UNM SCHOOL of ENGINEERING
Department of Computer Science

Program Optimization

CS 341: Intro. to Computer Architecture & Organization

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Today

- **Overview**
- Generally Useful Optimizations
  - Code motion/precomputation
  - Strength reduction
  - Sharing of common subexpressions
  - Removing unnecessary procedure calls
- Optimization Blockers
  - Procedure calls
  - Memory aliasing
- **Exploiting Instruction-Level Parallelism**
- Dealing with Conditionals

Performance Realities

- *There's more to performance than asymptotic complexity*

- Constant factors matter too!
  - Easily see 10:1 performance range depending on how code is written
  - Must optimize at multiple levels:
    - algorithm, data representations, procedures, and loops
- Must understand system to optimize performance
  - How programs are compiled and executed
  - How to measure program performance and identify bottlenecks
  - How to improve performance without destroying code modularity and generality

Optimizing Compilers

- Provide efficient mapping of program to machine
  - register allocation
  - code selection and ordering (scheduling)
  - dead code elimination
  - eliminating minor inefficiencies
- Don't (usually) improve asymptotic efficiency
  - up to programmer to select best overall algorithm
  - big-O savings are (often) more important than constant factors
    - but constant factors also matter
- Have difficulty overcoming "optimization blockers"
  - potential memory aliasing
  - potential procedure side-effects

Limitations of Optimizing Compilers

- Operate under fundamental constraint
  - Must not cause any change in program behavior
  - Often prevents it from making optimizations when would only affect behavior under pathological conditions.
- Behavior that may be obvious to the programmer can be obfuscated by languages and coding styles
  - e.g., Data ranges may be more limited than variable types suggest
- Most analysis is performed only within procedures
  - Whole-program analysis is too expensive in most cases
- Most analysis is based only on *static* information
  - Compiler has difficulty anticipating run-time inputs
- When in doubt, the compiler must be conservative

Generally Useful Optimizations

- Optimizations that you or the compiler should do regardless of processor / compiler
- Code Motion
  - Reduce frequency with which computation performed
    - If it will always produce same result
    - Especially moving code out of loop

```
void set_row(double *a, double *b,
    long i, long n)
{
    long j;
    for (j = 0; j < n; j++)
    a[n*i+j] = b[j];
}
```

```
long j;
int ni = n*i;
for (j = 0; j < n; j++)
  a[ni+j] = b[j];
```

Compiler-Generated Code Motion

```
void set_row(double *a, double *b,
    long i, long n)
{
    long j;
    for (j = 0; j < n; j++)
    a[n*i+j] = b[j];
}
```

```
long j;
long ni = n*i;
double *rowp = a+ni;
for (j = 0; j < n; j++)
 *rowp++ = b[j];
```

Where are the FP operations?

```
set_row:
    testq  %rcx, %rcx       # Test n
    jle  .L4              # If 0, goto done
    movq   %rcx, %rax       # rax = n
    imulq  %rdx, %rax       # rax *= i
    leaq   (%rdi,%rax,8), %rdx  # rowp = A + n*i*8
    movl   $0, %r8d             # j = 0
.L3:                          # loop:
    movq (%rsi,%r8,8), %rax # t = b[j]
    movq %rax, (%rdx)       # *rowp = t
    addq $1, %r8            # j++
    addq $8, %rdx           # rowp++
    cmpq %r8, %rcx        # Compare n:j
    jg   .L3              # If >, goto loop
.L4:                          # done:
    rep ; ret
```

Reduction in Strength

- Replace costly operation with simpler one
- Shift, add instead of multiply or divide

  `16*x      -->     x << 4`

  - Utility machine dependent
  - Depends on cost of multiply or divide instruction
    - On Intel Nehalem, integer multiply requires 3 CPU cycles
- Recognize sequence of products

```
for (i = 0; i < n; i++)
  for (j = 0; j < n; j++)
    a[n*i + j] = b[j];
```

```
int ni = 0;
for (i = 0; i < n; i++) {
  for (j = 0; j < n; j++)
    a[ni + j] = b[j];
  ni += n;
}
```

