

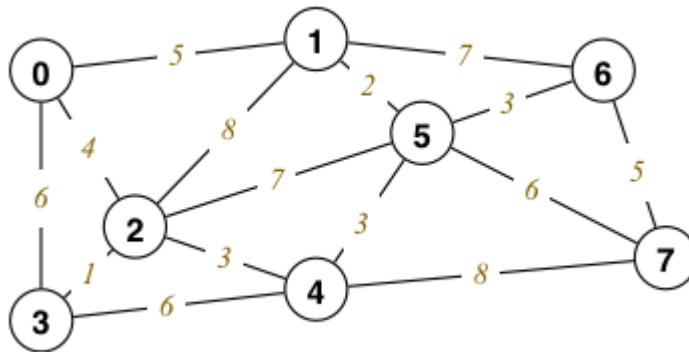
## 目录

1. Week 9 Exercises
  1. Dijkstra.txt
  2. Prim.txt
  3. Kruskal.txt
  4. eulerianCycle.c
  5. unreachable.c

# Week 9 Exercises

## Dijkstra.txt

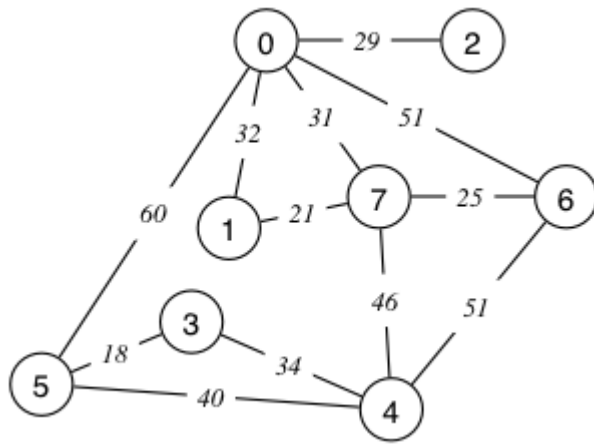
Dijkstra's algorithm computes the minimum path costs from the source node 0 to all the other vertices, resulting in the Shortest Path Tree (SPT). For the following graph:



- a. show the order that vertices are visited, and the path cost  $pacost[]$  of each of the vertices
  - (when faced with a choice, select the edge to the lowest vertex)
- b. there are 4 vertices that undergo *non-trivial edge relaxation*. What are they, and what is the reduction in cost for each vertex?
- c. draw the SPT

## Prim.txt

Prim's Algorithm generates a Minimum Spanning Tree (MST). For the following graph:



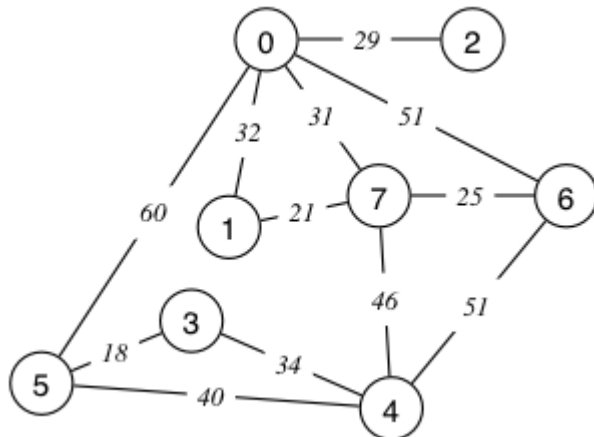
- a. by hand, use the sets *mst* and *rest* to build an MST one vertex at a time Do this in the form of a table, where the first 2 lines have been completed for you.

mst	rest vertex	cost
{0}	2	29
{0,2}	7	31
...	...	...

- b. draw the MST  
 c. how many edges in the MST?  
 d. what is the cost of the MST?

## Kruskal.txt

Kruskal's Algorithm also generates an MST. For the following graph:



- a. show which edges are considered at each step and ...  
     ◦ whether the edge gets accepted, or rejected (because it causes a cycle)  
 b. draw the MST.  
     ◦ is the MST the same of different to the MST generated by Prim's algorithm?

