

A POSITION ALLOCATION PROBLEM APPROACH TO THE BATTERY ELECTRIC BUS CHARGING PROBLEM

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Outline

Introduction

The Position Allocation Problem Approach With Linear Battery Dynamic

The Simulated Annealing Approach With Linear Battery Dynamics

TODO Summary and Questions

Introduction

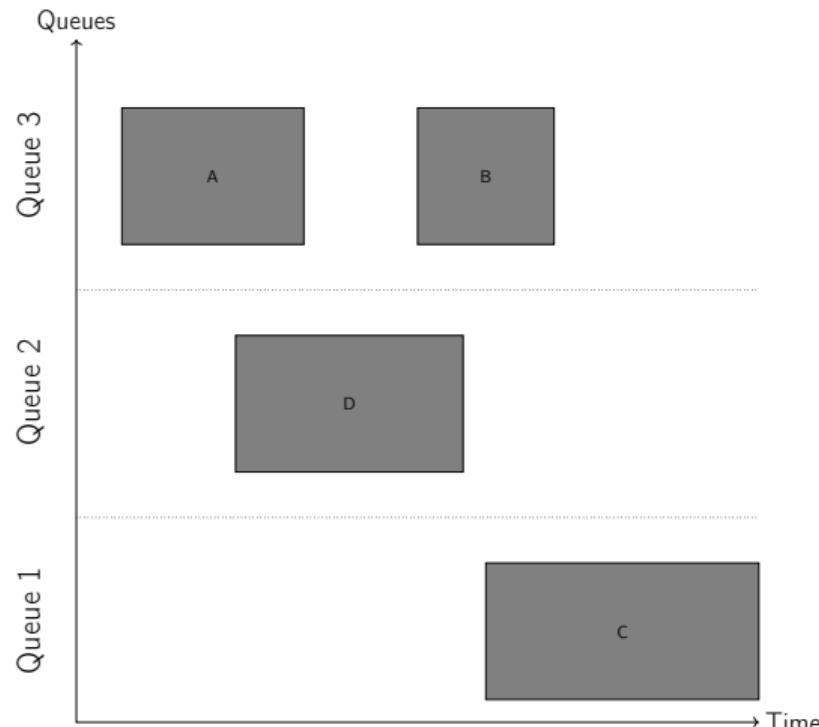
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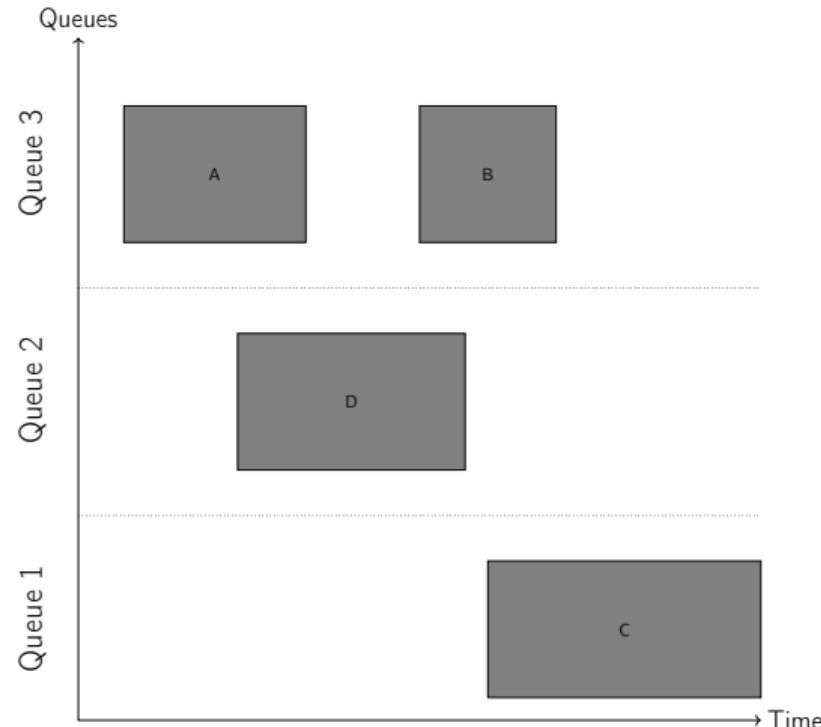
Problem Description

-	-	6:11 AM	6:16 AM	6:25 AM
6:30 AM	6:34 AM	6:41 AM	6:46 AM	6:55 AM
7:00 AM	7:04 AM	7:11 AM	7:16 AM	7:25 AM
7:30 AM	7:34 AM	7:41 AM	7:46 AM	7:55 AM
8:00 AM	8:04 AM	8:11 AM	8:16 AM	8:25 AM
8:30 AM	8:34 AM	8:41 AM	8:46 AM	8:55 AM
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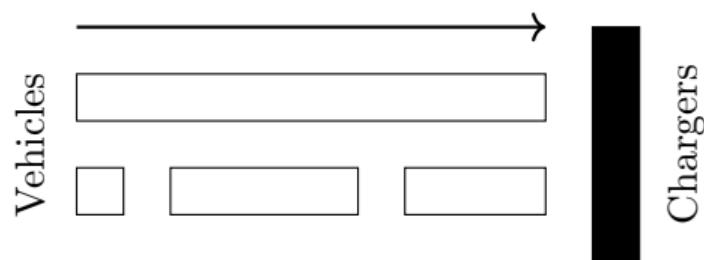
Overall Objective

- ▶ Create an operationally feasible schedule quickly
- ▶ Minimize cost
 - ▶ Demand cost
 - ▶ Consumption cost
- ▶ Battery dynamics
 - ▶ Linear
 - ▶ Non-linear
- ▶ Consider battery health
- ▶ Allow partial charging
- ▶ Multiple charger types
- ▶ Minimize total amount of chargers



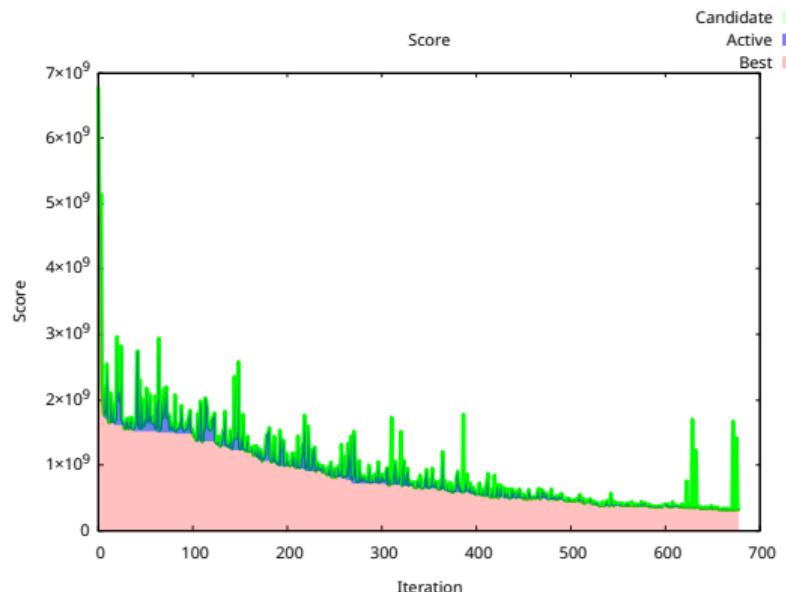
BEB Implementation of the PAP (BPAP)

- ▶ Minimize cost
 - ▶ Consumption cost
- ▶ Battery dynamics
 - ▶ Linear
- ▶ Consider battery health
- ▶ Allow partial charging
- ▶ Multiple charger types
- ▶ Minimize total amount of chargers



SA PAP with Linear Battery Dynamics

- ▶ Create an operationally feasible schedule quickly
- ▶ Minimize cost
 - ▶ Demand cost
 - ▶ Consumption cost
- ▶ Linear battery dynamics
- ▶ Consider battery health
- ▶ Allow partial charging
- ▶ Multiple charger types



Brief State Of The Art

Ref	Consumption	Demand	Linear	Non-linear	Charger Min.	Battery Health
[1]	✓		✓			
[2]	✓	✓	✓			
[3]	✓	✓	✓	✓		
[4]	✓	✓	✓			
[5]	✓	✓	✓			
[6]	✓		✓			
[7]	✓			✓	✓	✓
[8]	✓		✓			
[9]	✓		✓		✓	
BPAP	✓		✓		✓	✓
SA BPAP	✓	✓	✓	✓	*	✓

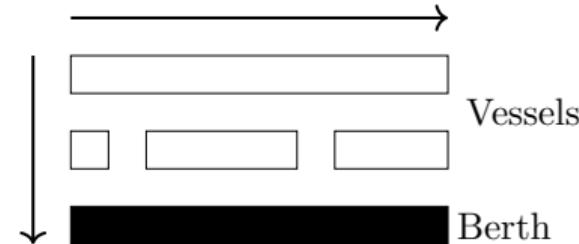
The Berth Allocation Problem¹



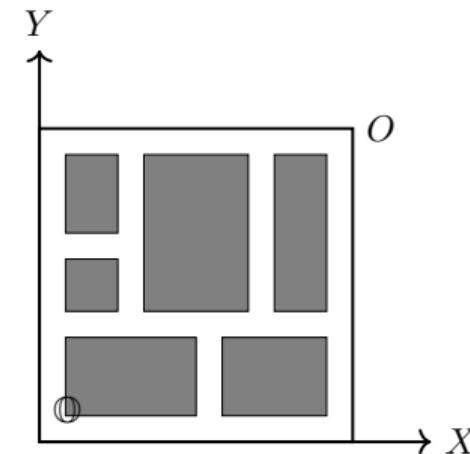
¹<https://www.mdpi.com/2077-1312/11/7/1280>

