AVL - trees:

Balanced 1887: guarateed O(legn) performances for insertion and deletion (search)

AVL-tree: fer any node, the height difference between its left and right subtreess is < 1.

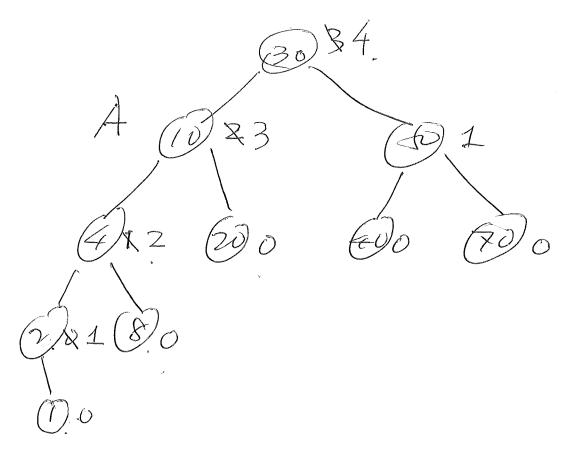
Reall that for an AVL-tree of n modes, ies height < 1.4 hyn

insertion in AVL

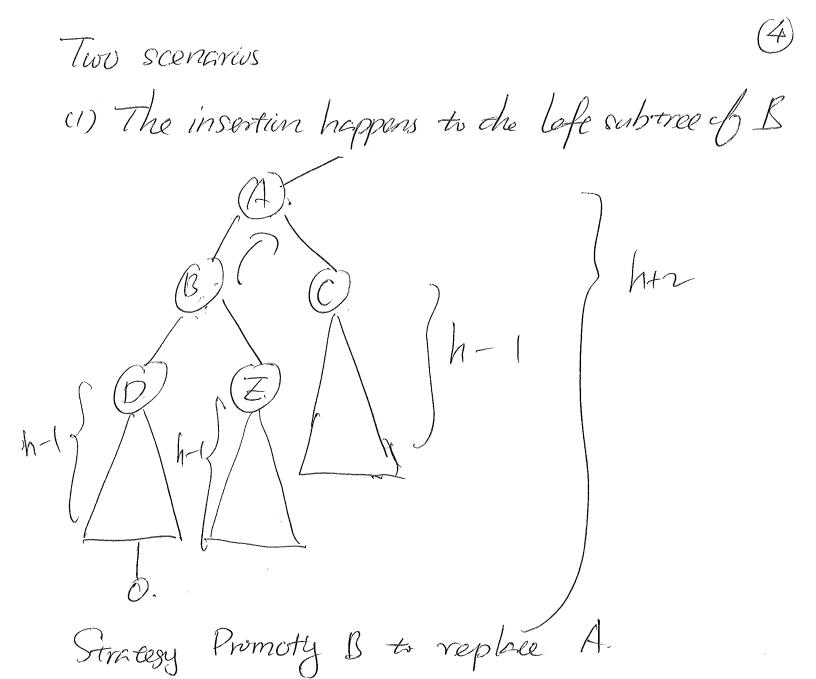
Da normal BST insertion

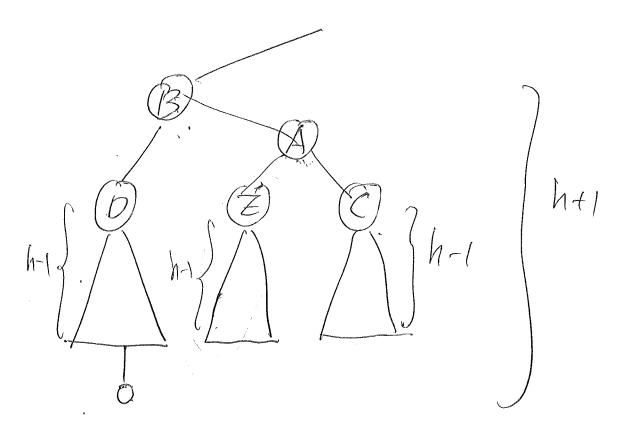
D if the normal insertion cause a violation of the AVL property, we will rebalance the tree with rotation (s)