Share Common Subexpressions

- Reuse portions of expressions
- Compilers not very good at exploiting arithmetic properties

```
/* Sum neighbors of i,j */
up =    val[(i-1)*n + j  ];
down =  val[(i+1)*n + j  ];
left =  val[i*n     + j-1];
right = val[i*n     + j+1];
sum = up + down + left + right;
```

```
long inj = i*n + j;
up =    val[inj - n];
down =  val[inj + n];
left =  val[inj - 1];
right = val[inj + 1];
sum = up + down + left + right;
```

3 multiplications: i*n, (i–1)*n, (i+1)*n

1 multiplication: i*n

```
leaq   1(%rsi), %rax  # i+1
leaq   -1(%rsi), %r8  # i-1
imulq  %rcx, %rsi     # i*n
imulq  %rcx, %rax     # (i+1)*n
imulq  %rcx, %r8      # (i-1)*n
addq   %rdx, %rsi     # i*n+j
addq   %rdx, %rax     # (i+1)*n+j
addq   %rdx, %r8      # (i-1)*n+j
```

```
imulq   %rcx, %rsi   # i*n
addq %rdx, %rsi   # i*n+j
movq %rsi, %rax   # i*n+j
subq %rcx, %rax   # i*n+j-n
leaq (%rsi,%rcx), %rcx # i*n+j+n
```

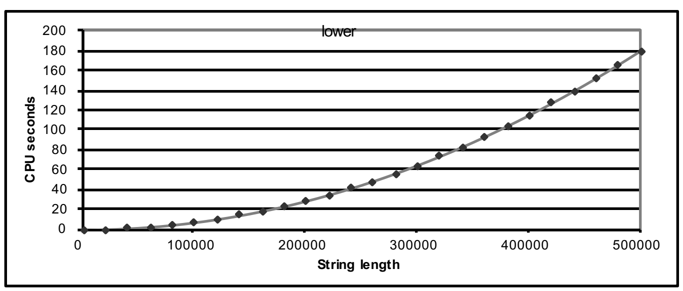
Optimization Blocker #1: Procedure Calls

- Procedure to Convert String to Lower Case

```
void lower(char *s)
{
  int i;
  for (i = 0; i < strlen(s); i++)
    if (s[i] >= 'A' && s[i] <= 'Z')
      s[i] -= ('A' - 'a');
}
```

Lower Case Conversion Performance

- Time quadruples when double string length
- Quadratic performance



Convert Loop To Goto Form

```
void lower(char *s)
{
   int i = 0;
   if (i >= strlen(s))
     goto done;
 loop:
   if (s[i] >= 'A' && s[i] <= 'Z')
       s[i] -= ('A' - 'a');
   i++;
   if (i < strlen(s))
     goto loop;
 done:
}
```

- `strlen` executed every iteration

Calling Strlen

```
/* My version of strlen */
size_t strlen(const char *s)
{
    size_t length = 0;
    while (*s != '\0') {
   s++;
   length++;
    }
    return length;
}
```

- Strlen performance
  - Only way to determine length of string is to scan its entire length, looking for null character.
- Overall performance, string of length N
  - N calls to strlen
  - Require times N, N-1, N-2, …, 1
  - Overall $O(N^2)$ performance

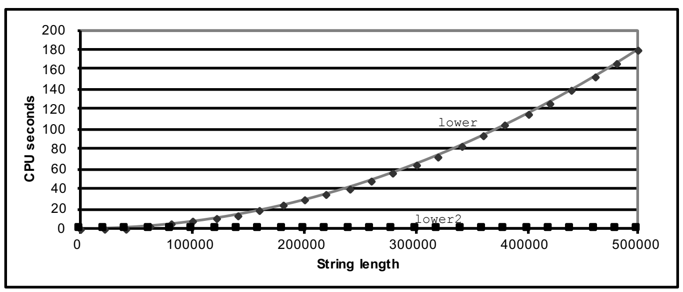
Improving Performance

```
void lower(char *s)
{
  int i;
  int len = strlen(s);
  for (i = 0; i < len; i++)
    if (s[i] >= 'A' && s[i] <= 'Z')
      s[i] -= ('A' - 'a');
}
```

- Move call to `strlen` outside of loop
- Since result does not change from one iteration to another
- Form of code motion

