March 31

Reminder Exam I Newt Thursday (Apr 7)

Min-cut Algerichm Cuhile there are De vertices, randomly chouse an edge and centrare. return all edges between the 2 remaining vortices tachots the probability that the a minest survives to the very end? $> \frac{n}{n} \cdot \frac{n}{n-1} \cdot \frac{n}{n}$ $\frac{2}{n(n-1)} =$

Suppose we run the algorithm (n) times, and each nen will produce a cut, we will return the sinclese cut out of these (n) runs. what's the probability that this returned out is a minat? Cohat's the pubability that the returnal cut is not a minut? $\left(1-\frac{1}{\binom{n}{2}}\right)^{\binom{n}{2}} \left(2-\frac{1}{\binom{n}{2}}\right)^{\binom{n}{2}} = e^{-1}$ Recall the Taylor expansion of $e^{x} = 1 + x + \frac{x}{L'} + \frac{x^{3}}{3!} +$

When x is small, $e^{x} \approx 1+x$, $e^{-x} \approx 1+(-x)$

= 1-x

 $\int_{n}^{\infty} \left(1 + \frac{1}{n}\right)^n = e$

For exam,

- 1) understand the difference hotween Les Veges and Minte Carly Algorithms
- D) Basie tail analysis technique

 Marker Inequality $Pr(X;S) \leq \frac{ZB}{S}$ Chebapher Inequality

A

Randomiced Linear Programmig

max $C_1x_1 + C_2x_2$ \leftarrow

s,t. anxi+ anx = bi,

auxi+ aux = br

 $G_{n_1} X_1 + G_{n_1} X_2 \leq b_n$

max X,+x- : E objective finction

Sit. $2x_i + x_i \le 4$ 7 linear

 $-X \leq 0$ Constraints

 $-\alpha \leq 0$

1 x = 4

Let
$$x = \begin{bmatrix} x_i \\ x_i \end{bmatrix}$$

$$C = \begin{bmatrix} C_1 \\ C_V \end{bmatrix}$$

$$= \frac{(C_1,C_2)(X_1)}{g_{ij}}$$

Let
$$A = \begin{cases} G_{11} G_{12} \\ G_{21} G_{22} \\ \vdots \\ G_{n1} G_{n2} \end{cases}$$

$$b = \begin{cases} b_{1} \\ k \\ \vdots \\ b_{n} \end{cases}$$

$$b = \begin{cases} b_1 \\ k \\ b_n \end{cases}$$

then the conservint becomes

$$\begin{array}{c}
\left(\frac{\partial u}{\partial u} \times i + \frac{\partial u}{\partial u} \times i + \frac{$$

The condensed firm

 $max c^T x$

 $S,t, Ax \leq b$

 $min - C^T x$

S,t. -Axz-b

maro $(1,1)\begin{pmatrix} \pi_1 \\ \pi_2 \end{pmatrix}$

Sie $\begin{bmatrix} 2 \\ 3 \\ \end{bmatrix}$ $\begin{bmatrix} x_i \\ 5 \end{bmatrix}$

max XI+XI

y,t, $2x_1+x_1 \leq 4$

3×1-2×15

mas x,+x+0.5,+0.52

 $S.t \qquad (2x_i + x_i) + S_i = 4$

 $(3x_1-2x_1)+S_v=5$

5,70

5720

max CTX siti Ax= b

min ctx Sit, Anzb

max ctx

 $S, \epsilon, Ax = h$

ous (d-b) cus (d+B) sin (a+B) sin (x-p)

complex numbers $\sqrt{-1} = i \propto j$.

a= Youb

b= YsinO

at ib = a rang + i rsing = r (as b+isinb)

