

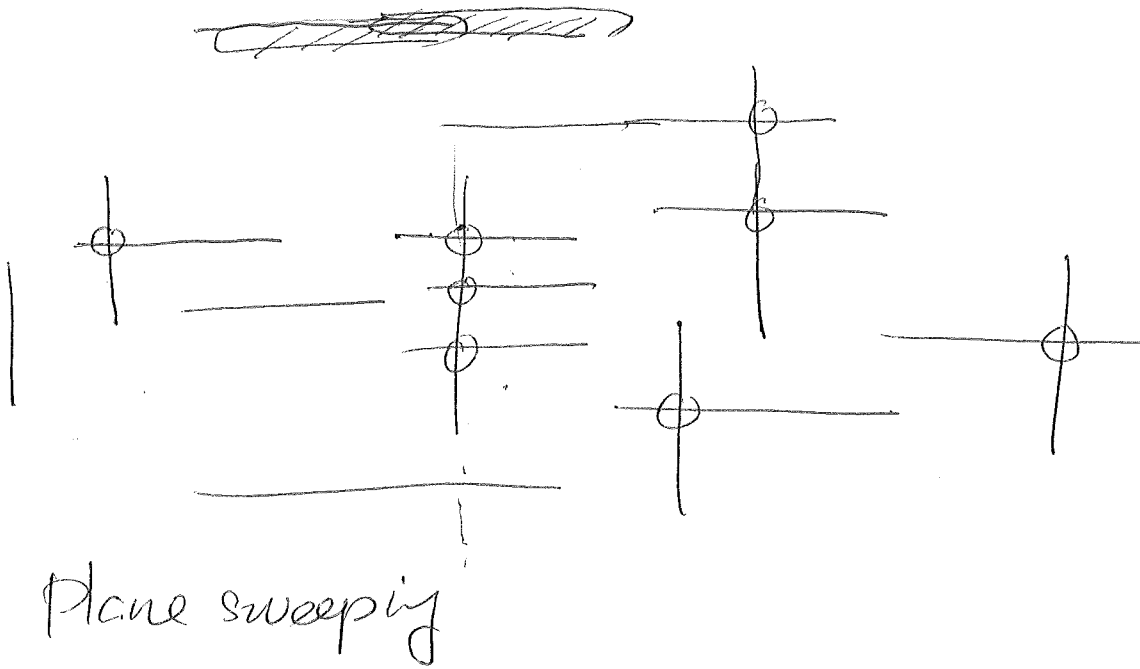
Feb 18, 2016

Exam 1 Thursday Feb 25.

Recall Line segment Intersection Problem

Given n vertical and n horizontal line segments

Find all intersections.



Plane sweeping

① Sort all vertical line segments based on their x -coordinates from left to right

② For each vertical line segment from left to right

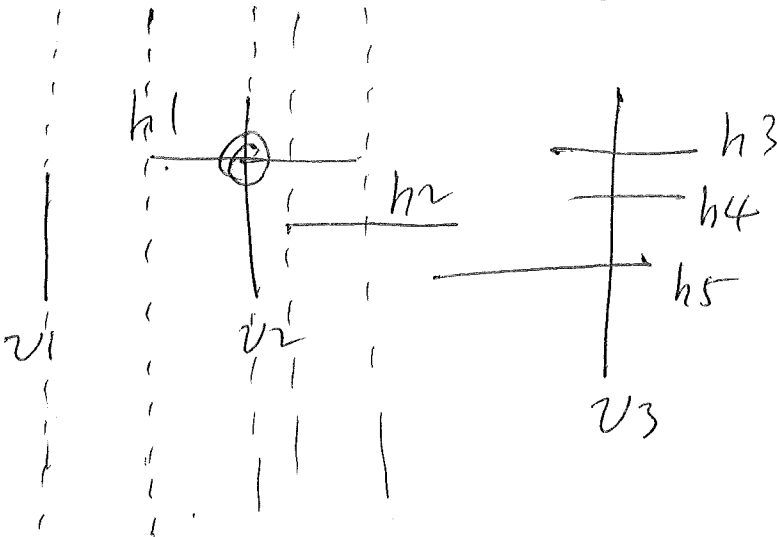
Find all intersections on it

(2)

Alg V2.