Lower Case Conversion Performance

- Time doubles when double string length
- Linear performance of lower2



Optimization Blocker: Procedure Calls

- *Why couldn't compiler move* `strlen` *out of inner loop?*
  - Procedure may have side effects
    - Alters global state each time called
  - Function may not return same value for given arguments
    - Depends on other parts of global state
    - Procedure `lower` could interact with `strlen`
- Warning:
  - Compiler treats procedure call as a black bo
  - Weak optimizations near them
- Remedies:
  - Use of `inline` functions
    - GCC does this with –O2
    - See web aside ASM:OPT
  - Do your own code motion

```
int lencnt = 0;
size_t strlen(const char *s)
{
    size_t length = 0;
    while (*s != '\0') {
   s++; length++;
    }
    lencnt += length;
    return length;
}
```

Memory Matters

```
/* Sum rows is of n X n matrix a
   and store in vector b  */
void sum_rows1(double *a, double *b, long n) {
    long i, j;
    for (i = 0; i < n; i++) {
      b[i] = 0;
      for (j = 0; j < n; j++)
        b[i] += a[i*n + j];
    }
}
```

```
# sum_rows1 inner loop
.L53:
     addsd    (%rcx), %xmm0          # FP add
     addq$8, %rcx
     decq%rax
     movsd    %xmm0, (%rsi,%r8,8)    # FP store
     jne .L53
```

- Code updates `b[i]` on every iteration
- Why couldn't compiler optimize this away?

Memory Aliasing

```
/* Sum rows is of n X n matrix a
   and store in vector b  */
void sum_rows1(double *a, double *b, long n) {
    long i, j;
    for (i = 0; i < n; i++) {
      b[i] = 0;
      for (j = 0; j < n; j++)
        b[i] += a[i*n + j];
    }
}
```

```
double A[9] =
  { 0,   1,   2,
    4,   8,  16},
   32,  64, 128};

double B[3] = A+3;

sum_rows1(A, B, 3);
```

Value of `B`:

```
init:   [4, 8, 16]
i = 0: [3, 8, 16]
i = 1: [3, 22, 16]
i = 2: [3, 22, 224]
```

- Code updates `b[i]` on every iteration
- Must consider possibility that these updates will affect program behavior

Removing Aliasing

```
/* Sum rows is of n X n matrix a
   and store in vector b  */
void sum_rows2(double *a, double *b, long n) {
    long i, j;
    for (i = 0; i < n; i++) {
      double val = 0;
      for (j = 0; j < n; j++)
        val += a[i*n + j];
       b[i] = val;
    }
}
```

```
# sum_rows2 inner loop
.L66:
     addsd    (%rcx), %xmm0   # FP Add
     addq$8, %rcx
     decq%rax
     jne .L66
```

- No need to store intermediate results

Optimization Blocker: Memory Aliasing

- Aliasing
  - Two different memory references specify single location
  - Easy to have happen in C
    - Since allowed to do address arithmetic
    - Direct access to storage structures
  - Get in habit of introducing local variables
    - Accumulating within loops
    - Your way of telling compiler not to check for aliasing

Program Optimizations

Optimizing Compilers (Recap)

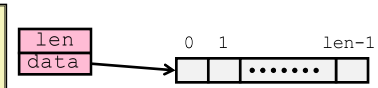
- Efficiently map program to machine
  - register allocation
  - code selection and ordering (scheduling)
  - dead code elimination
  - eliminating minor inefficiencies
- Don't (usually) improve asymptotic efficiency
  - Programmer should select best algorithm
  - big-O savings are (often) more important than constant factors
    - but constant factors also matter
- Limitations: Optimization blockers dependent on runtime behavior
  - Typically static, intra-procedural analysis
  - Safe optimizations only: original program behavior must be preserved
    - Memory aliasing: two pointers to overlapping memory locations
    - Stateful functions (with side effects): function output not solely dependent on inputs

Other (Potential Compiler) Optimizations

- Code motion
  - reduce execution frequency of an operation
- Strength reduction
  - substitute cheap operations for costly ones
- Function inlining
  - replace function call with function body
- Sharing common sub-expressions
  - a more complex variant of code motion
- Eliminating unnecessary memory references
  - use intermediates that compiler can store in registers

Benchmark Example: Vectors Data Type

```
/* data structure for vectors */
typedef struct{
    int len;
    double *data;
} vec;
```




```
/* retrieve vector element and store at val */
double get_vec_element(*v, idx, double *val)
{
    if (idx < 0 || idx >= v->len)
        return 0;

    *val = v->data[idx];

    return 1;
}
```