## eulerianCycle.c

- a. What determines whether a graph is Eulerian or not?  
 b. Write a C program that reads a graph, prints the graph, and determines whether an input graph is Eulerian or not.  
     ◦ if the graph is Eulerian, the program prints an Eulerian path

- you should start with vertex 0
- note that you may use the function *findEulerianCycle()* from the lecture on Graph Search Applications
  - if it is not Eulerian, the program prints the message *Not Eulerian*

For example,

- The graph:

```
#4
0 1 0 2 0 3 1 2 2 3
```

is not Eulerian (*can you see why?*). Using this as input, your program should output:

```
V=4, E=5
<0 1> <0 2> <0 3>
<1 0> <1 2>
<2 0> <2 1> <2 3>
<3 0> <3 2>
Not Eulerian
```

- In the above-named lecture I showed a 'concentric squares' graph (called *concsquares*):

```
#8
0 7 7 5 5 1 1 0
6 0 6 7
2 5 2 7
4 1 4 5
3 0 3 1
```

which is Eulerian, although I've labelled the vertices differently here. For this input your program should produce the output:

```
V=8, E=12
<0 1> <0 3> <0 6> <0 7>
<1 0> <1 3> <1 4> <1 5>
<2 5> <2 7>
<3 0> <3 1>
<4 1> <4 5>
<5 1> <5 2> <5 4> <5 7>
<6 0> <6 7>
<7 0> <7 2> <7 5> <7 6>
Eulerian cycle: 0 1 4 5 2 7 5 1 3 0 6 7 0
```

Draw *concsquares*, label it as given in the input file above, and check the cycle is indeed Eulerian.

- The function *findEulerCycle()* in the lecture notes does not handle disconnected graphs. In a disconnected Eulerian graph, each subgraph has an Eulerian cycle.
  - Modify this function to handle disconnected graphs.
  - With this change, your program should now work for the graph consisting of 2 disconnected triangles:

```
#6
0 1 0 2 1 2 3 4 3 5 4 5
```

It should now find 2 Eulerian paths:

```
V=6, E=6
<0 1> <0 2>
<1 0> <1 2>
```

```

<2 0> <2 1>
<3 4> <3 5>
<4 3> <4 5>
<5 3> <5 4>
Eulerian cycle: 0 1 2 0
Eulerian cycle: 3 4 5 3

```

## unreachable.c

Write a program that uses a *fixed-point computation* to find all the vertices in a graph that are unreachable from the start vertex (assume it to be 0). Note the following:

- the fixed-point computation should be iterative
- **you should not use recursion, or stacks or queues**

If a graph is disconnected:

- then those vertices not reachable (say vertices 8 and 9) should be output as follows:

```
Unreachable vertices = 8 9
```

If a graph is connected then all vertices are reachable and the output is :

```
Unreachable vertices = none
```

For example:

- Here is a graph that consists of 2 disconnected triangles:

```

#6
0 1 0 2 1 2 3 4 3 5 4 5

```

If the start vertex is 0, then the output should be:

```

V=6, E=6
<0 1> <0 2>
<1 0> <1 2>
<2 0> <2 1>
<3 4> <3 5>
<4 3> <4 5>
<5 3> <5 4>
Unreachable vertices = 3 4 5

```

because obviously the vertices in the second triangle are not reachable from the first.

- here is a connected graph:

```

#5
0 1 1 2 2 3 3 4 4 0
1 3 1 4
2 4

```

Starting at any vertex, the result should be:

```

V=5, E=8
<0 1> <0 4>
<1 0> <1 2> <1 3> <1 4>
<2 1> <2 3> <2 4>
<3 1> <3 2> <3 4>

```

```
<4 0> <4 1> <4 2> <4 3>  
Unreachable vertices = none
```

Week9Exercises (2019-08-22 14:04:43由AlbertNymeyer编辑)