Recall Bernolli Trials

$$\frac{1}{0} \qquad \qquad P \qquad \qquad P \qquad \qquad 0$$

$$\mu = p$$

$$6^2 = p$$

Binomial Distribution

n Bernoulli Trials what's the probability of

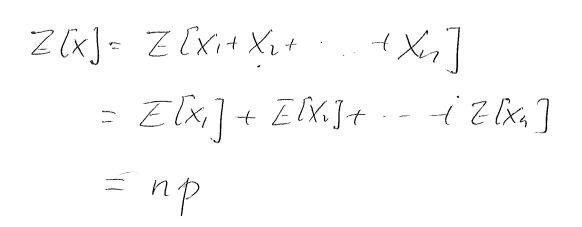
$$P(X=k) = \binom{n}{k} p^k g^{n-k} = \binom{n}{k} p^k (p)^{n-k}$$

$$X \sim 0, 1, 2, \ldots, n$$

$$\mathbb{E}[x]$$
 and $Var(x)$

Let Xj represent the oritaine of the jeth Bernoulli Trial

$$P(X_{j}=0) = 1-p$$
 $p(X_{j}=1)=p$



Example censider tossing a coin with 0.3 of head and 0.7 of tail.

Assume we toss the coin lov, over times how many heads should we expect?

100,000 x 0.3 = 30,000

Variance.

Var(X) = Var (Xi+Xr+ · - + Xn)

Reall the variance of the sum of independent v.v.

is the sum of the vairance

individual

Geometric Distribution

A Bernoulli Trial with a probability p of success. Let X be the v.v. representing the number of trials until the success occurs.

 $X \sim 1, 2, \cdots$ $P(X=k) = C_1 - p^{k-1}p$ O Check if the distribution is well-defined, $\sum_{j=1}^{\infty} P(X=j) = \lim_{n\to\infty} \left(\sum_{j=1}^{n} (1-p)^{j-1}p \right)$ $\sum_{j=1}^{n} (1-p)^{j-1}p = O(1-p)^{n}p + (1-p)^{n}p$ $+ (1-p)^{n}p$

 $= \frac{p - (1-p)^{n}p(1-p)}{(-(1-p))} \frac{p - 0}{v - \infty}$

$$Z[x] = \sum_{j=1}^{\infty} (pj^{j+}p \cdot j)$$

$$= (1-p)^{\circ}p \cdot 1 + (1+p)p \cdot 2 + (1+p)^{2}p \cdot 3 + (1+p)^{3}p \cdot 4 + \cdots$$

$$S = (1-p)^{\circ}p \cdot 1 + (1+p)p \cdot 2 + (1+p)^{2}p \cdot 3 + (1+p)^{3}p \cdot 4 + \cdots$$

$$PS = (1-p)^{\circ}p \cdot 1 + (1-p)^{\circ}p \cdot 2 + (1-p)^{3}p \cdot 3 + \cdots$$

$$PS = (1-p)^{\circ}p + (1-p) \cdot p + (1-p)^{2}p$$

$$= 1$$

$$S = \frac{1}{p}$$

$$Z[x] = \frac{1}{p}$$

Example: consider a coin with probability 0.2 of H

 $Var(X) = Z[X] - (Z[X])^T$ $Z(x^{2}) = \sum_{j=1}^{\infty} (-p)^{j-1} p \cdot j^{2} = \frac{(a+i)^{2} - a^{2}}{(a+i+a)^{2}}$ $= \frac{(a+i)^{2} - a^{2}}{(a+i+a)^{2}}$ $= \frac{(a+i)^{2} - a^{2}}{(a+i+a)^{2}}$ $S = (1-p)^{0}p \cdot 1 + (1-p)^{1} \cdot p \cdot 2 + (1-p)^{2} \cdot p \cdot 3 + (1-p)^{4} \cdot p \cdot 4 + (1-p)^{2} \cdot p \cdot 3 + (1-p)^{4} \cdot p \cdot 4 + (1 (1-p)^{1} \cdot p \cdot 1^{2} + (1-p)^{2} \cdot p \cdot 2^{2} + (1-p)^{3} \cdot p \cdot 3^{2} + \cdots$ pS = (1-p) p.18+ (1-p) p.(2.1+1)+(1-p).p(2.2+1)+ (1-p)3-p(2.3+1)+-(1-p) (2.2+1) = ((1-p)^0.p+ (1-p)'.p+ (1-p)^2.p+ (1-p)^p.+... (1-p) p.2.+ (1-p) p.2.2+(1-p) p.2.3-

$$= \sum_{j=1}^{\infty} (1-p)^{j-j} p$$

$$+ 2 \sum_{j=1}^{\infty} (1-p)^{j-j} p + 2 (1-p) \sum_{j=1}^{\infty} (1-p)^{j-j} p + 2 (1-p)^{j-j} p + 2$$

(Z)

Poisson Distribution

A r.v. X is subject to the Poisson Distribution
$$\int P(X=k) = \frac{\chi k e^{-\chi}}{k!} \qquad \chi = 0, 1, --$$

$$P(x-k) = \begin{cases} \frac{\lambda^{k}e^{-\lambda}}{k!} & \frac{\lambda^{k}e^{-\lambda}}{k!} \\ 0 & \text{is it well-defined?} \end{cases}$$

$$\begin{cases} \frac{\infty}{k} & \frac{\lambda^{k}e^{-\lambda}}{k!} & = e^{-\lambda} & \frac{\lambda^{k}e^{-\lambda}}{k!} \\ \frac{\infty}{k!} & = e^{-\lambda} & \frac{\lambda^{k}e^{-\lambda}}{k!} & = e^{-\lambda}e^{-\lambda} \end{cases}$$

$$= \begin{cases} \frac{\lambda^{k}e^{-\lambda}}{k!} & = e^{-\lambda} & \frac{\lambda^{k}e^{-\lambda}}{k!} & = e^{-\lambda}e^{-\lambda} \\ \frac{\lambda^{k}e^{-\lambda}}{k!} & = e^{-\lambda}e^{-\lambda}e^{-\lambda} \end{cases}$$

$$= \begin{cases} \frac{\lambda^{k}e^{-\lambda}}{k!} & = e^{-\lambda}e^{-\lambda}e^{-\lambda} \\ \frac{\lambda^{k}e^{-\lambda}}{k!} & = e^{-\lambda}e^{-\lambda}e^{-\lambda}e^{-\lambda} \end{cases}$$

$$= \begin{cases} \frac{\lambda^{k}e^{-\lambda}}{k!} & = e^{-\lambda}e^{\lambda}e^{-\lambda}e^{$$

$$V_{ar}(x) = \frac{2[x]}{(2[x])} (766)$$

$$= \frac{x^{k}e^{-x}}{k!} (k-1)+1)$$

$$= \frac{x^{k}e^{-x}}{(k-1)!} (k-1) + \frac{x^{k}e^{-x}}{(k-1)!}$$

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 $Var(x) = (x + \lambda) - x = \lambda$

Randomied Algorithm: an algorithm that make random choices during its execution.

It assume the on-demand availability of uniform random bits

Generally speeky: Two types.

Las Vegas Algorichm: guarantees convereness

No guarantee of running time

Monte Carlo Algorichm:

Suarantees running time

No guarantee of correreness

usually with a success pubability of $\propto p < 1$

applications.

- (1) Beating the adversary
- (2) Random Sampling
- (1) Hashig
- (4) Existence Prof.

$$X_{n+1} = (\alpha X_n + c) \mod (m)$$

$$X_0 = seed$$