

1. Induction on strings.

On page 302 of the book, the operation **rev** on strings is defined as follows:

$$\begin{aligned}\mathbf{rev}(\epsilon) &:= \epsilon \\ \mathbf{rev}(xa) &:= a\mathbf{rev}(x)\end{aligned}$$

for any  $x \in \Sigma^*$  and  $a \in \Sigma$ .

Show that, for any strings  $x, y \in \Sigma^*$ , we have  $\mathbf{rev}(xy) = \mathbf{rev}(y) \mathbf{rev}(x)$ . In other words, the reverse of  $xy$  is the reverse of  $y$  followed by the reverse of  $x$ .

**Base Case:**

Consider the case where  $|y| = 0$ . This is equivalent to  $y = \epsilon$ .

By definition, the null string  $\epsilon$  serves as an identity for string concatenation.

*Proof.*

$\mathbf{rev}(xy) = \mathbf{rev}(x\epsilon)$	by substituting $\epsilon$ in for $y$	
$= \mathbf{rev}(x)$	by definition of $\epsilon$	
$= \epsilon \mathbf{rev}(x)$	by definition of $\epsilon$	
$= \mathbf{rev}(\epsilon) \mathbf{rev}(x)$	by definition of <b>rev</b>	
$= \mathbf{rev}(y) \mathbf{rev}(x)$	by substituting $y$ in for $\epsilon$	□

**Induction Hypothesis:**

Assume that for  $|y| = k$ , where  $k \geq 1$ ,  $\mathbf{rev}(xy) = \mathbf{rev}(y) \mathbf{rev}(x)$ .

**Inductive Step:**

Let  $y \neq \epsilon$ . Therefore,  $y$  can be represented as the concatenation  $bc$  where  $b \in \Sigma^*$  and  $c \in \Sigma$ .

*Proof.*

$\mathbf{rev}(xy) = \mathbf{rev}(x(bc))$	by substituting $bc$ in for $y$	
$= \mathbf{rev}((xb)c)$	by associativity of concatenation	
$= c \mathbf{rev}(xb)$	by definition of <b>rev</b>	
$= c \mathbf{rev}(b) \mathbf{rev}(x)$	by Induction Hypothesis	
$= \mathbf{rev}(bc) \mathbf{rev}(x)$	by definition of <b>rev</b>	
$= \mathbf{rev}(y) \mathbf{rev}(x)$	by substituting $y$ in for $bc$	□