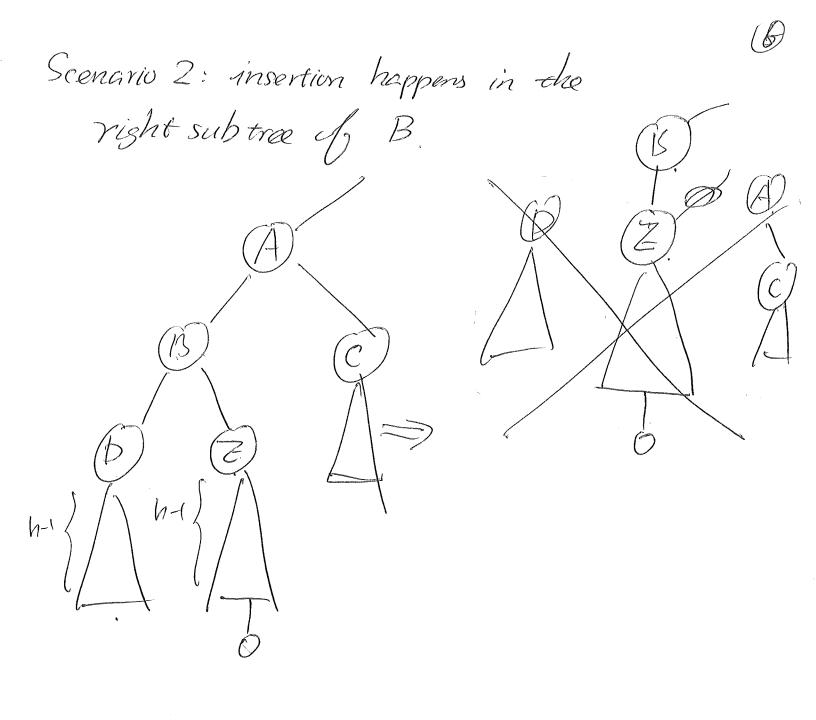
Let A be the node of the smallest subtree that is violating the AVL property after insertion.

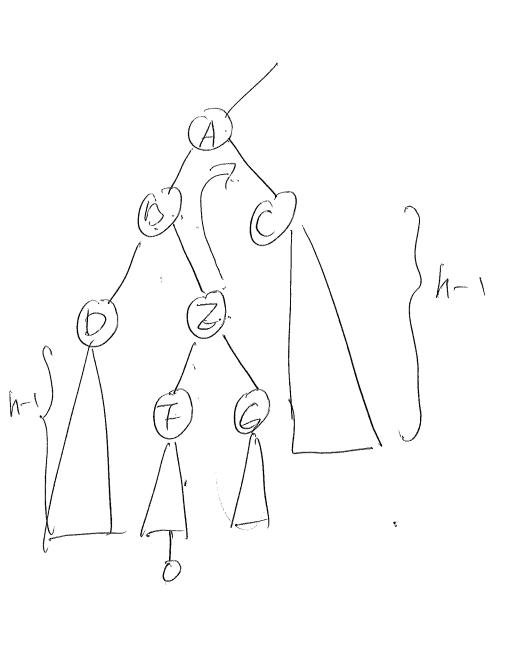


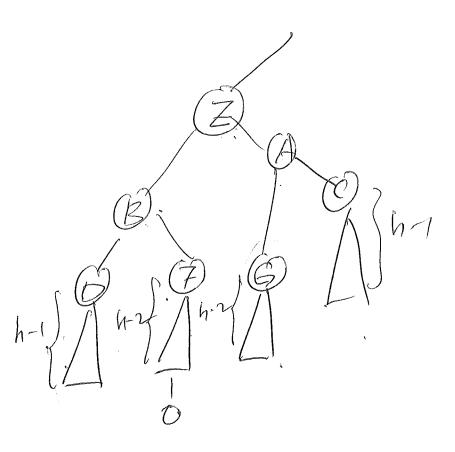
not of the smallest subtree violating AVL proposty Assume the insertion happens in the subtree routed at B inserted bey