Benchmark Computation

```
void combine1(vec_ptr v, data_t *dest)
{
    long int i;
    *dest = IDENT;

    for (i = 0; i < vec_length(v); i++) {
       data_t val;
       get_vec_element(v, i, &val);
       *dest = *dest OP val;
    }
}
```

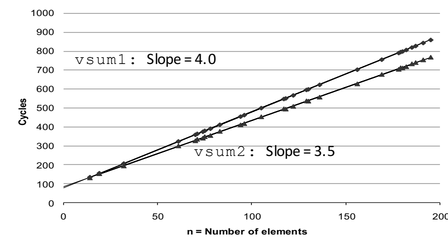
Compute sum or product of vector elements

Data Types
- Different declarations for `data_t`
- `int`
- `float`
- `double`

Operations
- Different `OP` and `IDENT`
  - `+ / 0`
  - `* / 1`

Cycles Per Element (CPE)

- Convenient expression of program performance for vectors or lists
- Length = n
- In our case: CPE = cycles per OP
- T = CPE*n + Overhead
  - CPE is slope of line



Benchmark Performance

```
void combine1(vec_ptr v, data_t *dest)
{
    long int i;
    *dest = IDENT;
    for (i = 0; i < vec_length(v); i++) {
       data_t val;
       get_vec_element(v, i, &val);
       *dest = *dest OP val;
    }
}
```

Compute sum or product of vector elements

Method	Integer		Double FP	
Operation	Add	Mult	Add	Mult
Combine1 unoptimized	29.0	29.2	27.4	27.9
Combine1 –O1	12.0	12.0	12.0	13.0

Basic Optimizations

```
void combine4(vec_ptr v, data_t *dest)
{
  int i;
  int length = vec_length(v);
  data_t *d = get_vec_start(v);
  data_t t = IDENT;
  for (i = 0; i < length; i++)
    t = t OP d[i];
  *dest = t;
}
```

- Move vec_length out of loop
- Avoid bounds check on each cycle
- Accumulate in temporary

Effect of Basic Optimizations

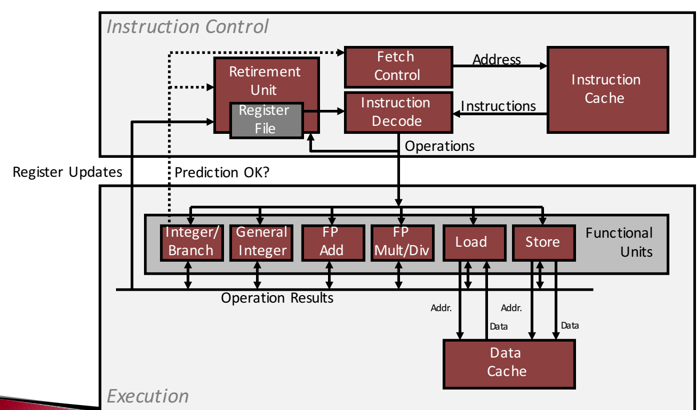
```
void combine4(vec_ptr v, data_t *dest)
{
  int i;
  int length = vec_length(v);
  data_t *d = get_vec_start(v);
  data_t t = IDENT;
  for (i = 0; i < length; i++)
    t = t OP d[i];
  *dest = t;
}
```

Method	Integer		Double FP	
Operation	Add	Mult	Add	Mult
Combine1 –O1	12.0	12.0	12.0	13.0
Combine4	2.0	3.0	3.0	5.0

Exploiting Instruction-Level Parallelism

- More complex optimizations require at least a basic understanding of modern processor design
  - Hardware can execute multiple instructions in parallel
- Performance limited by data dependencies
- Simple transformations can yield dramatic improvements
  - Compilers often cannot make these transformations
  - Lack of associativity and distributivity in floating-point arithmetic

Modern CPU Design



Superscalar Processor

- Definition: A superscalar processor can issue and execute *multiple instructions in one cycle*. The instructions are retrieved from a sequential instruction stream and are usually scheduled dynamically.
- Benefit: without programming effort, superscalar processor can take advantage of the *instruction level parallelism* that most programs have
- Most CPUs since about 1998 are superscalar.
- Intel: since Pentium Pro

Nehalem CPU

- Multiple instructions can execute in parallel
  1 load, with address computation
  1 store, with address computation
  2 simple integer (one may be branch)
  1 complex integer (multiply/divide)
  1 FP Multiply
  1 FP Add