Sort all vertical line segments from left to right and place them in an array $A[1..n]$

initialize a BST storing all horizontal line segments intersecting the current sweep line.



A

v_1	v_2	v_3
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$BST = \phi$

A

v_1	v_2	v_3
-----------------------------	-------	-------

$BST = \{h_1\}$

A

v_1	v_2	v_3
-----------------------------	-----------------------------	-------

upper intersection

A

v_1	v_2	v_3
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$BST = \{h_1, h_2\}$

Alg V3

(3)

Sort the leftends, rightends of horizontal line segments and the x-coordinates of the vertical line segments from left to right, and put those in an

event array $A[1..3n]$

For $j = 1$ to $3n$

~~case~~ ~~case~~

Switch $A[j]$:

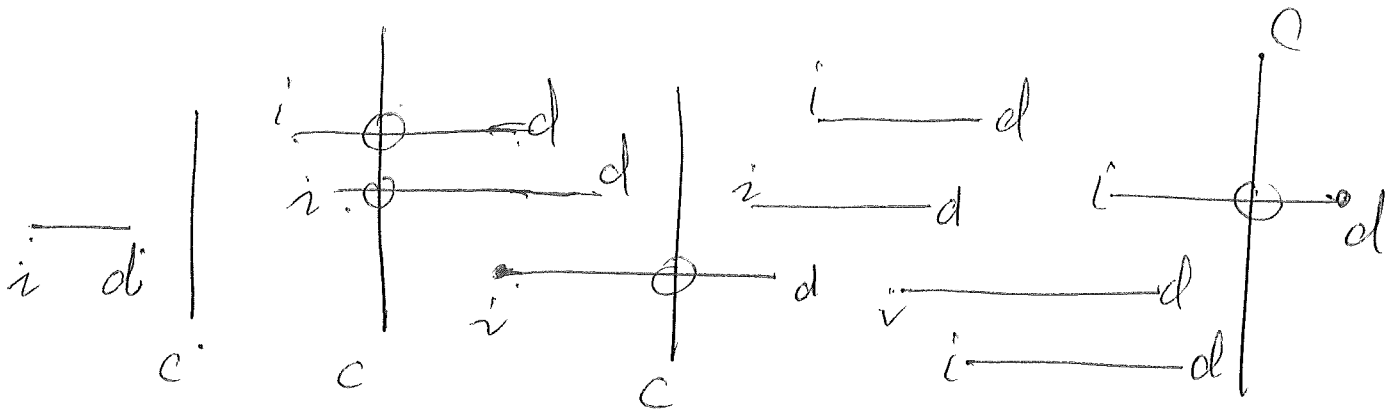
case 1 left end of a horizontal line segment
insert the corresponding line segment
into the BST