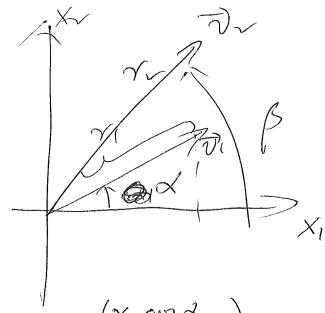
= Ye io

$$= (ac-bd) + i ((a+b)(r+d) - ac-bd)$$



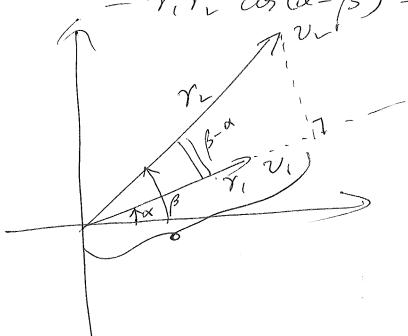
cosd+isind = 0cus B+ isinß = eiß (cosd+isind) (confs+isinf) = i (X+B) (cos & coss-sinxsins)+ i (sind our B + cord sin B) (cond cosp-sindsing)+ cus (2+/5)+ 2'sin(d+/s) r'(sidaß+coxsinß) cos (d+/s) = ond os/s - sin a sin/s sin (x+B) = sind our B+ on asin B cos (d+(-b)) = cos x cos (-b) - sind sin (-b) cus (d-b) = cosd ouß + sind sins

inner product



$$\overrightarrow{v}_{i} = \begin{pmatrix} \gamma_{i} \cos \lambda \\ \gamma_{i} \sin \lambda \end{pmatrix}$$

$$= r_1 r_1 \cos(\alpha - \beta) = r_1 (r_1 \cos(\beta - \alpha))$$



 $\vec{v}_i, \vec{v}_i = 1$

No Avi

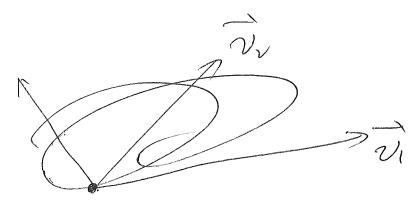
1. VV.

(νινι) = γιγι cos (β-α) = γι cos (β-α)

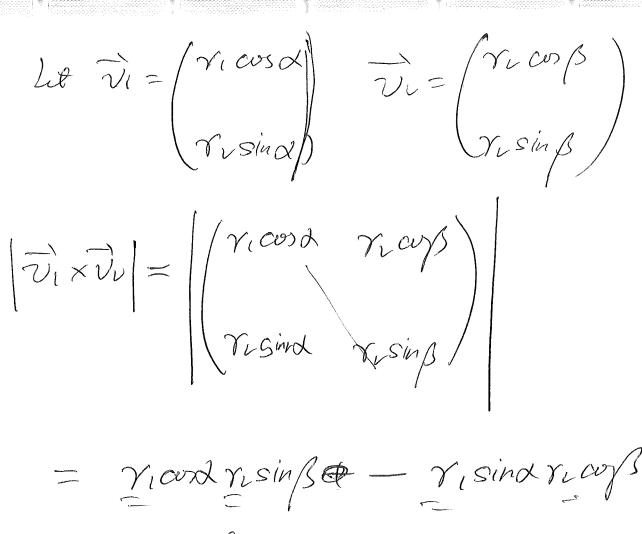
 $\langle v_1, v_v \rangle = |v_i| \cdot |v_v| \cos Q$

whon 0 = The evi, vi) =

Cross produce.







= $\gamma_i \cos \lambda \gamma_i \sin \beta - \gamma_i \sin \lambda \gamma_i \cos \beta$ = $\gamma_i \gamma_i \left(\cos \lambda \sin \beta - \sin \lambda \cos \beta \right)$ = $\gamma_i \gamma_i \left(\beta - \lambda \right) = -\gamma_i \gamma_i \sin (\lambda - \beta)$

Gradient Search



$$f(x) = x^{2} + 1$$

$$f(x) = 2$$

$$f(x) = 2$$

$$f(x) = 2$$

$$f(x) = 3$$

$$f$$

$$f(x)$$

$$\nabla f = \begin{cases} \frac{\partial f}{\partial x} \\ \frac{\partial f}{\partial x} \end{cases}$$

$$f(x+\partial x) = f(x) + \langle \nabla f, \Delta x \rangle$$

$$min f(x)$$

$$we want f(x+\partial x) < f(x)$$

$$\langle \nabla f, \Delta x \rangle < O$$

$$|\nabla f| \cdot |\Delta x| \cos \theta < C$$