- Some instructions take > 1 cycle, but can be pipelined

Instruction	*Latency*	*Cycles/Issue*
Load / Store	4	1
Integer Multiply	3	1
Integer/Long Divide	**11--21**	**11--21**
Single/Double FP Multiply	4/5	1
Single/Double FP Add	3	1
Single/Double FP Divide	**10--23**	**10--23**

x86-64 Compilation of Combine4

- Inner Loop (Case: Integer Multiply)

```
.L519:                          # Loop:
  imull  (%rax,%rdx,4), %ecx    # t = t * d[i]
  addq   $1, %rdx               # i++
  cmpq   %rdx, %rbp             # Compare length:i
  jg     .L519                  # If >, goto Loop
```

Method	Integer		Double FP	
Operation	Add	Mult	Add	Mult
Combine4	2.0	3.0	3.0	5.0
Latency Bound	1.0	3.0	3.0	5.0

Loop Unrolling

- Increase number of elements computed per iteration
  - Reduces loop overhead
  - Exposes other optimization opportunities

Loop Unrolling

```
void unroll2a_combine(vec_ptr v, data_t *dest)
{
    int length = vec_length(v);
    int limit = length-1;
    data_t *d = get_vec_start(v);
    data_t x = IDENT;
    int i;
    /* Combine 2 elements at a time */
    for (i = 0; i < limit; i+=2) {
        x = (x OP d[i]) OP d[i+1];
    }
    /* Finish any remaining elements */
    for (; i < length; i++) {
        x = x OP d[i];
    }
    *dest = x;
}
```

- Perform 2x more useful work per iteration

Effect of Loop Unrolling

Method	Integer		Double FP	
Operation	Add	Mult	Add	Mult
Combine4	2.0	3.0	3.0	5.0
Unroll 2x	2.0	1.5	3.0	5.0
Latency Bound	1.0	3.0	3.0	5.0

- Helps integer multiply
  - below latency bound
  - Compiler does clever optimization
- Others don't improve. *Why?*
  - Still sequential dependency

```
x = (x OP d[i]) OP d[i+1];
```

Loop Unrolling with Reassociation

```
void unroll2aa_combine(vec_ptr v, data_t *dest)
{
    int length = vec_length(v);
    int limit = length-1;
    data_t *d = get_vec_start(v);
    data_t x = IDENT;
    int i;
    /* Combine 2 elements at a time */
    for (i = 0; i < limit; i+=2) {
        x = x OP (d[i] OP d[i+1]);
    }
    /* Finish any remaining elements */
    for (; i < length; i++) {
        x = x OP d[i];
    }
    *dest = x;
}
```

- Can this change the result of the computation?
- Yes, for FP. *Why?*

Compare to before

```
x = (x OP d[i]) OP d[i+1];
```

Effect of Reassociation

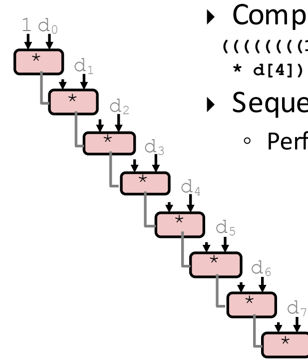
Method	Integer		Double FP	
Operation	Add	Mult	Add	Mult
Combine4	2.0	3.0	3.0	5.0
Unroll 2x	2.0	1.5	3.0	5.0
Unroll 2x, reassociate	2.0	1.5	1.5	3.0
Latency Bound	1.0	3.0	3.0	5.0
Throughput Bound	1.0	1.0	1.0	1.0

- Nearly 2x speedup for Int *, FP +, FP *
  - Reason: Breaks sequential dependency

```
x = x OP (d[i] OP d[i+1]);
```

  - Why is that? (next slide)

Combine4 = Serial Computation (OP = *)



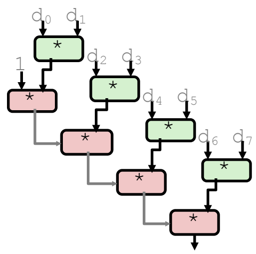
- Computation (length=8)

```
((((((((1 * d[0]) * d[1]) * d[2]) * d[3])
 * d[4]) * d[5]) * d[6]) * d[7])
```