The Berth Allocation Problem

- ▶ Vessels move down toward the quay
- ▶ Receive service
- ▶ Exit to the right



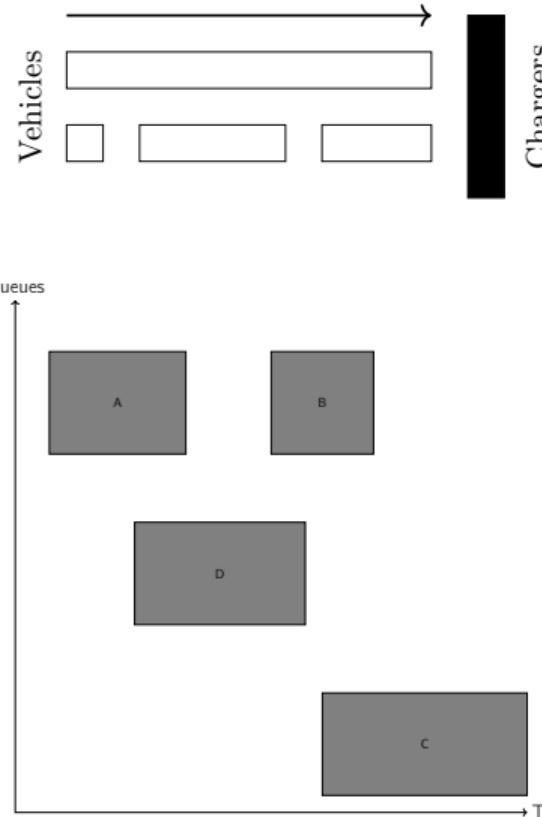
-
- ▶ A variant of the rectangle packing problem
 - ▶ Solves the problem of optimally assigning incoming vessels to be serviced



The Position Allocation Problem

PAP Behavior

- ▶ Service flow is left to right
 - ▶ Single charger type
 - ▶ All arrivals are considered unique
 - ▶ Service times are assumed to be known
-



Desired Behavior

- ▶ Discrete charger queues
- ▶ Multiple charger types
- ▶ Bus can have multiple visits
- ▶ Propagate battery SOC across visits

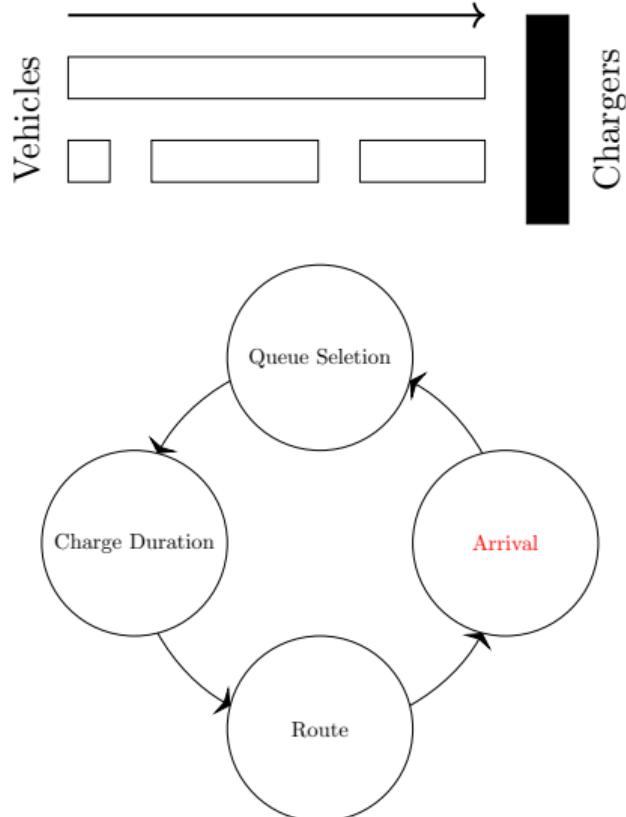
Introduction

The Position Allocation Problem Approach With Linear Battery Dynamic

The Simulated Annealing Approach With Linear Battery Dynamics

TODO Summary and Questions

Requirements For BEB Implementation



- ▶ Charges must be able to be tracked
- ▶ Service time is unknown
- ▶ Accommodate different charger rates
- ▶ Minimize charger count
- ▶ Minimize consumption cost
- ▶ Encourage slow charger use for battery health

Mixed Integer Linear Program



$$\min \sum_{i=1}^{n_V} \sum_{q=1}^{n_Q} (w_{iq} m_q + g_{iq} \epsilon_q)$$

$$u_j - u_i - s_i - (\sigma_{ij} - 1)T \geq 0$$

$$v_j - v_i - (\psi_{ij} - 1)n_Q \geq 1$$

$$\sigma_{ij} + \sigma_{ji} \leq 1$$

$$\psi_{ij} + \psi_{ji} \leq 1$$

$$\sigma_{ij} + \sigma_{ji} + \psi_{ij} + \psi_{ji} \geq 1$$

$$\sum_{q=1}^{n_Q} w_{iq} = 1$$

$$v_i = \sum_{q=1}^{n_Q} q w_{iq}$$

$$s_i + u_i = d_i$$

$$a_i \leq u_i \leq (T - s_i)$$

$$d_i \leq \tau_i$$

$$\eta_i + \sum_{q=1}^{n_Q} g_{iq} r_q - \Delta_i = \eta_{\gamma_i}$$

$$\eta_i + \sum_{q=1}^{n_Q} g_{iq} r_q - \Delta_i \geq \nu_{\Gamma_i} \kappa_{\Gamma_i}$$

$$\eta_i + \sum_{q=1}^{n_Q} g_{iq} r_q \leq \kappa_{\Gamma_i}$$

$$\eta_{\Gamma_b^0} = \alpha_{\Gamma_i} \kappa_{\Gamma_i}$$

$$\eta_{\Gamma_b^f} \geq \beta_{\Gamma_b} \kappa_{\Gamma_b}$$

Queuing Constraints



$$u_j - u_i - s_i - (\sigma_{ij} - 1)T \geq 0$$

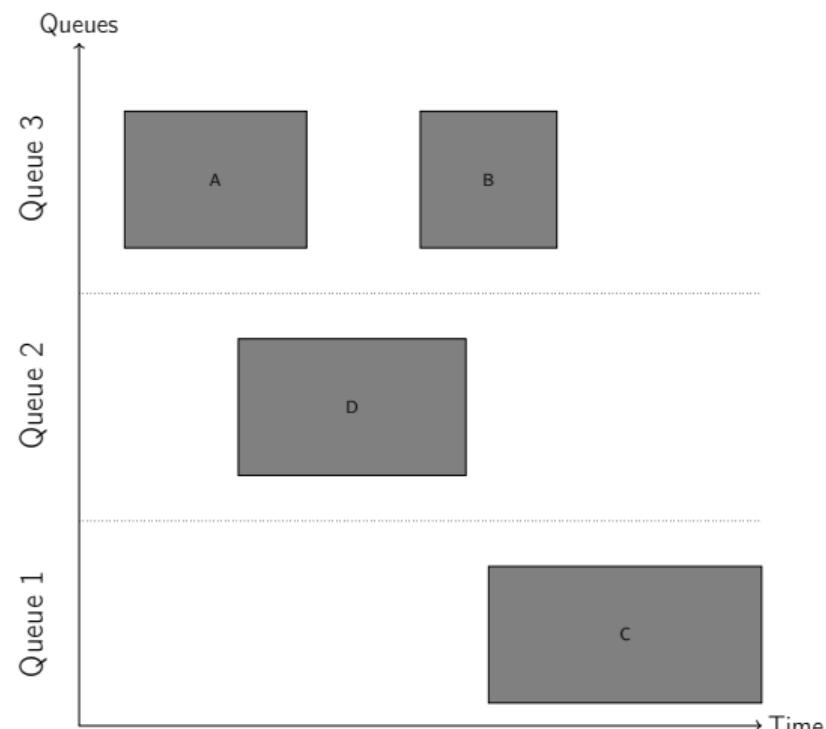
$$v_j - v_i - (\psi_{ij} - 1)n_Q \geq 1$$

$$\sigma_{ij} + \sigma_{ji} \leq 1$$

$$\psi_{ij} + \psi_{ji} \leq 1$$

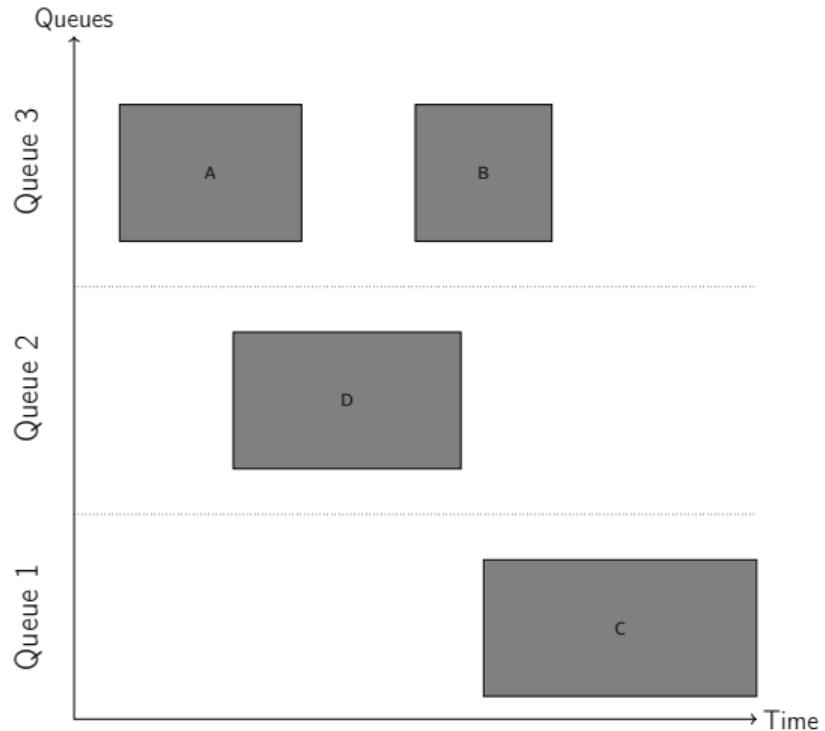
$$\sigma_{ij} + \sigma_{ji} + \psi_{ij} + \psi_{ji} \geq 1$$

- ▶ Used to ensure that gray rectangles do not overlap
- ▶ σ_{ij} establishes temporal ordering when active
- ▶ ψ_{ij} establishes spacial ordering when active