case 2 right end of a horizontal line
segment

delete it from the BST

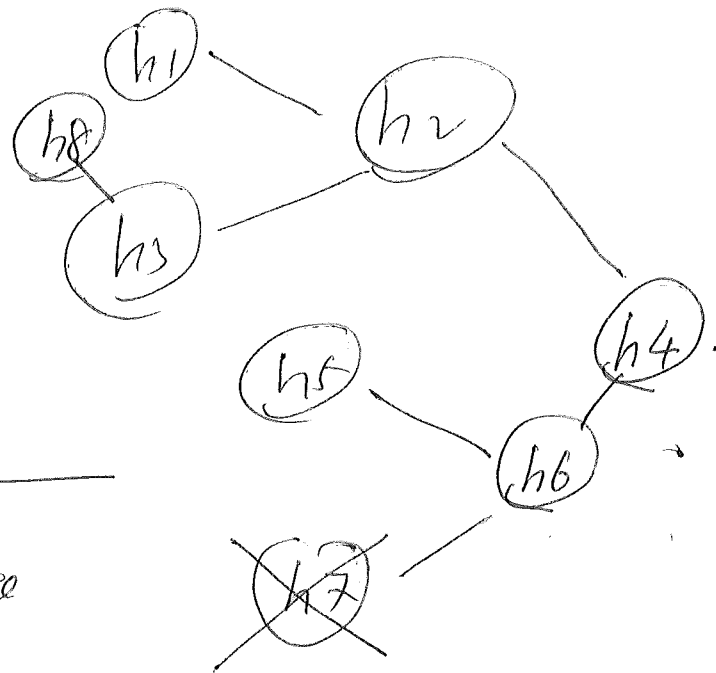
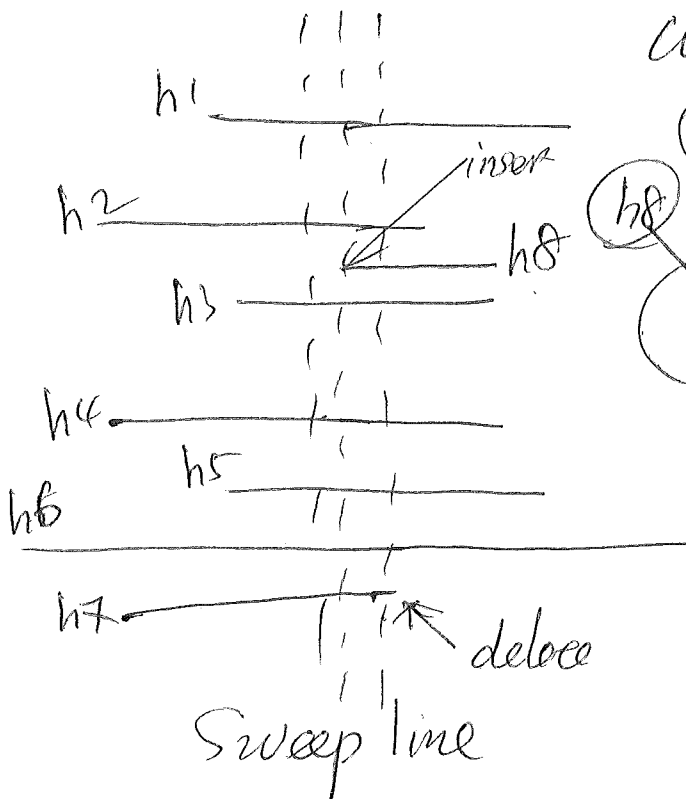
case 3 vertical line segment
output intersections

(4)



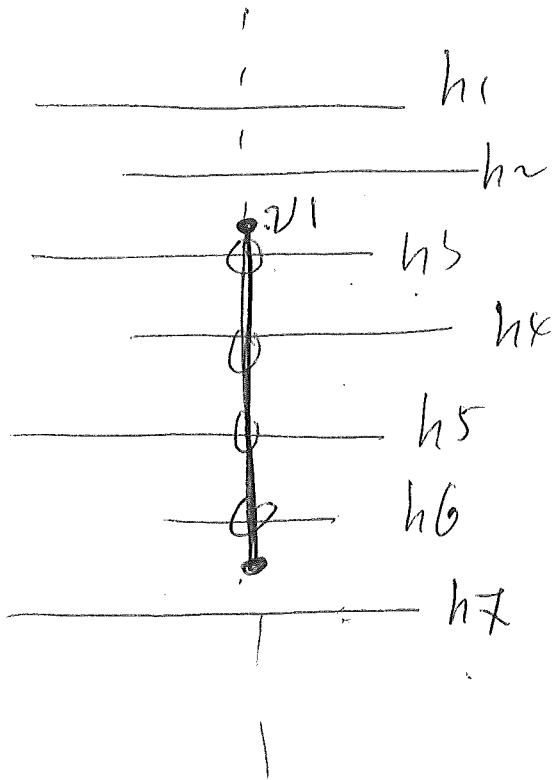
Q1 How BST is organised, i.e., how the horizontal line segments are stored in the BST

We will store based on y-coord.