- Sequential dependence
  - Performance: determined by latency of OP

Reassociated Computation

```
x = x OP (d[i] OP d[i+1]);
```



- What changed:
  - Ops in the next iteration can be started early (no dependency)
- Overall Performance
  - N elements, D cycles latency/op
  - Should be (N/2+1)*D cycles:
    CPE = D/2
  - Measured CPE slightly worse for FP mult

Loop Unrolling with Separate Accumulators

```
void unroll2a_combine(vec_ptr v, data_t *dest)
{
    int length = vec_length(v);
    int limit = length-1;
    data_t *d = get_vec_start(v);
    data_t x0 = IDENT;
    data_t x1 = IDENT;
    int i;
    /* Combine 2 elements at a time */
    for (i = 0; i < limit; i+=2) {
       x0 = x0 OP d[i];
       x1 = x1 OP d[i+1];
    }
    /* Finish any remaining elements */
    for (; i < length; i++) {
        x0 = x0 OP d[i];
    }
    *dest = x0 OP x1;
}
```

- Different form of reassociation

Effect of Separate Accumulators

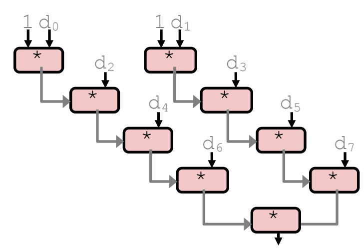
Method	Integer		Double FP	
Operation	Add	Mult	Add	Mult
Combine4	2.0	3.0	3.0	5.0
Unroll 2x	2.0	1.5	3.0	5.0
Unroll 2x, reassociate	2.0	1.5	1.5	3.0
Unroll 2x Parallel 2x	1.5	1.5	1.5	2.5
Latency Bound	1.0	3.0	3.0	5.0
Throughput Bound	1.0	1.0	1.0	1.0

- 2x speedup (over unroll2) for Int *, FP +, FP *
  - Breaks sequential dependency in a "cleaner," more obvious way

```
x0 = x0 OP d[i];
x1 = x1 OP d[i+1];
```

Separate Accumulators

```
x0 = x0 OP d[i];
x1 = x1 OP d[i+1];
```



- **What changed:**
  - Two independent "streams" of operations
- **Overall Performance**
  - N elements, D cycles latency/op
  - Should be (N/2+1)*D cycles:
    CPE = D/2
  - CPE matches prediction!

What Now?

Unrolling & Accumulating

- Idea
  - Can unroll to any degree L
  - Can accumulate K results in parallel
  - L must be multiple of K
- Limitations
  - Diminishing returns
    - Cannot go beyond throughput limitations of execution units
  - Large overhead for short lengths
    - Finish off iterations sequentially

Unrolling & Accumulating: Double *

- Case
  - Intel Nehelam (Shark machines)
  - Double FP Multiplication

FP *	Unrolling Factor L							
K	1	2	3	4	6	8	10	12
1	5.00	5.00	5.00	5.00	5.00	5.00		
2		2.50		2.50		2.50		
3			1.67					
4				1.25		1.25		
6					1.00			1.19
8						1.02		
10							1.01	
12								1.00

Accumulators

Unrolling & Accumulating: Int +

- Case
  - Intel Nehelam (Shark machines)
  - Integer addition

FP *	Unrolling Factor L							
K	1	2	3	4	6	8	10	12
1	2.00	2.00	1.00	1.01	1.02	1.03		
2		1.50		1.26		1.03		
3			1.00					
4				1.00		1.24		
6					1.00			1.02
8						1.03		
10							1.01	
12								1.09

Accumulators

Achievable Performance

Method	Integer		Double FP	
Operation	Add	Mult	Add	Mult
Scalar Optimum	1.00	1.00	1.00	1.00
Latency Bound	1.00	3.00	3.00	5.00
Throughput Bound	1.00	1.00	1.00	1.00

- Limited only by throughput of functional units
- Up to 29X improvement over original, unoptimized code

Using Vector Instructions

Method	Integer		Double FP	
Operation	Add	Mult	Add	Mult
Scalar Optimum	1.00	1.00	1.00	1.00
Vector Optimum	0.25	0.53	0.53	0.57
Latency Bound	1.00	3.00	3.00	5.00
Throughput Bound	1.00	1.00	1.00	1.00
Vec Throughput Bound	0.25	0.50	0.50	0.50

- Make use of SSE Instructions
  - Parallel operations on multiple data elements
  - See Web Aside OPT:SIMD on CS:APP web page

What About Branches?