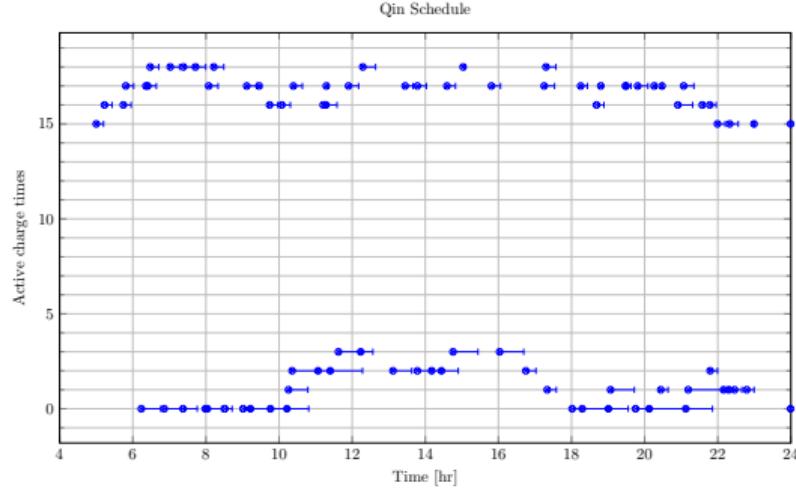
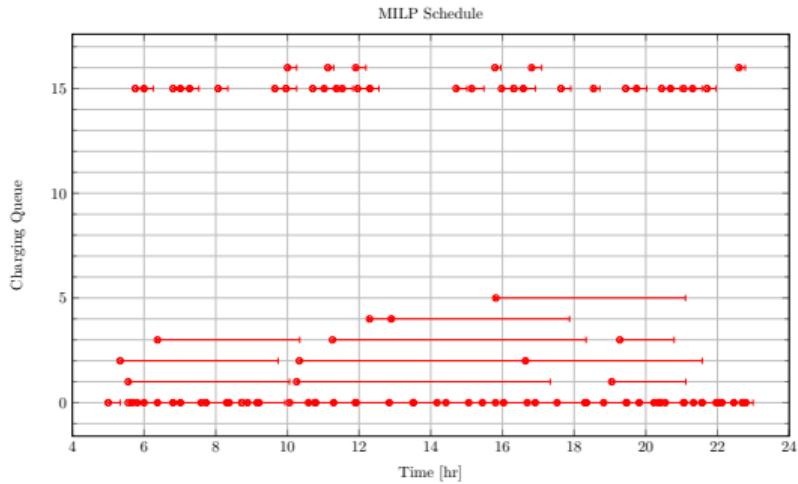


Parameters

Parameter	Value
Execution	2 hr
T	24 hr
n_V	338
n_A	35
α_i	90%
ν_i	25%
β_i	70%
m_q	1000q



Schedules



- ▶ Slow chargers utilized more frequently
- ▶ Fast chargers utilized sparingly

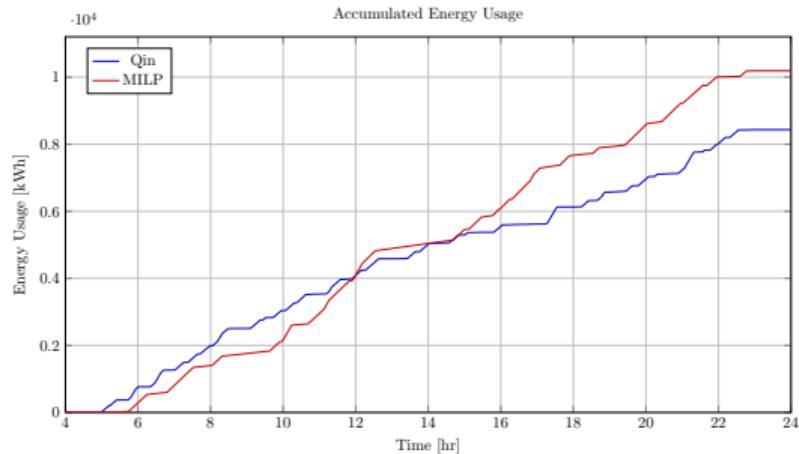
- ▶ Slow chargers used for short durations
- ▶ Extensive use of fast chargers

SOC And Energy Use



	MILP [kWh]	Qin [kWh]
Mean	265.873	355.93
Min	97.04	0.000
Max	388	368.354

- ▶ SOC of Qin allowed to drop to 0
- ▶ PAP maintained minimum SOC and final SOC



- ▶ PAP accumulated energy is larger due to minimum SOC constraints

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The Position Allocation Problem Approach With Linear Battery Dynamic

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Objective Function And Constraints



$$J(\mathbb{I}) = z_d p_d + \sum_{i=1}^{n_V} \left[\epsilon_{q_i} r_{q_i} + z_p \phi_i (\eta_i - \nu_{b_i} \kappa_{b_i}) + z_c r_{q_i} s_i \right]$$

► p_d : Demand cost

- $p_{T_p, h} = \frac{1}{T_p} \sum_{k=h-\frac{T_p}{dt}+1}^h p_k dt$
- $p_{max} = \max_{k \in [h-\frac{T_p}{dt}+1, h]} p_{T_p, h}$
- $p_d = \max(p_{fix}, p_{max})$

$$s_i = d_i - u_i$$

$$a_i \leq u_i \leq d_i \leq e_i \leq T$$

$$\eta_{\xi_i} = \eta_i + r_{q_i} s_i - \Delta_i$$

$$\kappa_{b_i} \geq \eta_i + r_{q_i} s_i$$

► $\epsilon_{q_i} r_{q_i}$: Assignment Cost

► $z_p \phi_i (\eta_i - \nu_{b_i} \kappa_{b_i})$: Penalty Function

► $z_c r_{q_i} s_i$: Consumption Cost

$$u_j - d_i - (\sigma_{ij} - 1)T \geq 0$$

$$q_j - q_i - 1 - (\psi_{ij} - 1)Q \geq 0$$

$$\sigma_{ij} + \sigma_{ji} \leq 1$$

$$\psi_{ij} + \psi_{ji} \leq 1$$

$$\sigma_{ij} + \sigma_{ji} + \psi_{ij} + \psi_{ji} \geq 1$$

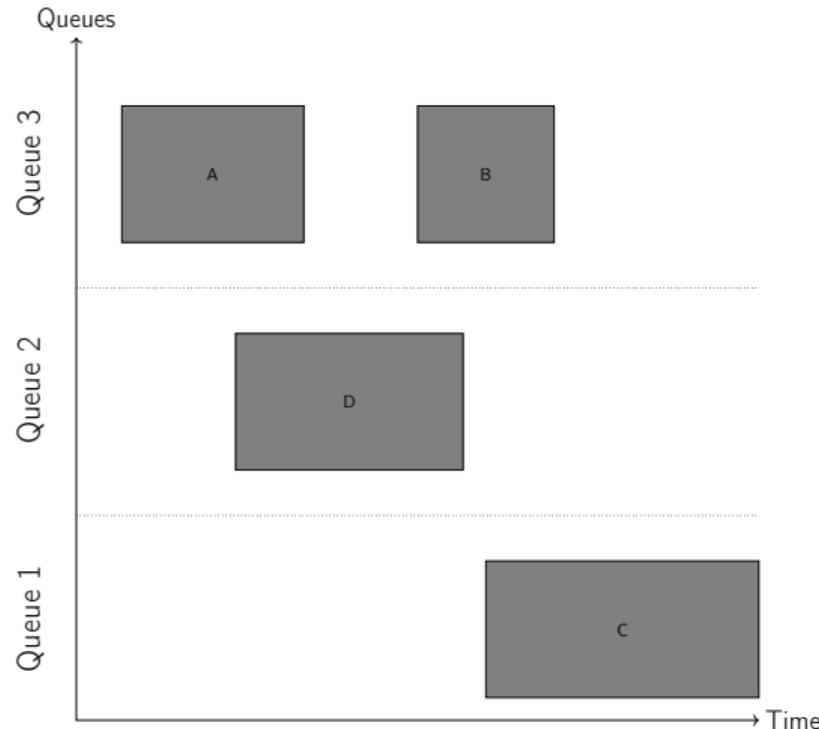
Parameters

NOTES:

- ▶ Results are different from what is presented in the thesis
- ▶ BPAP model does not take demand cost into consideration

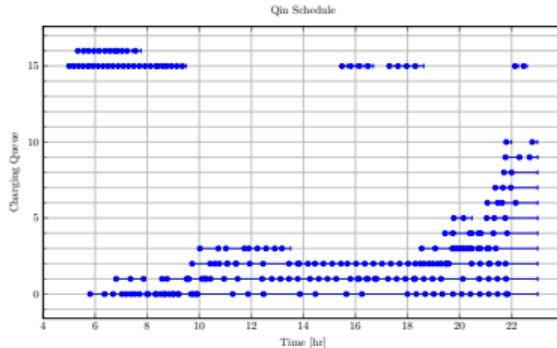
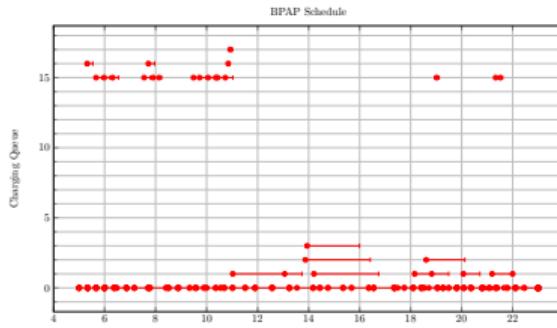
Model	Execution Time [s]
BPAP	1900
Quick	1533
Heuristic	1916