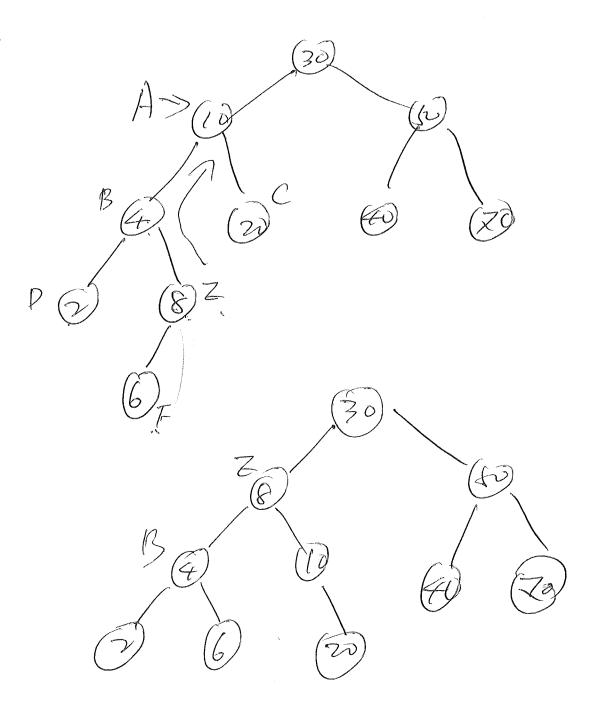








Double Rotation

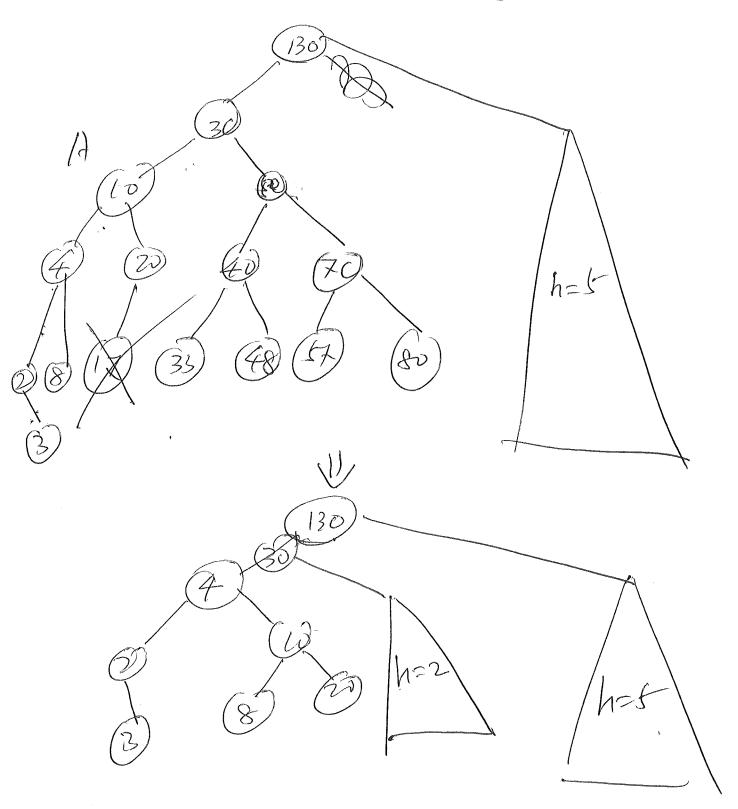




Deletion in AVL-tree

1 Normal deletion as in a 1887

1 vehalance if necessary



| | - |
|----|------------|
| 1 | 1 |
| 1/ | ' / |
| K | Market Co. |

Selection Algorichm.

Given an array A [1. n] unsorted when h= 1 and an integer k, $1 \le k \le n$ (median selection) Find the k smallest number in A

Naive Solution:

O Sort A (nlyn)

(1) Output Alk) Reduce the running time to O(n)

- @ Partition A into groups of 5.
- 1) Find the median of each group
 - 3) Put all the medians into a set S!
- (4) Recursively finel the median & of S
- (1) Use x as the pivot and partition A is to $S_1 < x < S_1$
- 6) If $k \leq |S_1|$ recursively find search S_2 else verisively search S_2

A= (8) 23 21 1,5, 7, 15, 16, 11, 2 (8) 10, 11,5 12, 13,3,8 (8) 15,1 2, 18, 22 (8) (8) 22

 $S = \{21, 11, 17, 11.59\}$ $k' = \emptyset 3$ X = 11.5

 $S_{i}=\{1,5,7,11,2,4,10,11,5,3,8,5,3$

Sn= {25, 23, 21, 13, 14, 20, 17, \$18, 12, 13 15, 1, 223

 $|S_1|=15$ $|S_2|=12$ recursisely search (S, k=7) if k=20

recursily seach (Sr, k-1511)

$$T(n) = n + T(\frac{\ln n}{5}) + T(mex[1s,1,1s,4])$$

$$0 \le 0 \le 0 \ne 0 \le 0$$

$$0 \le 0 \le 0 \ne 0 \le 0$$

$$0 \le 0 \le 0 \ne 0 \le 0$$

$$0 \le 0 \le 0 \ne 0 \le 0$$

$$0 \le 0 \le 0 \ne 0 \le 0$$

$$|s_1| \ge \frac{1}{5} \cdot \frac{1}{2} \cdot 3 = \frac{3}{10} \ln |s_1| = \frac{3}$$

$$|S_{1}| > |S_{1}| = \frac{3}{10}$$
 $|S_{1}| + |S_{1}| = |S_{1}|$
 $|S_{1}| + |S_{1}| = |S_{1}|$
 $|S_{1}| + |S_{1}| = \frac{7n}{10}$

$$T(n) = n + T(\frac{h}{f}) + T(\frac{7n}{10})$$

$$= 7(n) = \Theta(n) \approx lon$$