- Challenge
  - Instruction Control Unit must work well ahead of Execution Unit to generate enough operations to keep EU busy

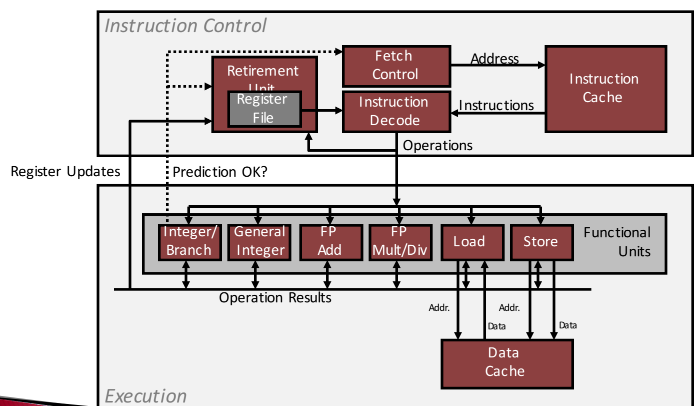
```
80489f3:  movl    $0x1,%ecx
80489f8:  xorl    %edx,%edx
80489fa:  cmpl    %esi,%edx
80489fc:  jnl     8048a25
80489fe:  movl    %esi,%esi
8048a00:  imull   (%eax,%edx,4),%ecx
```

Executing

How to continue?

  - When encounters conditional branch, cannot reliably determine where to continue fetching

Modern CPU Design



Branch Outcomes

- **When encounter conditional branch, cannot determine where to fetch**
  - Branch Taken: Transfer control to branch target
  - Branch Not-Taken: Continue with next instruction in sequence
- **Cannot resolve until outcome determined by branch/integer unit**

```
80489f3:  movl    $0x1,%ecx
80489f8:  xorl    %edx,%edx
80489fa:  cmpl    %esi,%edx
80489fc:  jnl     8048a25
80489fe:  movl    %esi,%esi
8048a00:  imull   (%eax,%edx,4),%ecx
```

Branch Not-Taken

Branch Taken

```
8048a25:  cmpl    %edi,%edx
8048a27:  jl      8048a20
8048a29:  movl    0xc(%ebp),%eax
8048a2c:  leal    0xffffffe8(%ebp),%esp
8048a2f:  movl    %ecx,(%eax)
```

Branch Prediction

- Idea
  - Guess which way branch will go
  - Begin executing instructions at predicted position
    - But don't actually modify register or memory data

```
80489f3: movl   $0x1,%ecx
80489f8: xorl   %edx,%edx
80489fa: cmpl   %esi,%edx
80489fc: jnl    8048a25
. . .
```

Predict Taken

```
8048a25: cmpl   %edi,%edx
8048a27: jl     8048a20
8048a29: movl   0xc(%ebp),%eax
8048a2c: leal   0xffffffe8(%ebp),%esp
8048a2f: movl   %ecx,(%eax)
```

Begin Execution

Branch Prediction Through Loop

```
80488b1:   movl    (%ecx,%edx,4),%eax
80488b4:   addl    %eax,(%edi)
80488b6:   incl    %edx
80488b7:   cmpl    %esi,%edx        i = 98
80488b9:   jl      80488b1

80488b1:   movl    (%ecx,%edx,4),%eax
80488b4:   addl    %eax,(%edi)
80488b6:   incl    %edx
80488b7:   cmpl    %esi,%edx        i = 99
80488b9:   jl      80488b1

80488b1:   movl    (%ecx,%edx,4),%eax
80488b4:   addl    %eax,(%edi)
80488b6:   incl    %edx
80488b7:   cmpl    %esi,%edx        i = 100
80488b9:   jl      80488b1

80488b1:   movl    (%ecx,%edx,4),%eax
80488b4:   addl    %eax,(%edi)
80488b6:   incl    %edx
80488b7:   cmpl    %esi,%edx        i = 101
80488b9:   jl      80488b1
```

Assume
vector length = 100

Predict Taken (OK)

Predict Taken (Oops)