- ▶ $T_0 = 90000$
- ▶ $\beta = 0.997$
- ▶ $|t| = 3797$
- ▶ $n_K = 500$

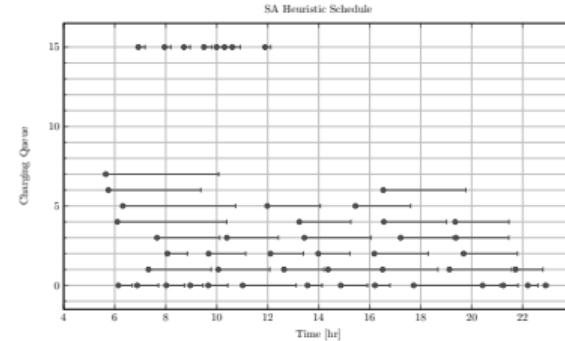
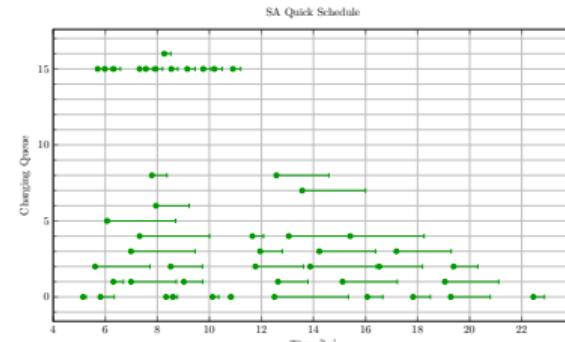


Schedule

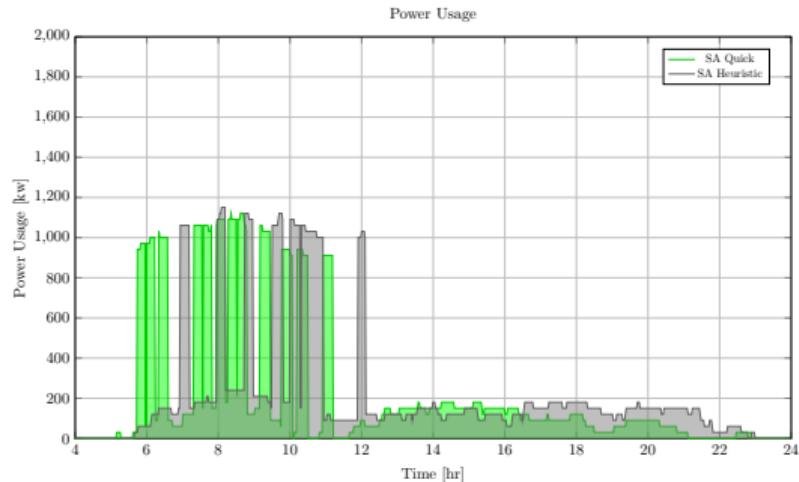
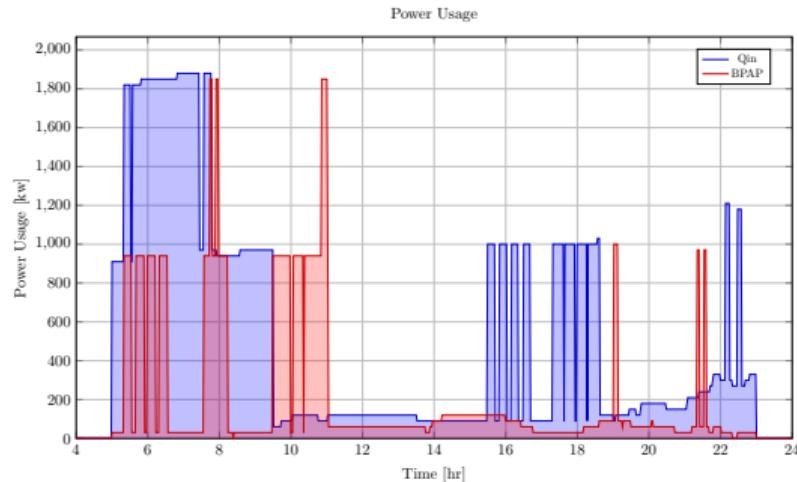
- Qin heavily utilizes fast charger
- BPAP emphasizes slow queues, but utilizes fast chargers readily



- SA techniques heavily utilize slow charging
- SA techniques utilize fast chargers more sparingly



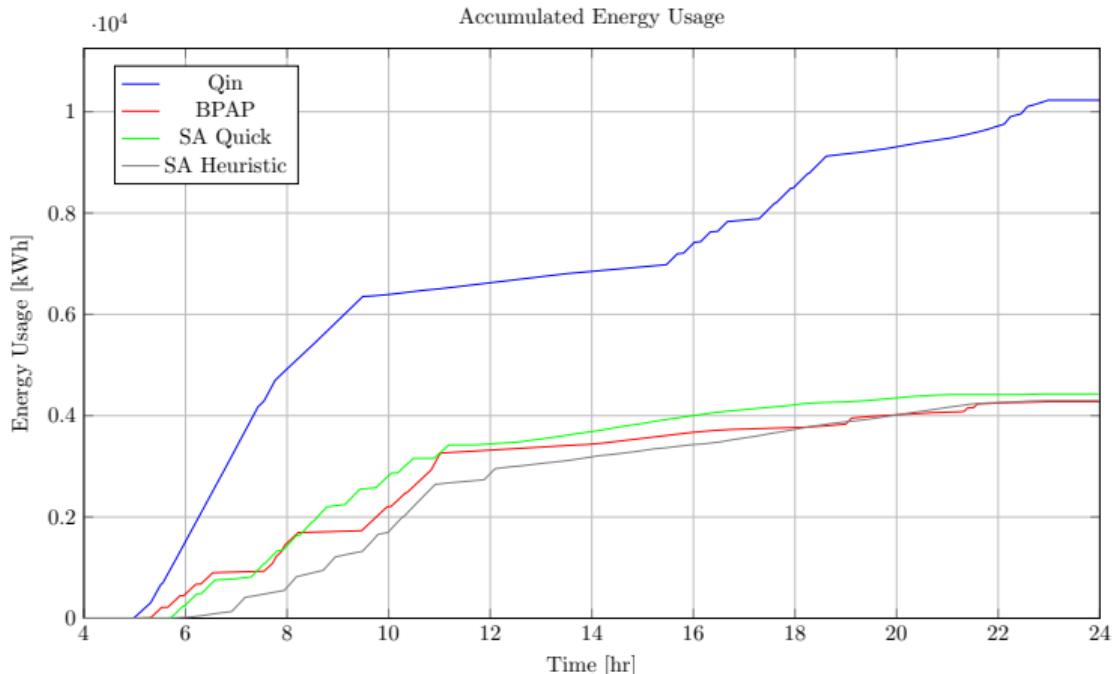
Power



- Heuristic SA maintained the lowest demand peak
- Quick SA peak demand comparable to the PAP peak

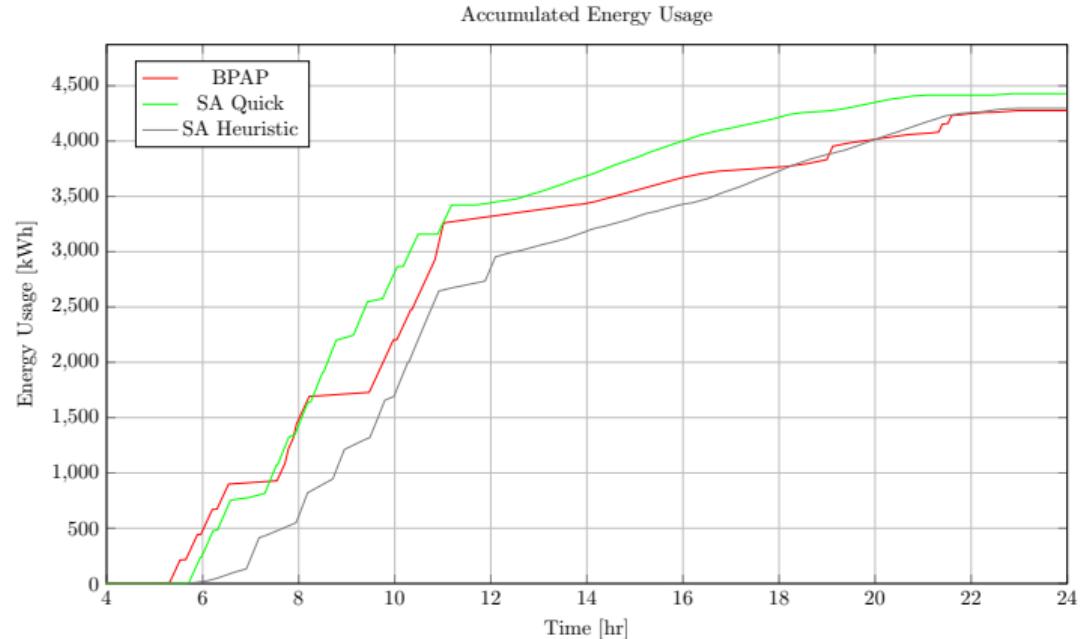
SOC and Energy

- ▶ SOC threshold deficits:
 - ▶ quick: 2.24 kWh
 - ▶ heuristic: 5.74 kWh
- ▶ Energy delta:
 - ▶ quick: 191.47 kWh
 - ▶ heuristic: 58.46 kWh
- ▶ Trade off of the lower peak demand



Scores and Energy

Schedule	Score
BPAP	18,500,000
Qin-Modified	34,578,526
Heuristic	11,673,937
Quick	11,234,577



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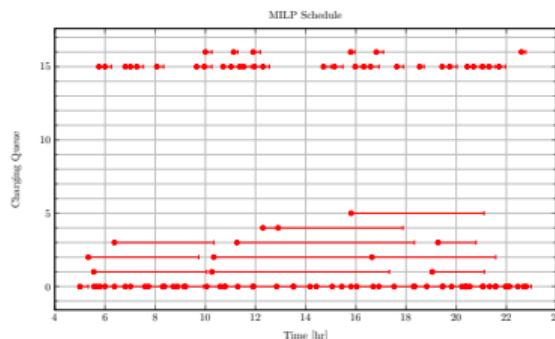
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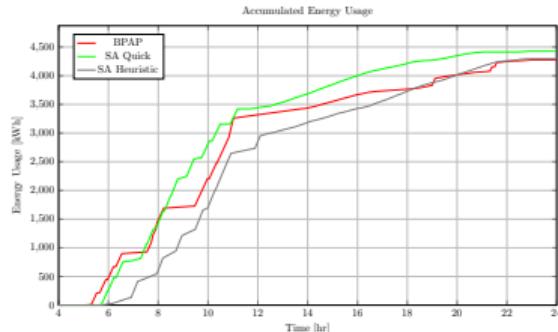
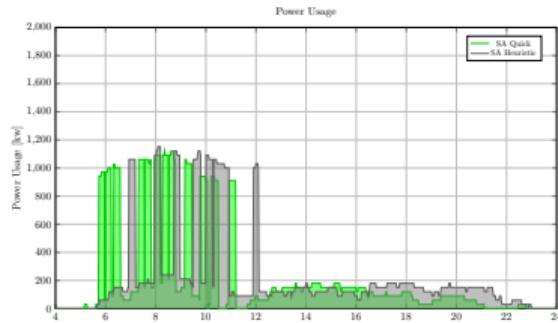
Summary

BPAP

	MILP [kWh]	Qin [kWh]
Mean	265.873	355.93
Min	97.04	0.000
Max	388	368.354



SA Techniques



Publications



College of Engineering
UtahStateUniversity

Questions?

