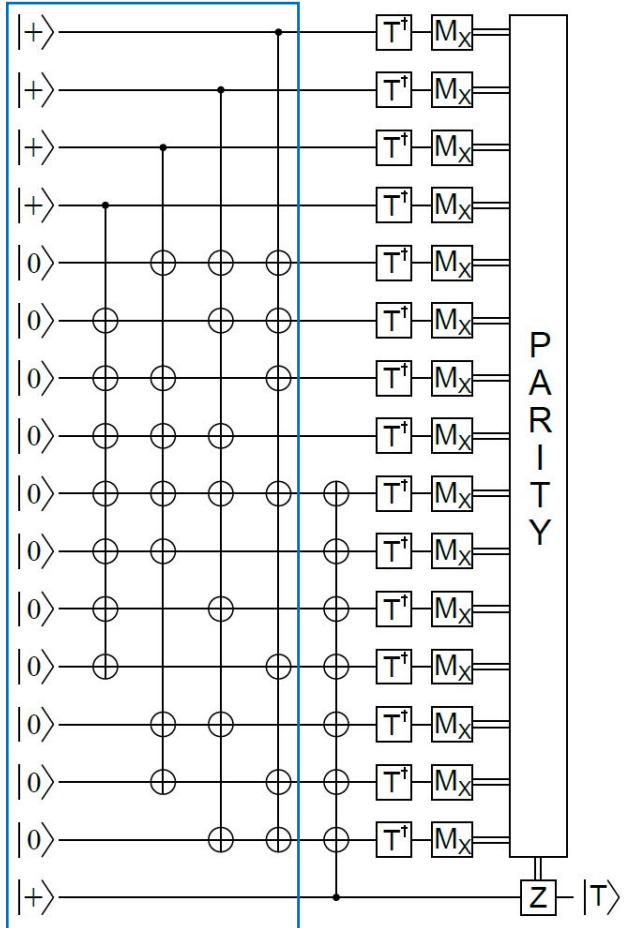
# State distillation



Measure multi-body X operator!



# State distillation



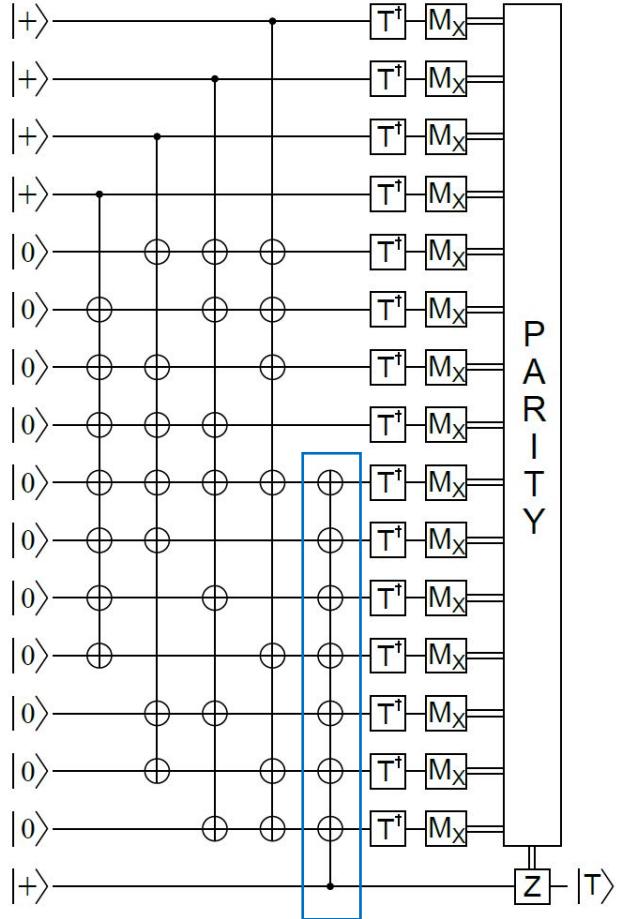
$$\frac{1}{\sqrt{2}} |0_L\rangle(|0\rangle + |1\rangle)$$