(5)

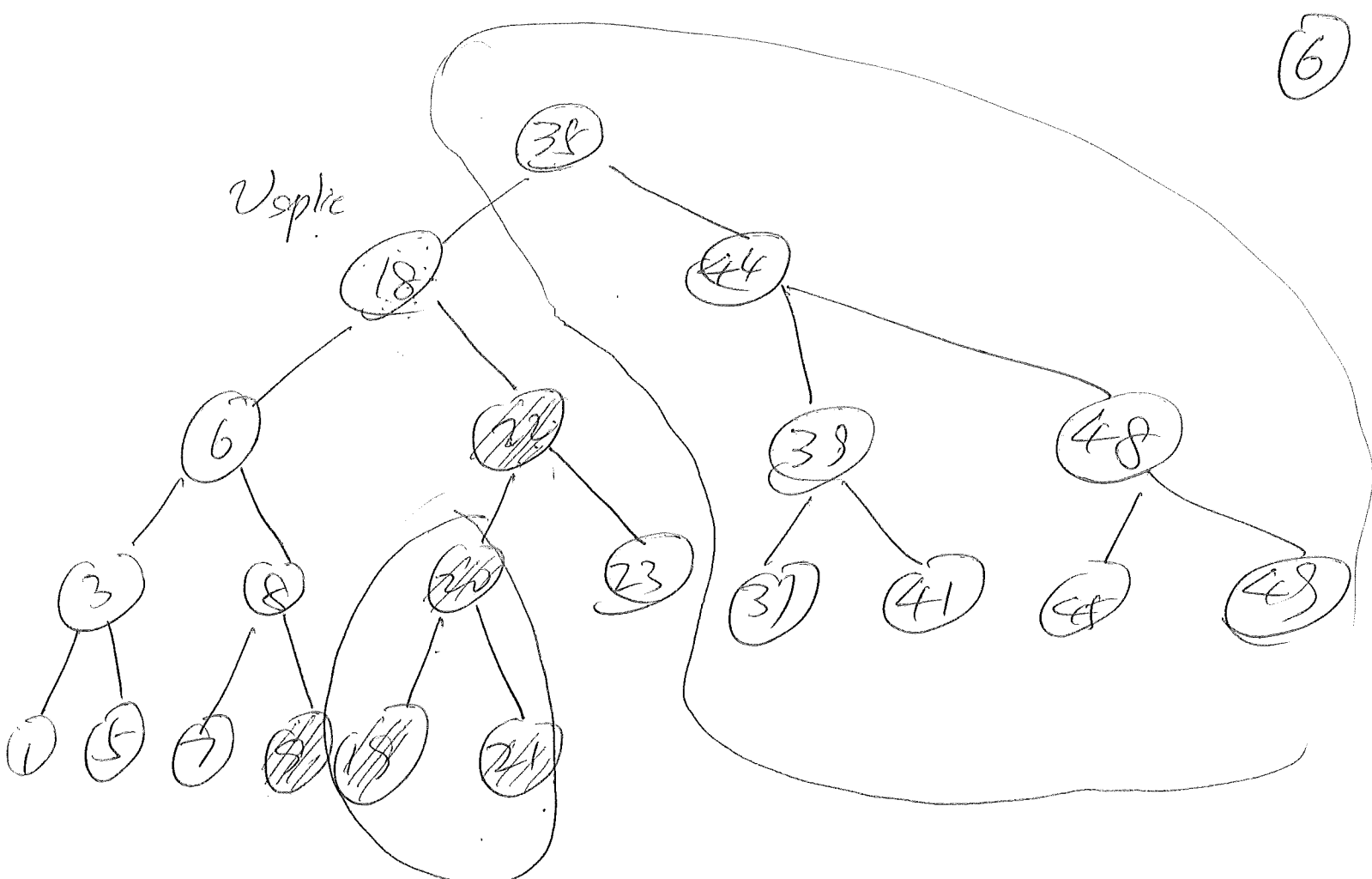
Q2 how to calculate intersections?