Appendix

Mixed Integer Linear Programming



$$\max J = \sum_j c_j x_j + \sum_k d_k y_k$$

$$\text{subject to } \sum_j a_{ij} x_j + \sum_k g_{ik} y_k \leq b_i \quad (i = 1, 2, \dots, m)$$

$$x_j \geq 0 \quad (j = 1, 2, \dots, n)$$

$$y_k \in \mathbb{Z}^+ \quad (k = 1, 2, \dots, n)$$

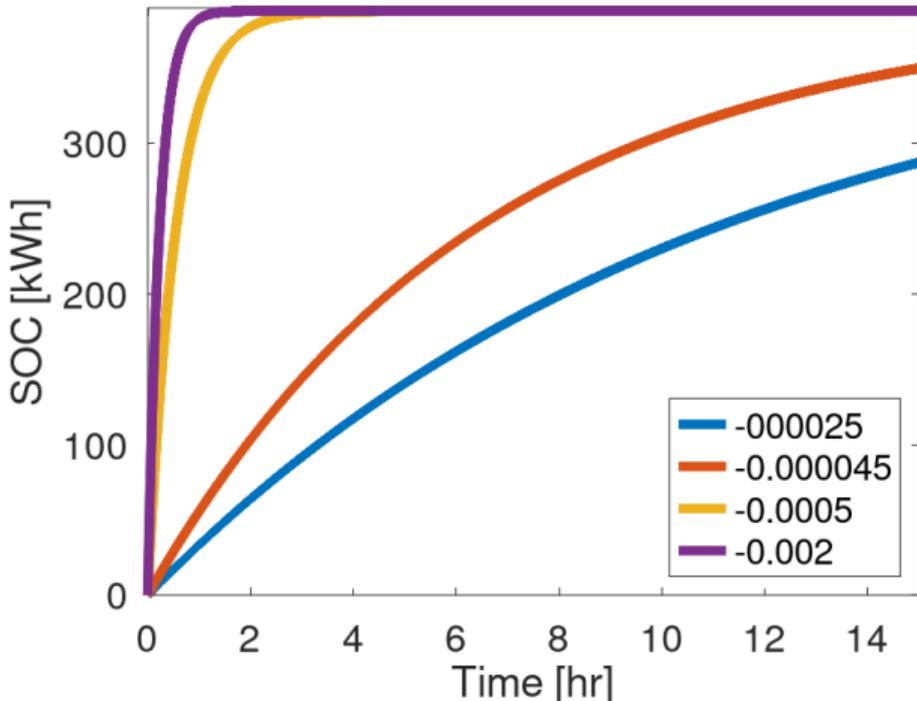
- ▶ J : Objective function
- ▶ $x_j \in \mathbb{R}$ and $y_k \in \mathbb{Z}^+$: Decision Variables
- ▶ $c_j, d_k, a_{ij}, g_{ik}, b_i \in \mathbb{R}$: Parameters

SA with Non-Linear Battery Dynamics

Introduction



- ▶ Higher fidelity in approximating charge at high SOC
- ▶ Implemented in SA for simplicity

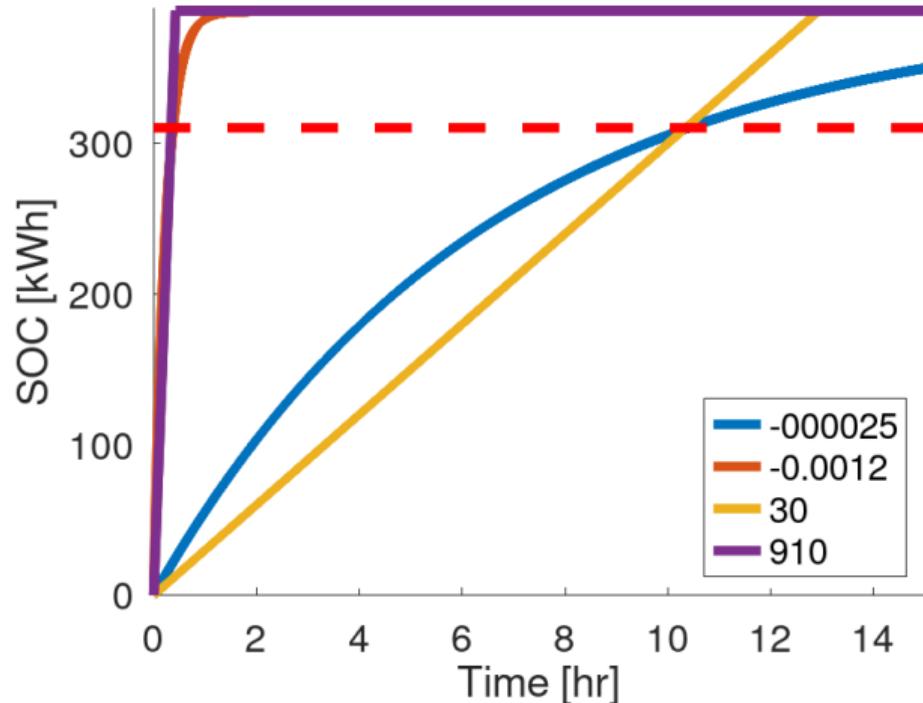


Non-Linear Battery Dynamics Model



$$\eta_{\xi_i} = \bar{a}_q \eta_i - \bar{b}_q \kappa b_i$$

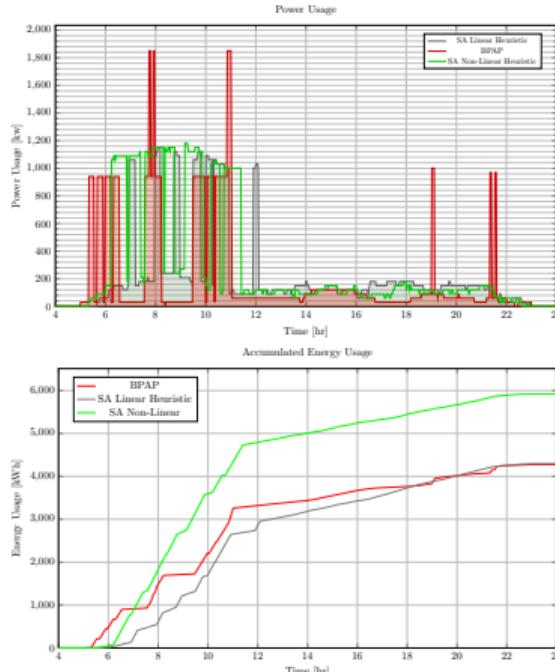
$$\bar{a}_q = e^{a_q dt} \quad \bar{b}_q = e^{a_q dt} - 1$$



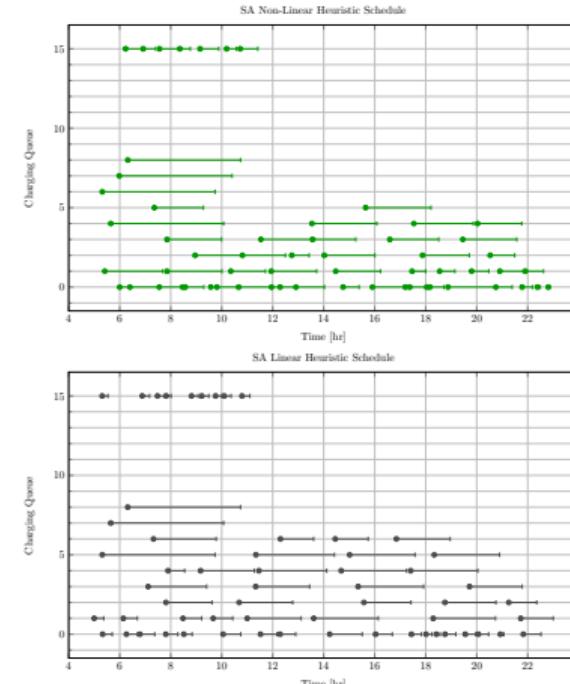
Results



- Minimized demand power
- Larger energy consumption due to fast charger duration



- Similar schedule output, longer fast charger durations due to nonlinear model
- Minimum SOC of 85 kWh



Simulated Annealing

Simulated Annealing²



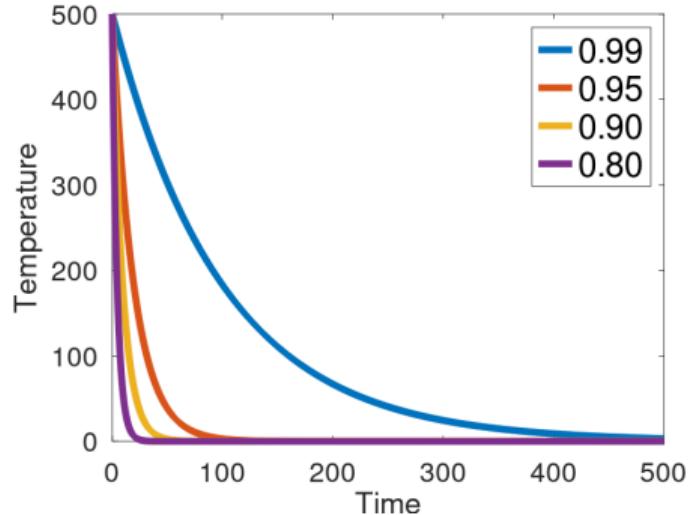
- ▶ A probabilistic technique for approximating the global optimum of a given function.
- ▶ Often applied to problems that contain many local solutions
- ▶ Three key components:
 - ▶ Cooling Schedule
 - ▶ Acceptance Criteria
 - ▶ Generation Mechanisms