$$\frac{1}{\sqrt{2}} (|0_L 0\rangle + |1_L 1\rangle)$$

$$\frac{1}{\sqrt{2}} (|0_L 0\rangle + e^{i\pi/4} |1_L 1\rangle)$$

$$\frac{1}{\sqrt{2}} (|0\rangle + (-1)^{M_X} e^{i\pi/4} |1\rangle)$$

# State distillation



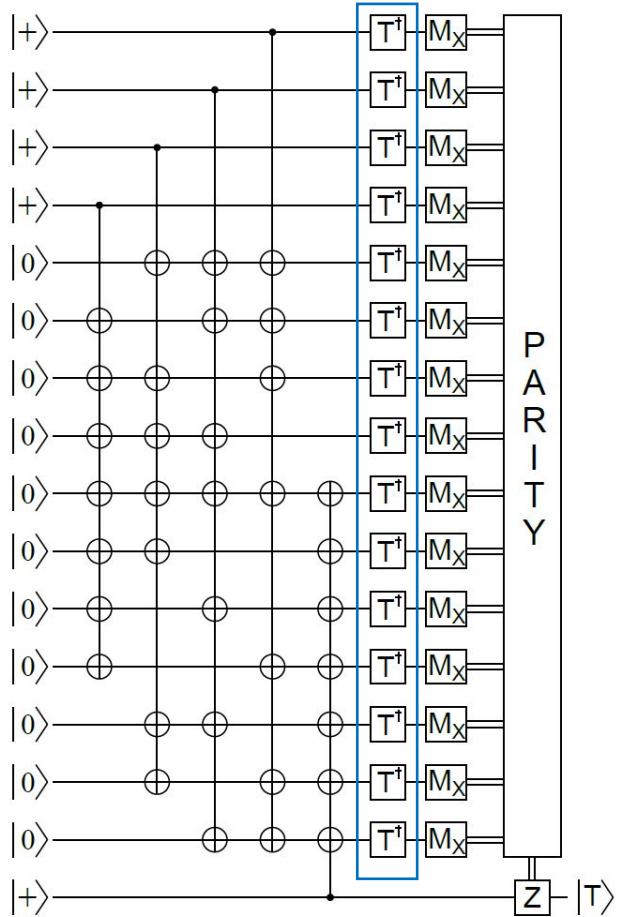
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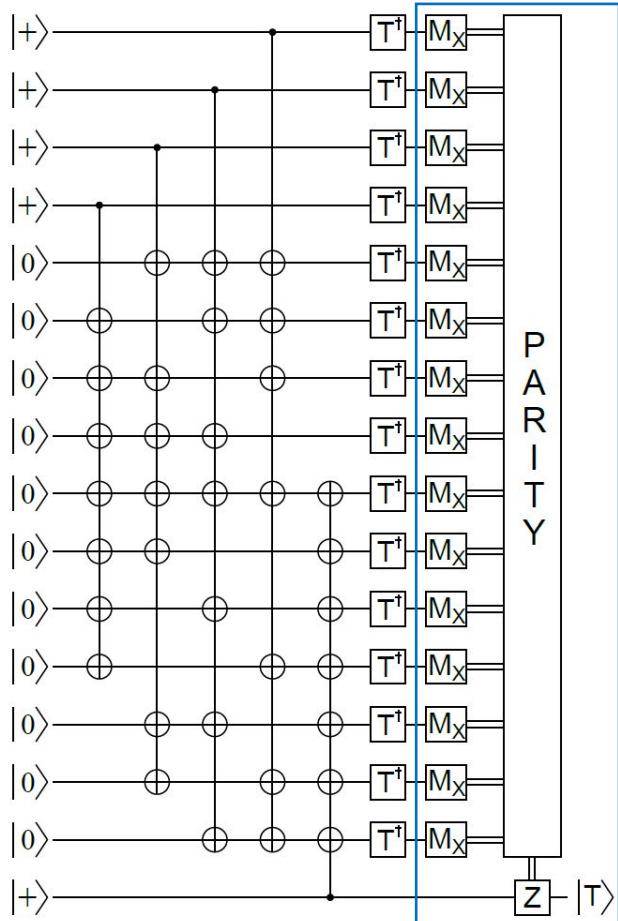
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# State distillation



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# State distillation