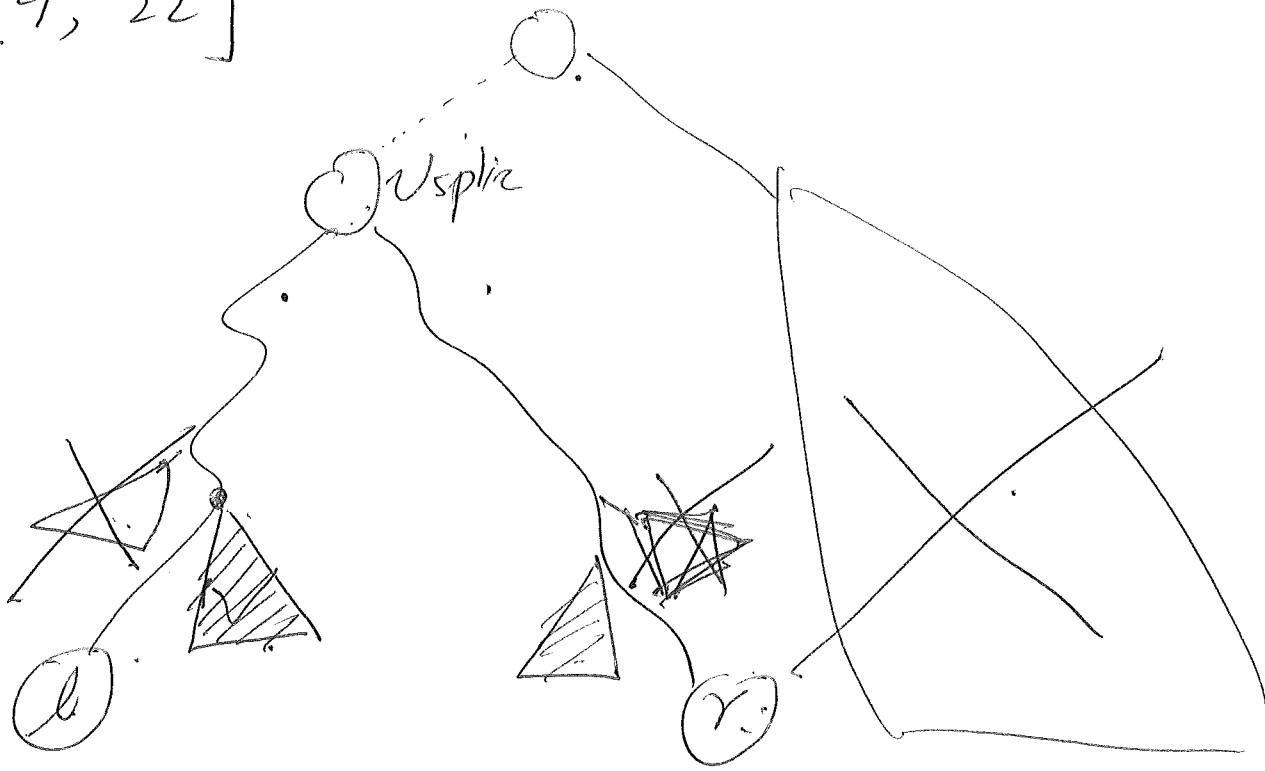
Observation: every horizontal line segments whose y -coordinate is within the y -range of the vertical line segment will generate an intersection

The Problem:

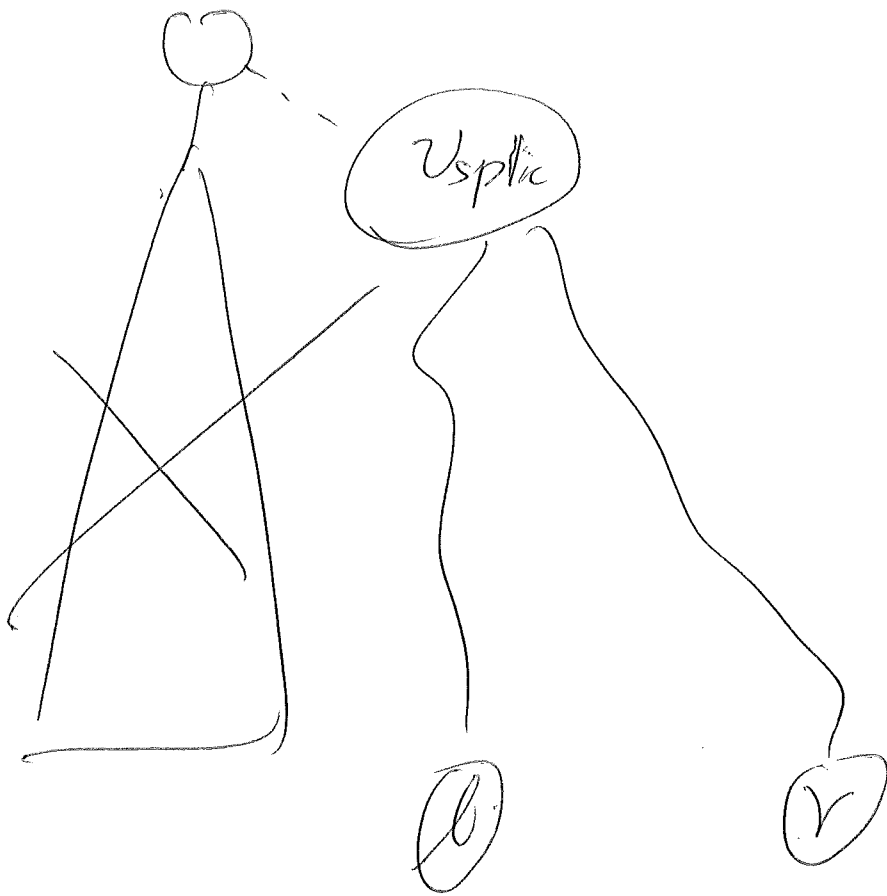
Given a BST. T and a query range $[l, r]$, output all ~~key~~ nodes of T whose key is within $[l, r]$



Query. [9, 22]



(7)



Range Query

input BST T , a query range $[l, r]$

output all nodes whose keys are within $[l, r]$

Find $Usplit$

Going down from $Usplit$ to (l) , output all right subtrees

Going down from $Usplit$ to (r) , output all left subtrees

For every node on the paths from $Usplit$ to (l) and (r) output it within $[l, r]$

Running Time

(8)

$O(\log n + k)$, where k is the number of nodes within the query range.

This type of running time is called output sensitive.

(9)

Sort the leftend points, rightend points and vertical line segments from left to right in an event array $A[1..3n]$

For $j=1$ to $3n$

switch $A[j]$

base y-coord

case 1: leftend point, insert into BST

case 2: right endpoint, delete from BST

case 3: vertical line segment

range query $[y_B, y_T]$ on the BST
output intersections

Running Time: $\underbrace{n \log n}_{\text{sort}} + \sum_{j=1}^{3n} (\log n + k_j)$

$$= n \log n + \left(\sum_{j=1}^{3n} \log n \right) + \left(\sum_{j=1}^{3n} k_j \right)$$

$$= O(n \log n + K)$$

where K is the number of intersections

$$A_{4 \times 6} (B_{6 \times 2} C_{2 \times 9})$$

$$(A_{4 \times 6} B_{6 \times 2}) C_{2 \times 9}$$

$$\textcircled{2} A_{4 \times 6} B_{6 \times 2} = (\quad)_{4 \times 2}$$

$$\begin{pmatrix} 0 & 0 & 0 & 0 & 0 & 0 \end{pmatrix} \begin{pmatrix} 0 \\ 0 \\ 6 \\ 0 \\ 0 \\ 0 \end{pmatrix} \quad \begin{matrix} \textcircled{4 \times 6 \times 2} \\ = 48 \end{matrix}$$

$$A_{4 \times 6} \underbrace{(B_{6 \times 2} C_{2 \times 9})}_{D_{6 \times 9}}$$

$$6 \times 2 \times 9 + 4 \times 6 \times 9 \\ = 324$$

$$\underbrace{(A_{4 \times 6} B_{6 \times 2})}_{D'_{4 \times 2}} C_{2 \times 9}$$