²https://en.wikipedia.org/wiki/Simulated_annealing

Cooling Schedule



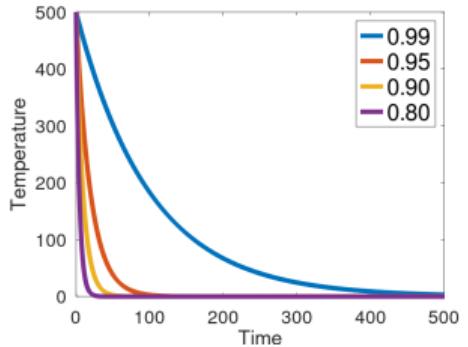
- ▶ The cooling equation models the rate at which the temperature decreases over time in the SA process.
- ▶ The temperature is high, SA encourages exploration. As the temperature decreases, exploitation is encouraged.



$$t_m = \beta t_{m-1}$$

Acceptance Criteria³

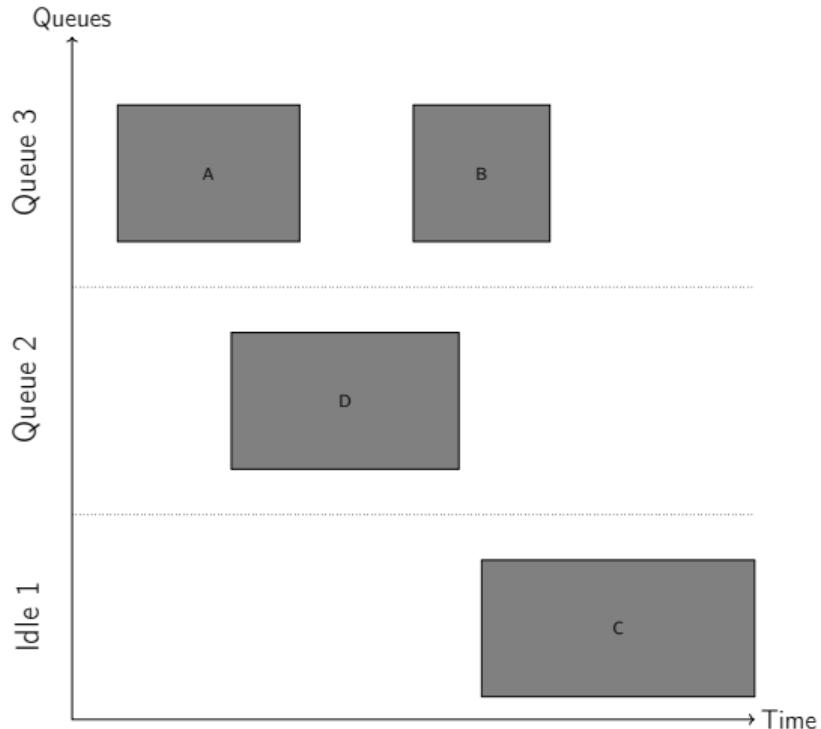
$$f(\mathbb{I}, \bar{\mathbb{I}}, t_m) = \begin{cases} 1 & \Delta E > 0 \\ e^{-\frac{\Delta E}{t_m}} & \text{otherwise} \end{cases}$$



³https://en.wikipedia.org/wiki/Simulated_annealing

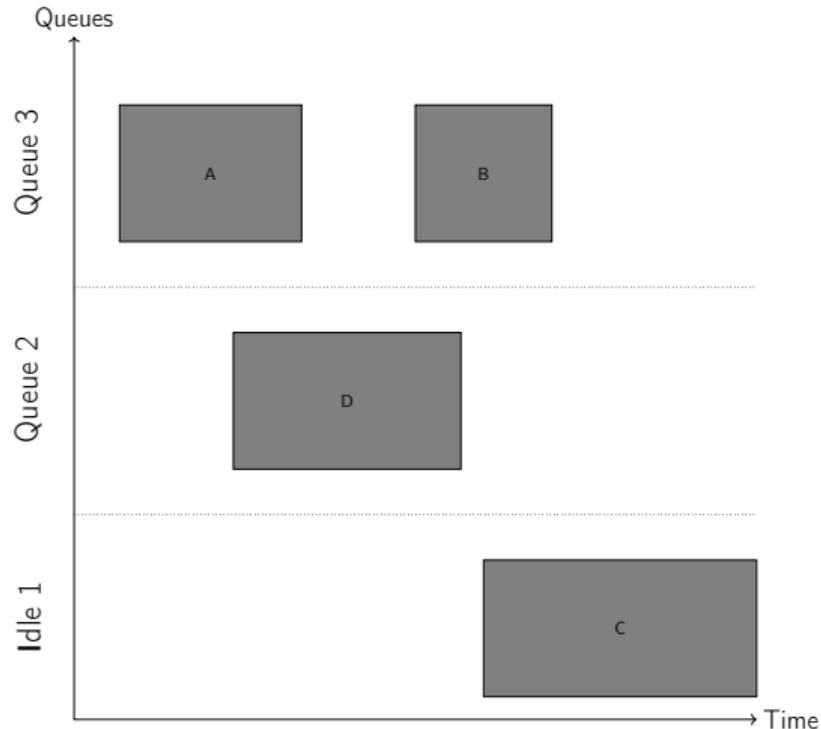
Generation Mechanisms - Primitive Functions

- ▶ New Visit: Move a bus from a wait queue to charge queue
- ▶ Slide Visit: Change the charge duration of a visit
- ▶ New Charger: Move a visit to a new charger
- ▶ Wait: Move a visit to its idle queue
- ▶ New Window: Execute Wait then New Visit primitives



Generation Mechanisms - Wrapper Functions

- ▶ Charge Schedule Generation:
Iterate through each visit and execute New Visit
- ▶ Perturb Schedule: Randomly execute one of the primitives with a weighted distribution



Algorithm 1: New visit algorithm

Algorithm: New Visit**Input:** \mathbb{S} **Output:** $\bar{\mathbb{S}}$ **1 begin**

```
2    $i \leftarrow \mathbb{S}_i;$                                 /* Extract visit index */
3    $a \leftarrow \mathbb{I}_{i.a};$                             /* Extract the arrival time for visit  $i$  */
4    $e \leftarrow \mathbb{I}_{i.e};$                             /* Extract the departure time for visit  $i$  */
5    $q \leftarrow \mathbb{I}_{i.q};$                             /* Extract the current charge queue for visit  $i$  */
6    $\bar{q} \leftarrow \mathcal{U}_Q;$                          /* Select a random charging queue with a uniform distribution */
7    $C \leftarrow \mathcal{U}_{\mathbb{C}_q};$                       /* Select a random time slice from  $\mathbb{C}_q$  */
8   if  $(\bar{C}, \bar{u}, \bar{d}) \leftarrow \text{findFreeTime}(C, i, q, a, e) \notin \emptyset$  then    /* If there is time available in  $C_q^j$  */
9     |   return  $(i, (\bar{q}, \bar{u}, \bar{d}), \bar{C})$                                 /* Return visit */
10  end
11  return  $(\emptyset);$                                 /* Return nothing */
12 end
```

Algorithm 2: Slide Visit Algorithm

Algorithm: Slide Visit

Input: \mathbb{S}

Output: $\bar{\mathbb{S}}$

```
1 begin
2    $(i, \mathbb{I}, \bar{\mathbb{C}}) \leftarrow \text{Purge}(\mathbb{S});$            /* Purge visit  $i$  from charger availability matrix */
3    $C \leftarrow \bar{\mathbb{C}}_{i,q_i};$                       /* Get the time availability of the purged visit */
4   /* If there is time available in  $C$  */
5   if  $(\bar{C}, \bar{u}, \bar{d}) \leftarrow \text{findFreeTime}(C, \mathbb{S}_i, \mathbb{I}_q, \mathbb{I}_{i.a}, \mathbb{I}_{i.e}) \notin \emptyset$  then
6     | return  $(i, \mathbb{I}, (\mathbb{I}_{i.q_i}, \bar{u}, \bar{d}), \bar{\mathbb{C}})$           /* Return updated visit */
7   end
8   return  $(\emptyset);$                                 /* Return nothing */
```

Algorithm 3: New Charger Algorithm

Algorithm: New Charger

Input: \mathbb{S}

Output: $\bar{\mathbb{S}}$

1 **begin**

```
2    $(i, \mathbb{I}, \bar{\mathbb{C}}) \leftarrow \text{Purge}(\mathbb{S});$            /* Purge visit  $i$  from charger availability matrix */
3    $q \leftarrow \mathcal{U}_Q;$            /* Select a random charging queue with a uniform distribution */
4   if  $(\bar{C}, \bar{u}, \bar{d}) \leftarrow \text{findFreeTime}(\bar{\mathbb{C}}_{i,q}, \mathbb{S}_i, \mathbb{I}_q, \mathbb{I}_{i,u}, \mathbb{I}_{i,d}) \notin \emptyset$  then /* If there is time available in
       $C_q$  */
      | /* Return visit, note  $u$  and  $d$  are the original initial/final charge times. */
      | return  $(i, \mathbb{I}, (q, \mathbb{I}_{i,u}, \mathbb{I}_{i,d}), \bar{\mathbb{C}})$ 
6   end
7   return  $(\emptyset);$            /* Return nothing */
8 end
```

Algorithm 4: Wait algorithm

Algorithm: Wait**Input:** \mathbb{S} **Output:** $\bar{\mathbb{S}}$

```
1 begin
2    $(i, \mathbb{I}, \bar{\mathbb{C}}) \leftarrow \text{Purge}(\mathbb{S});$            /* Purge visit  $i$  from charger availability matrix */
3    $\bar{\mathbb{C}}'_{\mathbb{I}_{i,r_i}} \leftarrow \mathbb{C}' \cup \{[\mathbb{I}_{i.b}, \mathbb{I}_{i.e}]\};$  /* Update the charger availability matrix for wait queue  $\bar{\mathbb{C}}_{i,q_i}$  */
4   return  $(i, \mathbb{I}, (\mathbb{I}_{i.b}, \mathbb{I}_{i.a}, \mathbb{I}_{i.e}), \bar{\mathbb{C}})$            /* Return visit */
5 end
```