Read invalid location

Executed

Fetched

Branch Misprediction Invalidation

```
80488b1:   movl    (%ecx,%edx,4),%eax
80488b4:   addl    %eax,(%edi)
80488b6:   incl    %edx
80488b7:   cmpl    %esi,%edx        i = 98
80488b9:   jl      80488b1

80488b1:   movl    (%ecx,%edx,4),%eax
80488b4:   addl    %eax,(%edi)
80488b6:   incl    %edx
80488b7:   cmpl    %esi,%edx        i = 99
80488b9:   jl      80488b1

80488b1:   movl    (%ecx,%edx,4),%eax
80488b4:   addl    %eax,(%edi)
80488b6:   incl    %edx
80488b7:   cmpl    %esi,%edx        i = 100
80488b9:   jl      80488b1

80488b1:   movl    (%ecx,%edx,4),%eax
80488b4:   addl    %eax,(%edi)
80488b6:   incl    %edx
80488b7:   cmpl    %esi,%edx        i = 101
80488b9:   jl      80488b1
```

Assume
vector length = 100

Predict Taken (OK)

Predict Taken (Oops)

Invalidate

Branch Misprediction Recovery

```
80488b1:   movl    (%ecx,%edx,4),%eax
80488b4:   addl    %eax,(%edi)
80488b6:   incl    %edx              i = 99
80488b7:   cmpl    %esi,%edx
80488b9:   jl      80488b1
80488bb:   leal    0xffffffe8(%ebp),%esp
80488be:   popl    %ebx
80488bf:   popl    %esi
80488c0:   popl    %edi
```

Definitely not taken

- Performance Cost
  - Multiple clock cycles on modern processor
  - Can be a major performance limiter

Effect of Branch Prediction

- Loops
  - Typically, only miss when hit loop end
- Checking code
  - Reliably predicts that error won't occur

```
void combine4b(vec_ptr v,
               data_t *dest)
{
    long int i;
    long int length = vec_length(v);
    data_t acc = IDENT;
    for (i = 0; i < length; i++) {
      if (i >= 0 && i < v->len) {
          acc = acc OP v->data[i];
      }
    }
    *dest = acc;
}
```

Method	Integer		Double FP	
Operation	Add	Mult	Add	Mult
Combine4	2.0	3.0	3.0	5.0
Combine4b	4.0	4.0	4.0	5.0

Getting High Performance

- Good compiler and flags
- Don't do anything stupid
  - Watch out for hidden algorithmic inefficiencies
  - Write compiler-friendly code
    - Watch out for optimization blockers: procedure calls & memory references
  - Look carefully at innermost loops (where most work is done)
- Tune code for machine
  - Exploit instruction-level parallelism
  - Avoid unpredictable branches
  - Make code cache friendly (Covered later in course)

GCC Optimizations

- -O1 (same as -O): reduce program size and time, quick optimizations
  - -fcompare-elim: identify instructions that compute processor flags, eliminate explicit comparison
  - -fcprop-registers: copy propagation to reduce dependencies
  - -fce: dead code elimination
  - -fdse: dead store elimination
  - -fdelayed-branch: exploit slot after branch delay
  - -fguess-branch-probability: based on heuristics and CFG
  - -fif-conversion: conditional jumps to branchless (cond. moves)equivalents -fif-conversion2
  - -ftree-ccp: conditional constant propagation
  - -ftree-dce: dead code elimination
  - -ftree-dominator-opts: dominator tree based optimizations (constant/copy propagation, redundancy elimination, ...) -fmerge-constants: across modules
  - -fipa-pure-const: identify pure or constant procedures
  - ...

GCC Optimizations (cont'd)

- -O2 (optimizations that don't bloat program)
  - -fcaller-saves
  - -fcrossjumping: unifies equivalent code
  - -fcse-*: common sub-expression elimination optimizations
  - -fexpensive-optimizations
  - -fsched-spec: speculative motion of non-loads instructions
  - -freorder-blocks: Reorder basic blocks: reduce taken branches and improve locality
  - -fstrict-aliasing: different types don't overlap
  - -freorder-functions
  - ...

GCC Optimizations (cont'd)

-O3: All optimizations

- -finline-functions
- -funswitch-loops: Move branches with loop invariant conditions out of the loop
- -fpredictive-commoning: reusing computations (especially memory loads and stores) performed in previous iterations of loops
- -fipa-cp-clone: Perform function cloning to make interprocedural constant propagation stronger

Other Considerations

- Insufficient registers (register spilling)
- Load/Store Costs
- Amdahl's Law: only speedups to bottlenecks matter!
- Profiling Tools
  - gprof