$$4 \times 6 \times 2 + 4 \times 2 \times 9 \\ = 48 + 72 \\ = 120$$

The Problem how to multiply

$$A_1 \times A_2 \times \dots \times A_n \quad ?$$

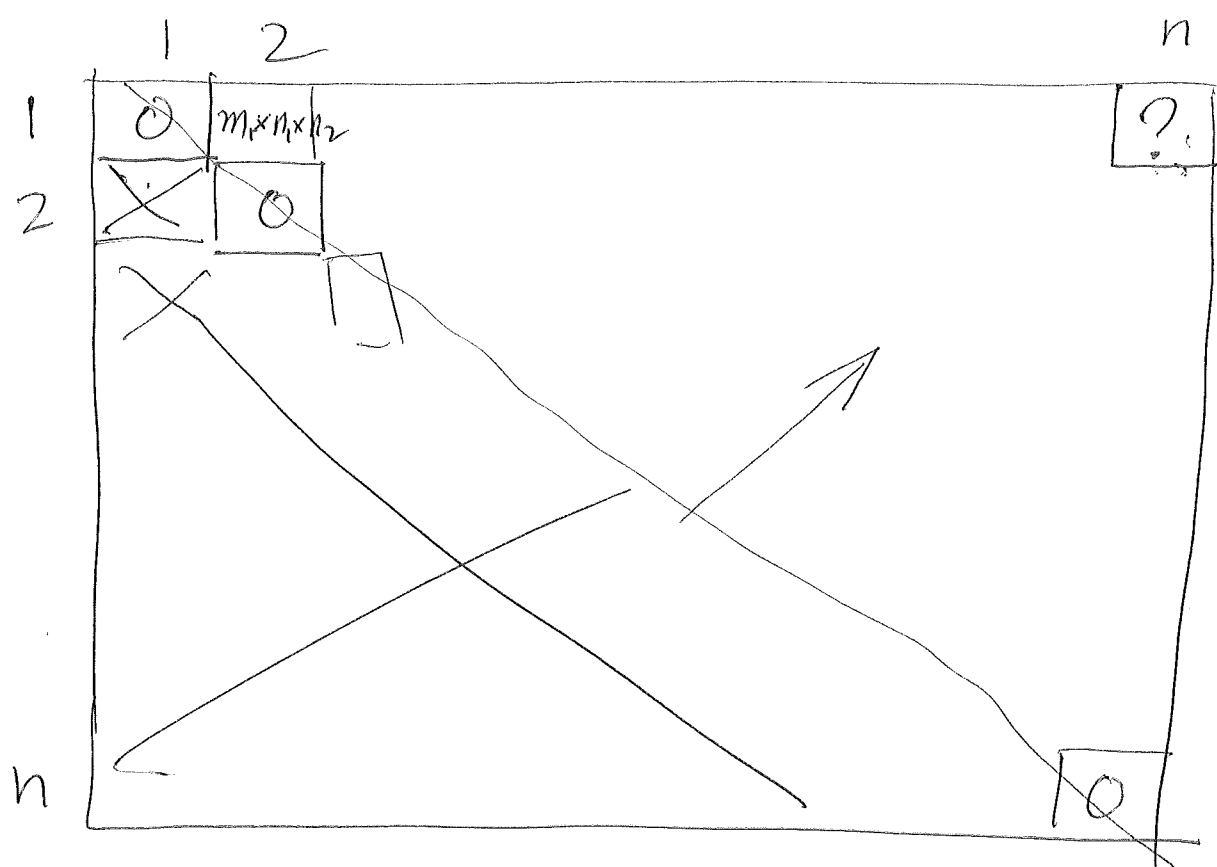
$$A_j \text{ is } m_j \times n_j$$

$$n_{j-1} = m_j$$

$$A_j \times A_{j+1}$$

$$m_j \times \underbrace{(n_j)}_{\underbrace{(m_{j+1})}} \times n_{j+1}$$

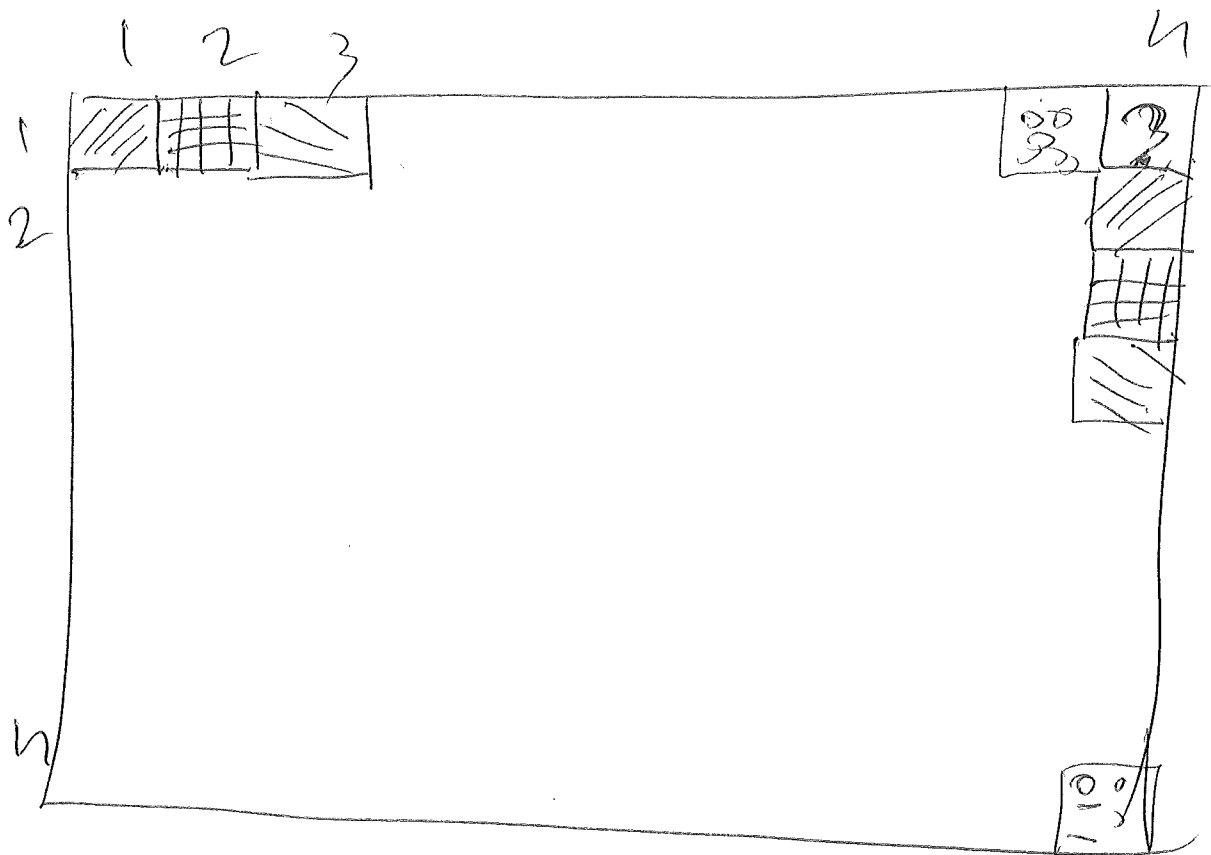
$$(A_1 \times \dots \times A_j) \cdot (A_{j+1} \times \dots \times A_n)$$



Each entry (i,j) means the optimal ways to multiply $A_i A_{i+1} \dots A_j$

- $(A_1 A_2 A_3 A_4) A_5$
- $(A_1 A_2 A_3) A_4 A_5$
- $(A_1 A_2) (A_3 A_4 A_5)$
- $A_1 (A_2 A_3 A_4 A_5)$

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A_1, A_2, \dots

A_n

$$O(n^3)$$