Algorithm 5: New window algorithm

Algorithm: New Window

Input: \mathbb{S}

Output: $\bar{\mathbb{S}}$

```
1 begin
2   |  $\bar{\mathbb{S}} \leftarrow \text{Wait}(\mathbb{S});$                                 /* Assign visit to its respective idle queue */
3   | if  $\bar{\mathbb{S}} \leftarrow \text{NewVisit}(\bar{\mathbb{S}}) \neq \emptyset$  then          /* Add visit  $i$  back in randomly */
4   |   | return  $\bar{\mathbb{S}}$                                               /* Return visit */
5   | end
6   | return  $(\emptyset);$                                          /* Return nothing */
7 end
```

Algorithm 6: Charge schedule generation algorithm

Algorithm: Candidate Solution Generator

Input: \mathbb{S}

Output: $\bar{\mathbb{S}}$

```
1 begin
2   | /* Select an unscheduled BEB visit from a randomly indexed set of visits */ 
3   | foreach  $\mathbb{I}_i \in \mathbb{I}$  do
4   |   |  $(i, \bar{\mathbb{I}}, \bar{\mathbb{C}}) \leftarrow \text{NewVisit}((\mathbb{I}_i, \mathbb{I}, \mathbb{C}))$ ;           /* Assign the bus to a charger */
5   |   end
6   | return  $(0, \bar{\mathbb{I}}, \bar{\mathbb{C}})$ 
7 end
```

Charge Schedule Perturbation



Algorithm 7: Perturb schedule algorithm

Algorithm: Perturb Schedule

Input: \mathbb{S}

Output: $\tilde{\mathbb{S}}$

```
1 begin
2      $p \leftarrow [\text{false}; n_A]$ ;                                /* Create vector to track priority routes */
3      $y^i \leftarrow [1.0; n_V]$ ;                                /* Create weight vector for index selection */
4     /* Loop through the visits in reverse order */                */
5     foreach  $\mathbb{I}_i \leftarrow \mathbb{I}_{|\mathbb{I}|} \text{ TO } \mathbb{I}_1$  do
6         /* If the current visit is part of a priority route */      */
7         if  $p_{\mathbb{I}_i.b} = \text{true}$  then
8              $y_{\mathbb{I}_i}^i = y_{\mathbb{I}_i,\xi}^i$ ;                            /* */
9         end
10        /* Else if the current visit's SOC does below the allowed threshold */ */
11        else if  $\mathbb{I}_{i.\eta} \leq \nu_{\mathbb{I}_i.b} \kappa_{\mathbb{I}_i.b}$  then
12             $p_{\mathbb{I}_i.b} = \text{true}$ ;                                /* Indicate the current BEB's routes are to be prioritized */
13             $y_{\mathbb{I}_i}^i = \kappa_{\mathbb{I}_i.b} (\nu_{\mathbb{I}_i.b} \kappa_{\mathbb{I}_i.b} - \mathbb{I}_{i.\eta})$ ; /* Calculate the weight of the current visit */
14        end
15    end
16     $\mathbb{I}_i \leftarrow \mathcal{W}_{\mathbb{I}}^{y^i}$ ;                                /* Select an index with a weighted distribution */
17     $i \leftarrow \mathbb{I}_i$ ;                                    /* Extract visit index */
18     $y^p \leftarrow [y_1^p, y_2^p, \dots]$ ;                      /* Define the weight of each primitive generator */
19     $PGF \leftarrow \mathcal{W}_{[1, n_G]}^{y^p}$ ;                  /* Select a generator function with weighted distribution */
20     $\tilde{\mathbb{S}} \leftarrow PGF((i, \mathbb{I}, \mathbb{C}))$ ;          /* Execute the generator function */
21    return  $(0, \mathbb{I}, \tilde{\mathbb{C}})$ 
22 end
```

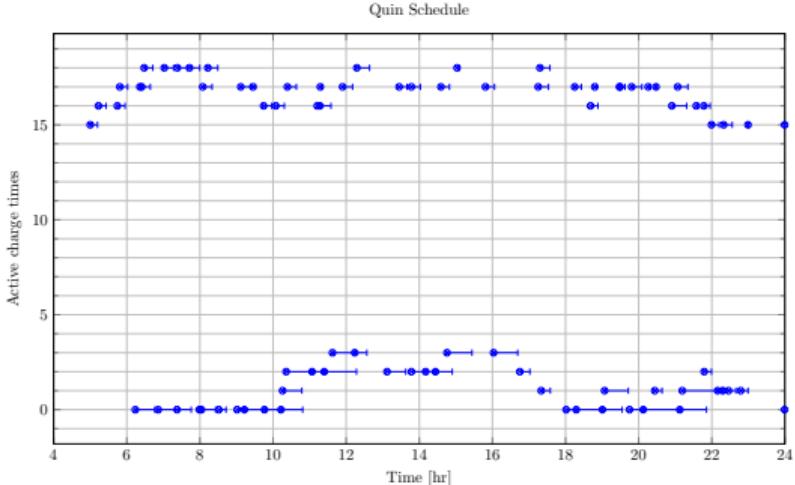
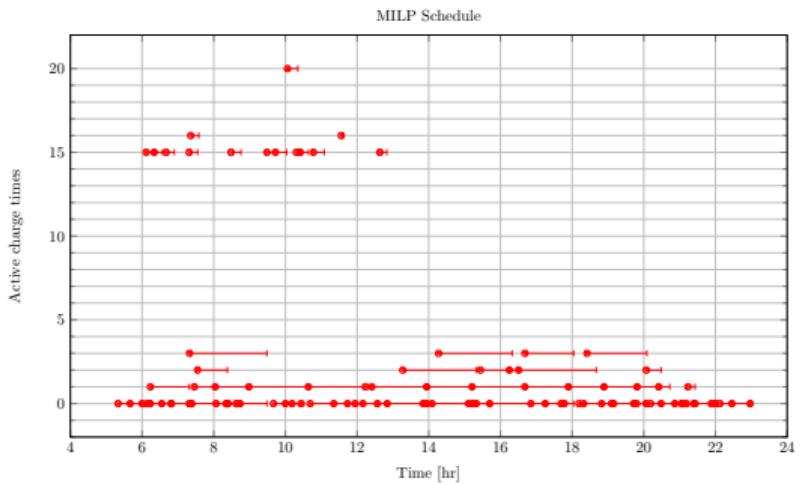
Thesis Results

Parameters

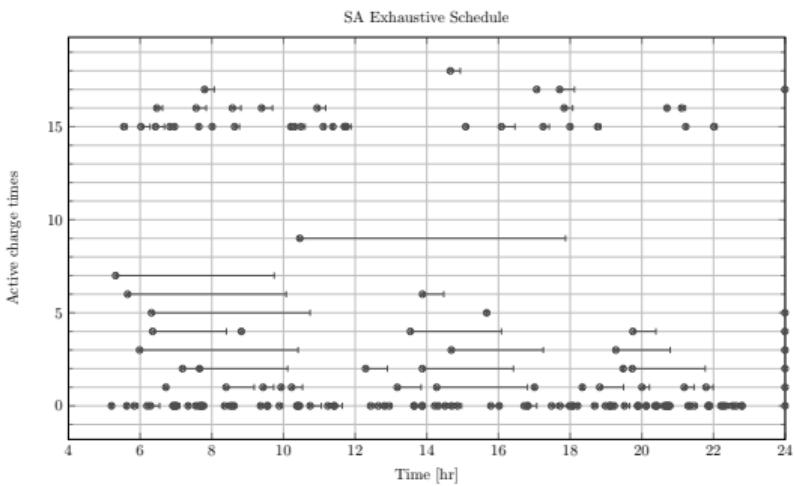
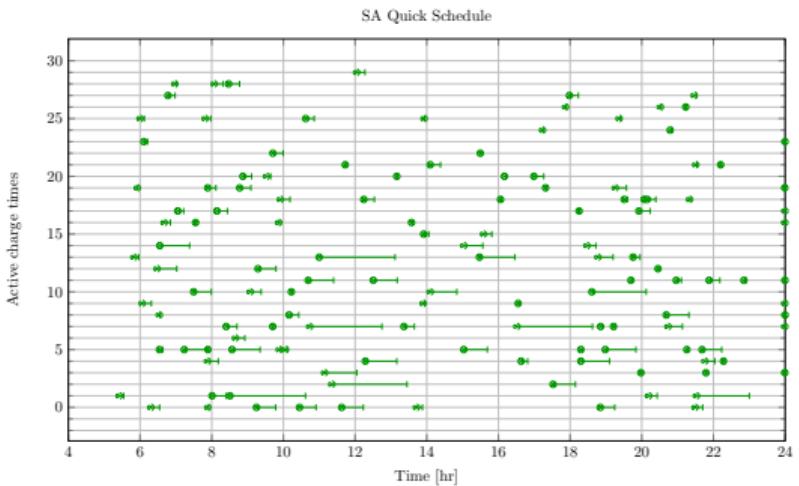
Model	Execution Time [s]	Iteration [s]
MILP	3600	N/A
Quick	2275.25	0.25
Heuristic	3640.4	0.4

- $T_0 = 99999$
- $\beta = 0.999$
- $|t| = 3797$
- $n_K = 500$

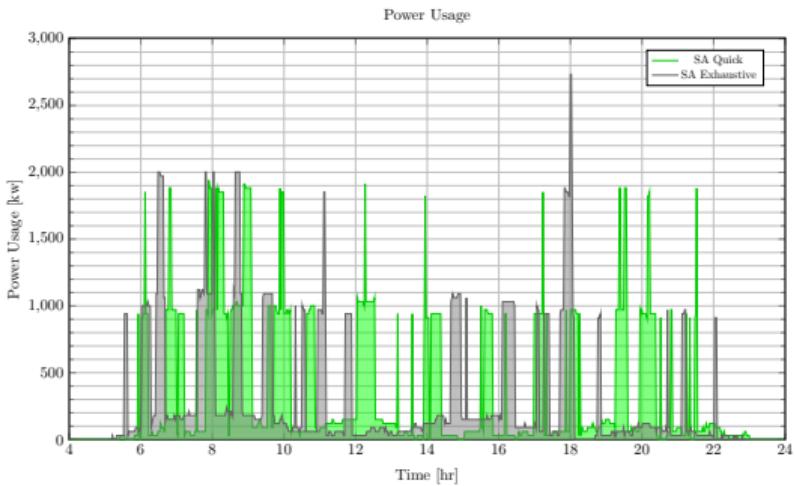
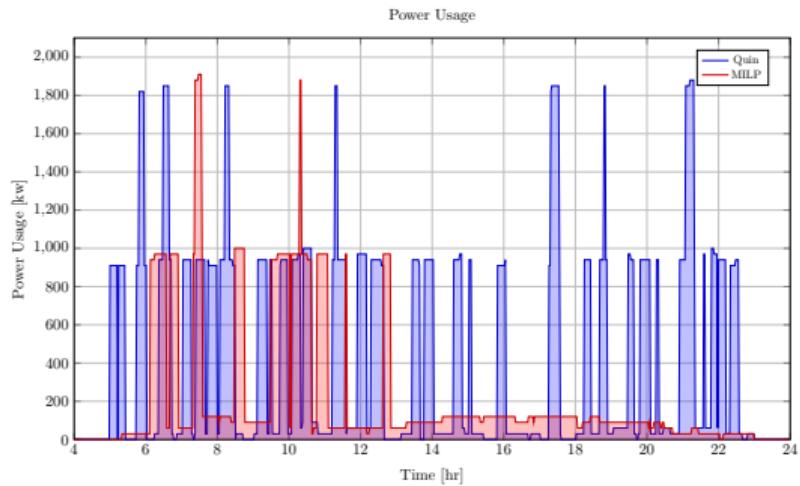
Schedule



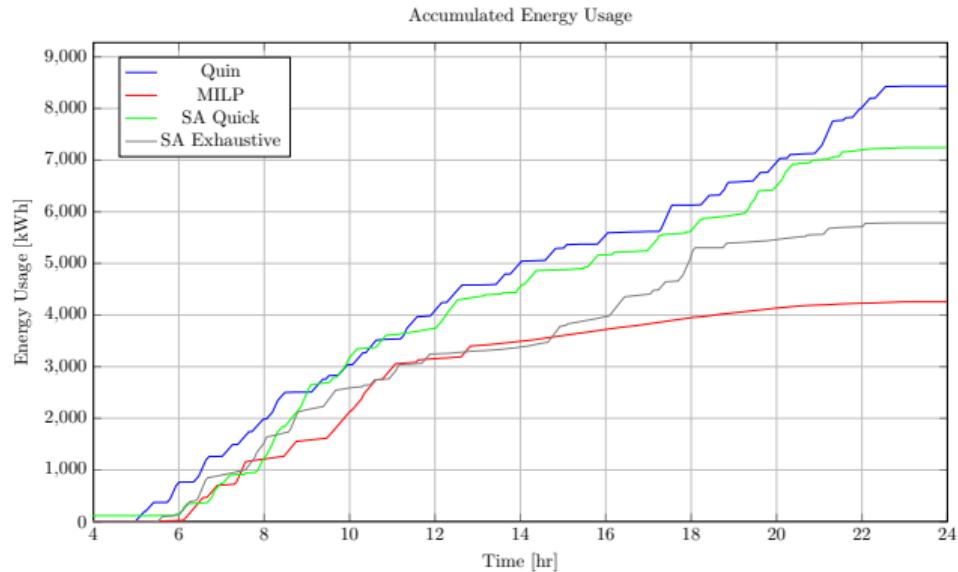
Schedule



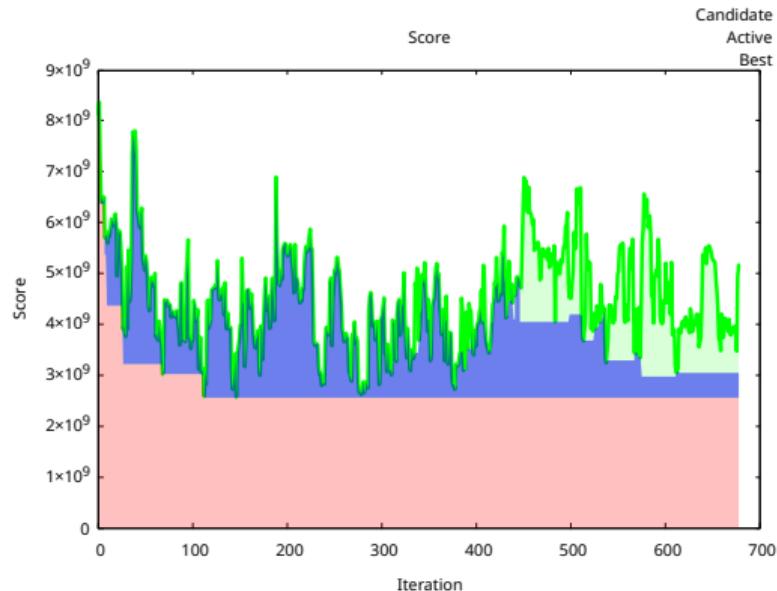
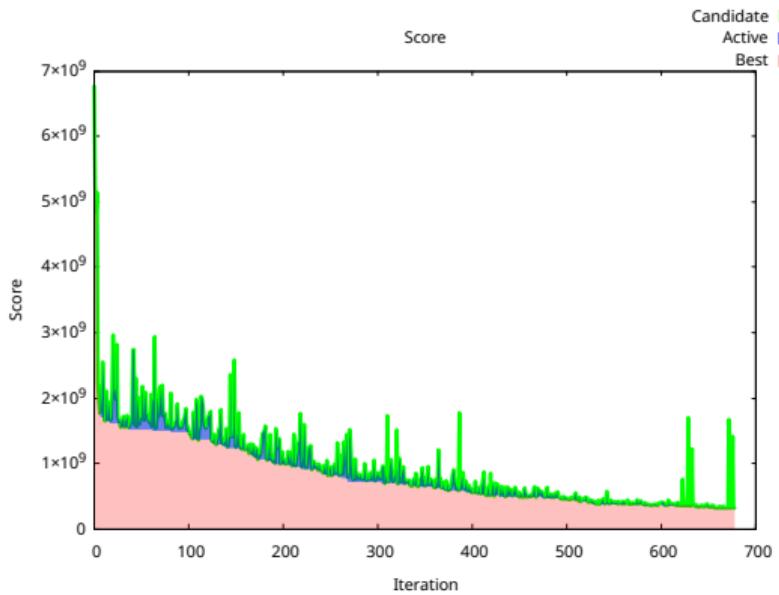
Power



Energy



How To Resolve This Problem?



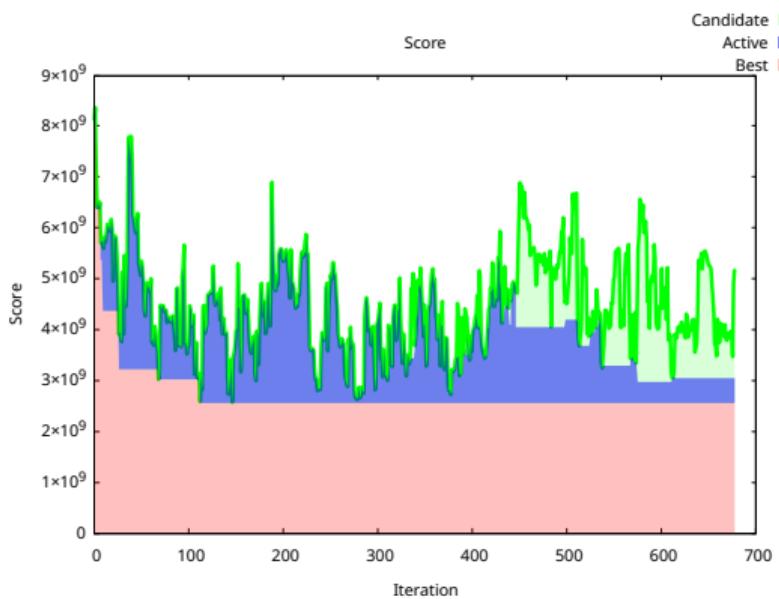
- ▶ Reverse search and weight the visit indices
- ▶ Be more aggressive in exploiting the best solution

- ▶ Candidate solutions diverge
- ▶ Hard time handling “difficult” routes

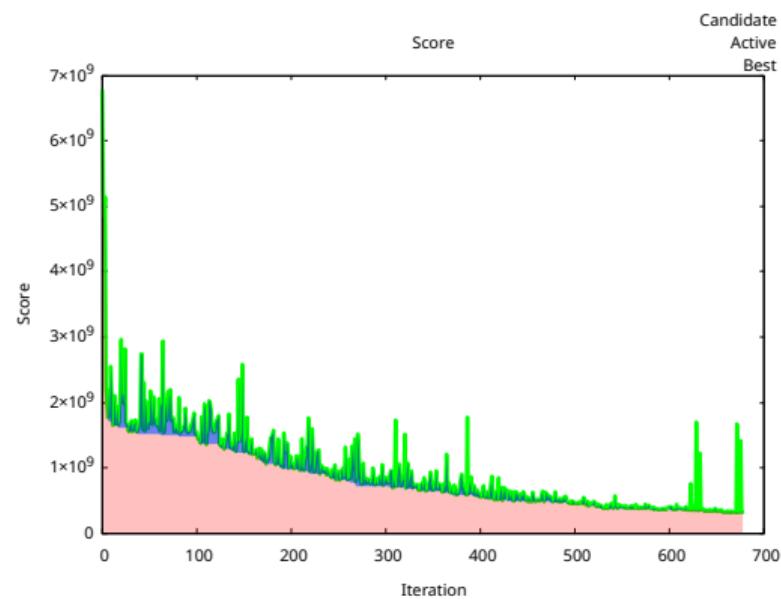
Score Convergence Comparison



Before Fix



After Fix



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